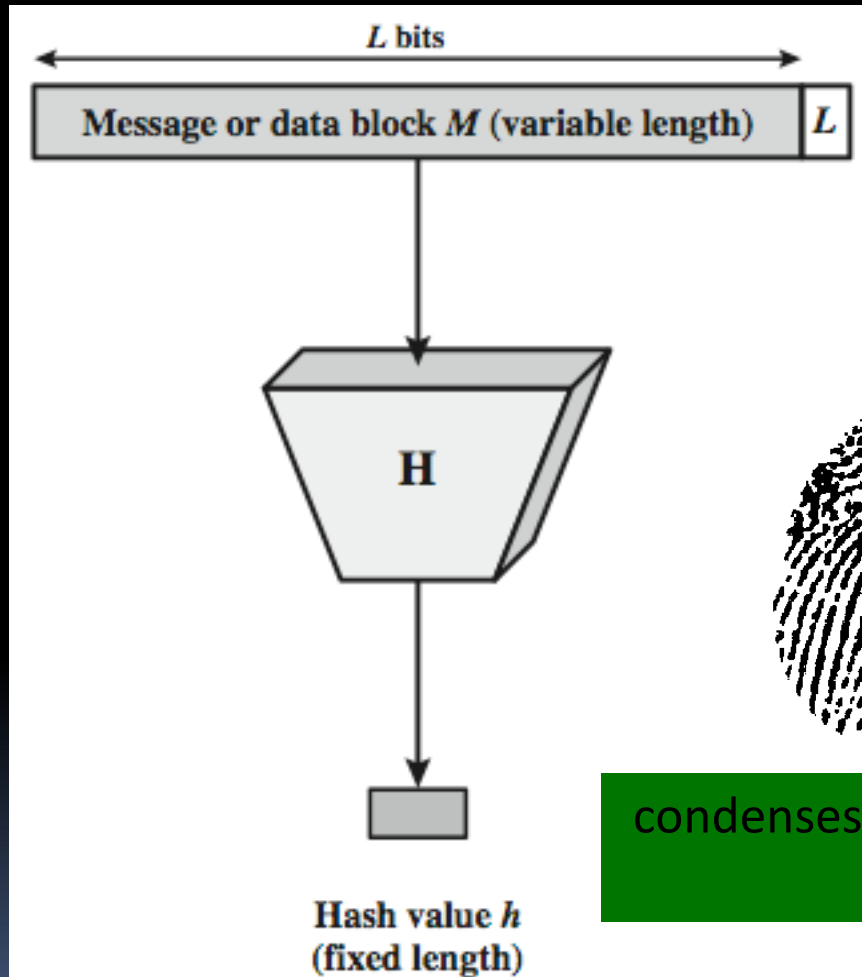




# **HASH AND MAC ALGORITHMS**

# Hash Function



- The hash value represents concisely the longer message
  - may called the *message digest*

