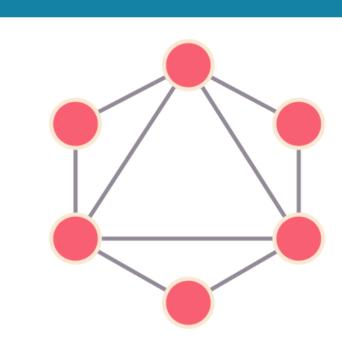
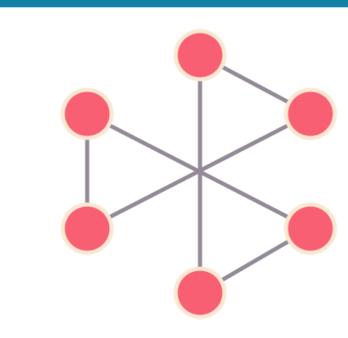
Path Planning for Dynamic Graphs using A* on GPU

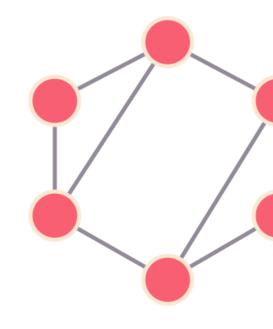
Abstract

- A* is one of the widely used path planning algorithms applied in a diverse set of problems in robotics and video games.
- Zhou and Zeng [1] proposed a parallel variant of A* for GPU, which keeps multiple priority queues to find the optimal path in a static graph (static A*).
- Here we present A* for dynamic graphs (dynamic A*) on GP-GPUs which achieves 2x-7x speedup than static A* on the SNAP dataset [2]

Dynamic Graphs



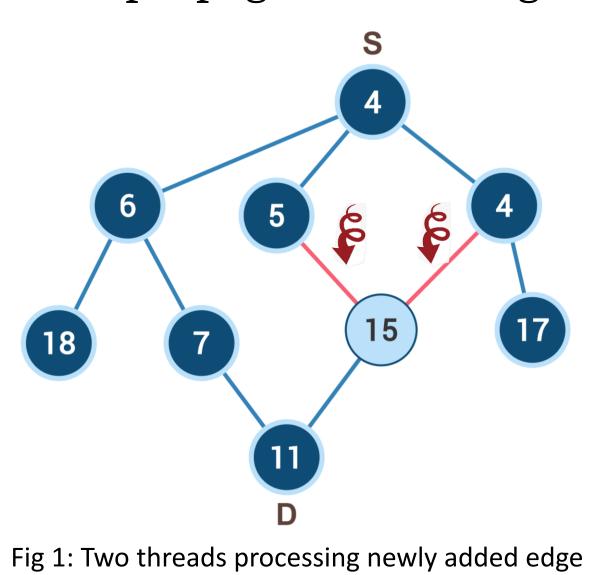




Time

Dynamic A*: Insertion of edges

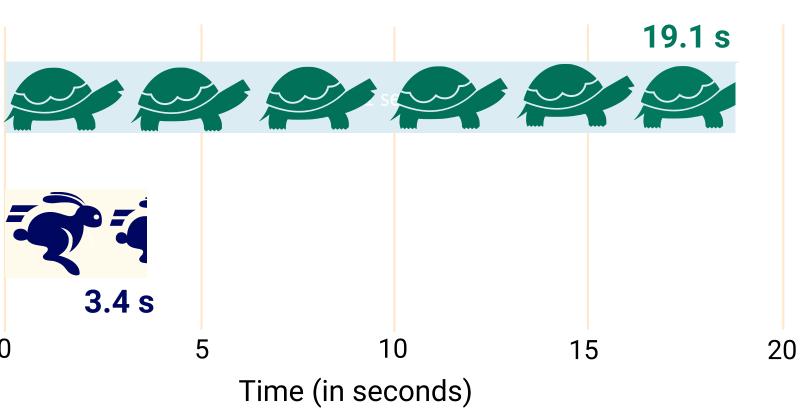
- Newly added edges can alter the optimal path.
- To find the new optimal path instead of executing A* from scratch, we propagate the change to the affected nodes of the graph.

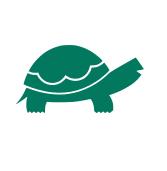


(shown in red)

Pseudocode!

