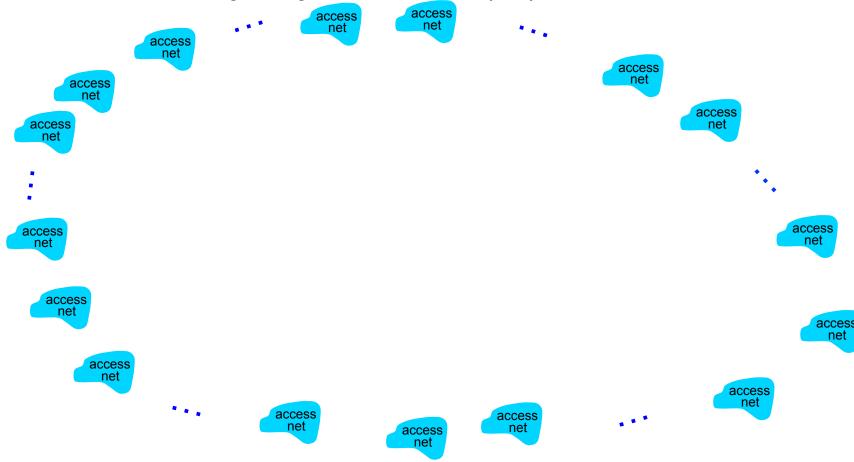
Chapter I: roadmap

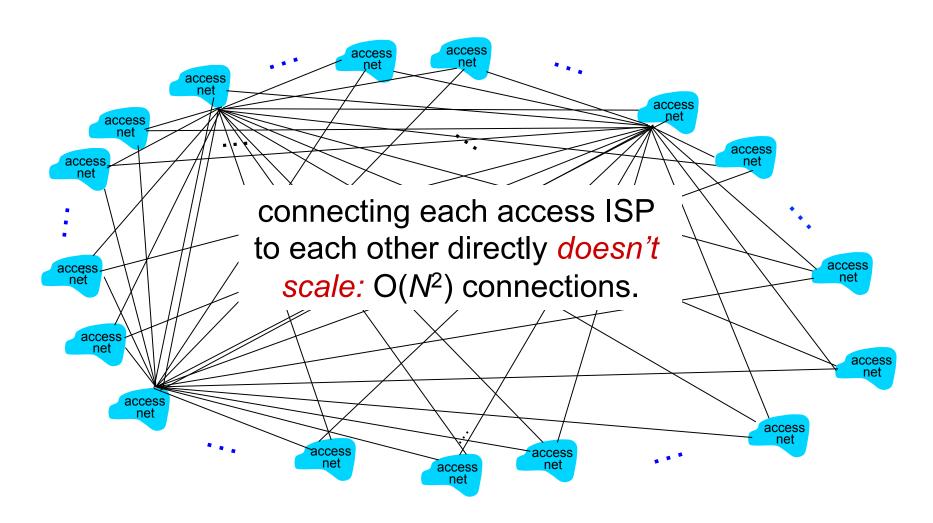
- I.I What is the Internet?
- I.2 Network edge
 - end systems, access networks, links
- 1.3 Network core
 - circuit switching vs packet switching,
 - network structure
- 1.4 Performance:
 - delay, loss and throughput
- 1.5 Protocol layers, service models
- 1.6 Networks under attack: security
- 1.7 History

- End systems connect to Internet via access ISPs (Internet Service Providers)
 - Residential, company and university ISPs
- Access ISPs in turn must be interconnected.
 - So that any two hosts can send packets to each other
- Resulting network of networks is very complex
 - Evolution was driven by economics and national policies
- Let's take a stepwise approach to describe current Internet structure