■ A message digest is as a "digital fingerprint" of the original document

condenses arbitrary message to fixed size

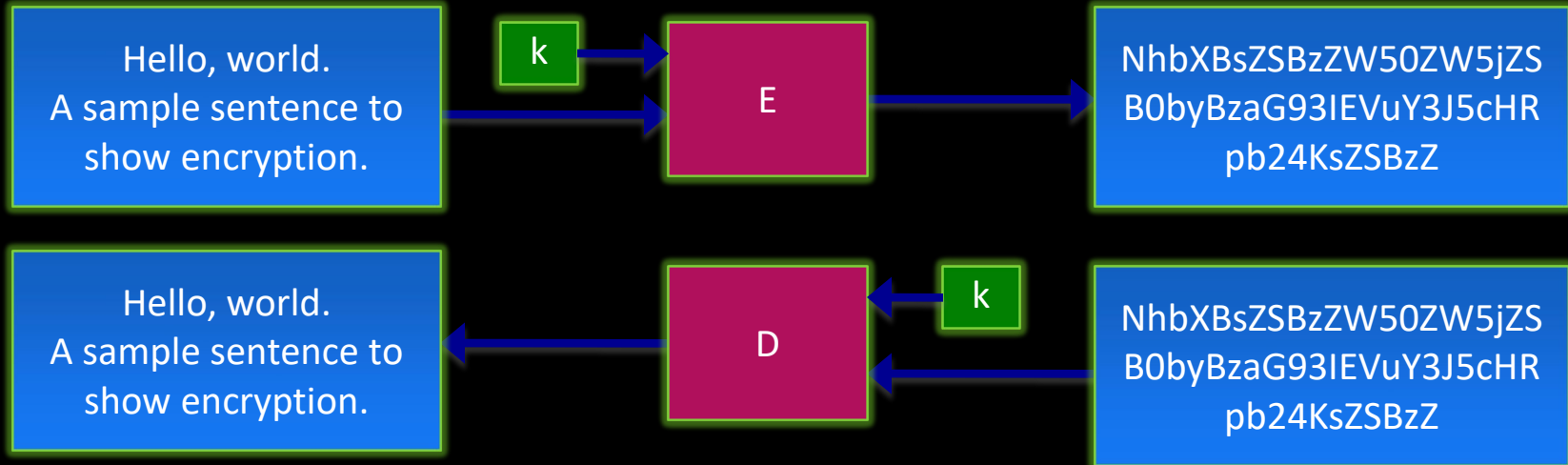
$$h = H(M)$$

# Chewing functions

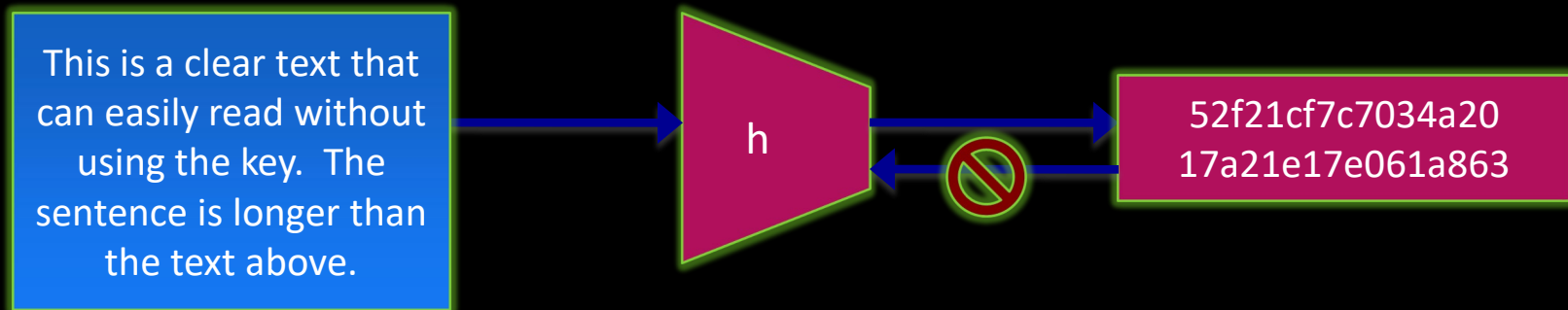
- ▶ Hashing function as “chewing” or “digest” function



# Hashing V.S. Encryption




- Encryption is two way, and requires a key to encrypt/decrypt



- Hashing is one-way. There is no 'de-hashing'



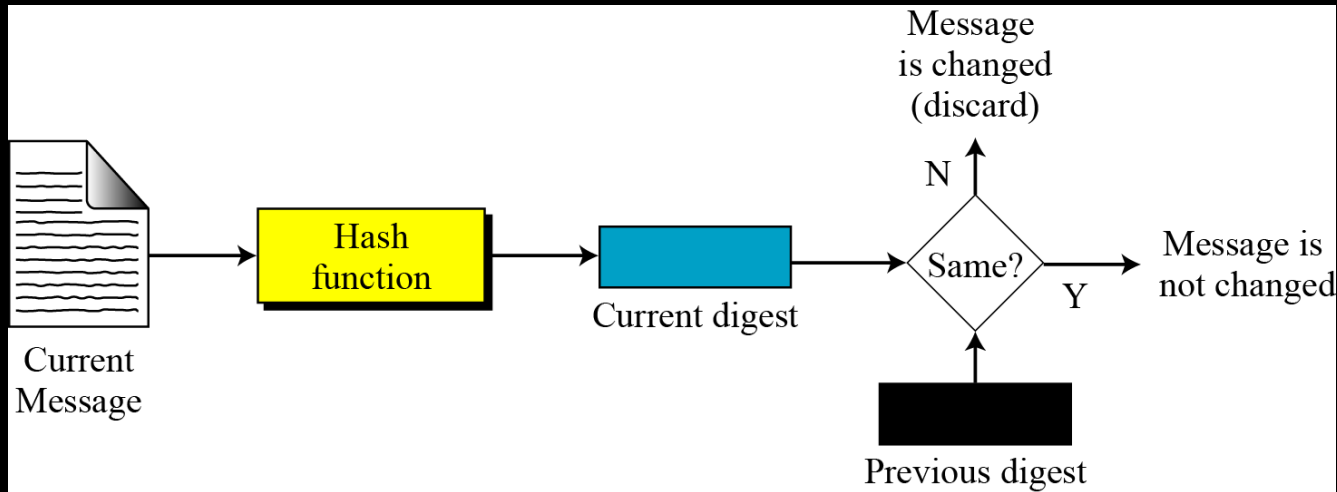
# Motivation for Hash Algorithms

- Intuition
    - Limitation on non-cryptographic checksum
    - Very possible to construct a message that matches the checksum
  - Goal
    - Design a code where the original message can not be inferred based on its checksum
    - such that an accidental or intentional change to the message will change the hash value
- 

# Hash Function Applications

- Used Alone
  - Fingerprint -- file integrity verification, public key fingerprint
  - Password storage (one-way encryption)
- Combined with encryption functions
  - Hash based Message Authentication Code (HMAC)
    - protects both a message's integrity and confidentiality
  - Digital signature
    - Ensuring Non-repudiation
    - Encrypt hash with private (signing) key and verify with public (verification) key

