



Theory of Computation

**Regular Expression**

**Lecture 7**

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Operators

$$\Sigma = \{a, b\}$$

\* : Kleene Star

+ : Kleene +

$$\Sigma^* = (a, b)^*$$

$$\left\{ \underbrace{\epsilon}_{0}, \underbrace{a, b}_{1}, \underbrace{aa, ab, ba}_{2}, \dots \right\}$$

$$\Sigma^+ = (a, b)^+$$

$$\{a, b, aa, ab, ba, bb, aaa, \dots\}$$

Unary

+ : OR

$$L_1 : \{ab, ba\}, L_2 : \{ba, bb\}$$

$$L_1 + L_2 : \{ab, ba, bb\}$$

• : Concatenation

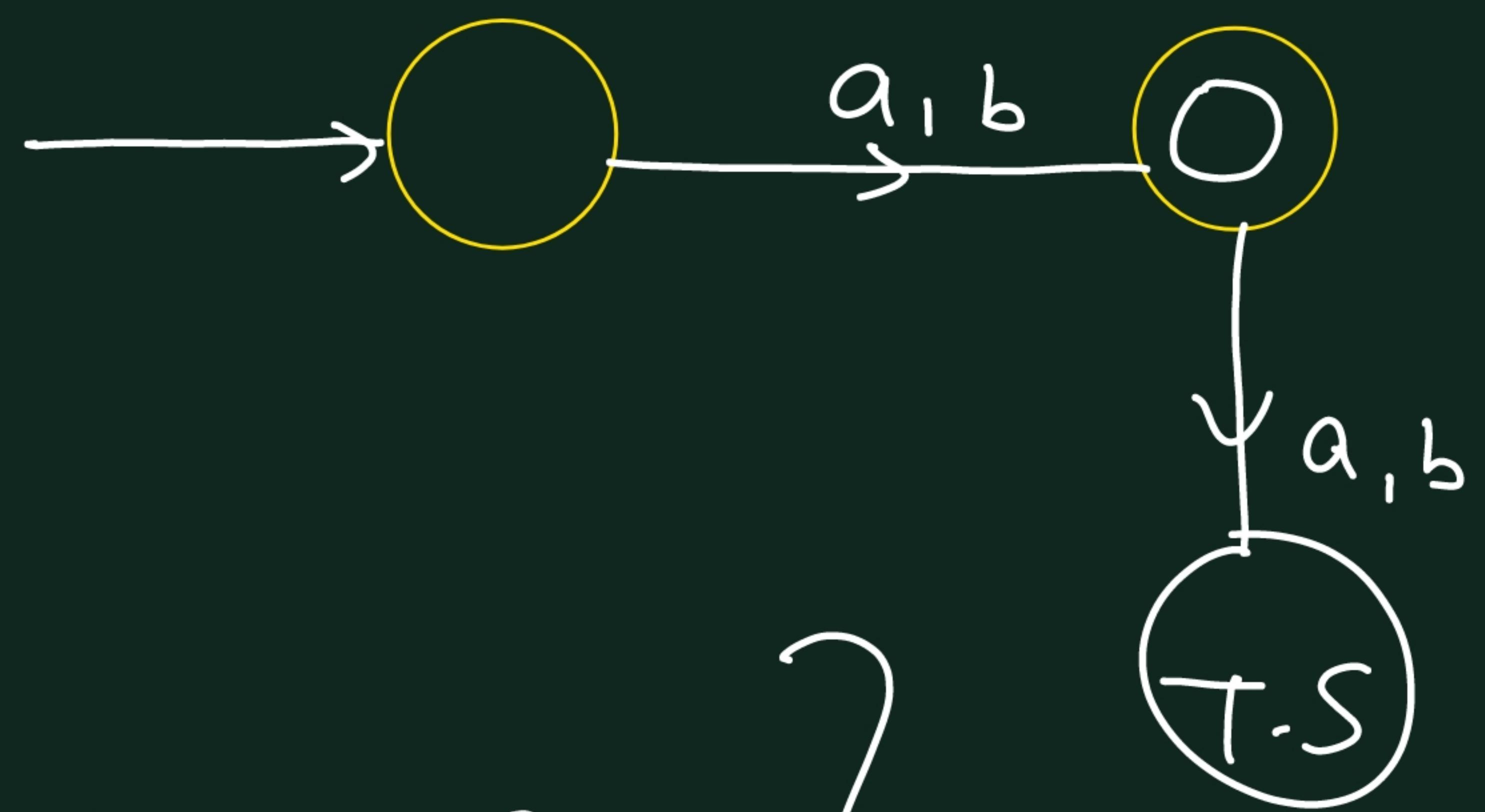
$$ab \cdot b = abb$$

$$a \cdot (b + a) = \{ab + aa\}$$

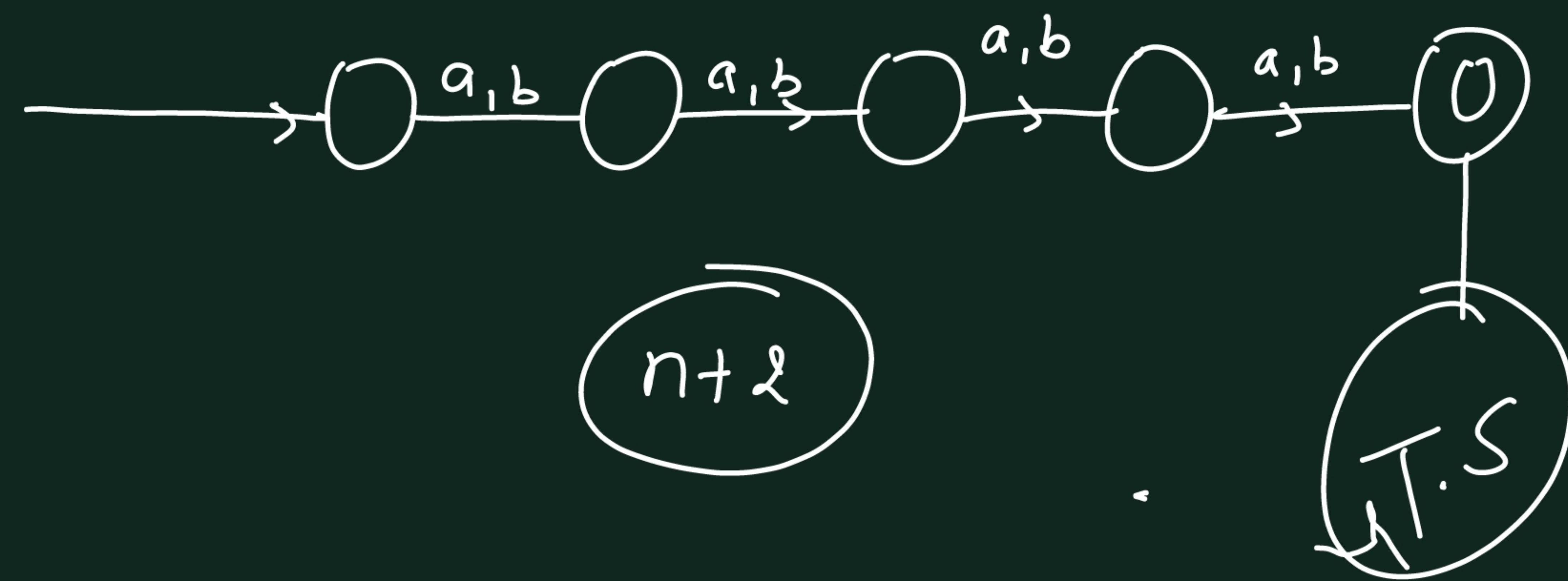
Binary

$$\Sigma = \{a, b\}$$

L 1: String length = 4



$$L_2 = \frac{a^3 b^3}{\{}} \quad | \quad a^{111} b^{111} \quad | \quad a^n b^n$$