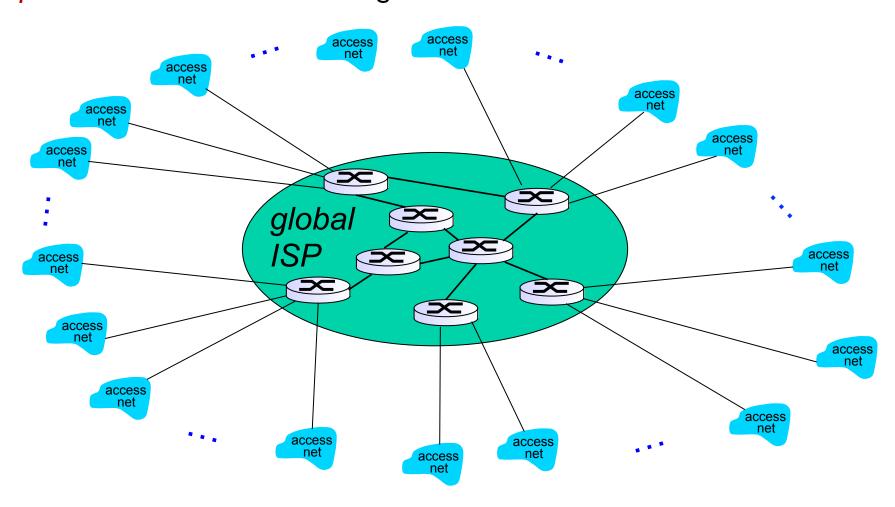
- End systems connect to Internet via access ISPs:
 - residential, company and university ISPs
- Question: given millions of access ISPs, how to connect them together?
 - Considerations: engineering but also economics+policy



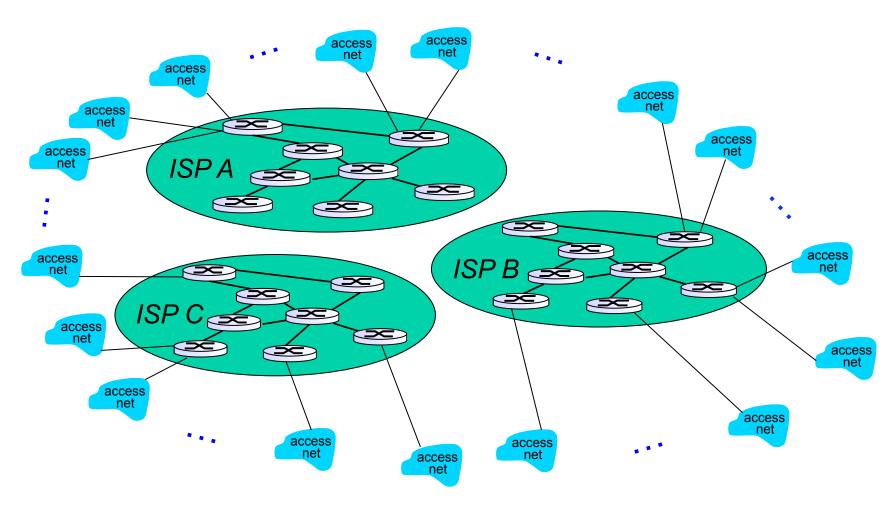
Option: connect each access ISP to every other access ISP?



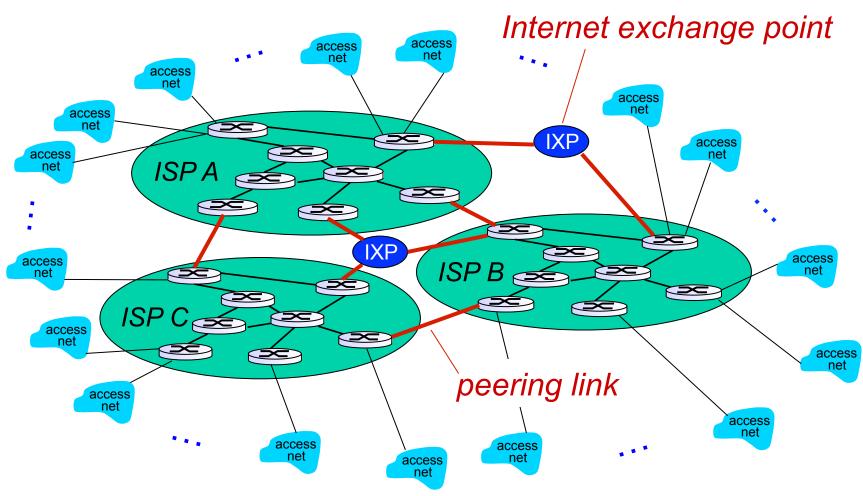
Option: connect each access ISP to a global transit ISP? Customer and provider ISPs have economic agreement.