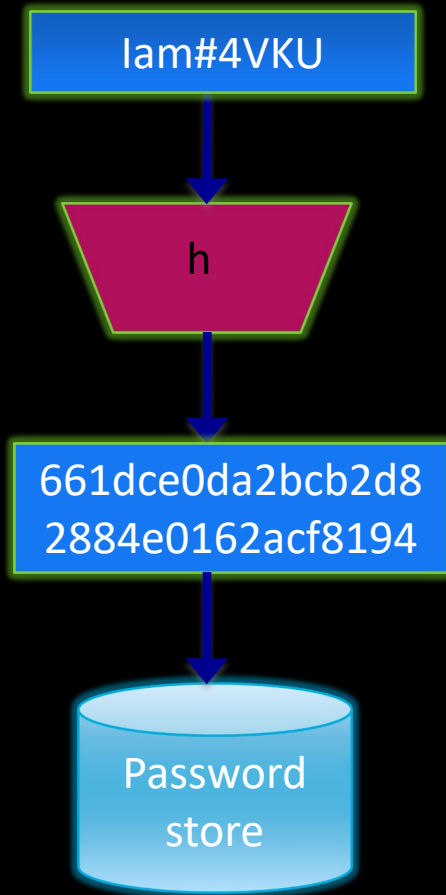
# Integrity



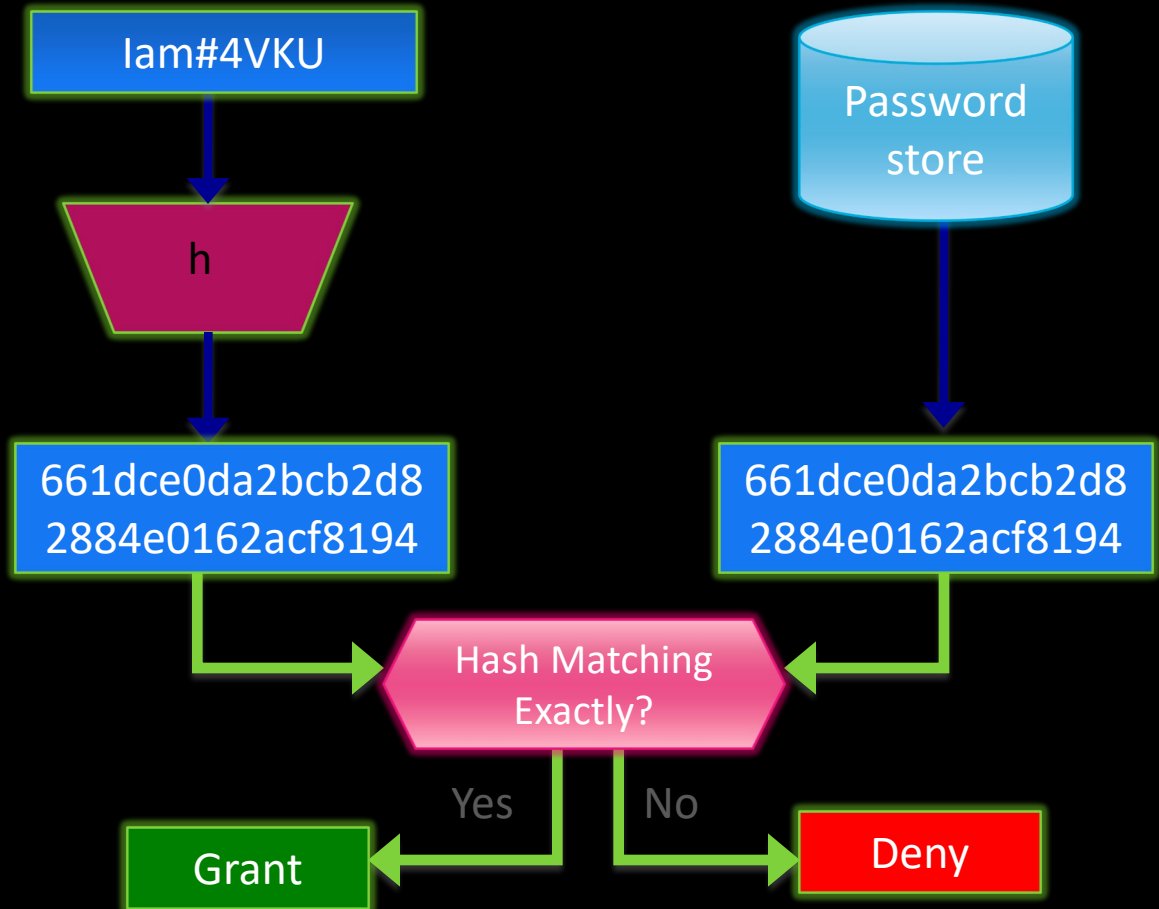
- to create a one-way password file
  - store hash of password not actual password
- for intrusion detection and virus detection
  - keep & check hash of files on system

