

EECS 114:

Engineering Data Structures and Algorithms

Lecture 6

Instructor: Ryan Rusich

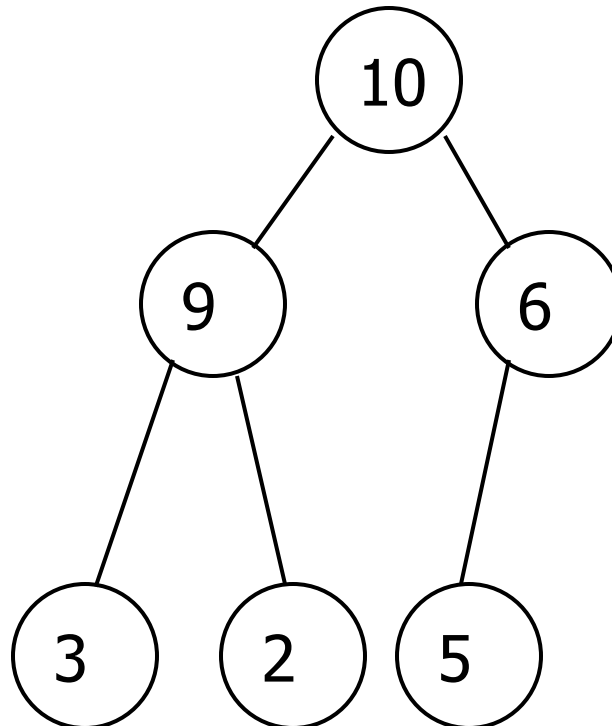
E-mail: rusichr@uci.edu

Office: EH 2204

The Henry Samueli School of Engineering
Electrical Engineering and Computer Science
University of California, Irvine

(Binary) Heaps

- Similar to a Binary Search Tree (BST)
- Heap is sorted in a weaker sense than BST
- Tree is a complete tree
 - Bottom level may not be filled but fills from left to right

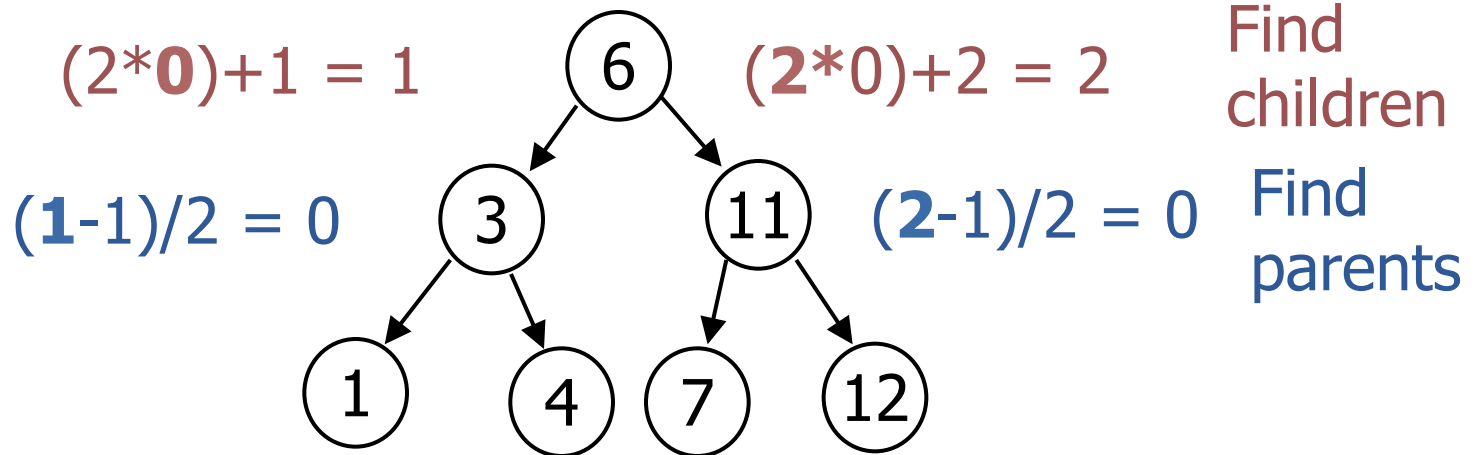


(Binary) Heaps

- Binary trees are easily stored in an array
 - Array with index started at 0
 - Parent = $(i-1)/2$
 - Left Child = $2i + 1$
 - Right Child = $2i + 2$
 - Array with index started at 1 (from CLRS)
 - Parent = $i/2$
 - Left Child = $2i$
 - Right Child = $2i + 1$

Binary Tree – Array

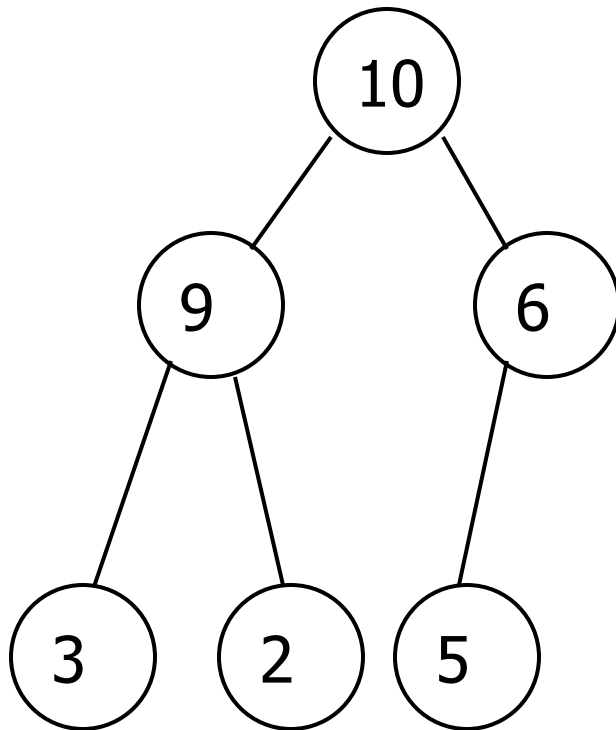
0	1	2	3	4	5	6	7
6	3	11	1	4	7	12	



Heap Order Properties

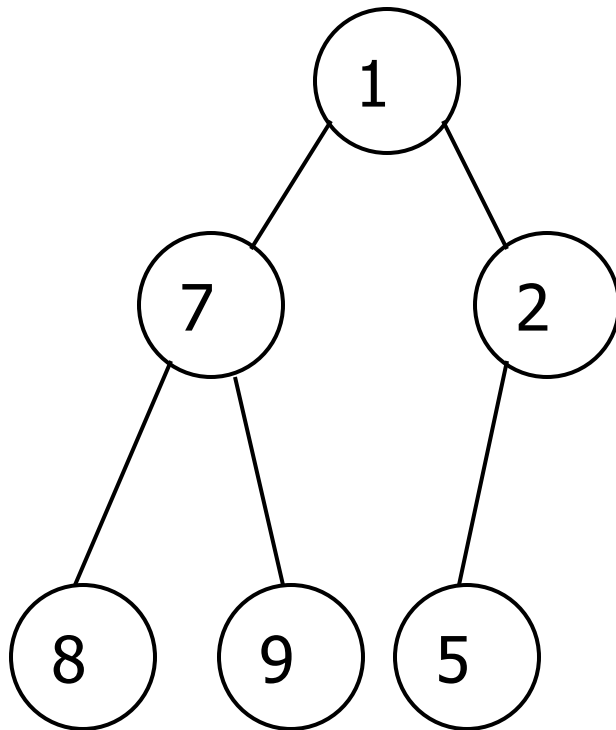
- Min Heap
 - For each node X , $X.key > (X.parent).key$
 - Smallest number is at the root
- Max Heap
 - For each node X , $X.key < (X.parent).key$
 - Largest number is at the root

Max Heap



10
9
6
3
2
5

Min Heap



1
7
2
8
9
5

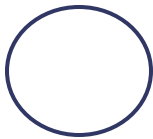
Heap - Insert

- Create a hole (empty node) in the next available complete tree location
 - Remember to fill bottom layer from left to right
- If the item can be inserted into the hole without violation of the heap property, insert item
- Otherwise, copy the hole's parent item into the hole then trickle up

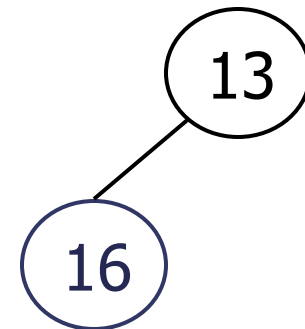
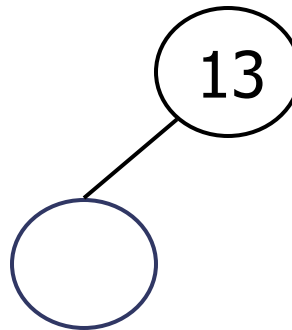
Min Heap - Insert

13, 16, 19, 14, 23, 17, 6

Insert 13



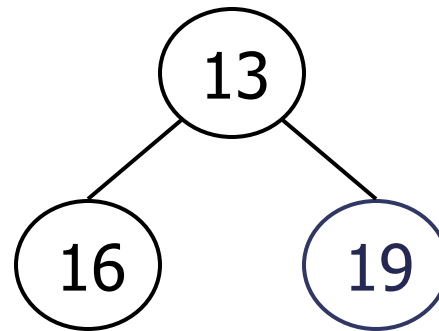
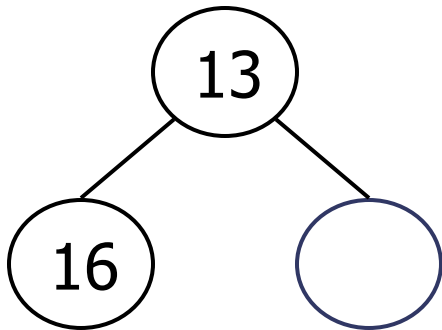
Insert 16



Min Heap - Insert

13, 16, 19, 14, 23, 17, 6

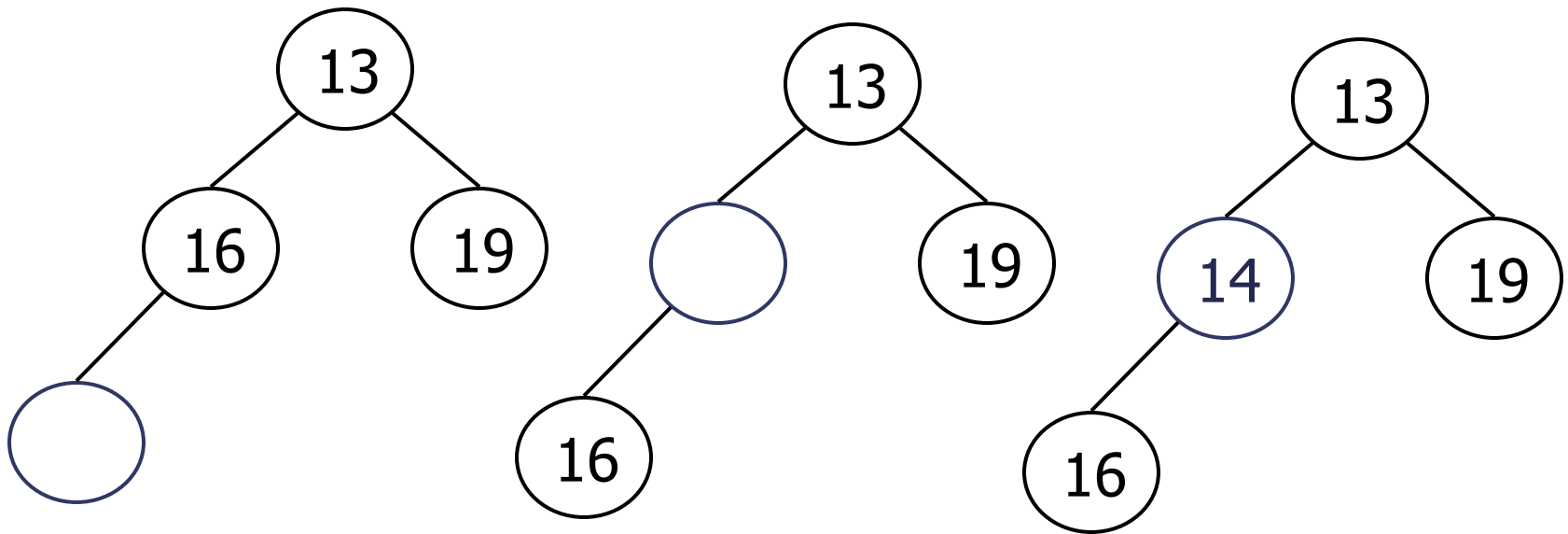
Insert 19



Min Heap - Insert

13, 16, 19, 14, 23, 17, 6

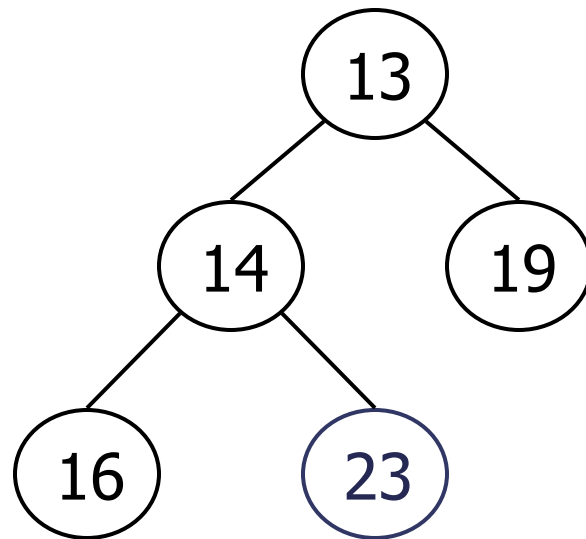
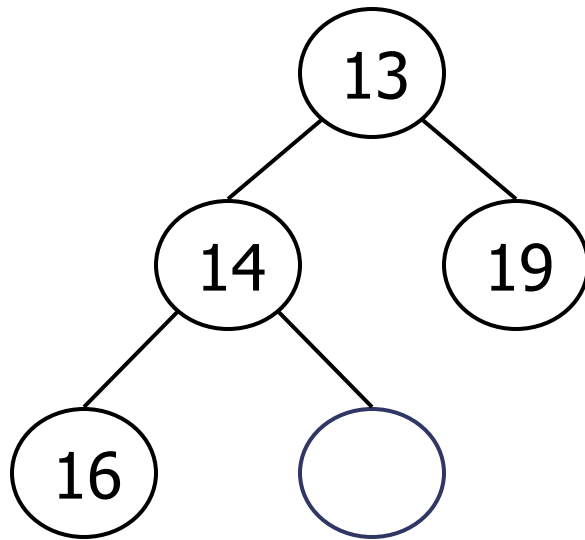
Insert 14