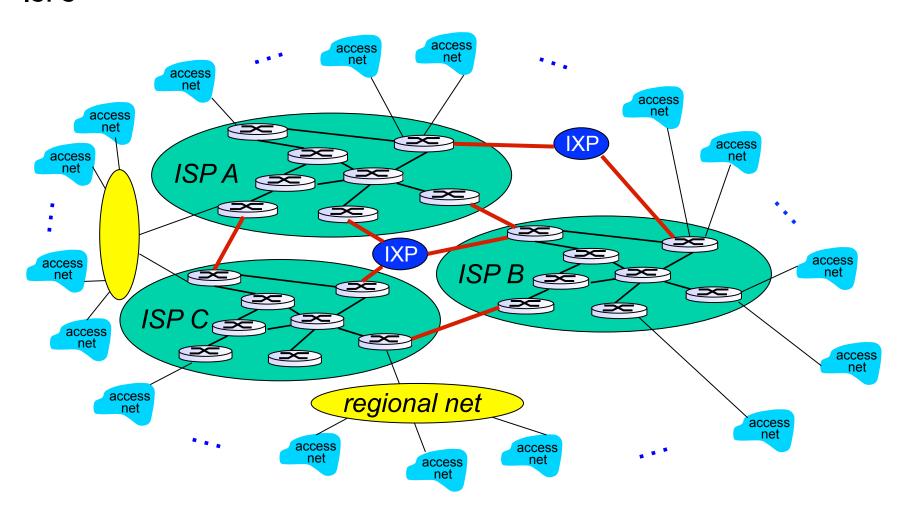
But if one global ISP is viable business, there will be competitors



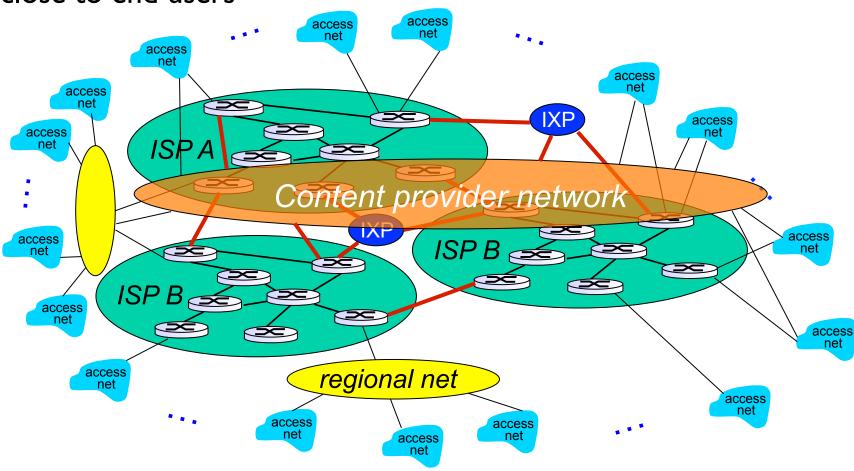
But if one global ISP is viable business, there will be competitors which must be interconnected