# Password Verification

## Store Hashing Password

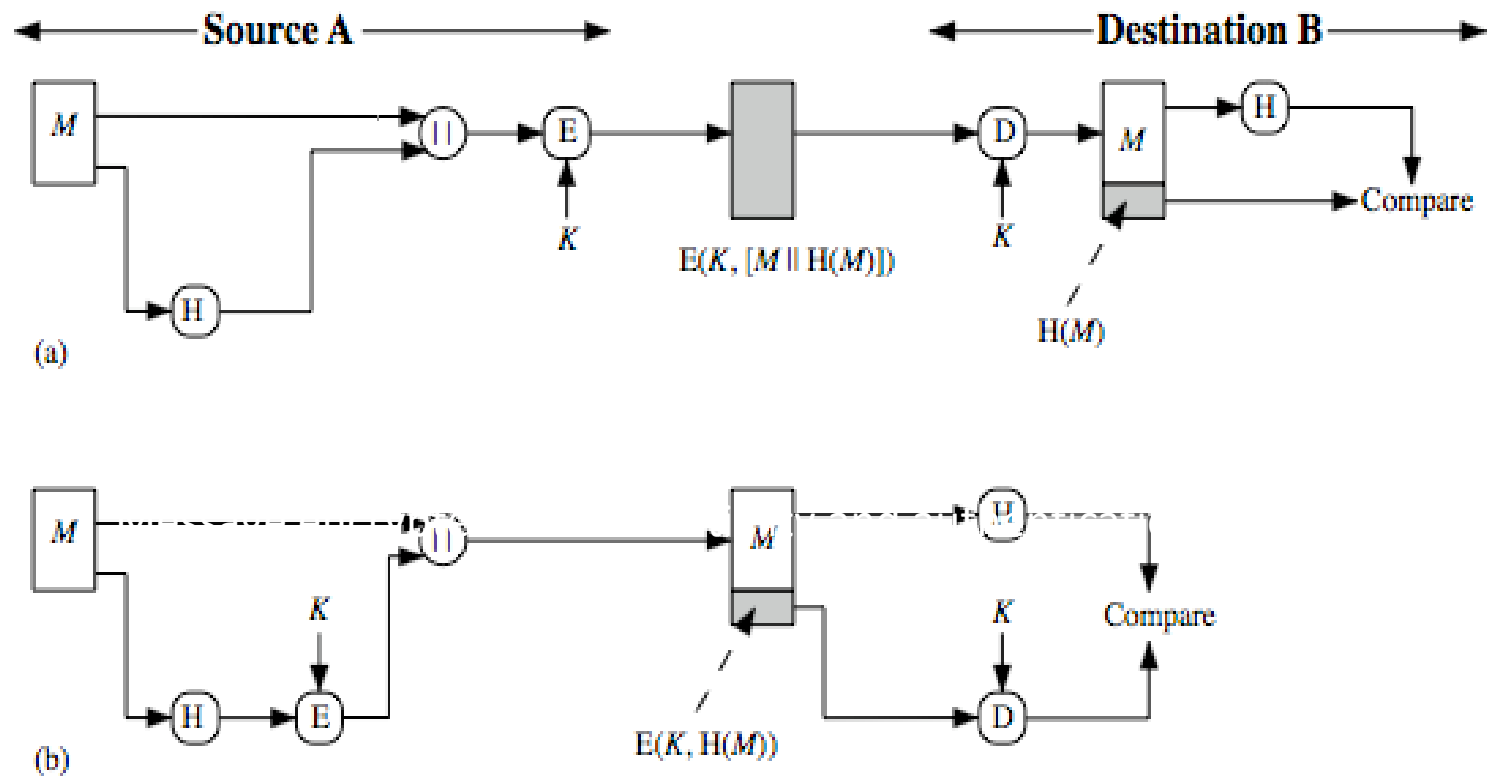


## Verification an input password against the stored hash



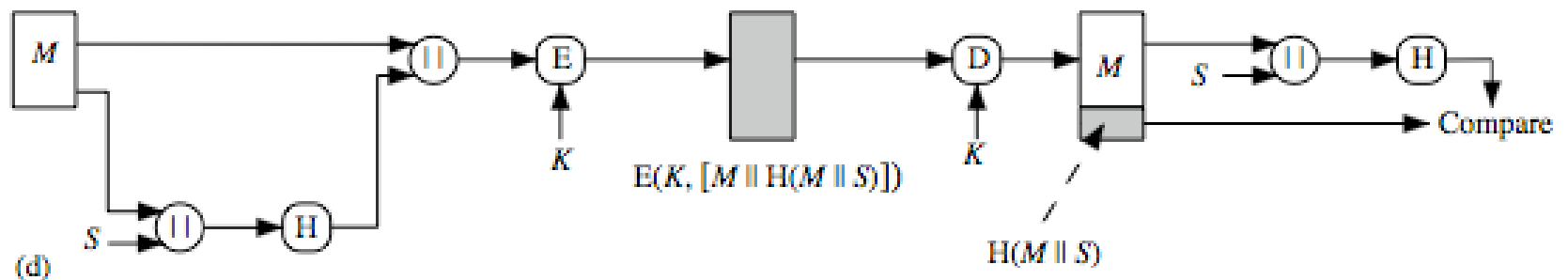
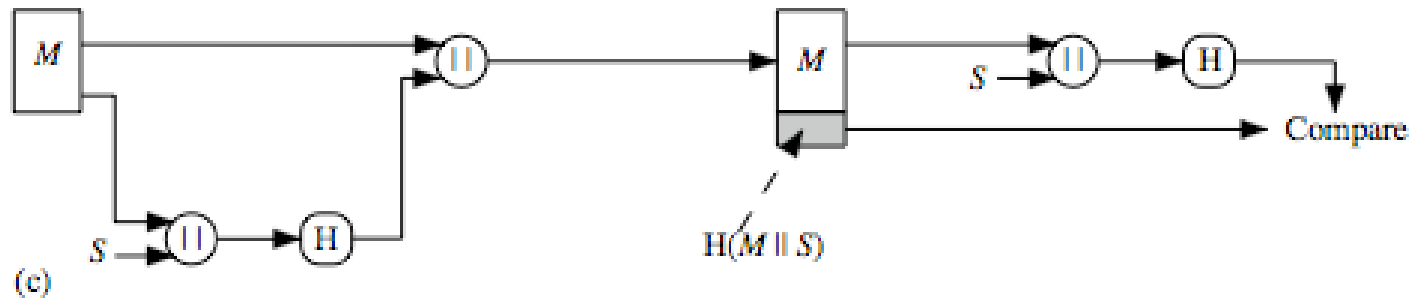


