

Secant Method

```
function x = mySecantRoot(f, x0, x1, N)
    y0 = f(x0);
    y1 = f(x1);
    for i=1:n
        x = x1 - (x1 - x0)* y1/ (y1 - y0);
        y = f(x);
        x0 = x1;
        y0 = y1;
        x1 = x;
        y1 = y;
    end
end
```

Regula Falsi Method (False position)- combination of bisection and secant

Inputs: a, b, MaxIter, Tolerance

Initialize: iteration = 0, xold = b (to start with)

while (iteration <= MaxIter)

 compute f(a)

 compute f(b)

$$x = a - ((f(a)*(b-a)) / (f(b) - f(a)))$$

 Test convergence:

 if ($|x - xold|/|x|$) < Tolerance then

 output x

 goto stop

 else

 output (iteration, a, b, x)

 xold = x

 compute f(x)

 if $f(a) * f(x) > 0$

 a = x

 else

 b = x

 end

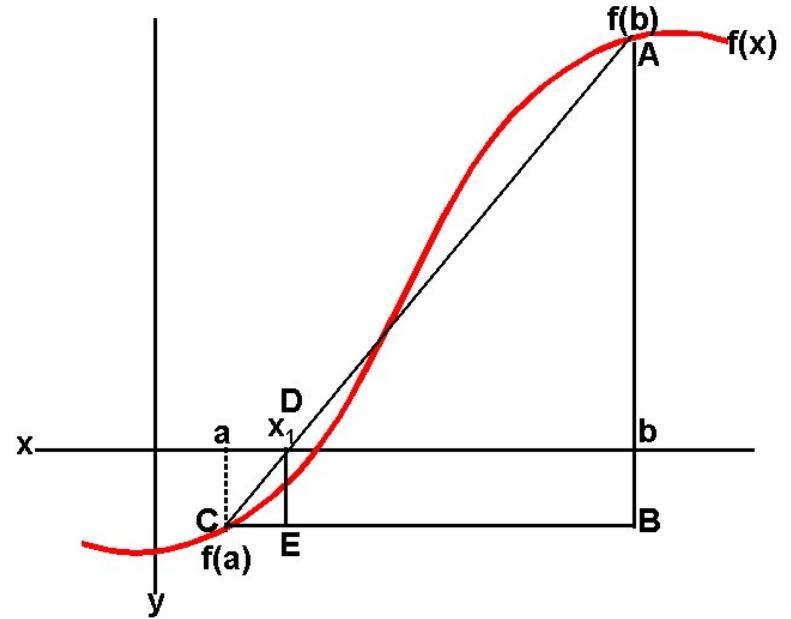
end

end

stop

Regula-falsi method (Algorithm)

1. Find points a and b so that $a < b$ and $f(a) f(b) < 0$
2. Take the interval $[a, b]$, determine the next value x_1
3. If $f(x_1) = 0$ then x_1 is the exact root



4. else if $f(x_1) f(b) < 0$ then let $a = x_1$
5. else if $f(a) f(x_1) < 0$ then let $b = x_1$
6. Repeat steps 2-5 until $f(x_i) \leq \text{Tolerance}$

Symbolic computations in Matlab

| `syms x y` (Declare variables used to be symbolic)

| `fx = cos(x)*sin(y) + 5*x^2*y` (Define function)

| `f = sin(x)+3*x^2`

| `subs(f, pi)` (to find specific value of the function)

| `g = exp(-y^2)`

| `h = compose(g,f)` (Define $h(x) = g(f(x))$)

| `k = f * g`

| `subs(k, [x,y], [0,1])`

| `f1=diff(f)`

| `k1x = diff(k,x)`

| `F=int(f), Fd=int(f,0,2*pi)`

Symbolic computation

Plot a symbolic function of one variable:

```
|ezplot(f), ezplot(g,-10,10)
```

Format a symbolic function of two variables

```
|ezsurf(k)
```

Simple algebra

```
poly = (x - 3)^5
```

```
polyex = expand(poly)
```

```
polysi = simplify(polyex)
```

```
solve(f), solve(g)
```

Linear Algebra

$$A = \begin{bmatrix} 1 & 3 & -2 & 5 \\ -1 & -1 & 5 & 4 \\ 0 & 1 & -9 & 0 \end{bmatrix}$$

$$u = [1 \ 1 \ 1 \ 1]$$

$$z = A * u$$

$$B = \begin{bmatrix} 3 & 2 & 1 \\ 7 & 6 & 5 \\ 4 & 3 & 2 \end{bmatrix}$$

$$B*B \text{ vs. } B.* B$$

$$B^3 \text{ vs. } B.^3$$

Inverse of a matrix

$$C = \text{randn}(5,5)$$

$$\text{inv}(C)$$

$$\text{Identity matrix } I = \text{eye}(3) \text{ (3x3 identity matrix)}$$

Norm of a matrix

If A is a $m \times n$ matrix, the norm of A is defined as

$$\|A\| = \max_{\|v\|=1} \|Av\|$$

Maximum is taken over all vectors with unit length. Norm of a matrix is defined as the largest factor by which it stretches or contracts a unit vector.

$$\|Av\| \leq \|A\| \|v\|$$

`norm(A)` norm of identity matrix is one.