Min Heap - Insert

13, 16, 19, 14, 23, 17, 6

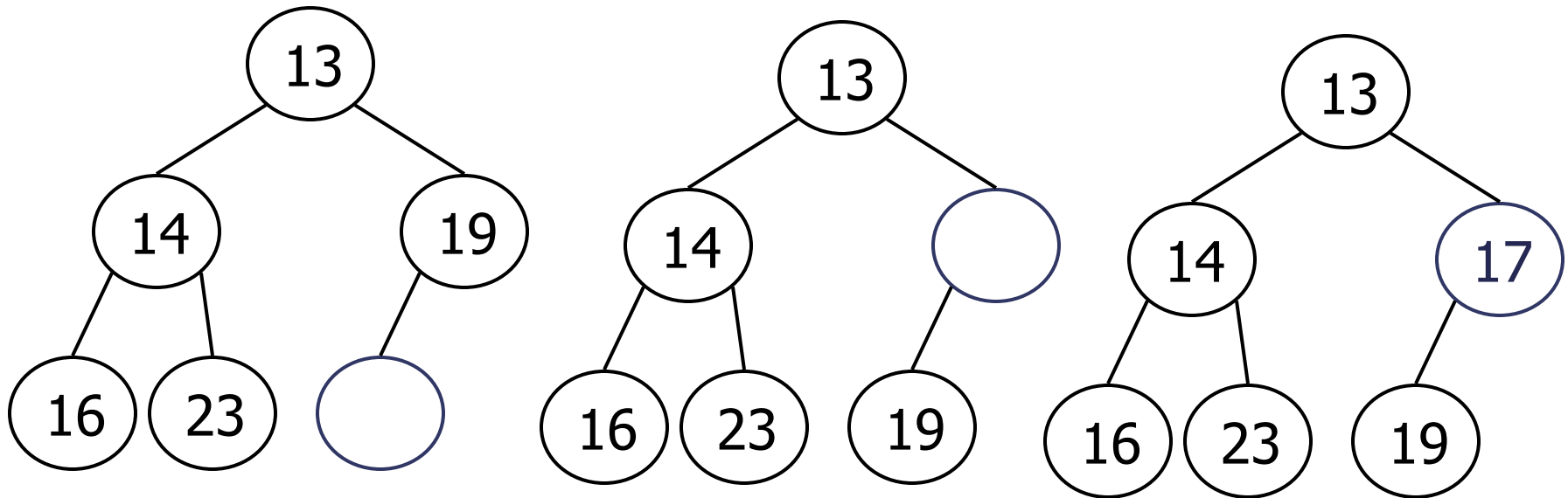
Insert 23



Min Heap - Insert

13, 16, 19, 14, 23, 17, 6

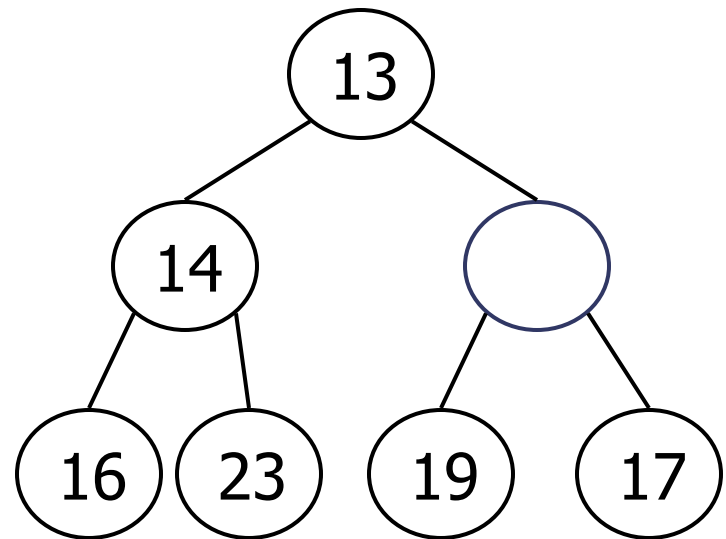
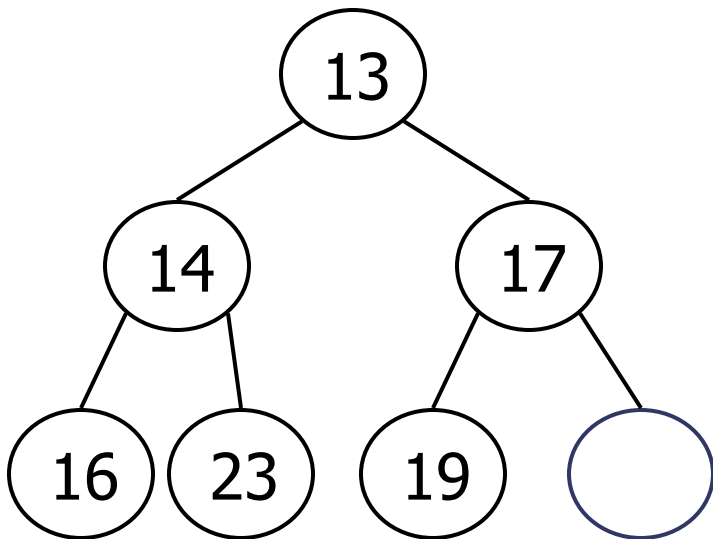
Insert 17



Min Heap - Insert

13, 16, 19, 14, 23, 17, 6

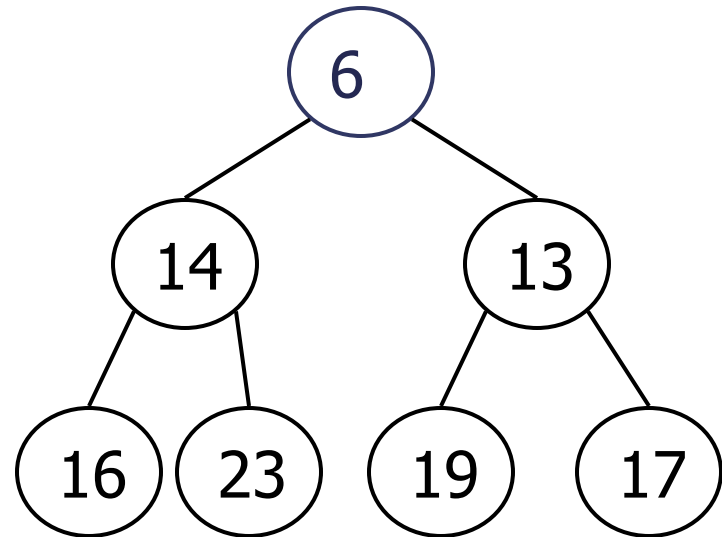
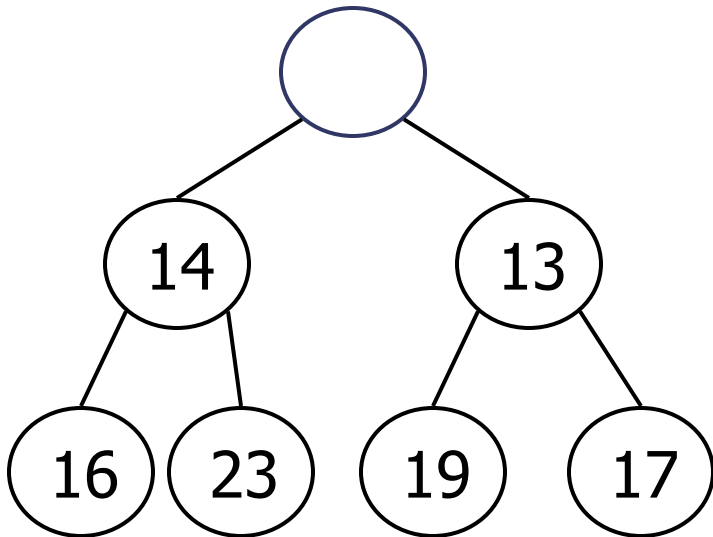
Insert 6



Min Heap - Insert

13, 16, 19, 14, 23, 17, 6

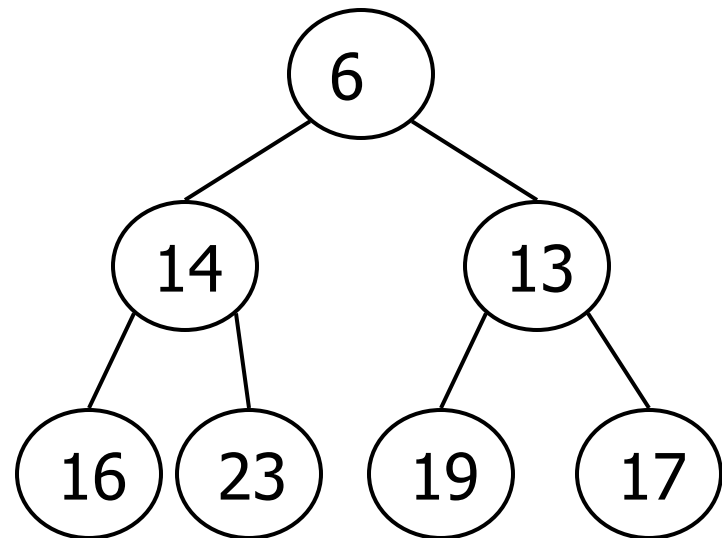
Insert 6



Min Heap – Extract Min

6, 14, 13, 16, 23, 19, 17

6,14,13,16 23,19,17
13,14,17,16,23,19
14,16,17,19,23
16,19,17,23
17,19,23
19,23
23



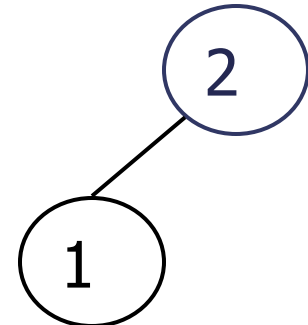
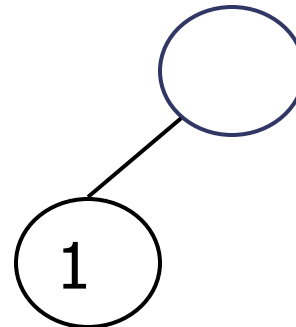
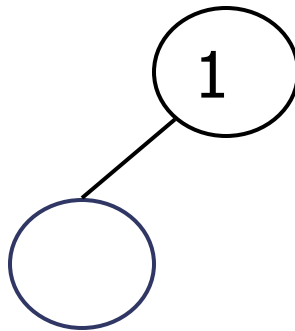
Max Heap - Insert

1, 2, 3, 4, 5, 6

Insert 1



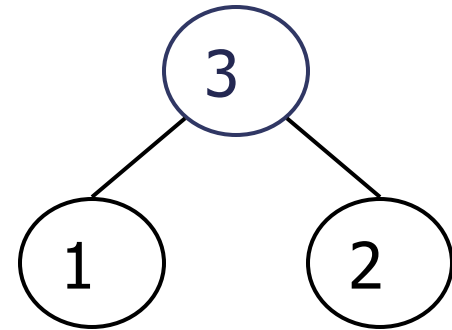
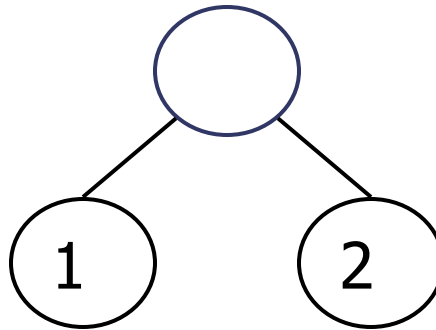
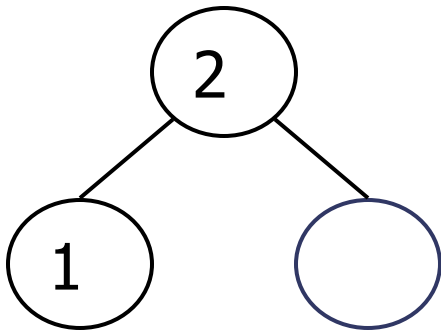
Insert 2