# Hash Function Usages (I)

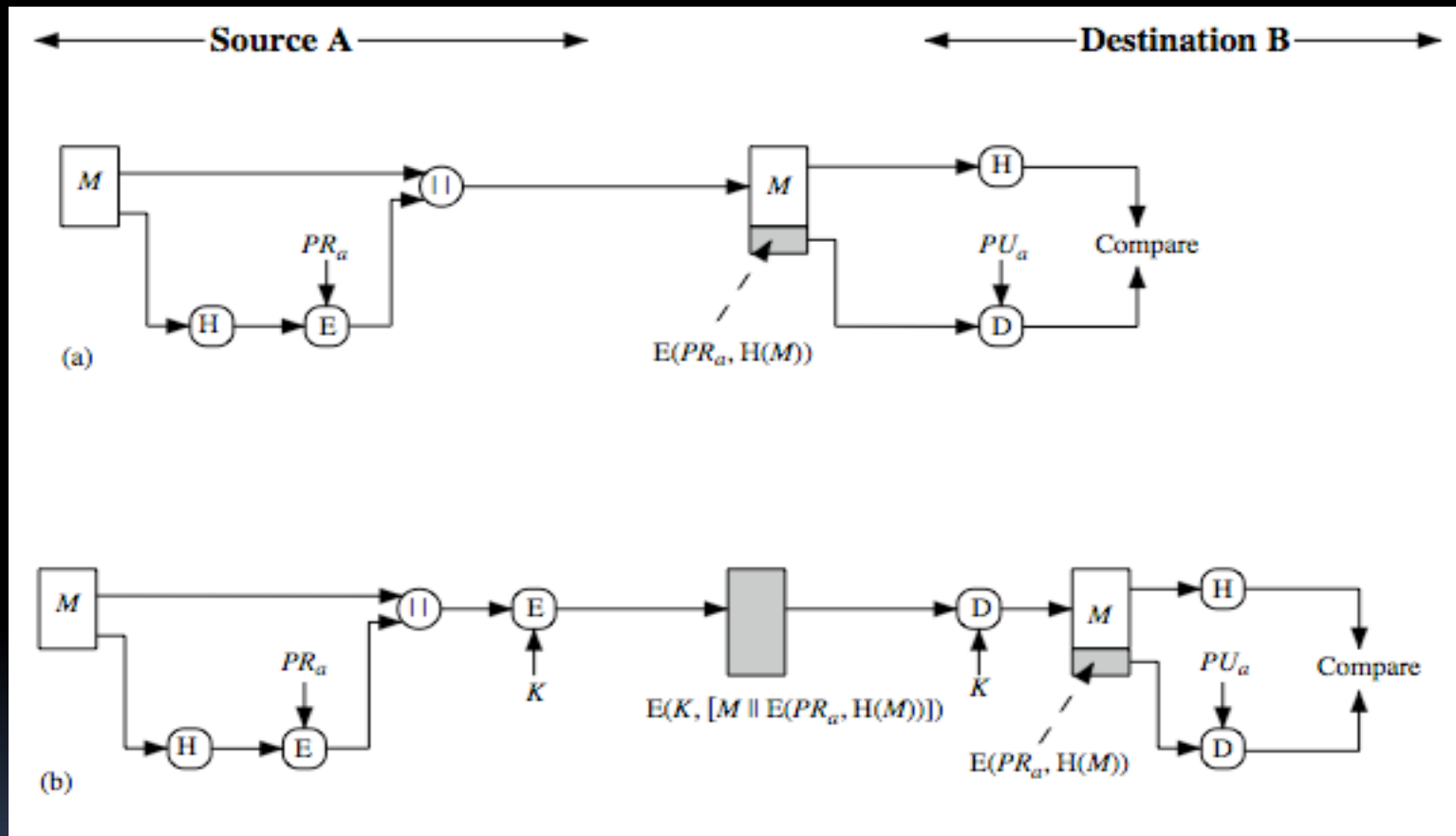


Message unencrypted: Authentication

# Hash Function Usages (II)




# Hash Function Usages (III)



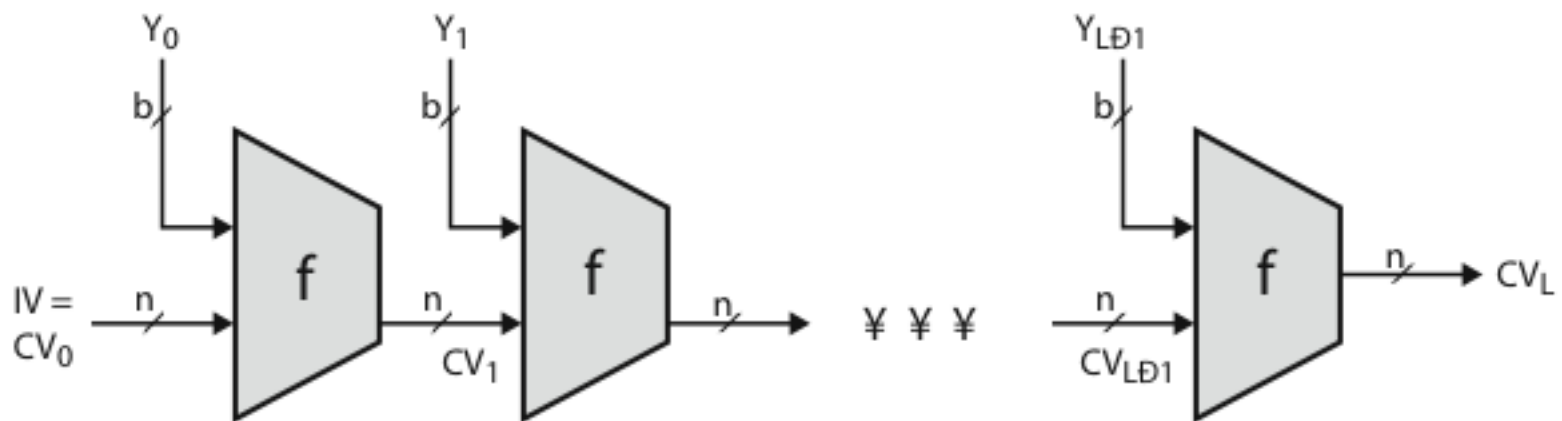
Authentication, digital signature, confidentiality



# Hash and MAC Algorithms

- Hash Functions
    - condense arbitrary size message to fixed size
    - by processing message in blocks
    - through some compression function
    - either custom or block cipher based
  - Message Authentication Code (MAC)
    - fixed sized authenticator for some message
    - to provide authentication for message
    - by using block cipher mode or hash function
- 

# Hash Algorithm Structure




$IV$  = Initial value  
 $CV_i$  = chaining variable  
 $Y_i$  =  $i$ th input block  
 $f$  = compression algorithm

$L$  = number of input blocks  
 $n$  = length of hash code  
 $b$  = length of input block




# Secure Hash Algorithm

- SHA originally designed by NIST & NSA in 1993
  - was revised in 1995 as SHA-1
  - US standard for use with DSA signature scheme
    - standard is FIPS 180-1 1995, also Internet RFC3174
    - nb. the algorithm is SHA, the standard is SHS
  - based on design of MD<sub>4</sub> with key differences
  - produces 160-bit hash values
  - recent 2005 results on security of SHA-1 have raised concerns on its use in future applications
- 