$$Q_n \left\{ ab + ba \right\}$$

# Regex : Node JS

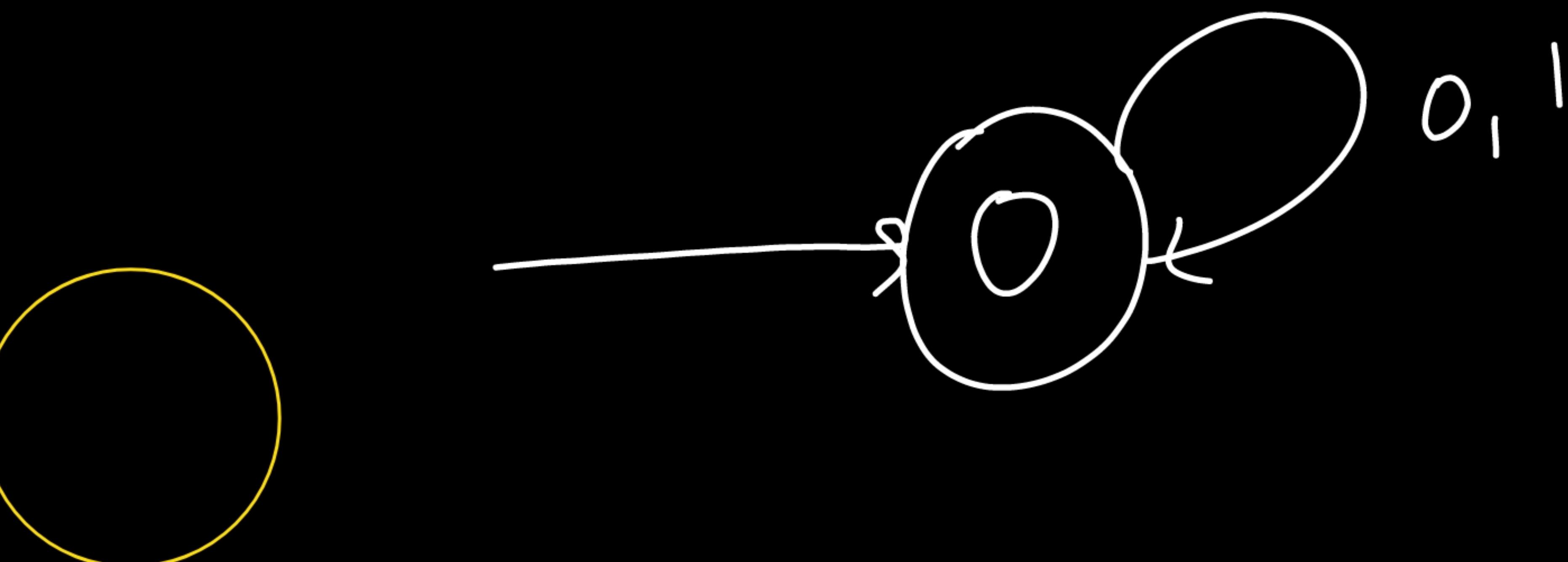