`size(A)` dimensions of A ; `det(A)` determinant; `max(A)` maximum of each column; `min(A)` minimum of each column; `sum(A)` sum of each column; `mean(A)` average of each column; `diag(A)` diagonal of each column; A' transpose of A .

$$\begin{aligned}
 A_{11}x_1 + A_{12}x_2 + \dots + A_{1n}x_n &= b_1 \\
 A_{21}x_1 + A_{22}x_2 + \dots + A_{2n}x_n &= b_2 \\
 A_{31}x_1 + A_{32}x_2 + \dots + A_{3n}x_n &= b_3 \\
 &\vdots \\
 A_{n1}x_1 + A_{n2}x_2 + \dots + A_{nn}x_n &= b_n
 \end{aligned}$$

Matrix notation

$$\begin{bmatrix} A_{11} & A_{12} & \cdots & A_{1n} \\ A_{21} & A_{22} & \cdots & A_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ A_{n1} & A_{n2} & \cdots & A_{nn} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_n \end{bmatrix}$$

Augmented coefficient matrix

$$[A|\mathbf{b}] = \left[\begin{array}{cccc|c} A_{11} & A_{12} & \cdots & A_{1n} & b_1 \\ A_{21} & A_{22} & \cdots & A_{2n} & b_2 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ A_{n1} & A_{n2} & \cdots & A_{nn} & b_n \end{array} \right]$$

Uniqueness of Solution

A system of n linear equations in n unknowns has a unique solution, provided that the determinant of the coefficient matrix is nonsingular ($\det(A) \neq 0$).

The rows and columns of a nonsingular (invertible) matrix are linearly independent i.e. no row or column can be expressed as a linear combination of other rows or columns.

When the coefficient matrix is singular, the equations have an infinite number of solutions or no solutions at all.

$$2x + y = 3; 4x + 2y = 6 \text{ (infinite solutions)}$$

$$2x + y = 3; 4x + 2y = 0 \text{ (no solutions)}$$

Ill-Conditioning

$$\|A\| = \sqrt{\sum_{i=1}^n \sum_{j=1}^n A_{ij}^2}, \quad \|A\| = \max_{1 \leq i \leq n} \sum_{j=1}^n |A_{ij}|$$

$$\text{cond}(A) = \|A\| \|A^{-1}\|$$

When $\text{cond}(A) \approx 1$, the matrix is well-conditioned.

Condition number increases with the degree of ill-conditioning.

$$2x + y = 3; \quad 2x + 1.001y = 0 \qquad \text{Solutions: } x = 1501.5, \quad y = -3000$$

$\det(A) = |A| = 0.002$ is much smaller than the coefficients (Ill – conditioned)

$$2x + y = 3; \quad 2x + 1.002y = 0 \qquad \text{Solutions: } x = 751.5, \quad y = -1500$$

0.1% change in coefficient produces 100% change in solution.

$$\begin{array}{rcl}
 E_1: & a_{11}x_1 + a_{12}x_2 + \cdots + a_{1n}x_n & = b_1 \\
 E_2: & a_{21}x_1 + a_{22}x_2 + \cdots + a_{2n}x_n & = b_2 \\
 & \vdots & \\
 E_n: & a_{n1}x_1 + a_{n2}x_2 + \cdots + a_{nn}x_n & = b_n
 \end{array}$$

The coefficients are a_{ij} ($i, j = 1, 2, \dots, n$) and b_i . We need to find unknowns x_1, x_2, \dots, x_n .

Three operations to simplify linear system:

- Equation E_i can be multiplied by any nonzero constant λ with the resulting equation replacing the original one. $(\lambda E_i) \rightarrow (E_i)$
- Equation E_j can be multiplied by any constant λ and added to equation E_i with the resulting equation used in place of E_i . $(E_i + \lambda E_j) \rightarrow (E_i)$
- Equations E_i and E_j can be transposed in order. $(E_i) \leftrightarrow (E_j)$.