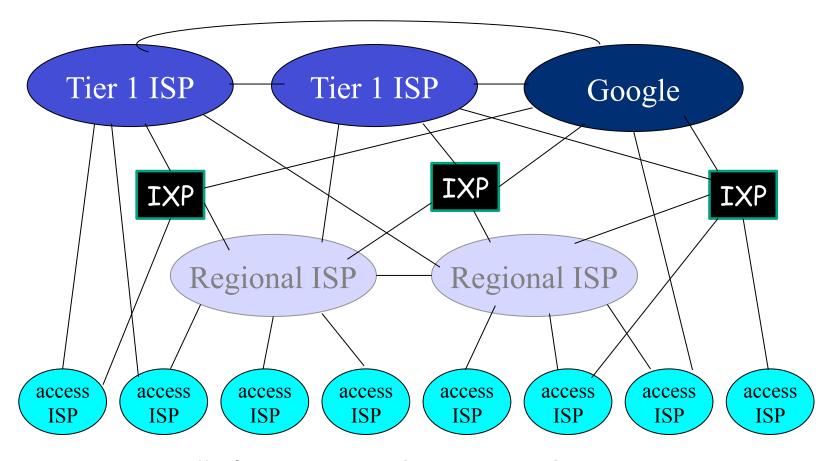


... and regional networks may arise to connect access nets to ISPS



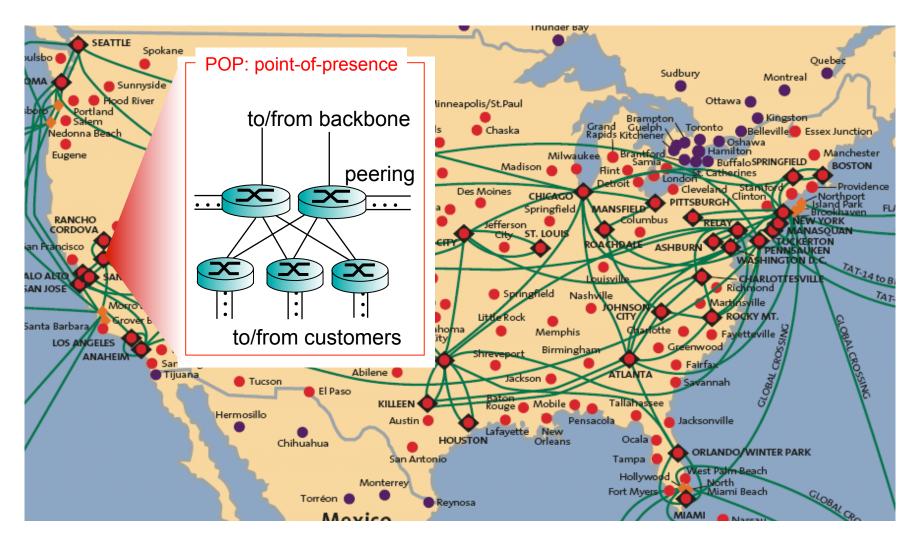
... and content provider networks (e.g., Google, Microsoft, Akamai) may run their own network, to bring services, content close to end users





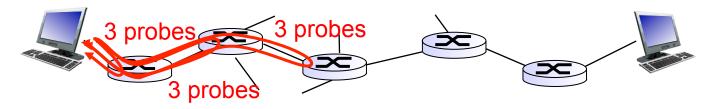
- □ at center: small # of well-connected large networks
 - * "tier-I" commercial ISPs (e.g., Level 3, Sprint, AT&T, NTT), national & international coverage
 - content provider network (e.g, Google): private network that connects it data centers to Internet, often bypassing tier-I, regional ISPs
 Introduction 1-89

Tier-I ISP: e.g., Sprint



"Real" Internet delays and routes

- what do "real" Internet delay & loss look like?
- traceroute program: provides delay measurement from source to router along endend Internet path towards destination. For all i:
 - sends three packets that will reach router i on path towards destination
 - router i will return packets to sender
 - sender times interval between transmission and reply.



"Real" Internet delays and routes

traceroute: from gaia.cs.umass.edu to www.eurecom.fr

```
Three delay measurements from
                                           gaia.cs.umass.edu to cs-gw.cs.umass.edu
I cs-gw (128.119.240.254) I ms I ms 2 ms
2 border l-rt-fa5-1-0.gw.umass.edu (128.119.3.145) l ms l ms 2 ms
3 cht-vbns.gw.umass.edu (128.119.3.130) 6 ms 5 ms 5 ms
4 jn1-at1-0-0-19.wor.vbns.net (204.147.132.129) 16 ms 11 ms 13 ms
5 jn I-so 7-0-0-0. wae. vbns.net (204. I 47. I 36. I 36) 21 ms I 8 ms I 8 ms
6 abilene-vbns.abilene.ucaid.edu (198.32.11.9) 22 ms 18 ms 22 ms
7 nycm-wash.abilene.ucaid.edu (198.32.8.46) 22 ms 22 ms 22 ms
                                                                         trans-oceanic
8 62.40.103.253 (62.40.103.253) 104 ms 109 ms 106 ms
                                                                         link
9 de2-1.de1.de.geant.net (62.40.96.129) 109 ms 102 ms 104 ms
10 de.fr1.fr.geant.net (62.40.96.50) 113 ms 121 ms 114 ms
II renater-gw.frI.fr.geant.net (62.40.103.54) II2 ms II4 ms II2 ms
12 nio-n2.cssi.renater.fr (193.51.206.13) 1 | 1 ms | 114 ms | 116 ms
13 nice.cssi.renater.fr (195.220.98.102) 123 ms 125 ms 124 ms
14 r3t2-nice.cssi.renater.fr (195.220.98.110) 126 ms 126 ms 124 ms
15 eurecom-valbonne.r3t2.ft.net (193.48.50.54) 135 ms 128 ms 133 ms
16 194.214.211.25 (194.214.211.25) 126 ms 128 ms 126 ms
17 ***
                     *means no response (probe lost, router not replying)
18 ***
    fantasia.eurecom.fr (193.55.113.142) 132 ms 128 ms 136 ms
```