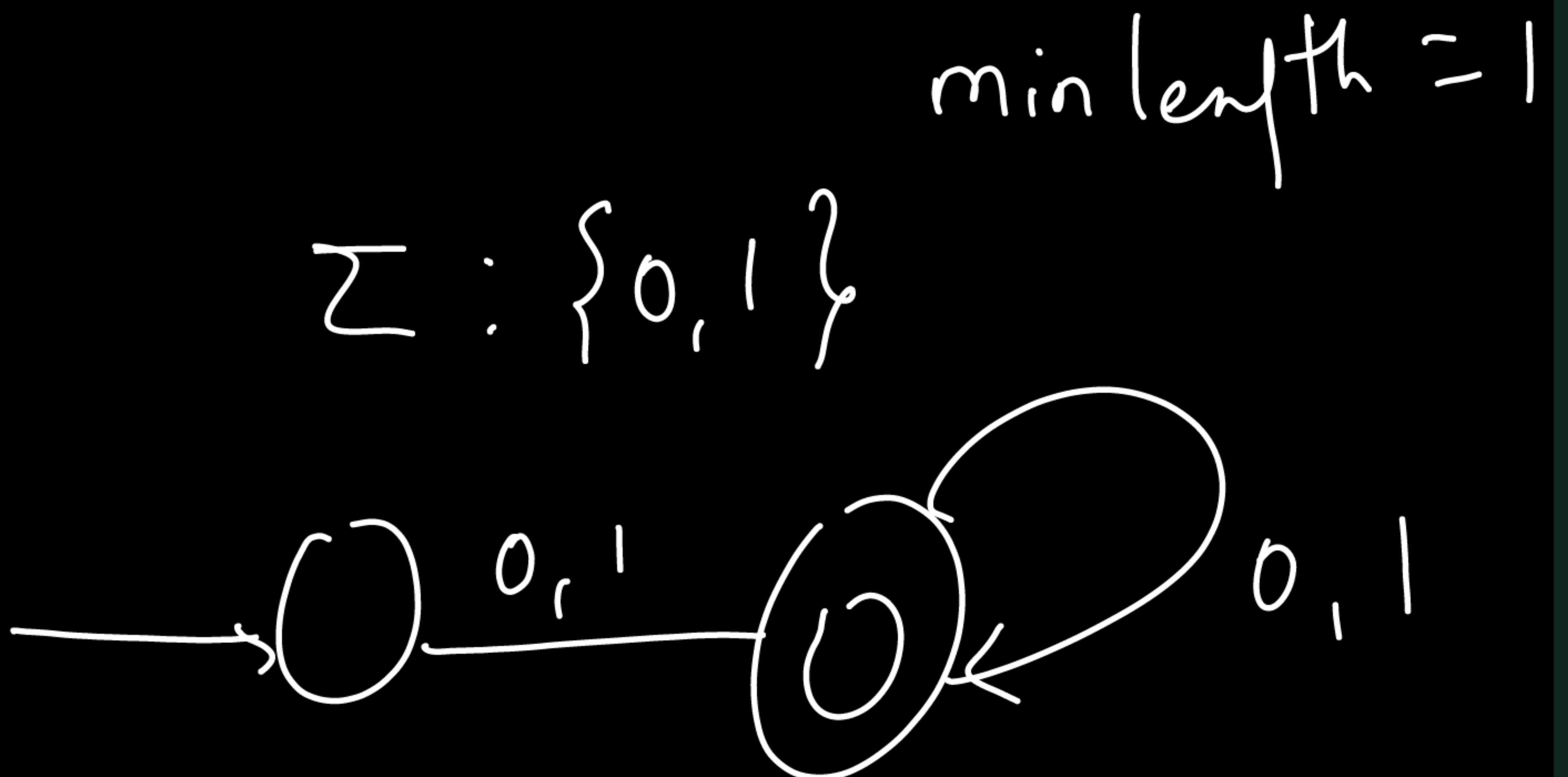
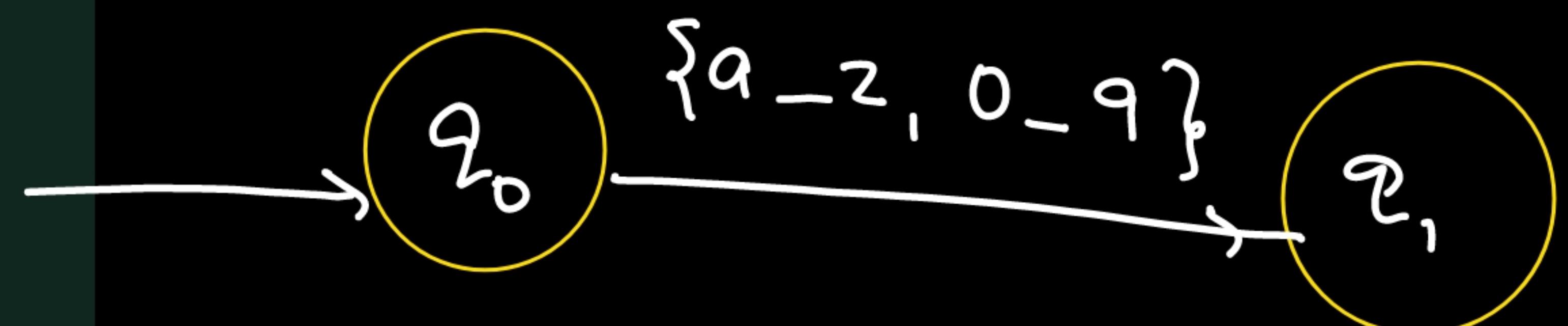
## Password Rule:

Password must:

1. Contain **only lowercase letters and digits**
2. Have **minimum length = 8**

Alphabet:

$$\Sigma = \{ a-z, 0-9 \}$$



| State | valid char [a-z0-9] | invalid |
|-------|---------------------|---------|
| → q0  | q1                  | dead    |
| q1    | q2                  | dead    |
| q2    | q3                  | dead    |
| q3    | q4                  | dead    |
| q4    | q5                  | dead    |
| q5    | q6                  | dead    |
| q6    | q7                  | dead    |
| q7    | q8                  | dead    |
| *q8   | q8                  | dead    |
| dead  | dead                | dead    |

} Real  
Scenario

## Symbols Reminder

- + → Union / OR
- . → Concatenation
- $\epsilon$  → Empty string
- $\emptyset$  → Empty language
- R, S, T → Regular expressions

## Properties of + (OR / Union)

### Identity

Identity element does **not change** the expression

$$R + \emptyset = R$$

.

Explanation:

Union with an empty language gives the same language.

$$\boxed{L_1 = L_2}$$

$$L_1 = \{aa, ab\} \cup \{ba, bb\}$$

### Commutative

Order does **not matter**

$$R + S = S + R$$

$$L_2 = \{ba, bb\} \cup \{aa, ab\}$$

Example:

$$a + b = b + a$$

## Associative

Grouping does **not** matter

$$(R+S)+T=R+(S+T)$$



Example:

$$(a + b) + c = a + (b + c)$$

$$\left\{ \begin{array}{l} R \rightarrow L \\ L \rightarrow R \end{array} \right\}$$

## Idempotent

Repeating does not change result

$$R+R=R$$

Example:

$$a + a = a$$

$$\left\{ ab, ba \right\} \vee \left\{ ab, ba \right\}$$

Proj.

$$\left\{ \begin{array}{c} a+b+c \\ \hline \Downarrow \\ ALU \end{array} \right\}$$

### Dominator (Absorbing Element)

$\Sigma^*$  dominates union

$$R + \Sigma^* = \Sigma^*$$

Because  $\Sigma^*$  already contains all strings.