- 1. For edges(u, v) inserted, add node v to update_list, if $f(v)_{new} < f(v)_{old}$.
- 2. While update_list not empty:
- a. Extract node n from update_list.
 - b. For each child of n:
 - i. lock(child)
 - ii. if f(child)_{new}< f(child)_{old}, add child
 - to update_list. iii. unlock(child)
- f(v): cost from source to v + heuristics value of v





Static A*

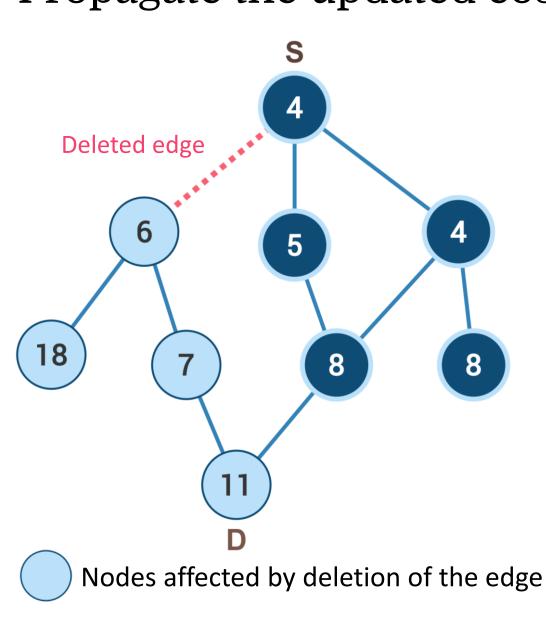


Dynamic A* (Insertions)

Fig 2: Execution time of static A* and dynamic A* (for only insertions) on graph Wiki-Talk

Dynamic A*: Deletion of edges

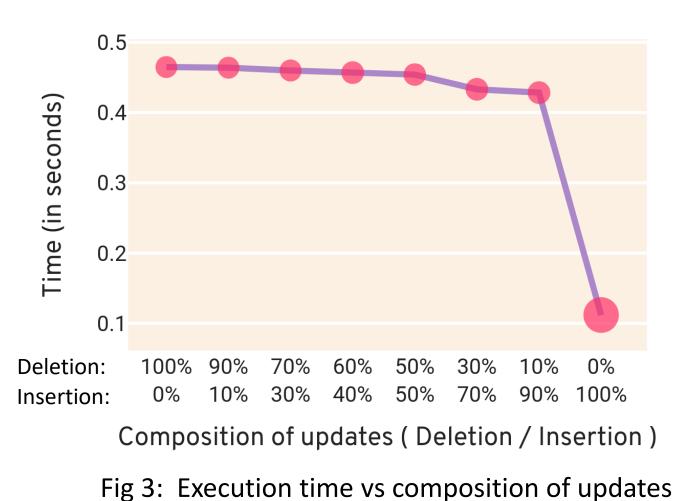
- Deleting only that edge which belongs to the optimal path can create a new optimal path.
- For all such affected nodes recompute the cost and select the neighbour with the least cost.
- Propagate the updated cost to all the affected nodes.

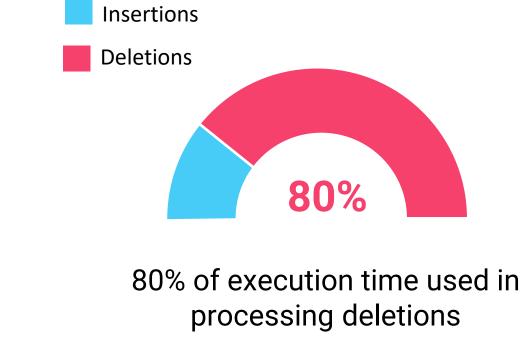


Pseudocode!

- For each deleted edge u → v:
 a. Compute f(v) from the neighbours of v.
 b. Add v to update_list.
- 2. While update_list is not empty:
- a. Extract node n from update_list.
- **b.** For each child of **n** such that
- optimal_parent(child) = n:
 i. If f(child) > f(child)_{new} then compute cost of child from each of its
- neighbour and add it to update_list.