Max Heap - Insert

1, 2, 3, 4, 5, 6

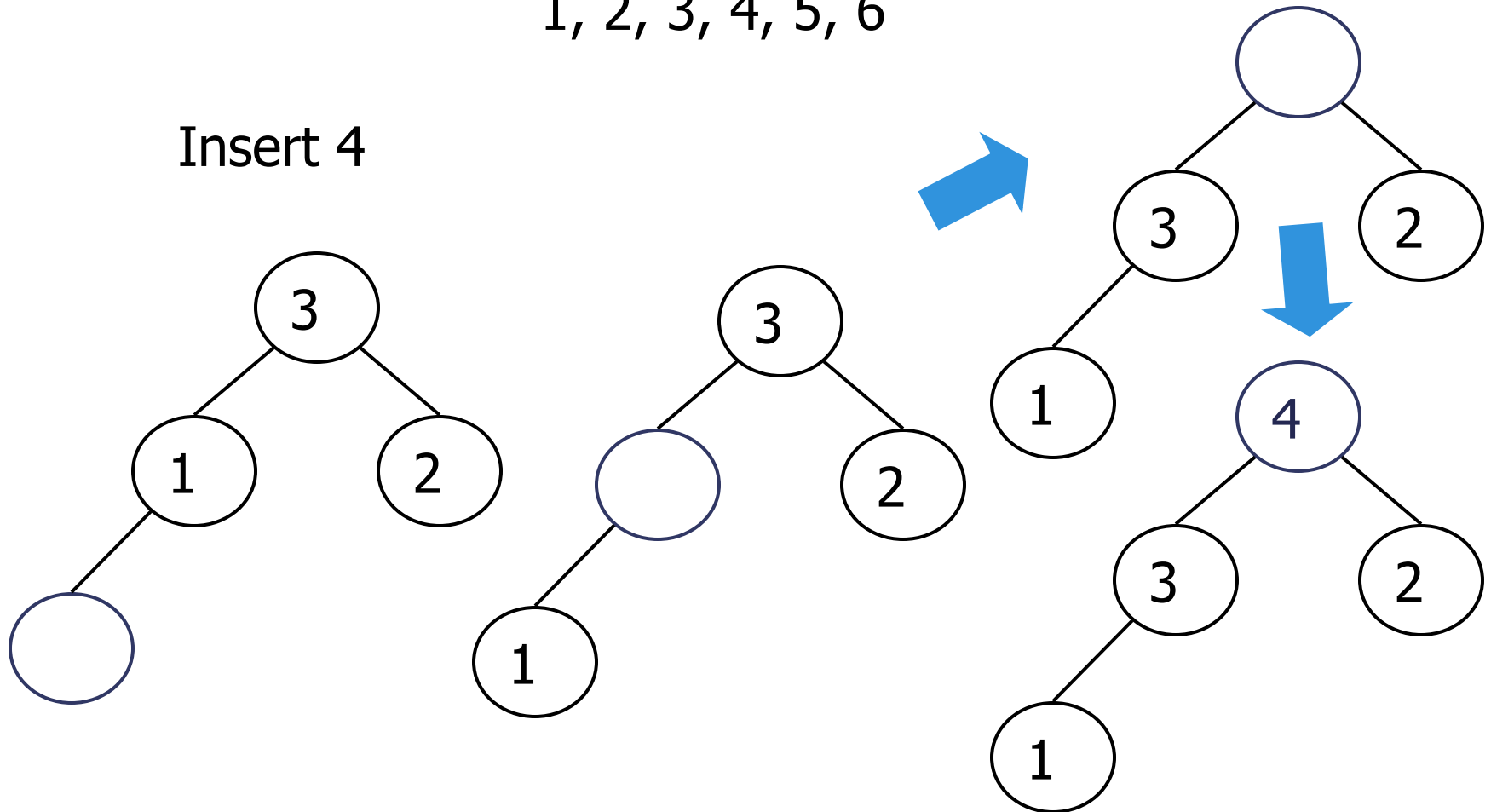
Insert 3



Max Heap - Insert

1, 2, 3, 4, 5, 6

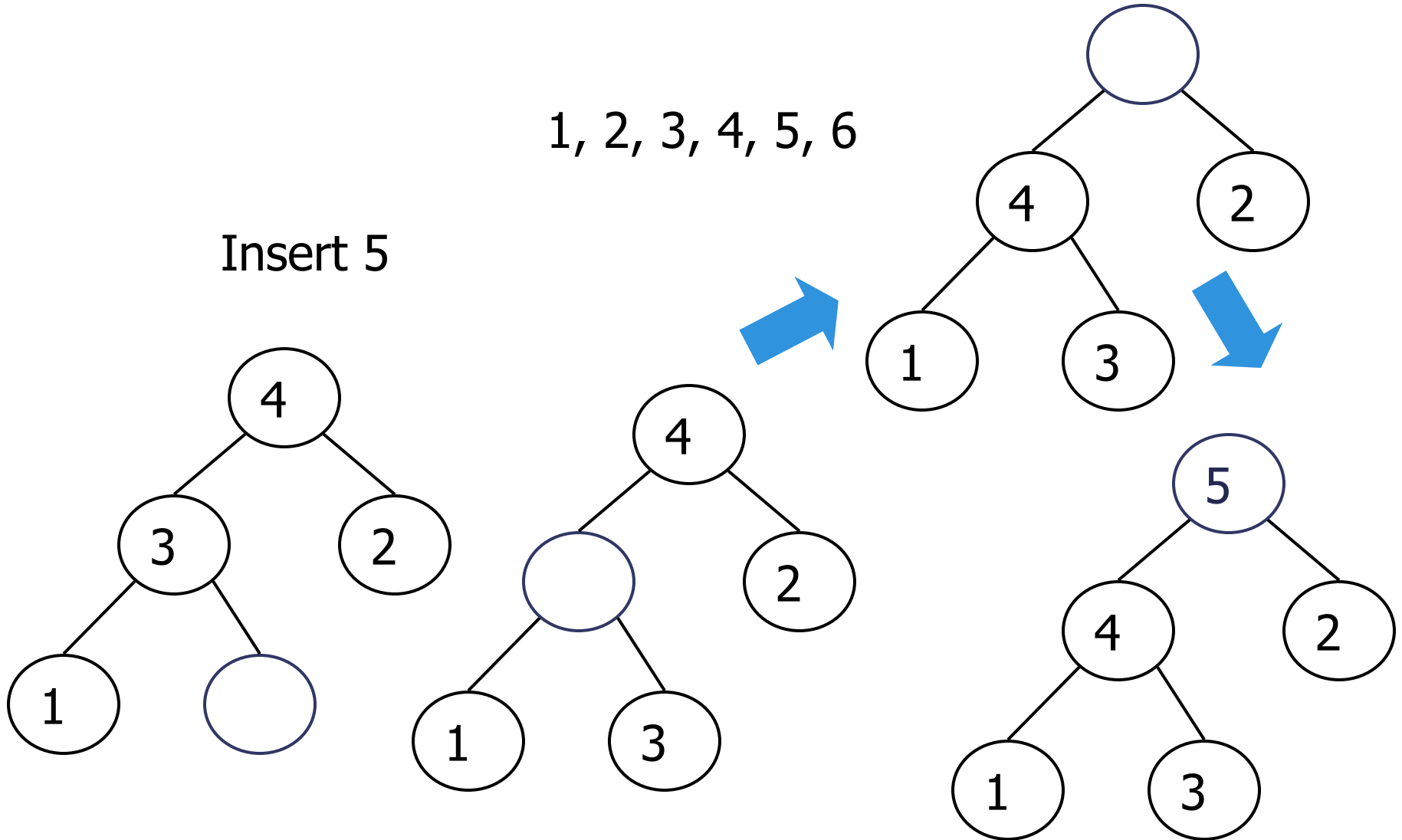
Insert 4



Max Heap - Insert

1, 2, 3, 4, 5, 6

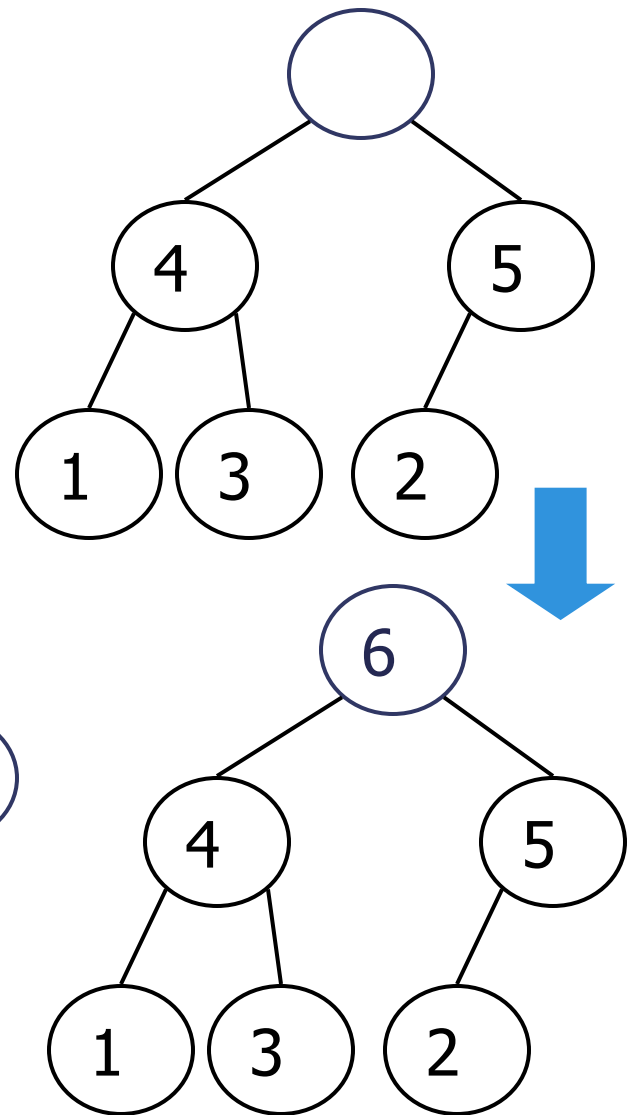
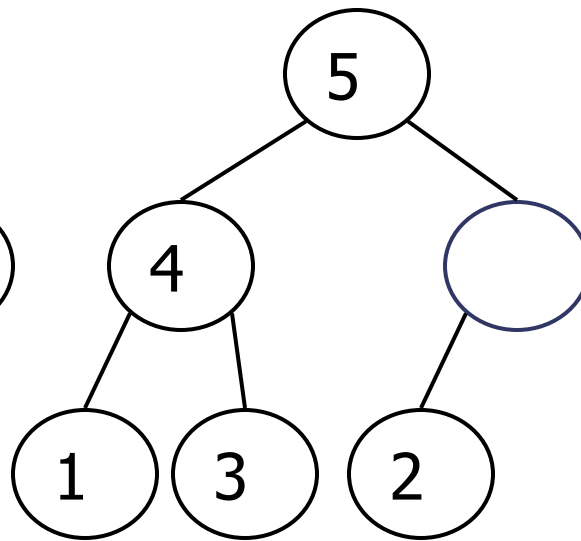
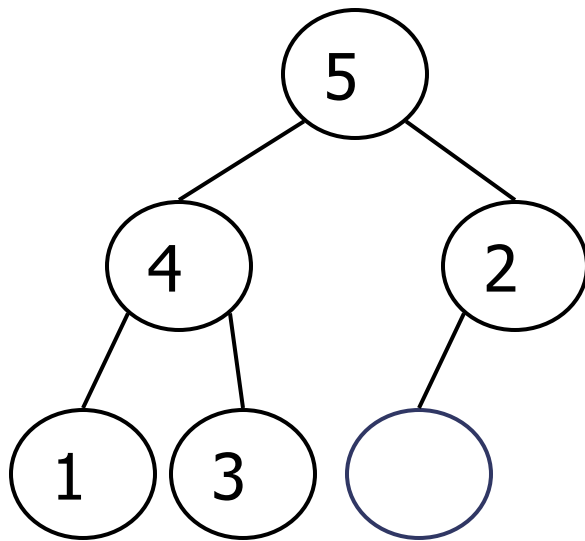
Insert 5



Max Heap - Insert

1, 2, 3, 4, 5, 6

Insert 6



Min Heap Insert

Running time = $O(\log n)$

```
void Min-Heap-Insert ( itemtype item )  
    if heap full  
        throw heap full exception  
    // Trickle up  
    x = numNodes  
    while (x > 0 && arr[(x-1)/2] > item)  
        arr[x] = arr[(x-1)/2]  
        x = (x-1)/2  
    arr[x] = item  
    numNodes++
```

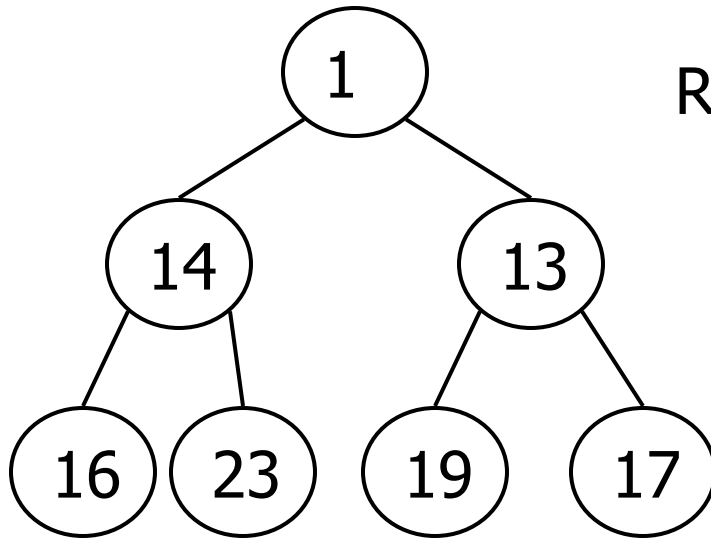
Heap Find Min/Max

- Finding the node with the highest priority is easy - it is just at the root
 - Running time = $O(1)$

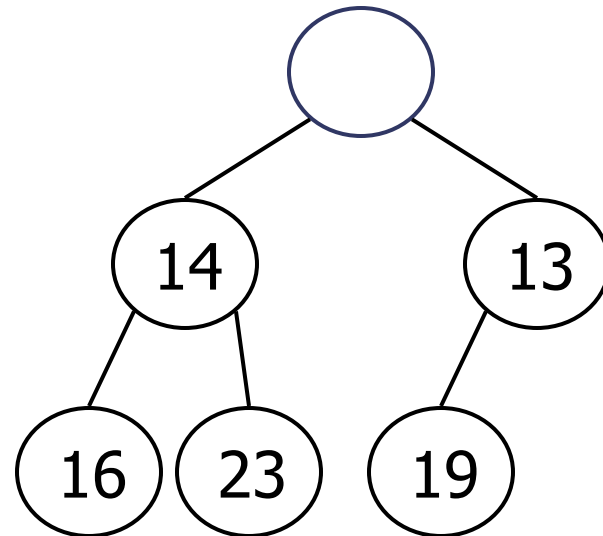
Heap – Remove (Extract Min or Max)

- Remove highest priority item
 - Makes a hole at the root
 - Want to remain a complete tree, so attempt to place last item in the heap into the hole
 - If item can be placed in hole without violation of the heap property, then done
 - Otherwise, trickle down
 - Pick the child with the highest priority

Min Heap – Remove

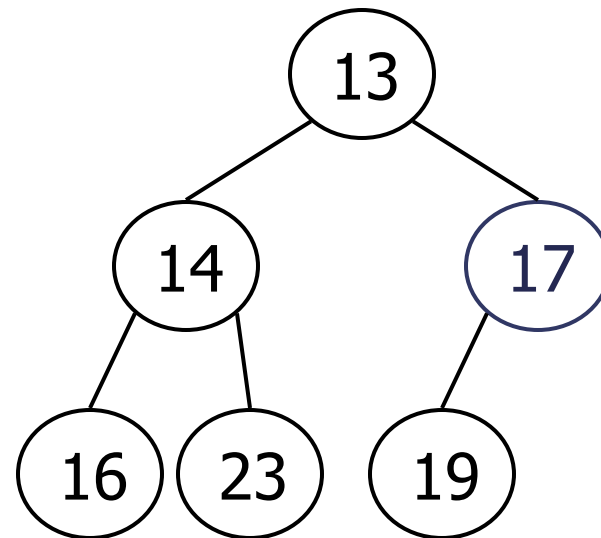
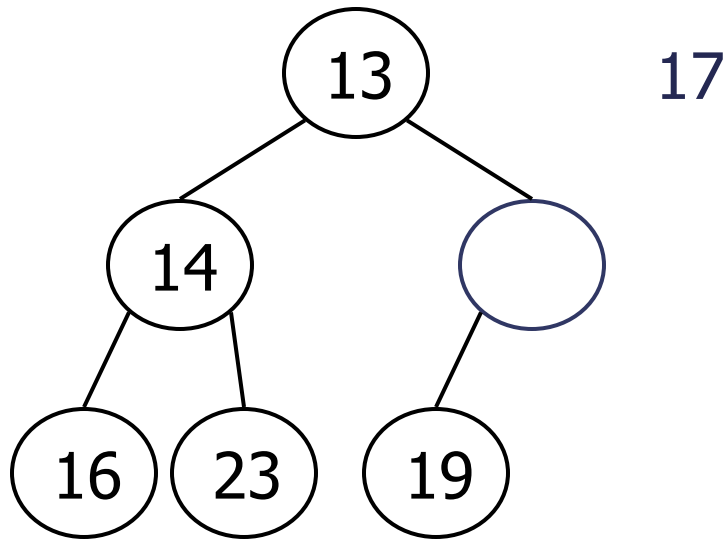