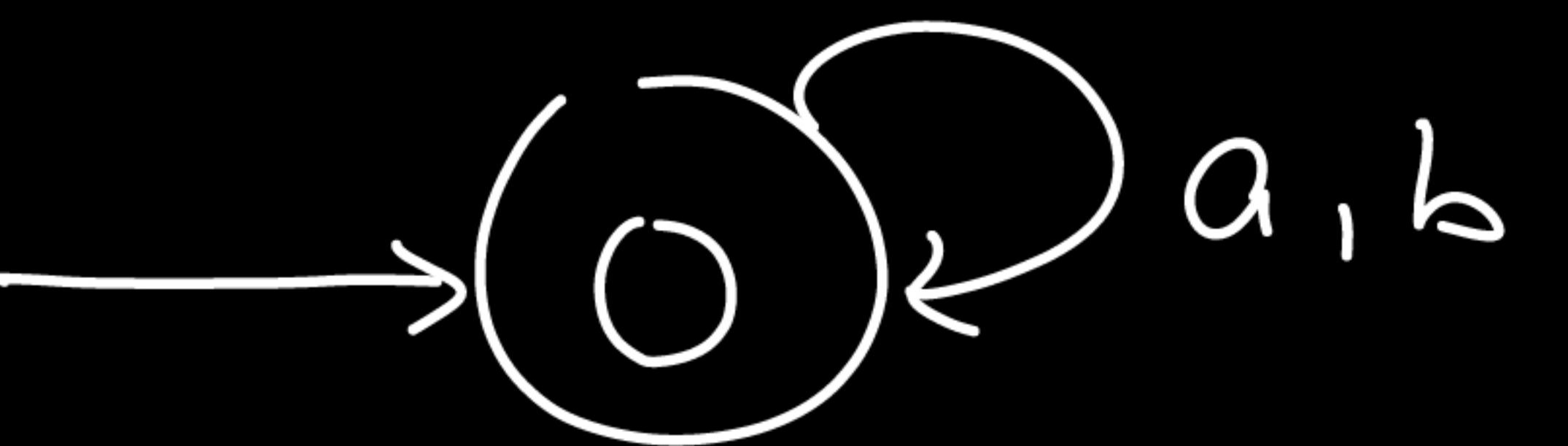
$$\Sigma = \{a, b\}$$

$$\Sigma^* = \{\epsilon, a, b, aa\dots\}$$

### Annihilator

NO annihilator exists for +

Because nothing makes union empty except both being empty.



$$L_1 \cup \{ \} = L_1$$

$\uparrow$

## Properties of . (Concatenation)

### Identity

Identity element keeps expression same

$$R.\epsilon = \epsilon.R = R$$

$$\{ab\}.\in = ab$$

Explanation:

Concatenating with empty string does nothing.

### Associative

Grouping does **not matter**

$$(R.S).T = R.(S.T)$$

Example:

$$(ab)c = a(bc)$$



## Commutative (NOT TRUE)

Order **matters**

$R.S \neq S.R$

Example:

$ab \neq ba$



$$\begin{array}{c} a.b \neq b.a \\ \hline a+b = b+a \end{array}$$

$$\left\{ \begin{array}{l} R.L : \text{Reg Exp} : R.G : \text{DFA} \\ \text{NF.} \end{array} \right.$$

## Annihilator

Element that destroys everything

$R.\emptyset = \emptyset.R = \emptyset$



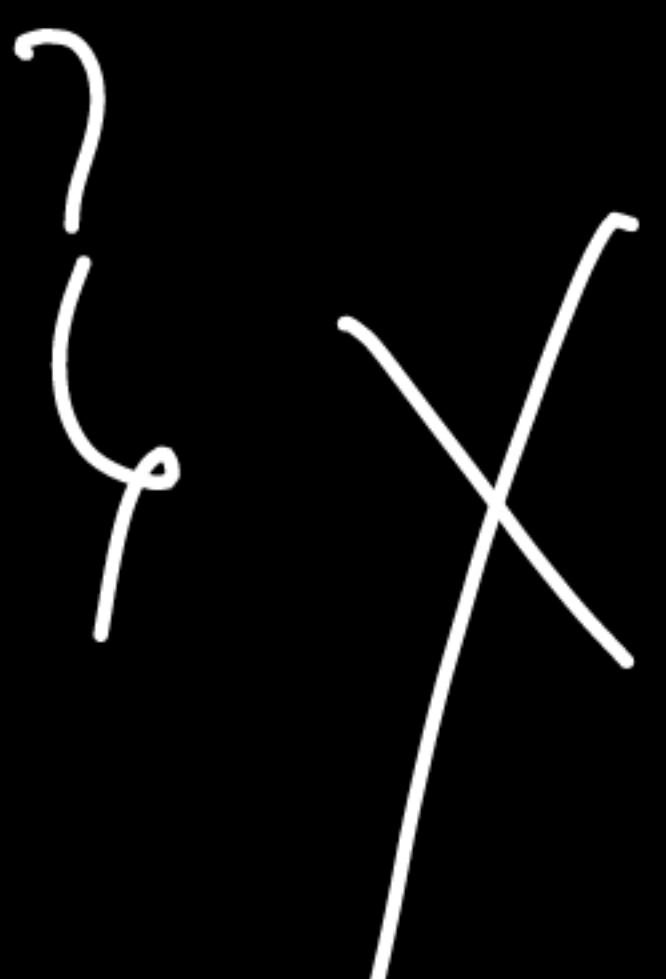
Explanation:

No strings available to concatenate.