# Revised Secure Hash Standard

- NIST issued revision FIPS 180-2 in 2002
  - adds 3 additional versions of SHA
    - ▣ SHA-256, SHA-384, SHA-512
  - designed for compatibility with increased security provided by the AES cipher
  - structure & detail is similar to SHA-1
  - hence analysis should be similar
  - but security levels are rather higher
- 

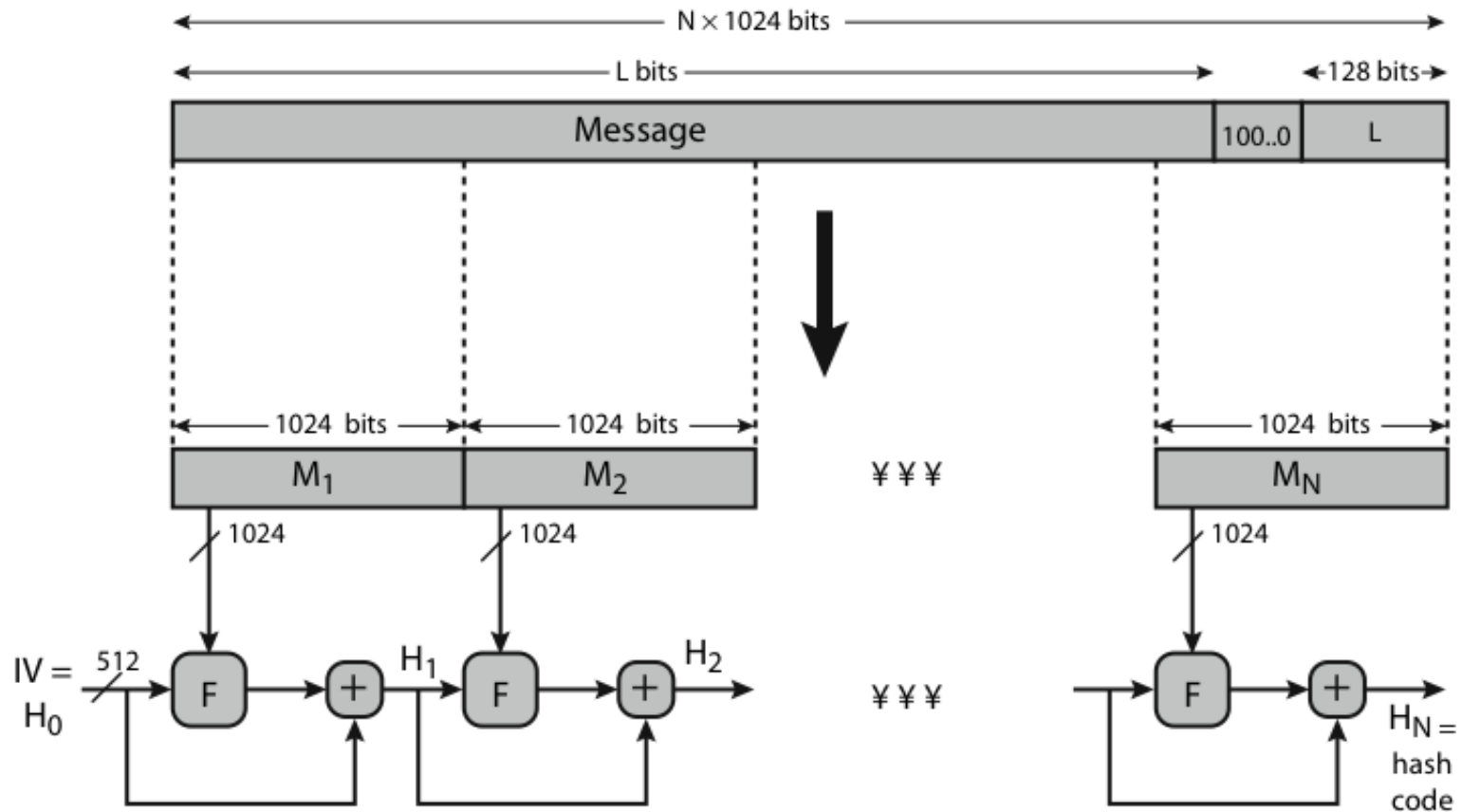


# SHA-512

- Step 1: Append padding bits
  - Step 2: Append length
  - Step 3: Initialize hash buffer
  - Step 4: Process the message in 1024-bit (128-word) blocks, which forms the heart of the algorithm
  - Step 5: Output the final state value as the resulting hash
- 