More traceroute examples

17 * * *

```
Athina-Markopoulous-MacBook-Air-2:web athina$ traceroute gaia.cs.umass.edu
traceroute to gaia.cs.umass.edu (128.119.245.12), 64 hops max, 52 byte packets
1 192.168.0.1 (192.168.0.1) 1.546 ms 0.247 ms 0.618 ms
2 10.75.152.1 (10.75.152.1) 6.692 ms 7.636 ms 7.960 ms
3 ip68-4-11-142.oc.oc.cox.net (68.4.11.142) 14.459 ms 8.553 ms 7.195 ms
4 ip68-4-11-12.oc.oc.cox.net (68.4.11.12) 26.645 ms 8.021 ms 7.628 ms
5 langbprj01-ae2.0.rd.la.cox.net (68.1.0.136) 8.452 ms 9.666 ms 9.866 ms
6 lsan0.tr-cps.internet2.edu (206.223.123.199) 9.921 ms 11.280 ms 10.319 ms
7 xe-0-2-0.0.sttl0.tr-cps.internet2.edu (64.57.20.223) 34.258 ms 34.985 ms 34.152 ms
8 xe-1-1-0.0.chic0.tr-cps.internet2.edu (64.57.20.169) 87.729 ms 88.794 ms 88.357 ms
9 198.71.47.62 (198.71.47.62) 112.777 ms 111.206 ms 111.823 ms
10 192.5.89.21 (192.5.89.21) 98.682 ms 98.637 ms 97.918 ms
11 nox300gw1-peer--207-210-142-242.nox.org (207.210.142.242) 98.917 ms 100.046 ms 99.847 ms
12 core1-rt-xe-0-0-0.gw.umass.edu (192.80.83.101) 100.235 ms 99.676 ms 98.883 ms
13 lgrc-rt-106-8-po-10.gw.umass.edu (128.119.0.233) 100.358 ms 100.123 ms 99.548 ms
14 128.119.3.32 (128.119.3.32) 101.201 ms 100.145 ms 99.514 ms
15 nscs1bbs1.cs.umass.edu (128.119.240.253) 101.423 ms 100.951 ms 103.533 ms
```

More traceroute examples

 $At hin a-Markopoulous-MacBook-Air-2: web \ at hin a \$\ traceroute\ nestor. calit 2.uci.edu$

traceroute to nestor.calit2.uci.edu (128.195.185.109), 64 hops max, 52 byte packets

- 1 302-wism-v1970.ucinet.uci.edu (169.234.0.1) 59.342 ms 27.768 ms 90.426 ms
- 2 cs1-core--302-wism.ucinet.uci.edu (128.195.249.129) 1.325 ms 1.361 ms 1.254 ms
- 3 325--cs1-core.ucinet.uci.edu (128.195.249.190) 1.051 ms 1.113 ms 1.035 ms
- 4 nestor.calit2.uci.edu (128.195.185.109) 1.019 ms 1.053 ms 1.030 ms

Athina-Markopoulous-MacBook-Air-2:web athina\$ traceroute www.google.com

traceroute: Warning: www.google.com has multiple addresses; using 74.125.224.178

traceroute to www.google.com (74.125.224.178), 64 hops max, 52 byte packets

- 1 302-wism-vl970.ucinet.uci.edu (169.234.0.1) 2.994 ms 0.930 ms 1.018 ms
- 2 cs1-core--302-wism.ucinet.uci.edu (128.195.249.129) 1.036 ms 1.003 ms 0.966 ms
- 3 kazad-dum-v1123.ucinet.uci.edu (128.200.2.222) 1.286 ms 1.088 ms 1.234 ms
- 4 dc-tus-agg2--uci-ge-1.cenic.net (137.164.24.49) 2.816 ms 2.306 ms 1.700 ms
- 5 riv-core1--tus-agg2-10ge-2.cenic.net (137.164.47.79) 9.473 ms 6.731 ms 13.066 ms
- 6 dc-lax-core1--riv-core1-10ge-2.cenic.net (137.164.46.57) 12.061 ms 9.388 ms 11.659 ms
- 7 72.14.223.85 (72.14.223.85) 7.697 ms 7.713 ms 7.751 ms
- 8 64.233.174.238 (64.233.174.238) 7.777 ms 9.668 ms 8.567 ms
- 9 72.14.236.11 (72.14.236.11) 8.098 ms 8.008 ms 8.128 ms
- 10 lax02s01-in-f18.1e100.net (74.125.224.178) 8.130 ms 7.953 ms 8.101 ms

