Gaurav Kuwar CSC 332 LAB (Spring 2022) May 5, 2022 Lab 5: Process Synchronization

Task 2 - Report

I used 3 semaphores to solve the issues.

semMutex – is used to lock a transaction and prevent race condition by making the portion of code modifying balance.txt and/or attemps.txt one transction.

semAtms – used to make sure 2 processes are running at once, since there are only 2 ATM machines in the bank.

semNoBalance – main purpose is to prevent negative balance. This is used to check if there is enough balance to be withdrawn by a son process.

This section of code is the first part of the dad process, where the dad process simply reads the balance from balance.txt. I use the P(semAtms) to wait for dad process to get access to the ATM. But, then the dad process sleeps, and at this time since other processes need access to the ATM, so I use V(semAtms) to stop dad process from access to ATM. After the sleep dad process waits again for access to the ATM.

```
P(semAtms); // Dad waits for access to ATM again
P(semMutex); // lock this transction / atomic
fp1 = fopen("balance.txt", "r+"); //Dad successfully got hold of the ATM.
fscanf(fp1, "%d", &bal2);
fseek(fp1,0L,0); //Dad will now deposit the money. And update the current balance
bal3 = bal2 + DepositAmount; // bal2 is saved as previous balance
fprintf(fp1, "%d \n", bal3);
fclose(fp1);
printf("Dad writes new balance = %d \n", bal3);
if (bal2 < WithdrawAmount && bal3 >= WithdrawAmount)
  V(semNoBalance);
if (bal2 < (2*WithdrawAmount) && bal3 >= (2*WithdrawAmount))
  V(semNoBalance);
V(semMutex); // unlock
printf("Dad will deposit %d more time\n",N-i); //Dad depostited the money.
V(semAtms); // Dad loses access to ATM
sleep(rand()%(sleepTime*2)+1); /* Dad will wait some time for requesting to see
```

This is the second part of the dad process. Here, the dad process gets access to the ATM again. Then, I used P(semMutex) so the atomic transaction can begin. I read the balance again, since it could have changed in the time dad process was asleep. I add the deposit amount to balance, but I have a bal3 integer, so I can store the balance before deposit, in bal2. I then use 2 if statements, which use V(semNoBalance) to signal the son processes waiting for enough balance. The second if statement is for situations where both sons are waiting for enough balance. I then use V(semMutex) to unlock the transaction and V(semAtms) so that the Dad process can lose access to ATM.

```
P(semAtms);

fp3 = fopen("attempt.txt" , "r+"); //Son_1 successfully got hold of the ATM.

fscanf(fp3, "%d", &N_Att); // Son_1 Checks if he has more than 0 attempt remaining.

printf("Attempt remaining: %d.\n", N_Att);

if(N_Att == 0)

{

fclose(fp3);

flag = TRUE;

}

else

fp2 = fopen("balance.txt", "r+"); //Son_1 reads the balance.

fscanf(fp2,"%d", &bal2);

printf("SON_1 reads balance. Available Balance: %d \n", bal2);

printf("SON_1 wants to withdraw money. \n"); //And if balance is greater than Withdraw amount
```

```
printf("SON_1 reads balance. Available Balance: %d \n", bal2);
printf("SON_1 wants to withdraw money. \n"); //And if balance is greater than Withdraw am
if (bal2 < WithdrawAmount) {</pre>
 V(semAtms); // son loses access to ATM
  P(semNoBalance); // waits for dad to put balance
  P(semAtms); // son gains access to ATM again
P(semMutex); // lock this transction / atomic
fp2 = fopen("balance.txt", "r+");
fscanf(fp2,"%d", &bal2);
fseek(fp2,0L, 0);
bal2 -=WithdrawAmount;
fprintf(fp2,"%d\n", bal2);
fclose(fp2);
printf("SON_1 withdrawed %d. New Balance: %d \n",WithdrawAmount, bal2);
fp3 = fopen("attempt.txt" , "r+");
fscanf(fp3, "%d", &N_Att);
fseek(fp3,0L, 0); //SON_1 will write the number of attempt remaining in the attampt.txt
N_Att -=1;
fprintf(fp3, "%d\n", N_Att);
fclose(fp3);
printf("Number of attempts remaining:%d \n", N_Att);
V(semMutex); // unlock
```

This is a section of the son 1 process code (but the same idea applies to son 2 process as well). First, the son process starts with P(semAtms) which gives the process access to ATM, and the iteration of the code exits with V(semAtms) before sleeping just like with the dad process. The if statement on line 172 is used to prevent negative balance. It checks if the bal2 is less than the withdrawamount, and if it is it will stop access to ATM, then it will call P(semNoBalance), at this point, since there isn't enough balance semNoBalance would be 0, so the son process will now wait for the dad process to add enough balance and call V(semNoBalance), as describe in the previous code. Once, there is enough balance, the process can continue, it gains access back to ATM (waits first). Then it waits for semMutex, to start its transaction.

In this transaction, balance is read again since it could have changed, and the withdraw actions is performed. I also read attempt.txt again, since it could have changed too. Finally, unlock the transction with V(semMutex). semMutex semaphore is used to prevent race conditions, since balance.txt and attempt.txt are both being modified in this transaction.