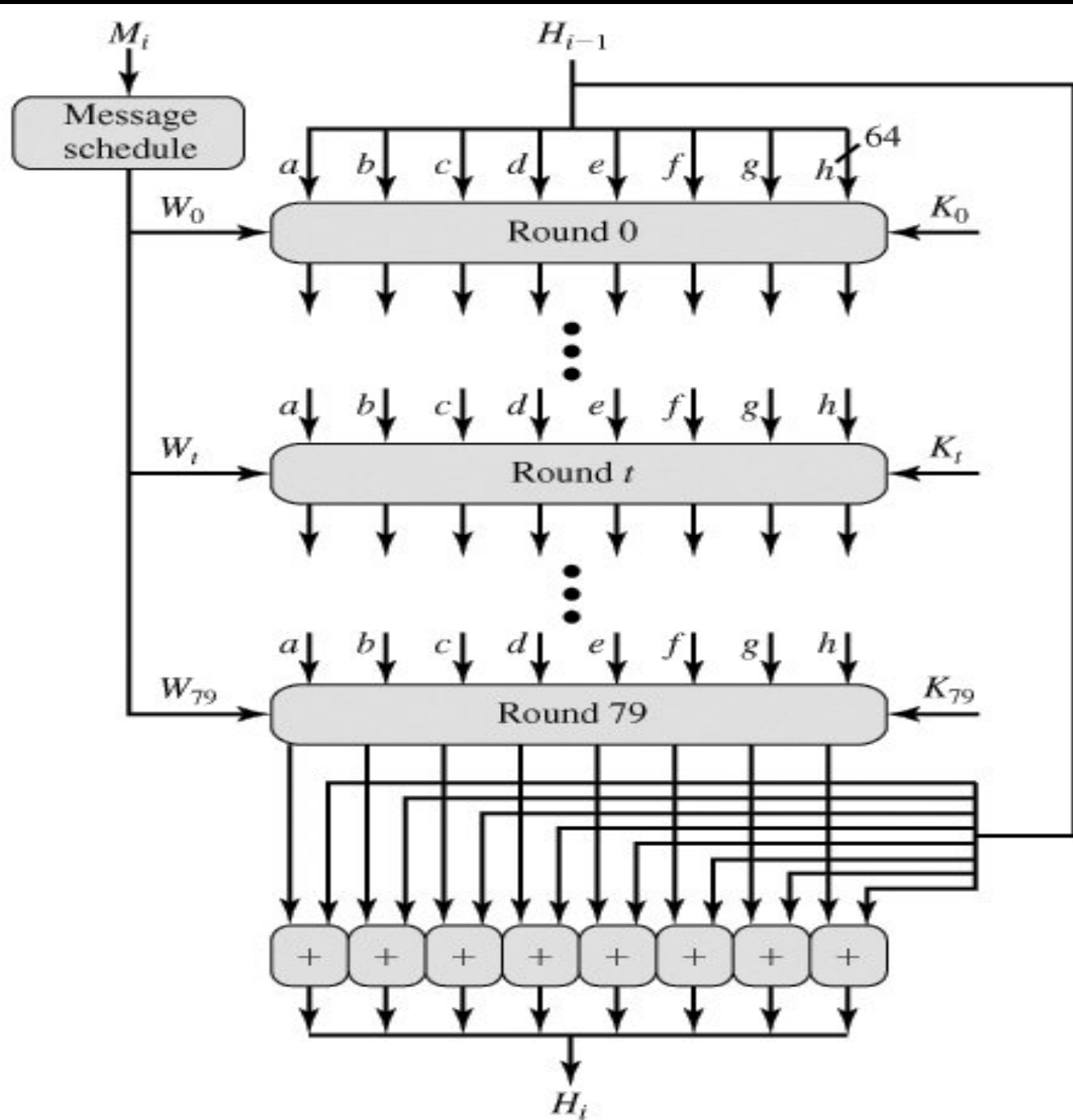
# SHA-512 Overview



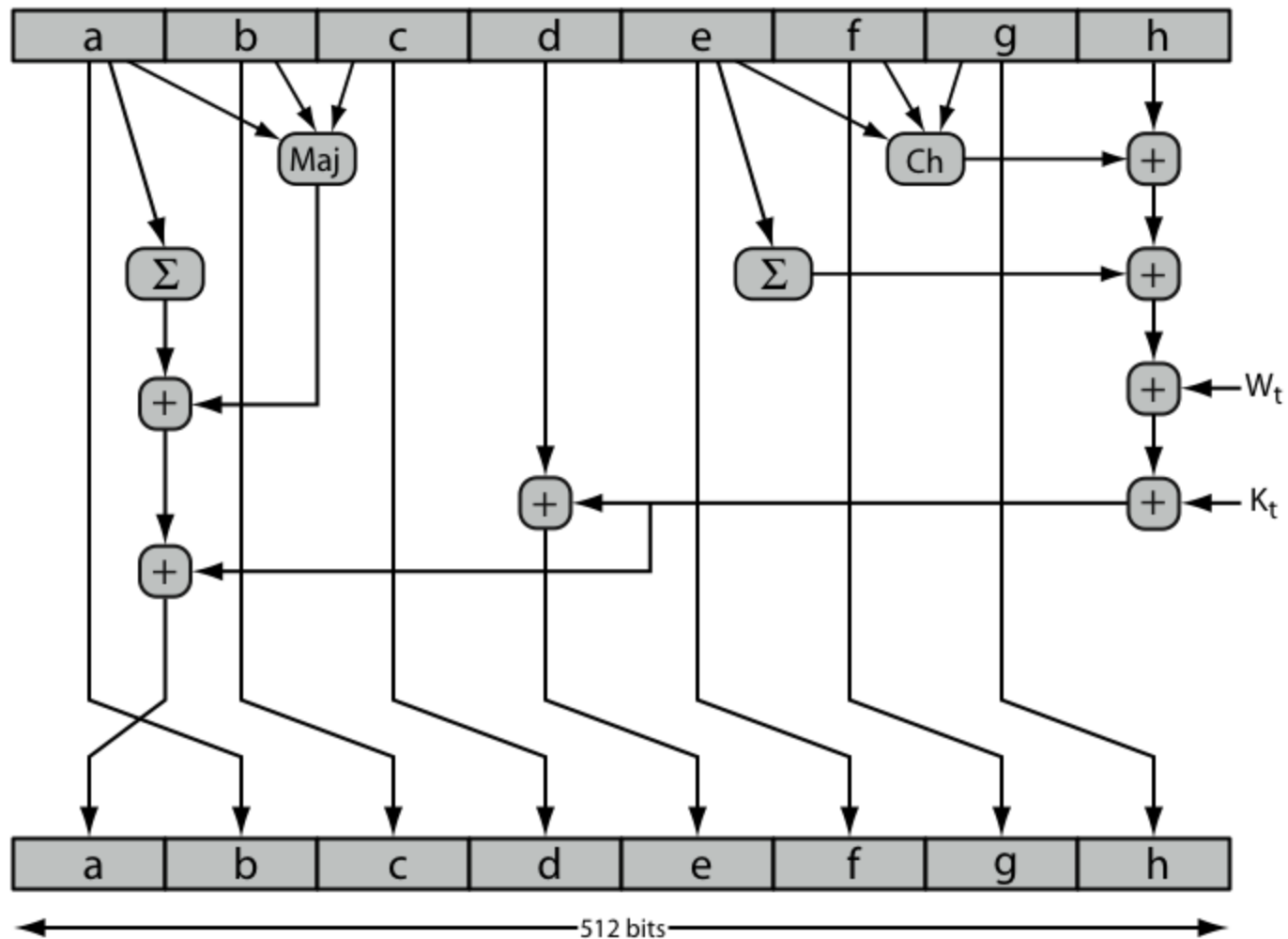
$+$  = word-by-word addition mod  $2^{64}$