Remove 1

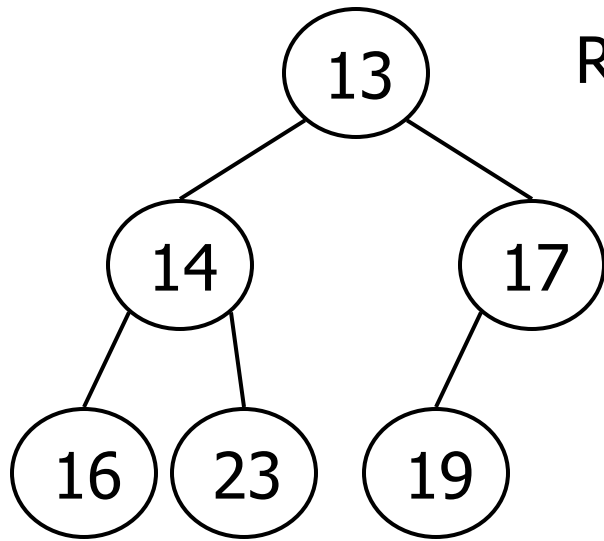


17

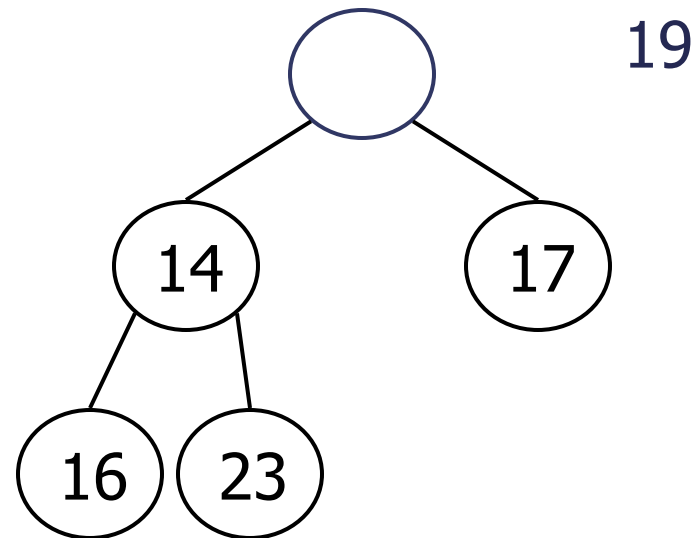
Min Heap - Remove



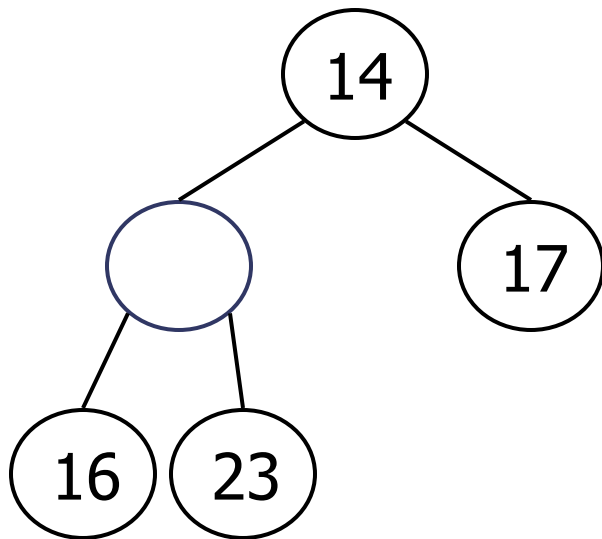
Min Heap - Remove



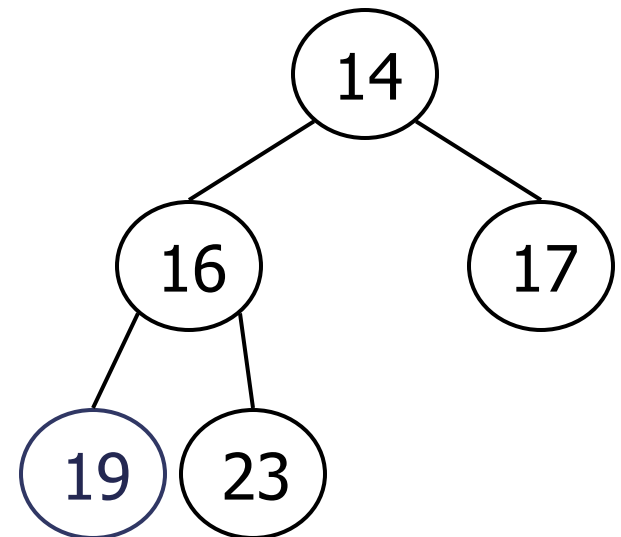
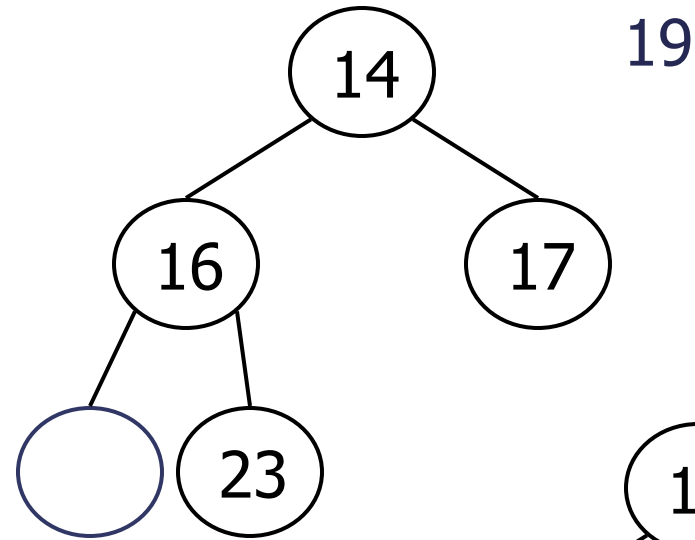
Remove 13



Min Heap - Remove

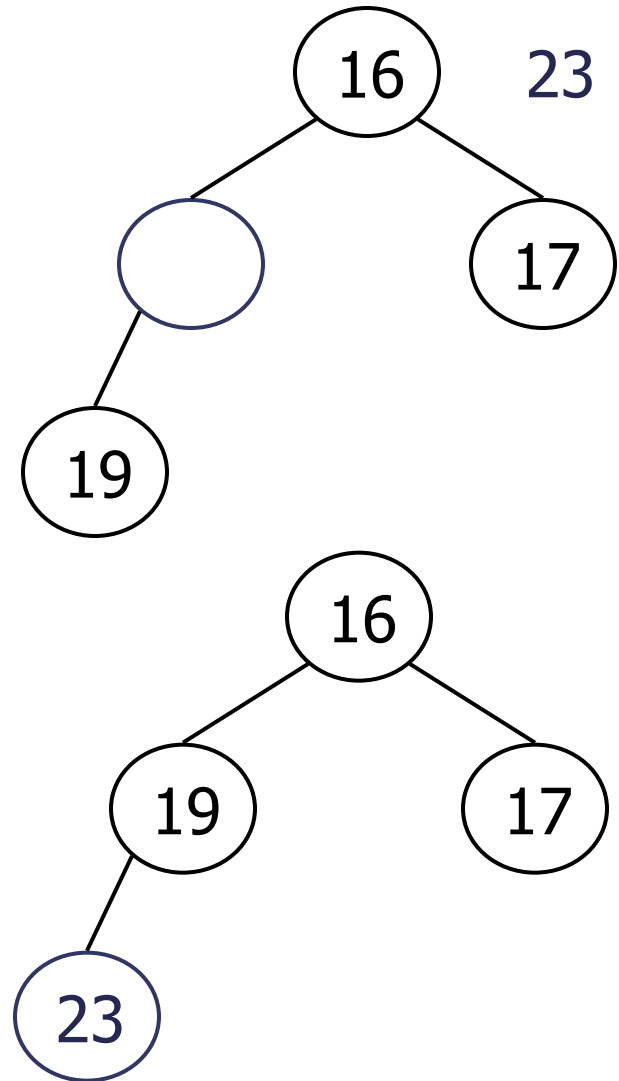
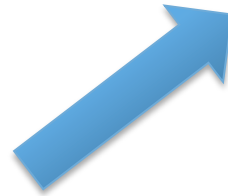
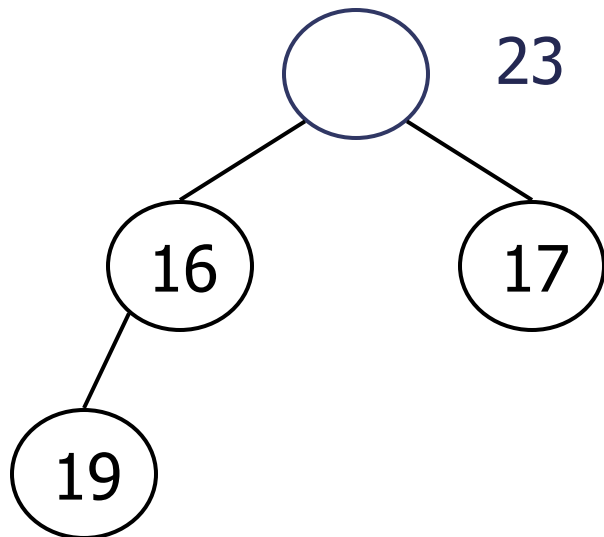
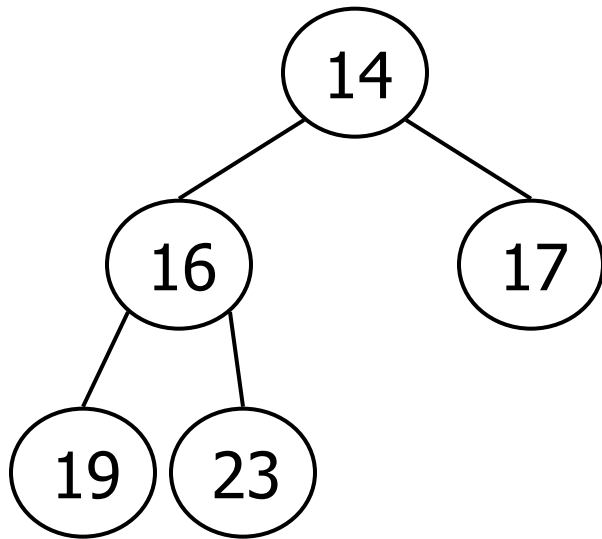


19

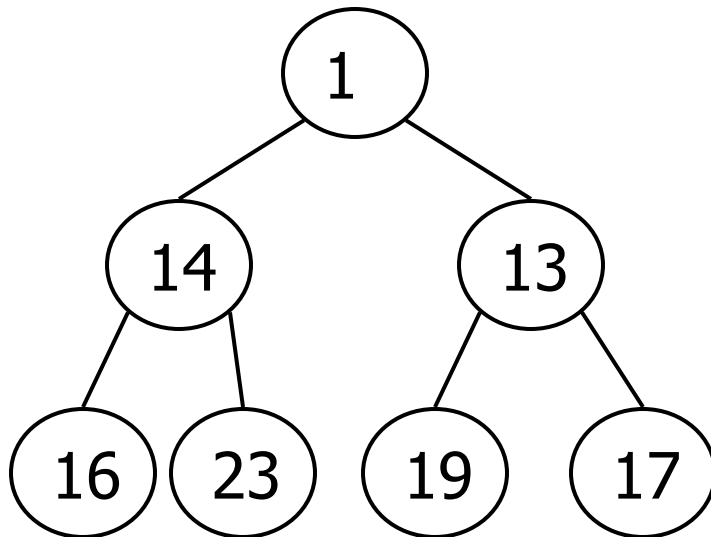


Min Heap - Remove

Remove 14



Min Heap - Remove



1,14,13,16,23,19,17
13,14,17,16,23,19
14,16,17,19,23
16,19,17,23
17,19,23
19,23
23