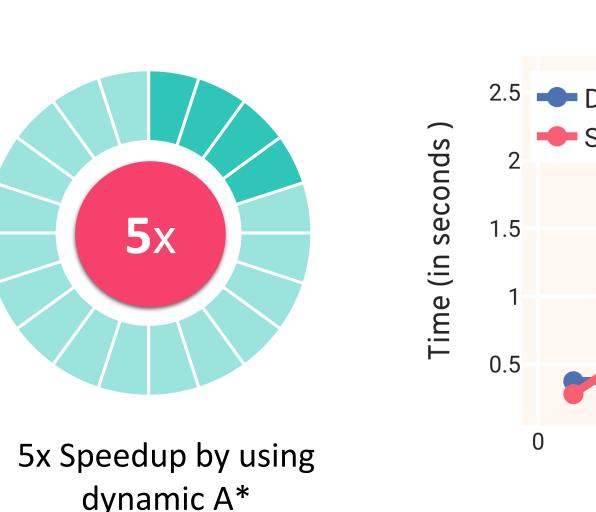
optimal_parent(v): neighbour of node v with least f(v)





Dynamic A*: Fully Dynamic

- The update contains both insertion and deletion of edges.
- Propagate insertions and deletions of edges separately.
- Performs better than re-executing static A* algorithm after each update.



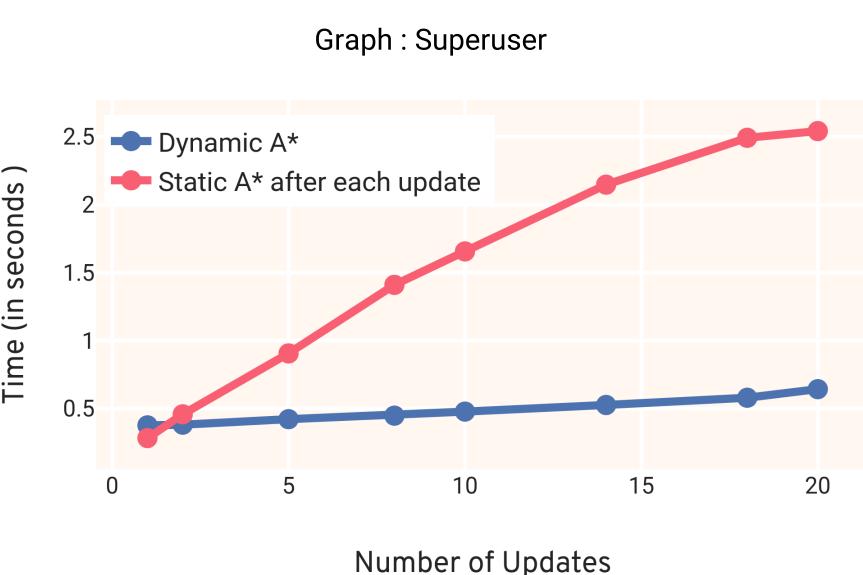


Fig 4: Execution time vs number of updates In the graph

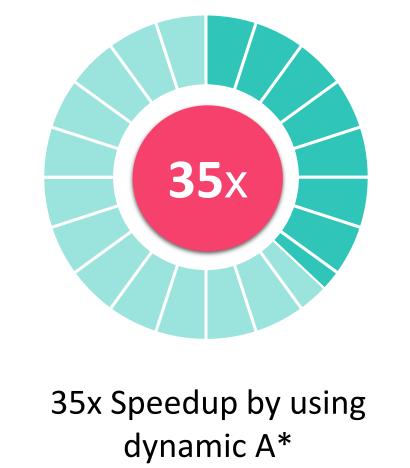
Results

The below table shows execution time (in seconds) and speedup of dynamic A* compared to re-executing static A* every time.

No.	Graph	Edges	Queries	Dynamic A*	Static A*	Speedup
1	Live Journal	34,681,189	10	6.01	33.93	5x
2	Wiki Talk	7,833,140	10	12.24	24.84	2x
3	Ask Ubuntu	964,437	10	0.25	1.31	5x
4	YouTube	2,987,624	10	0.81	5.78	7x
5	Math Overflow	506,550	10	0.09	0.67	7x
6	Live Journal	34,681,189	100	11.41	424.06	37x