# SHA-512 Compression Function

- heart of the algorithm
- processing message in 1024-bit blocks
- consists of 80 rounds
  - updating a 512-bit buffer
  - using a 64-bit value  $W_t$  derived from the current message block
  - and a round constant based on cube root of first 80 prime numbers



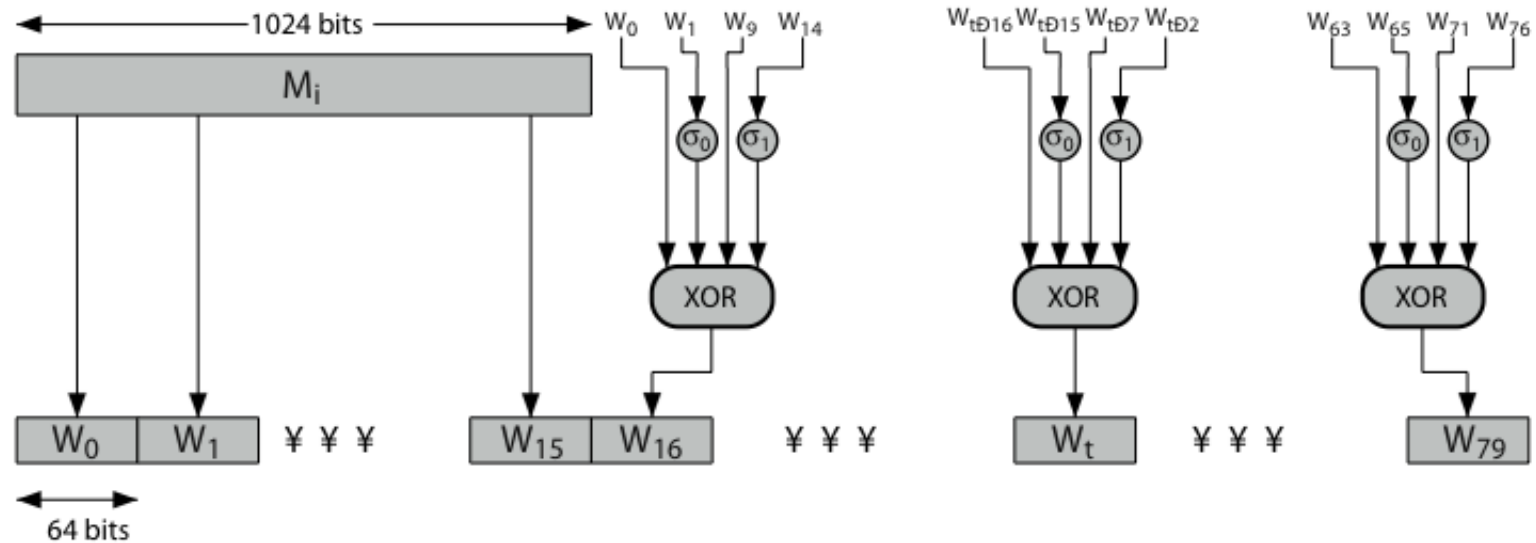
# SHA-512 Round Function



# Function Elements

- $\text{Ch}(e,f,g) = (e \text{ AND } f) \text{ XOR } (\text{NOT } e \text{ AND } g)$
- $\text{Maj}(a,b,c) = (a \text{ AND } b) \text{ XOR } (a \text{ AND } c) \text{ XOR } (b \text{ AND } c)$
- $\Sigma(a) = \text{ROTR}(a,28) \text{ XOR } \text{ROTR}(a,34) \text{ XOR } \text{ROTR}(a,39)$
- $\Sigma(e) = \text{ROTR}(e,14) \text{ XOR } \text{ROTR}(e,18) \text{ XOR } \text{ROTR}(e,41)$
- $+$  = addition modulo  $2^{64}$ 
  - $K_t$  = a 64-bit additive constant
  - $W_t$  = a 64-bit word derived from the current 512-bit input block.

# SHA-512 Round Function



# Keyed Hash Functions as MACs

- want a MAC based on a hash function
  - because hash functions are generally faster
  - code for crypto hash functions widely available
- hash includes a key along with message
- original proposal:  
$$\text{KeyedHash} = \text{Hash}(\text{Key} \parallel \text{Message})$$
  - some weaknesses were found with this
- eventually led to development of HMAC

# HMAC

- specified as Internet standard RFC2104
- uses hash function on the message:
$$\text{HMAC}_K = \text{Hash}[(K^+ \text{ XOR } \text{opad}) \parallel \text{Hash}[(K^+ \text{ XOR } \text{ipad}) \parallel M]]$$
- where  $K^+$  is the key padded out to size
- and opad, ipad are specified padding constants
- overhead is just 3 more hash calculations than the message needs alone
- any hash function can be used
  - eg. MD5, SHA-1, RIPEMD-160, Whirlpool



# HMAC Overview