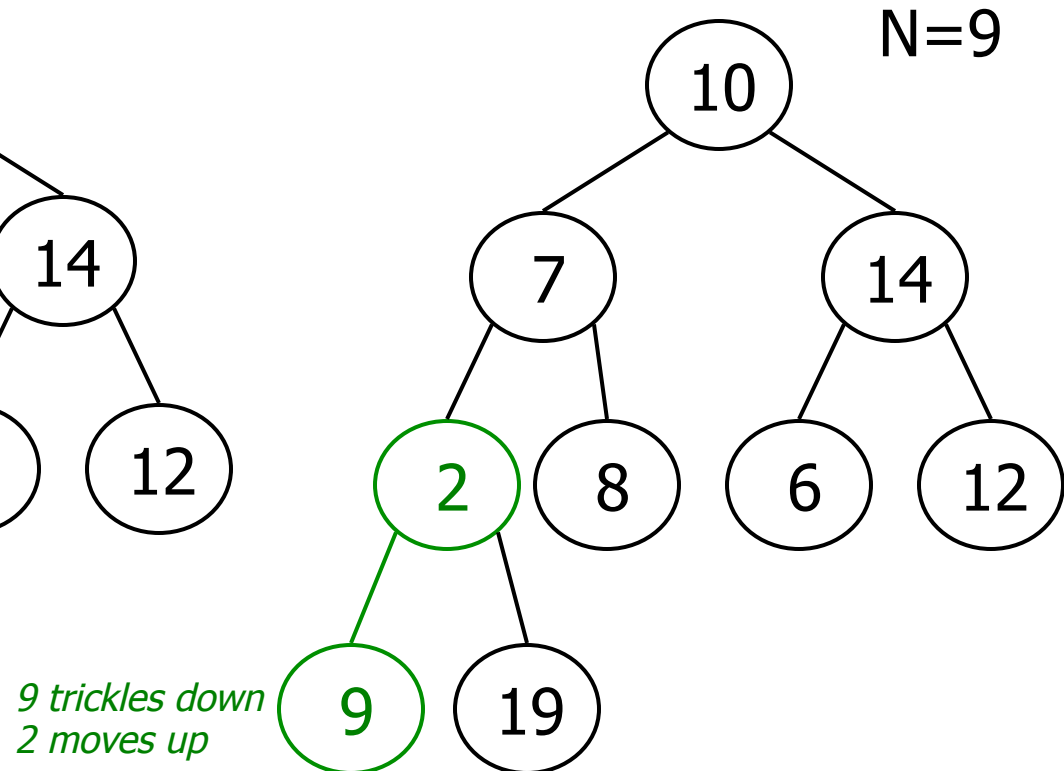
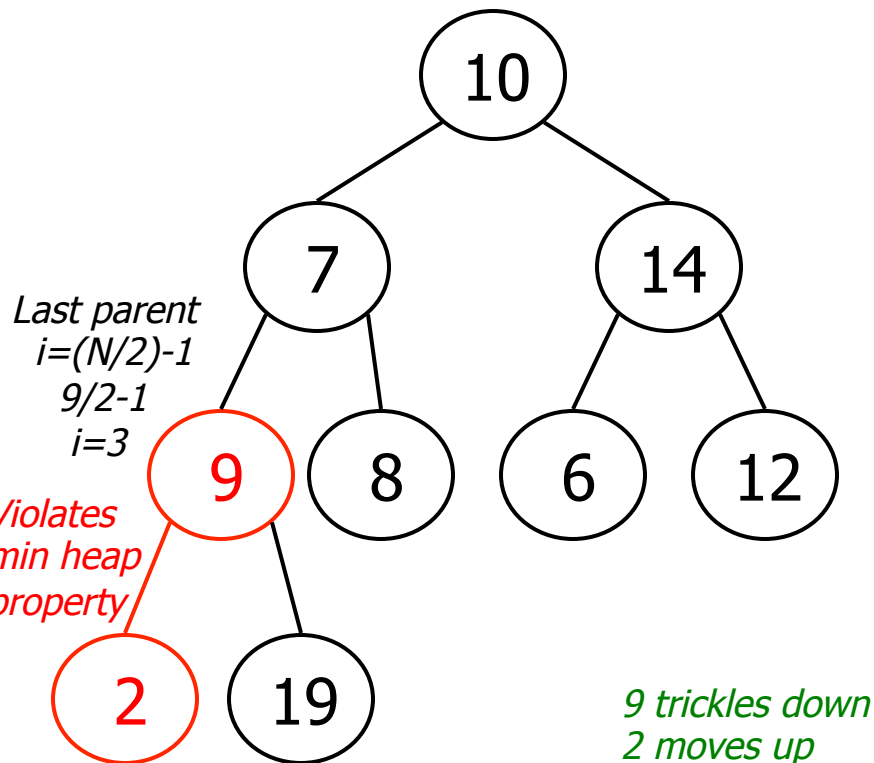
Building a Heap

- Naïve method
 - Perform N insert operations - $O(n \log n)$
- Better solution
 - Given an array of unordered items:

for (int i=(N/2)-1; i >= 0; i--)
Trickle down

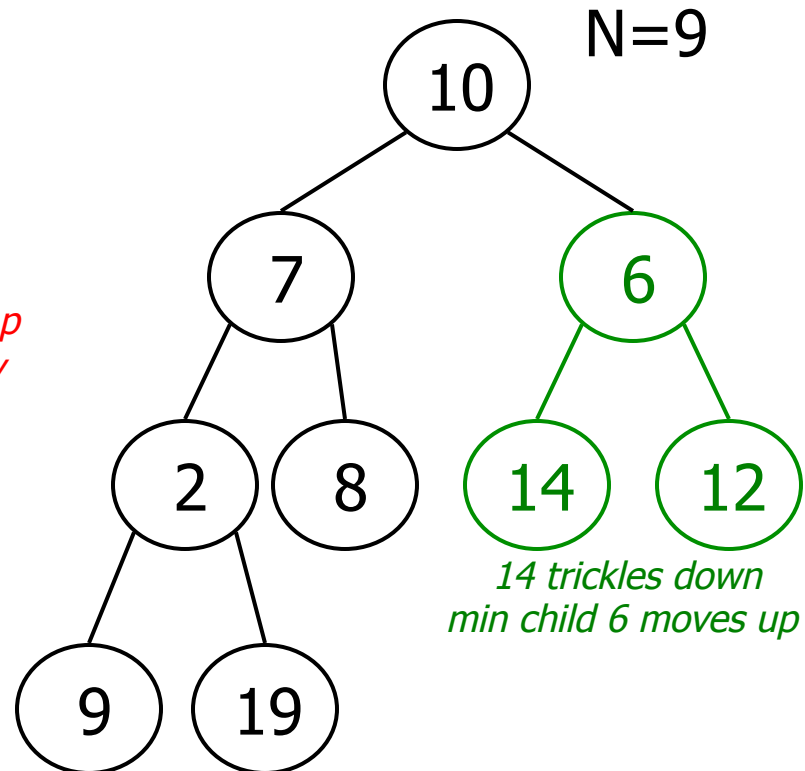
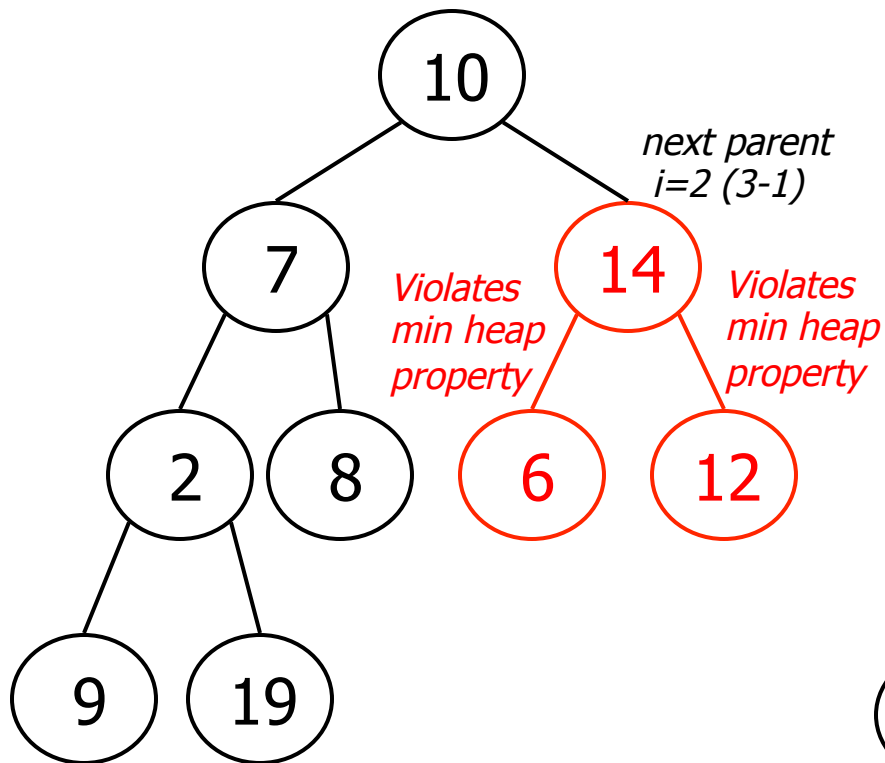
Building a Min Heap

i=0	i=1	i=2	i=3	i=4	i=5	i=6	i=7	i=8
10	7	14	9	8	6	12	2	19



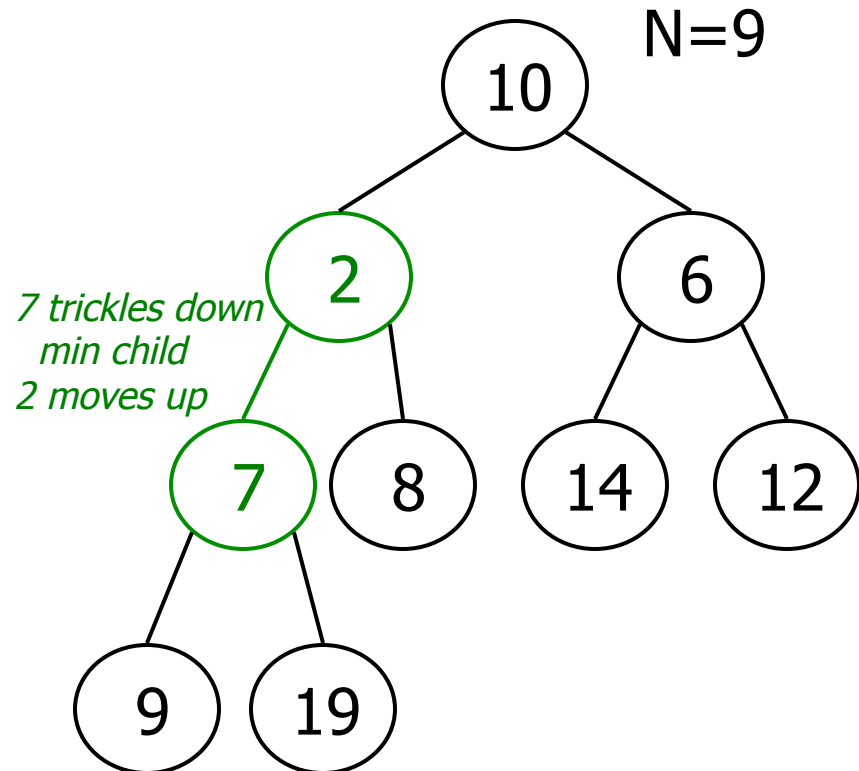
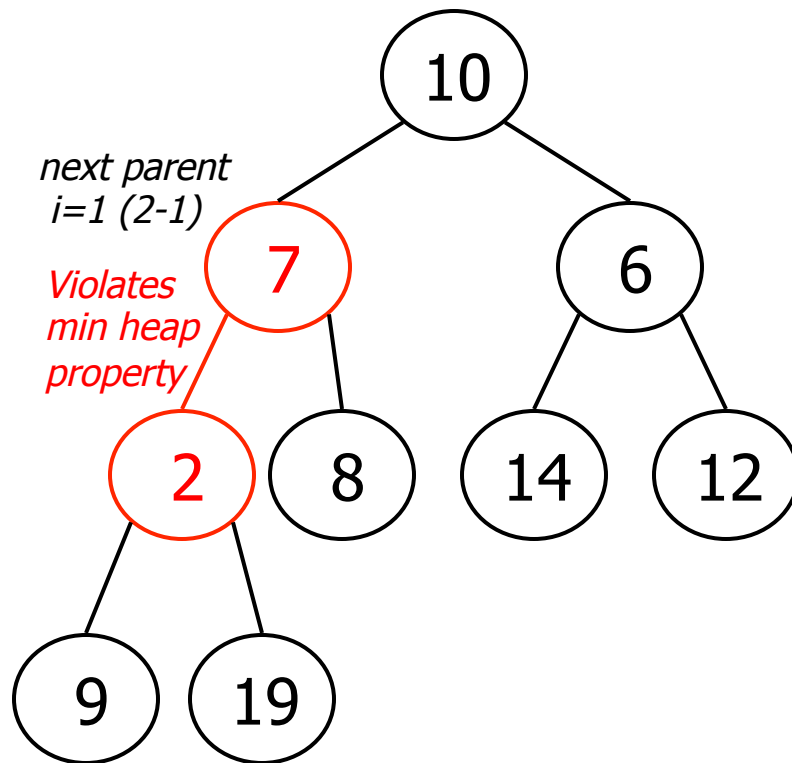
Building a Min Heap

i=0	i=1	i=2	i=3	i=4	i=5	i=6	i=7	i=8
10	7	14	2	8	6	12	9	19



Building a Min Heap

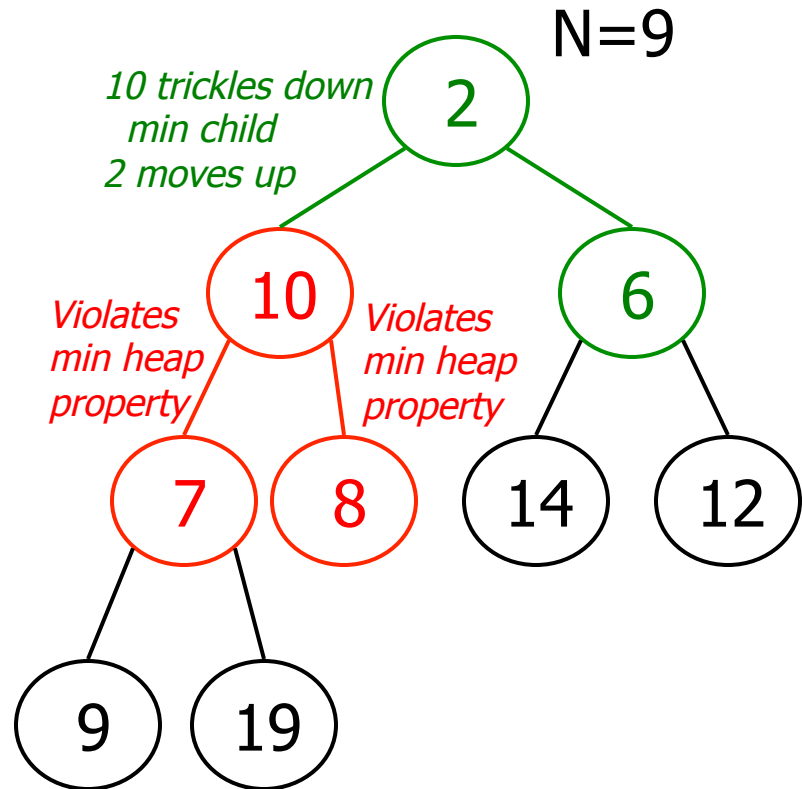
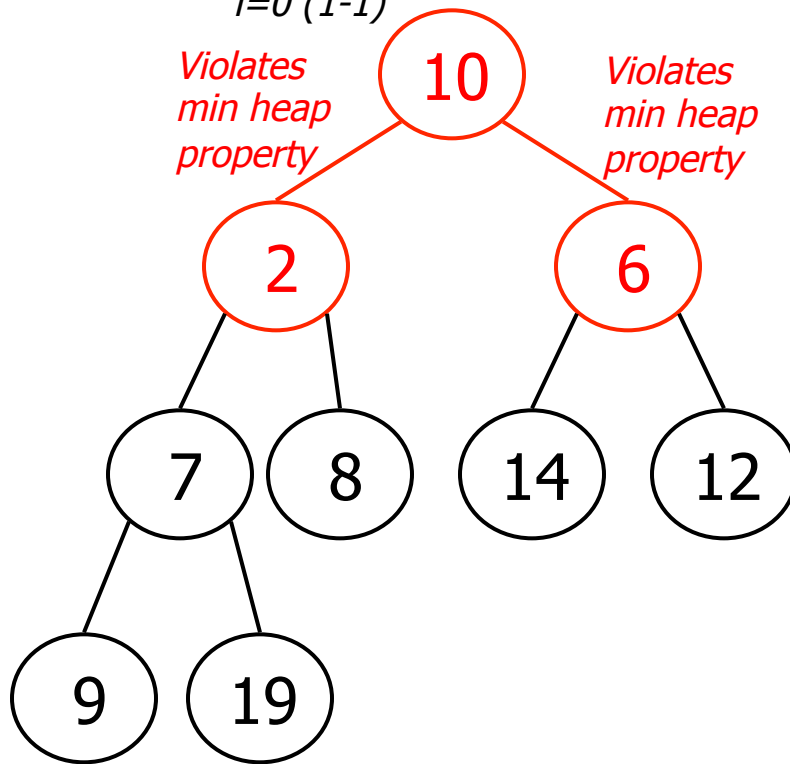
i=0	i=1	i=2	i=3	i=4	i=5	i=6	i=7	i=8
10	7	6	2	8	14	12	9	19