Applications

1. We have applied dynamic A* on energy efficient routing protocol (EERP) and achieved 35x speedup from static A*.



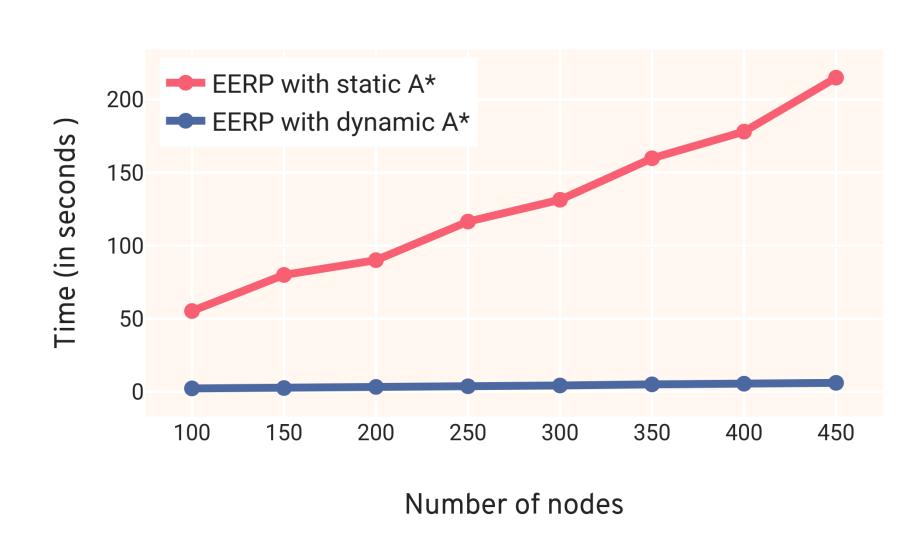


Fig 5: Comparison b/w static A* and dynamic A* on EERP algorithm

2. On applying dynamic A* for pathfinding in the maze, we achieved 8x speedup.

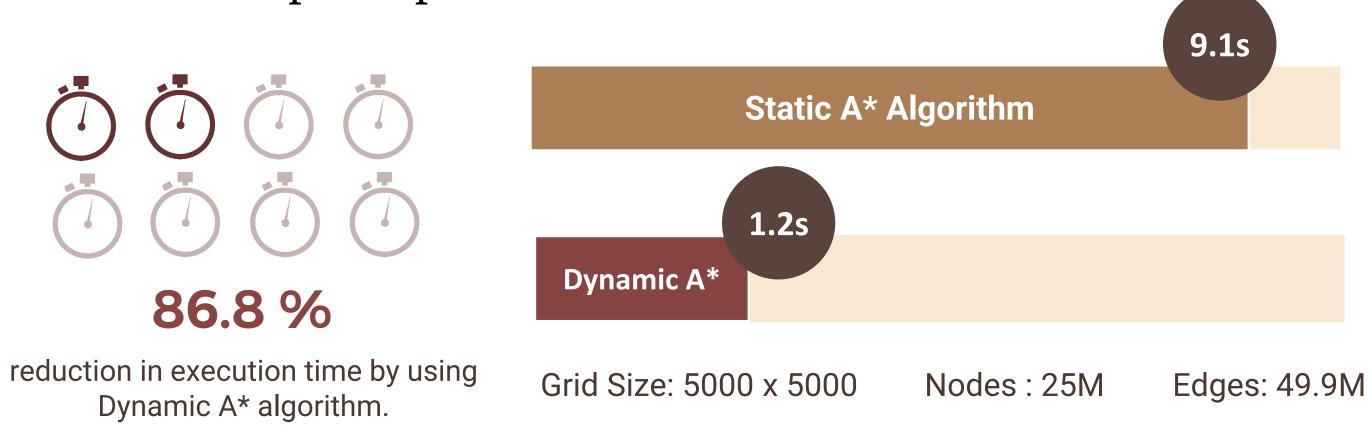


Fig 6: Comparison b/w static A* and dynamic A* on solving 5000 x 5000 maze

Time in seconds

References

- 1. Yichao Zhou and Jianyang Zeng. "Massively Parallel A* Search on GPU". In: Twenty-Ninth AAAI Conference on Artificial Intelligence (2015).
- 2. Jure Leskovec and Andrej Krevl. SNAP Datasets: Stanford Large Network Dataset Collection. http://snap.stanford.edu/data. June 2014.
- Pseudocode described above is to give a basic idea of the algorithm. It does not cover all the cases, for more information please refer to GitHub.