$$\underline{a \cdot \Sigma^*}$$

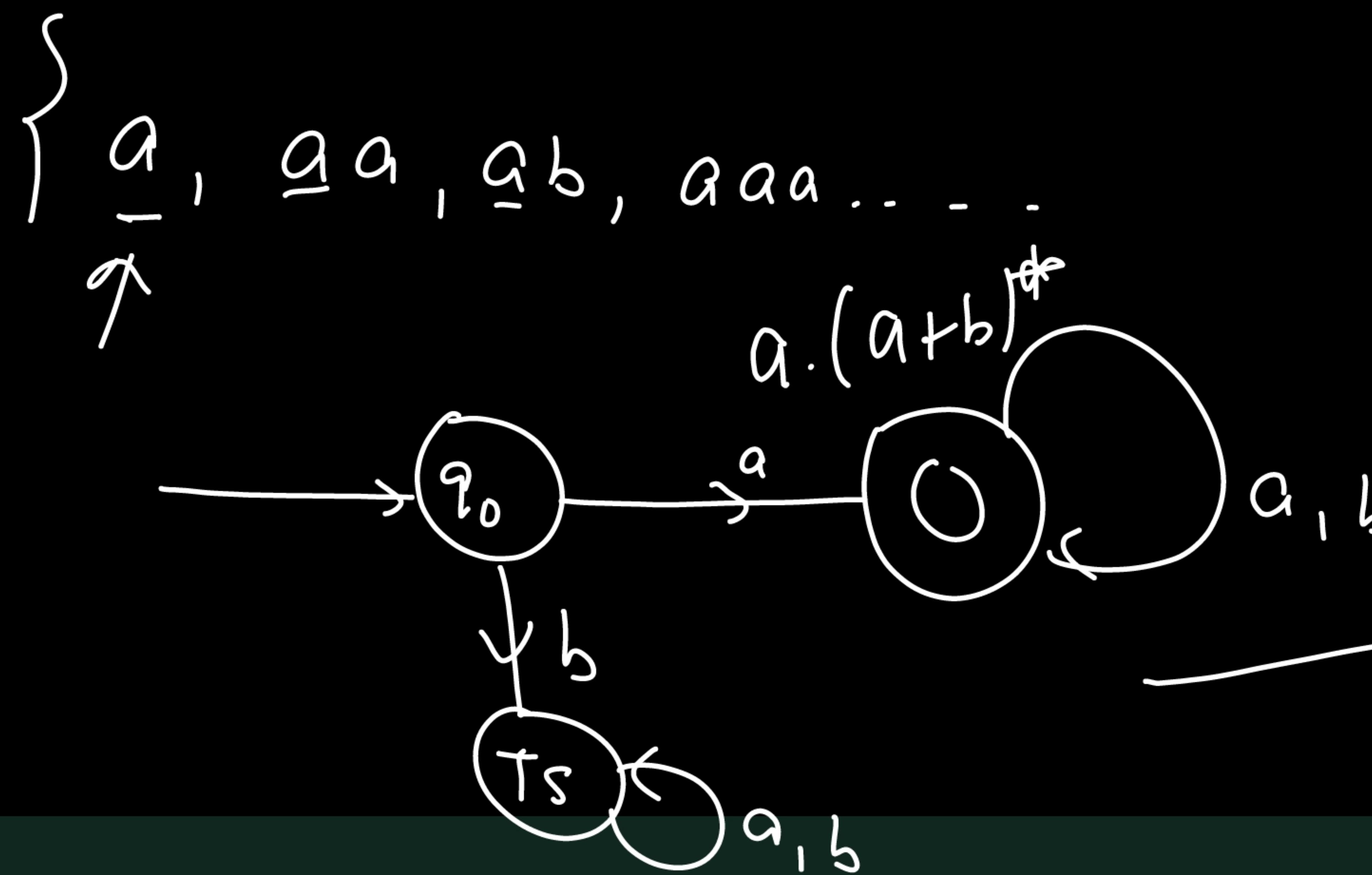
## Dominator

No dominator for concatenation



RegEx:  $a(a+b)^*$

$a \cdot \{ \in, a, b, aa, bb, ab, ba, \dots \}$



RegEx



Thank You