Building a Min Heap

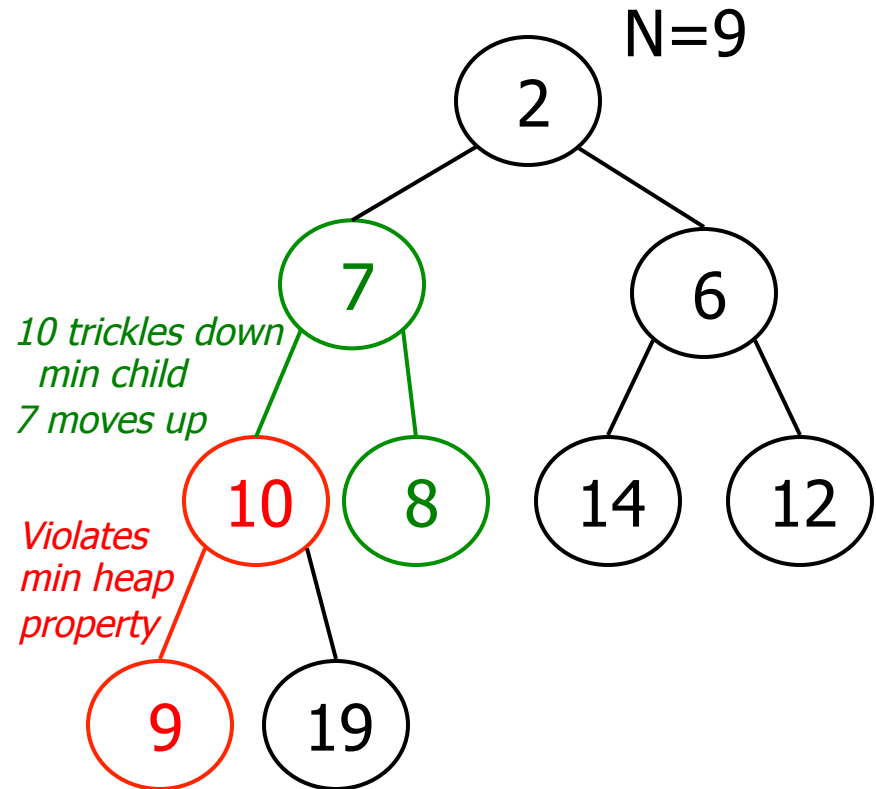
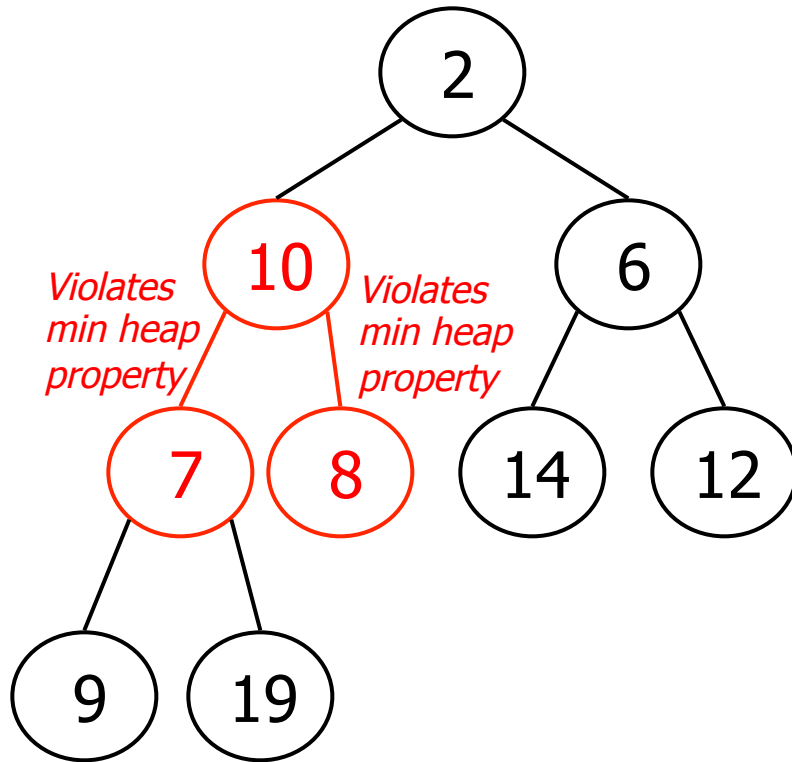
i=0	i=1	i=2	i=3	i=4	i=5	i=6	i=7	i=8
10	2	6	7	8	14	12	9	19

next parent
i=0 (1-1)



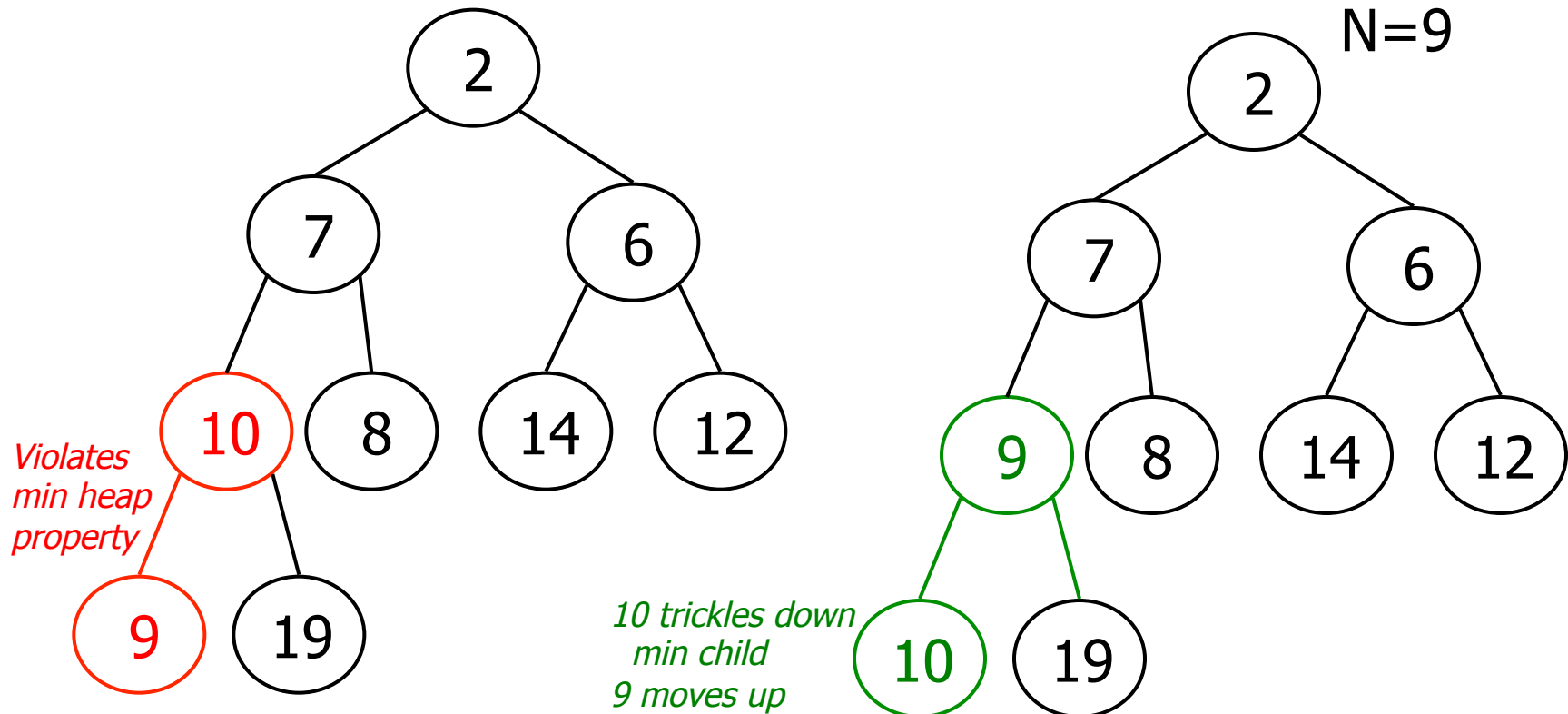
Building a Min Heap

i=0	i=1	i=2	i=3	i=4	i=5	i=6	i=7	i=8
2	10	6	7	8	14	12	9	19



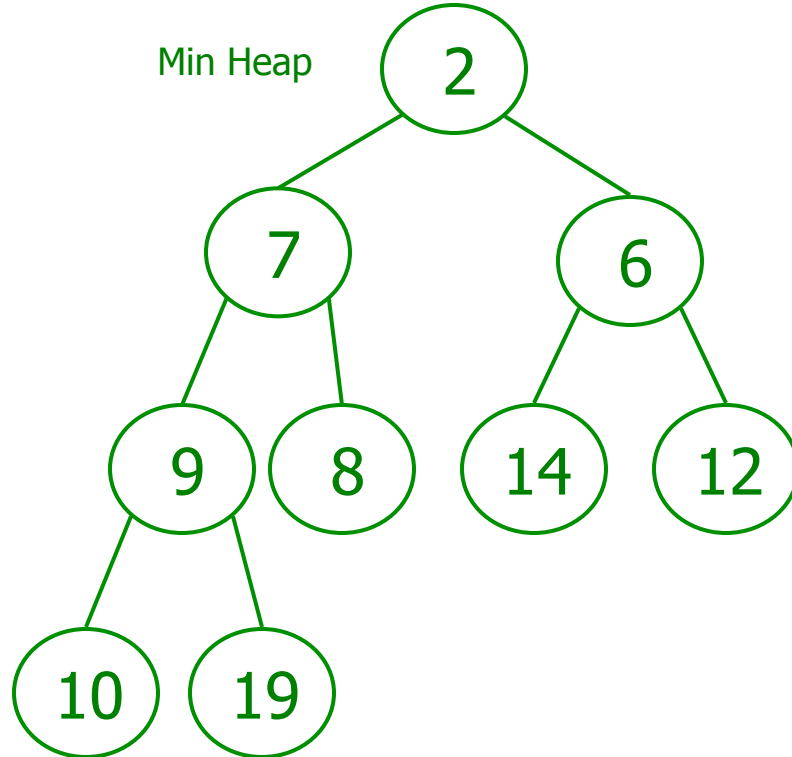
Building a Min Heap

i=0	i=1	i=2	i=3	i=4	i=5	i=6	i=7	i=8
2	7	6	10	8	14	12	9	19



Building a Min Heap

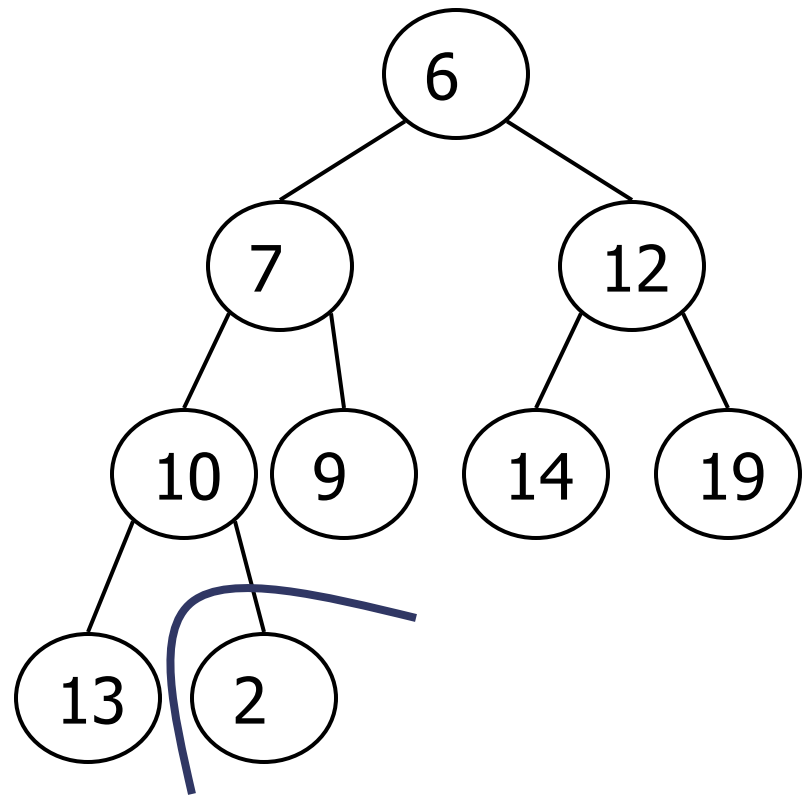
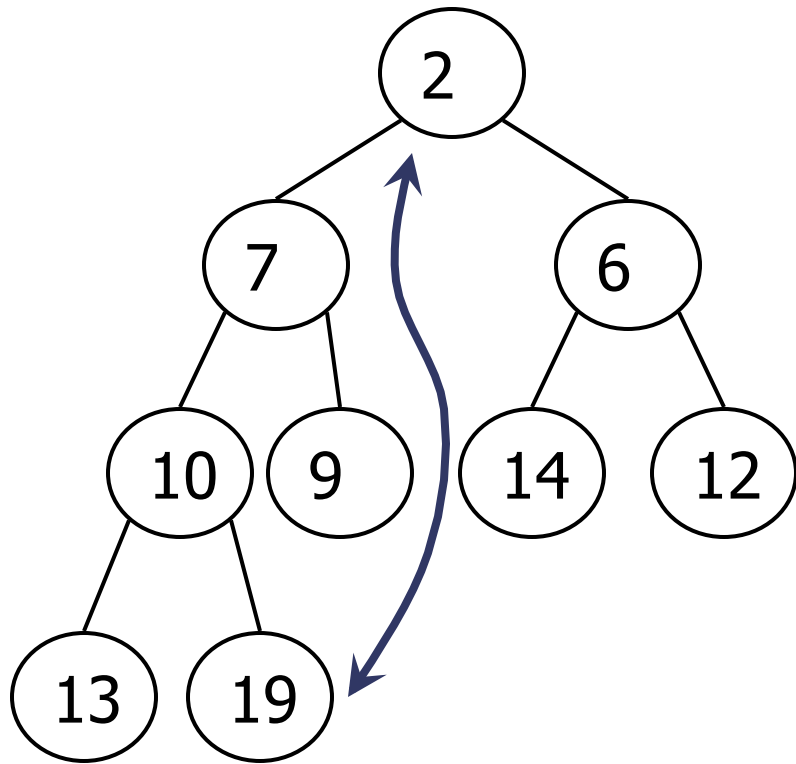
i=0	i=1	i=2	i=3	i=4	i=5	i=6	i=7	i=8
2	7	6	9	8	14	12	10	19