Athina-Markopoulous-MacBook-Air-2:web athina\$

Chapter I: roadmap

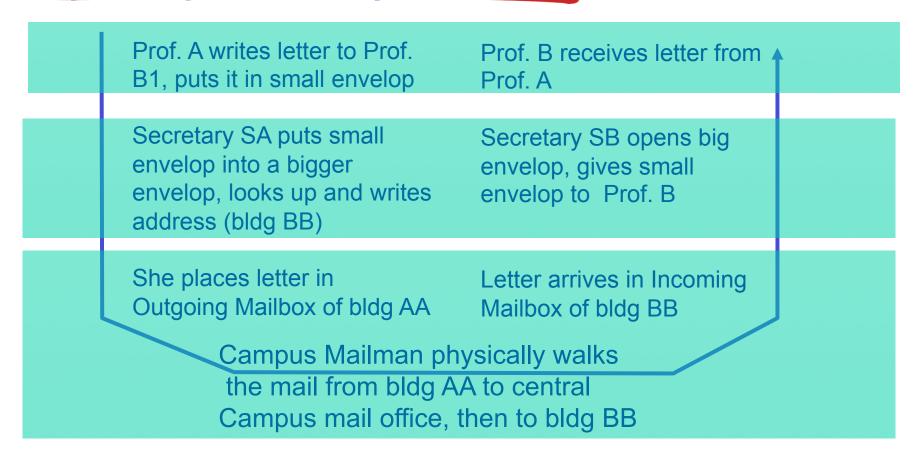
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Dealing with Scale and Complexity

- ☐ The Internet is highly complex, in terms of ...
 - Its sheer size (number of components)
 - The number of tasks it needs to manage: routing, congestion control, packet reordering, connection establishment ...
 - Diversity of Components, Topology, Functionality,
- ☐ Tasks are structured through modularization
 - Divide and Conquer
 - Break down into smaller pieces
 - Create different functional layers



Layering of Campus Mail



layers: each layer implements a service

- via its own internal-layer actions
- relying on services provided by layer below

Why layering?

dealing with complex systems:

- explicit structure allows identification, relationship of complex system's pieces
 - layered reference model for discussion
- modularization eases maintenance, updating of system
 - change of implementation of layer's service transparent to rest of system
 - e.g., change in secretary, secretary's routine, mailman's routine, doesn't affect rest of system
- layering considered harmful?
 - duplication of functionality?
 - cross-layer optimization?

Internet protocol stack

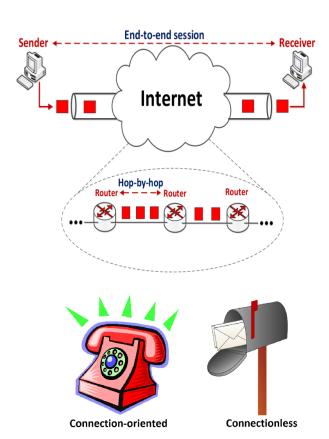
network medium (fiber, wireless,...)

bits on the wire

Application layer: 5 What we interact with: **Application** ■ Transport layer: **Transport** process-process data transfer 4 (TCP/UDP) End-to-end management: TCP, UDP ■ Network layer: **Network** routing of datagrams from source to destination (IP) IP, routing protocols 2 Link □ Link Layer: data transfer between neighboring network elements 1 **Physical** Ethernet, 802.111 (WiFi), PPP physical:

Transport & Network

- "Thin waist"
 - Other layers: evolved dramatically
 - TCP/IP: Stayed mostly the same
- Transport layer
 - End-to-end
 - Connection-oriented in TCP
- Network layer
 - Hop-by-hop
 - Connectionless in IP



ISO/OSI reference model

- presentation: allow applications to interpret meaning of data, e.g., encryption, compression, machinespecific conventions
- session: synchronization, checkpointing, recovery of data exchange
- Internet stack "missing" these layers!
 - these services, if needed, must be implemented in application

application
presentation
session
transport
network
link
physical

Headers **Application** Message Data **Encapsulation:** augment data Segment Data **Transport** Header with headers Payload: actual content **Packet** Network Data Header Header Header: identification and control information Frame Link Data Header Header Header

- ☐ This creates overhead. Why bother?
 - Allows us to distinguish messages
 - e.g.: Layer 3 header contains src & dst IP address of next hop
 - Essential for packet switching

"Understanding" Layers

- Different network elements process up to different layers
- End-hosts
 - Server, computer
 - Process all 5
- Routers
 - ❖ Up to 3
 - Have IP addresses
- Switches
 - ❖ Up to 2
 - Do not have or process IP