Illustration

$$E_1: \quad x_1 + x_2 + 0x_3 + 3x_4 = 4$$

$$E_2: \quad 2x_1 + x_2 - x_3 + x_4 = 1$$

$$E_3: \quad 3x_1 - x_2 - x_3 + 2x_4 = -3$$

$$E_4: \quad -x_1 + 2x_2 + 3x_3 - x_4 = 4$$

Augmented matrix

$$\left[\begin{array}{cccc|c} 1 & 1 & 0 & 3 & 4 \\ 2 & 1 & -1 & 1 & 1 \\ 3 & -1 & -1 & 2 & -3 \\ -1 & 2 & 3 & -1 & 4 \end{array} \right]$$

Use E_1 to eliminate x_1 from E_2, E_3 and E_4

Do

$$(E_2 - 2E_1) \rightarrow (E_2); (E_3 - 3E_1) \rightarrow (E_3); (E_4 + E_1) \rightarrow (E_4)$$

Augmented matrix

$$R_2 \rightarrow R_2 - 2R_1; R_3 \rightarrow R_3 - 3R_1; R_4 \rightarrow R_4 - (-R_1)$$

$$\left[\begin{array}{cccc|c} 1 & 1 & 0 & 3 & 4 \\ 2 & 1 & -1 & 1 & 1 \\ 3 & -1 & -1 & 2 & -3 \\ -1 & 2 & 3 & -1 & 4 \end{array} \right] \rightarrow \left[\begin{array}{cccc|c} 1 & 1 & 0 & 3 & 4 \\ 0 & -1 & -1 & -5 & -7 \\ 0 & -4 & -1 & -7 & -15 \\ 0 & 3 & 3 & 2 & 8 \end{array} \right]$$

$$R_3 \rightarrow R_3 - 4R_2; R_4 \rightarrow R_4 - (-3R_2)$$

$$\left[\begin{array}{cccc|c} 1 & 1 & 0 & 3 & 4 \\ 0 & -1 & -1 & -5 & -7 \\ 0 & 0 & 3 & 13 & 13 \\ 0 & 0 & 0 & -13 & -13 \end{array} \right]$$

Steps

1. If $a_{11} \neq 0$, we perform operations $(R_j - (a_{j1}/a_{11})R_1) \rightarrow (R_j)$ for each $j = 2, 3, \dots, n$ to eliminate coefficient of x_1 in each of the j rows.

2. Sequential procedure for $i = 2, 3, \dots, n - 1$:

$(R_j - (a_{ji}/a_{jj})R_i) \rightarrow (R_j)$ for each $j = i + 1, i + 2, \dots, n$

Backward substitution

$$\left[\begin{array}{ccc|c} a_{11} & a_{12} & a_{13} & b_1 \\ a_{21} & a_{22} & a_{23} & b_2 \\ a_{31} & a_{32} & a_{33} & b_3 \end{array} \right] \rightarrow \left[\begin{array}{ccc|c} a_{11} & a_{12} & a_{13} & b_1 \\ & a_{22}^{(1)} & a_{23}^{(1)} & b_2^{(1)} \\ & & a_{33}^{(2)} & b_3^{(2)} \end{array} \right]$$

$$x_3 = b_3^{(2)} / a_{33}^{(2)}$$

$$x_2 = (b_2^{(1)} - a_{23}^{(1)} x_3) / a_{22}^{(1)}$$

$$x_1 = (b_1 - a_{12} x_2 - a_{13} x_3) / a_{11}$$

General form

$$x_n = b_n^{(n-1)} / a_{nn}^{(n-1)}$$

$$x_i = \left(b_i^{(i-1)} - \sum_{j=i+1}^n a_{ij}^{(i-1)} x_j \right) / a_{ii}^{(i-1)}$$

$$\text{for } i = n-1, n-2, \dots, 1$$

Algorithm (Gaussian elimination with backward substitution)

Input: number of equations (unknowns) n ; augmented matrix $A = [a_{ij}]$, $1 \leq i \leq n$ and $1 \leq j \leq n + 1$

Output: solution x_1, x_2, \dots, x_n or “no unique solution”

1. For $i = 1, \dots, n - 1$ do steps 2 – 4:

2. Find the smallest integer p with $i \leq p \leq n$ and $a_{pi} \neq 0$.

 If no p found, then OUTPUT(no unique solution exists), STOP

3. If $p \neq i$ then perform $(R_p) \leftrightarrow (R_i)$

4. For $j = i + 1, \dots, n$ do:

$$m_{ji} = a_{ji} / a_{ii};$$

$$(R_j - m_{ji}R_i) \rightarrow (R_j)$$

5. If $a_{nn} = 0$ then OUTPUT(no unique solution exists)

6. Set $x_n = a_{n,n+1} / a_{nn}$

7. For $i = n - 1, \dots, 1$ set $x_i = [a_{i,n+1} - \sum_{j=i+1}^n a_{ij}x_j] / a_{ii}$

8. OUTPUT(x_1, \dots, x_n), STOP