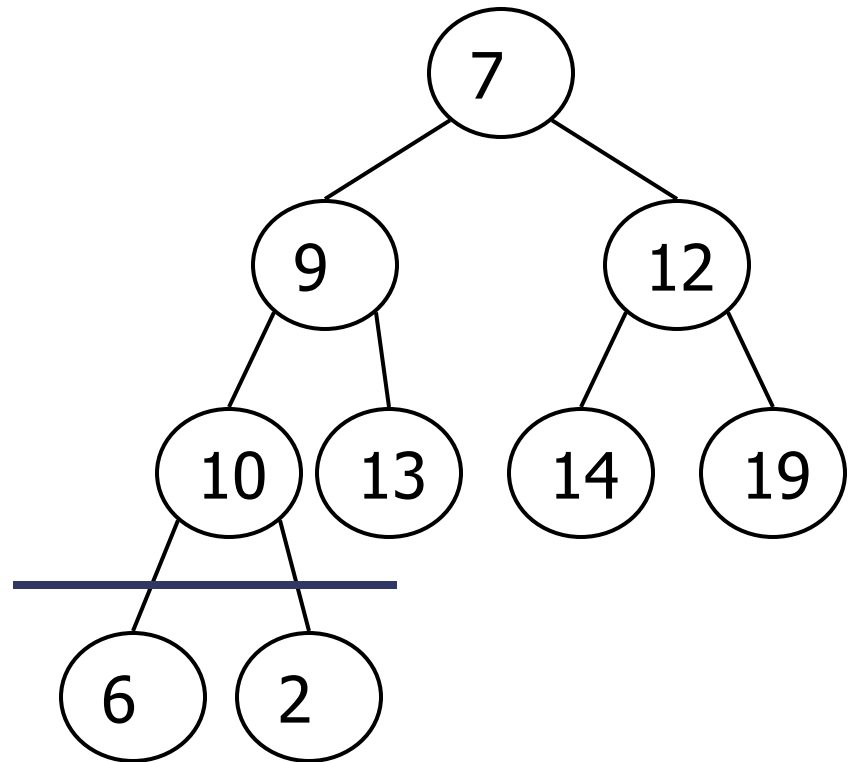
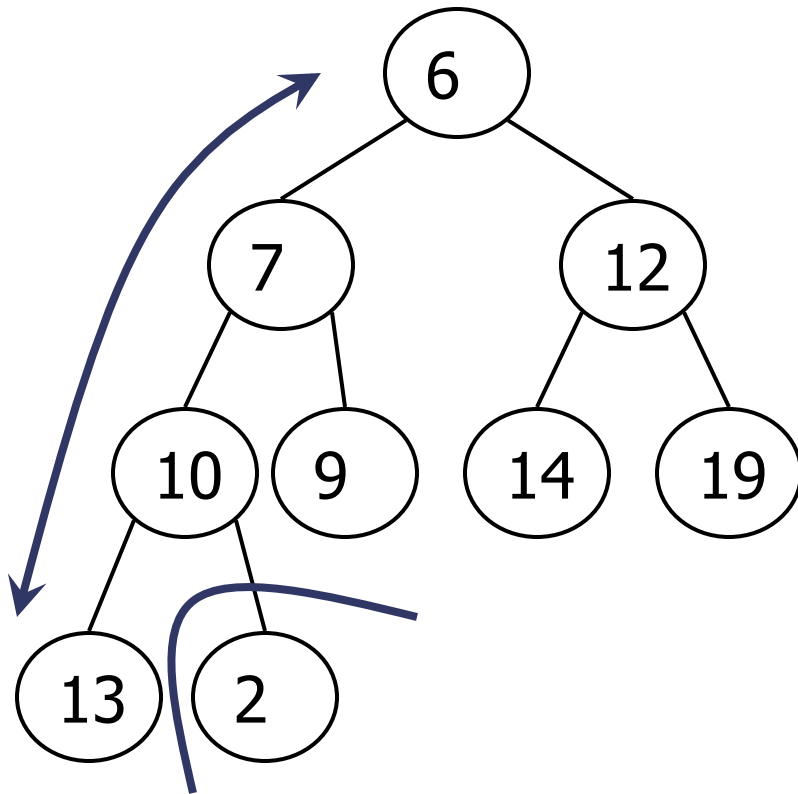
Heapsort

- Build heap for the opposite of what you want
 - Max heap for ascending order
 - Min heap for descending order
- Take root and place in last array position, then think of array as 1 smaller
- Trickle down from root to rebuild heap

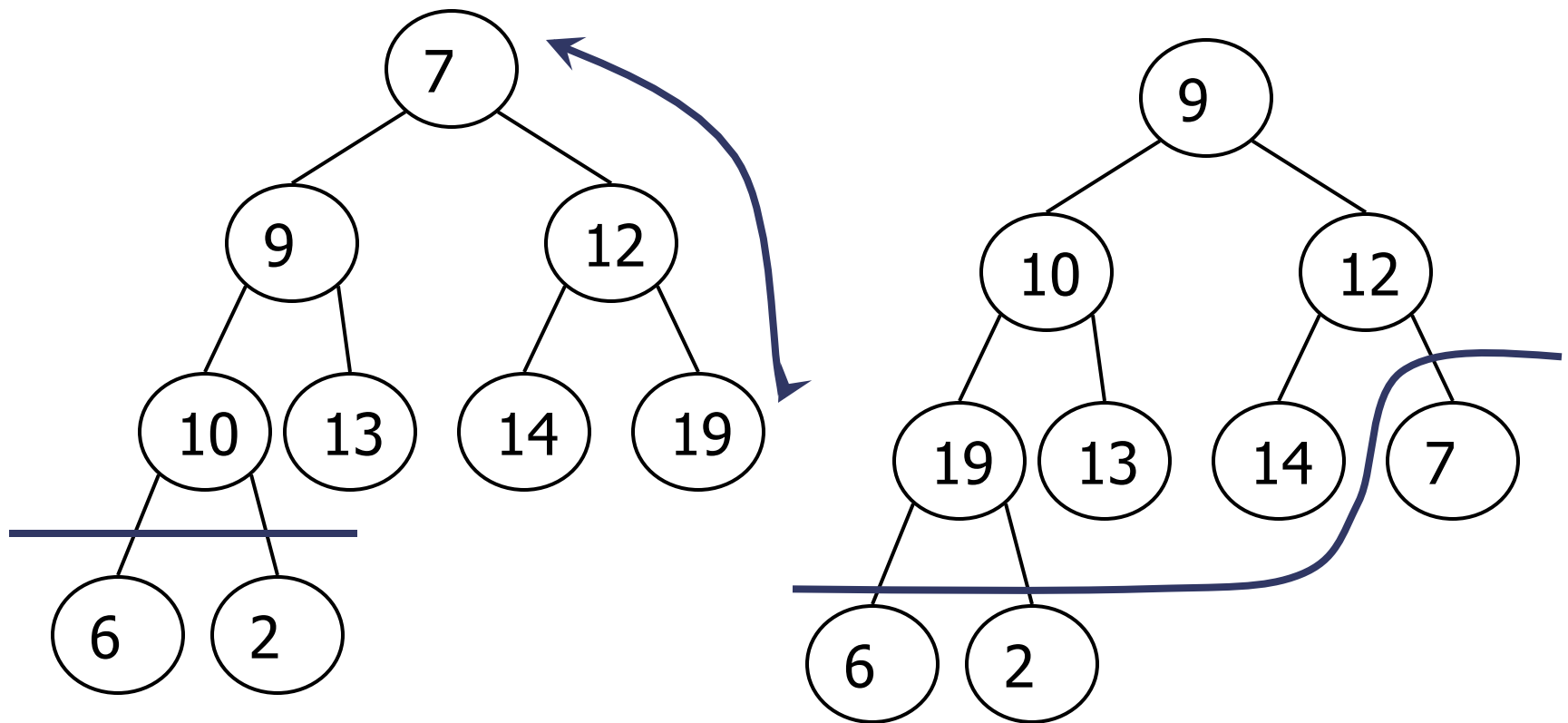
Heapsort - Sort in descending order



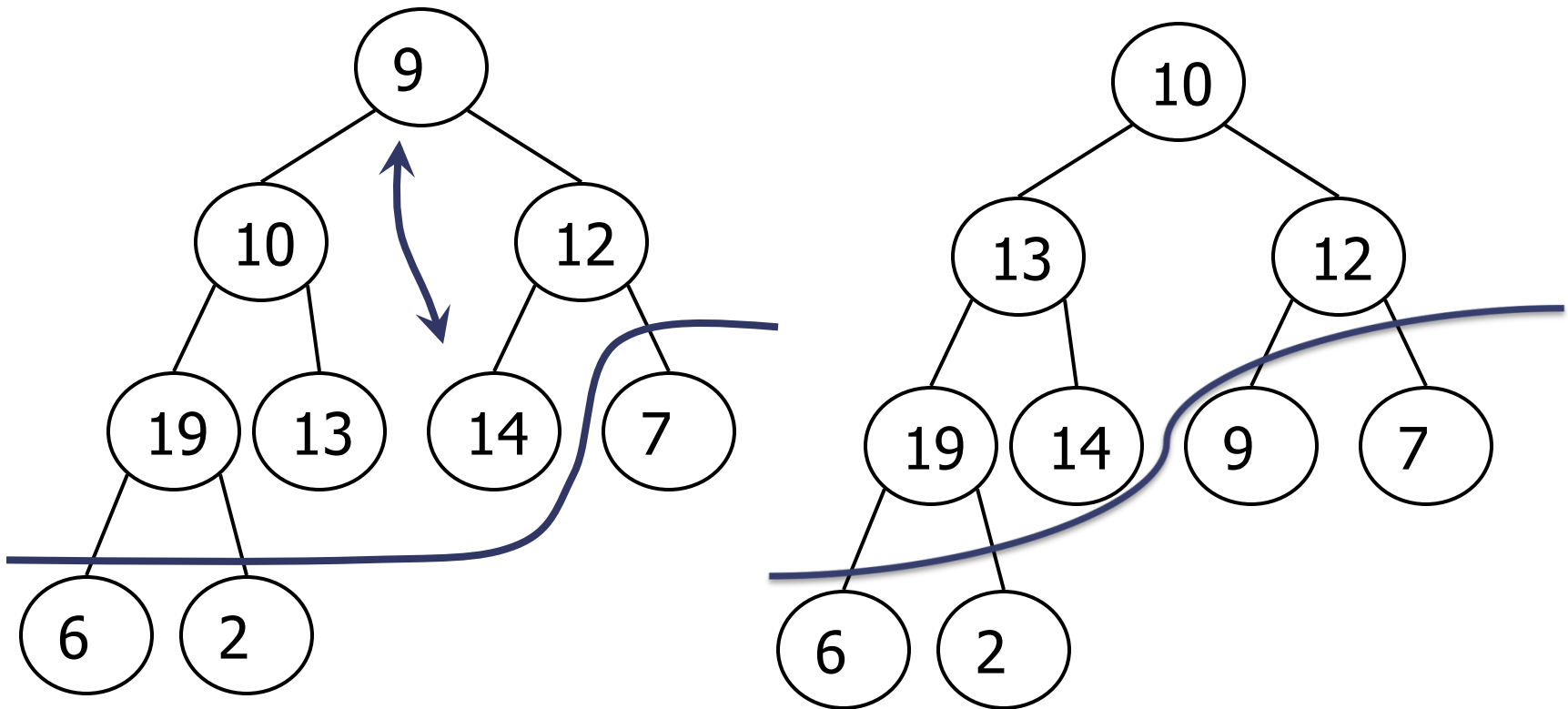
Heapsort - Sort in descending order



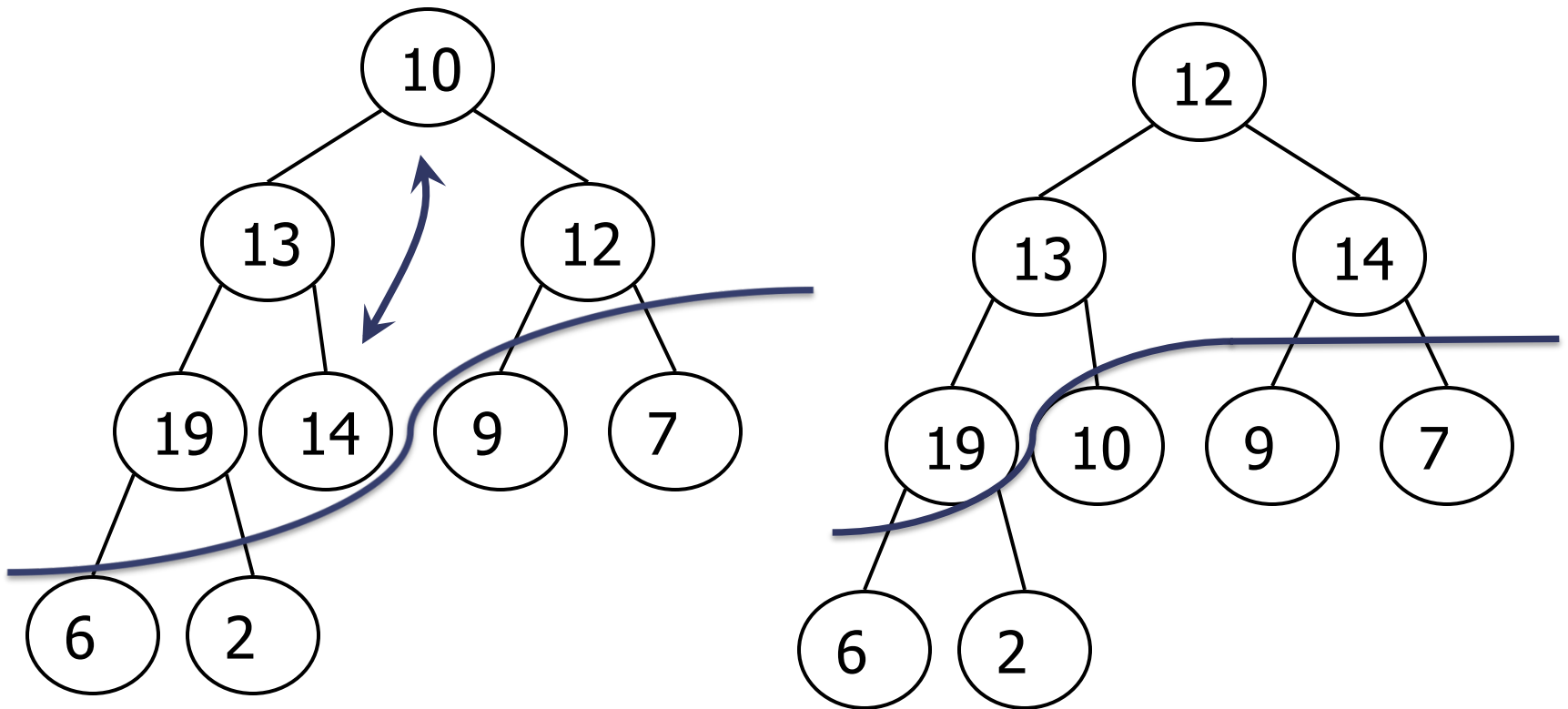
Heapsort - Sort in descending order



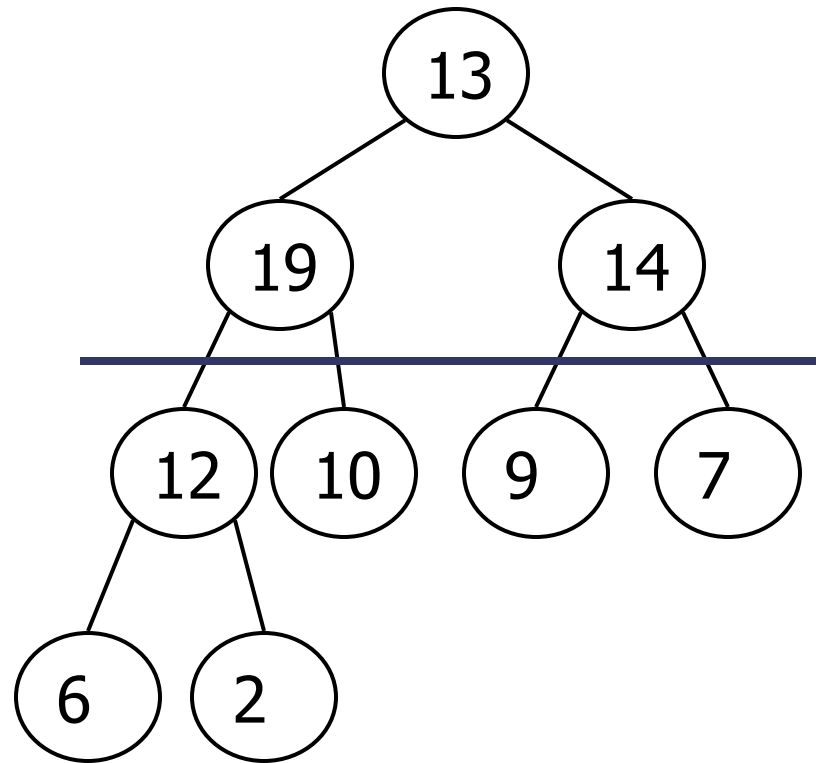
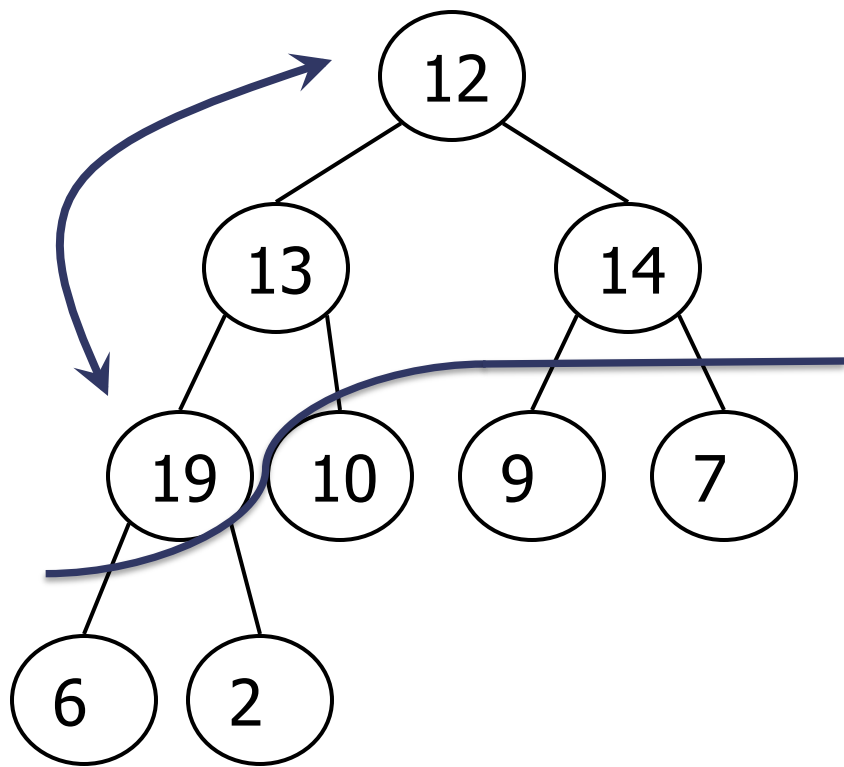
Heapsort - Sort in descending order



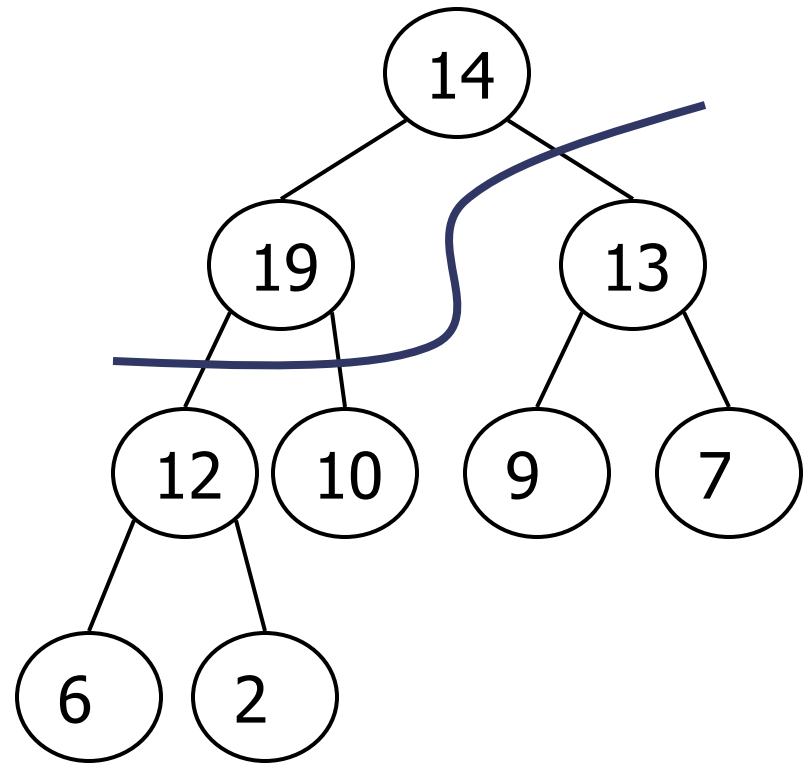
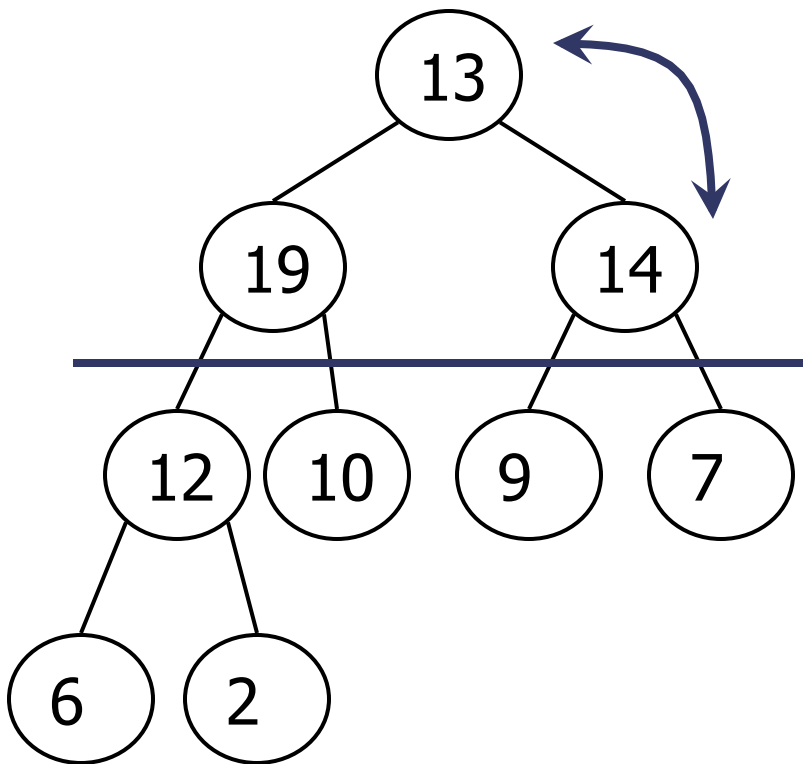
Heapsort - Sort in descending order



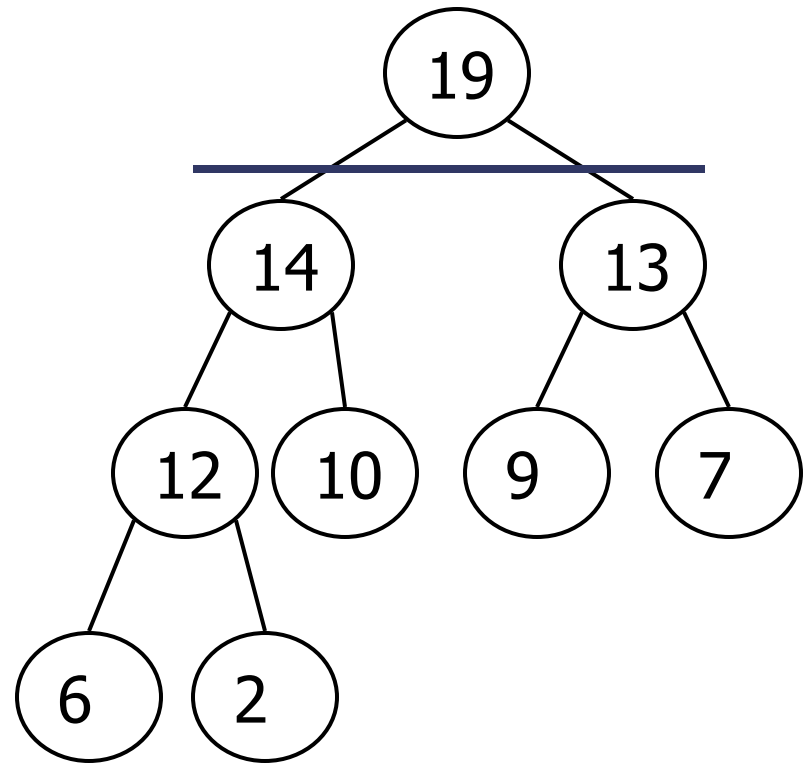
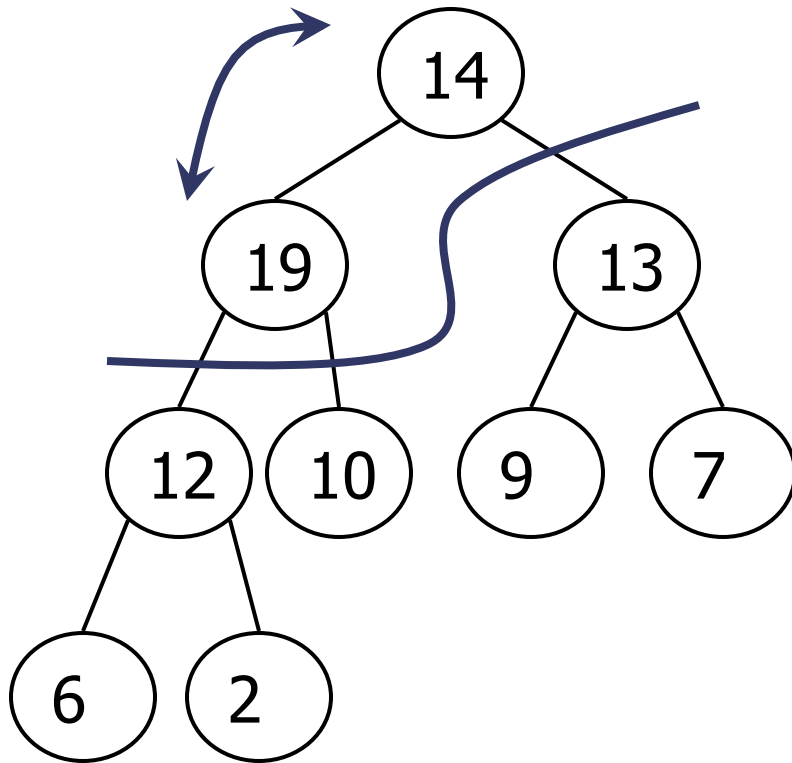
Heapsort - Sort in descending order



Heapsort - Sort in descending order



Heapsort - Sort in descending order



Heapsort Run Time

- Build step - $O(n)$
- Sorting step - $O(n \log n)$
- Heapsort - $O(n) + O(n \log n) = O(n \log n)$