

ded 2

# Min<sup>m</sup> cost climbing stairs

by recursion

$$f(n) \Rightarrow f(n+1) + f(n+2)$$

or

$$f(n) \Rightarrow f(n-1) + f(n-2)$$

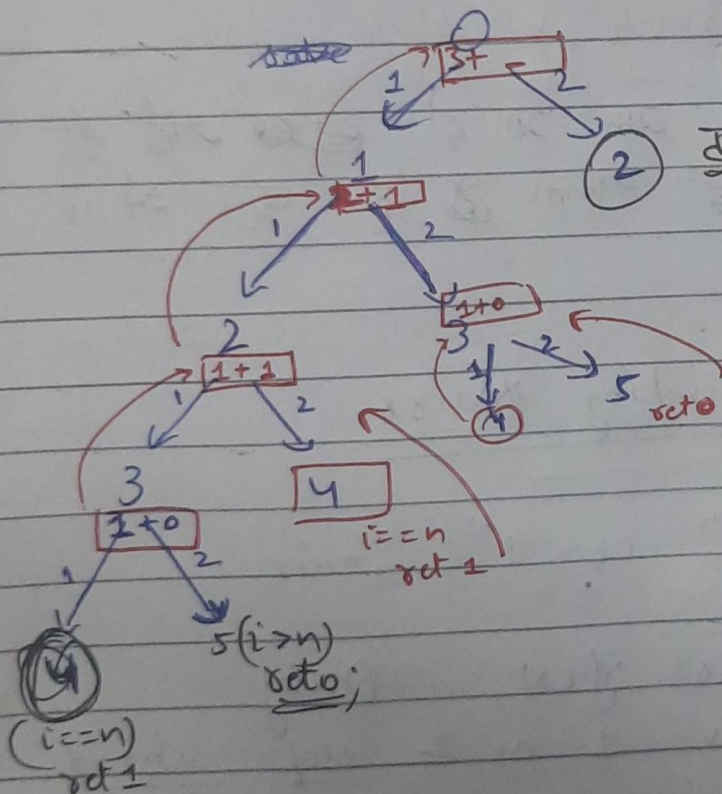
dry run

n=4

if (n==1) ret 1;

if (n<1) ret 0;

ret solve(n,i+1) +  
solve(n,i+2);



इसका answer तो ही पता है  
निकाल चुका है तो  
भाई फिर DP लीजो ॥

"Overlapping sub-problems"



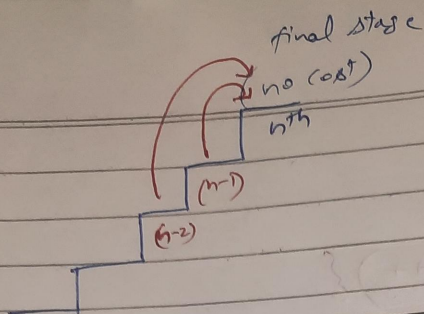
Q. If we write  $f(n)$   ~~$f(k)$~~ ; then  
it means <sup>any</sup> function should total possible  
way to go from  $k^{\text{th}}$  stair to  $n^{\text{th}}$  stairs  
by 1 or 2 step.

If need to find no. of way to  
go from  $0^{\text{th}}$  stair  $n^{\text{th}}$  stair  
तो मैं इस problem को एक subproblem  
में तोड़ सकता हूँ ॥

→ So can say, it will be equal to no. of  
way to reach ~~from~~  $n^{\text{th}}$  stair from  $1^{\text{st}}$  stair  
plus ways to reach  $n^{\text{th}}$  stairs from  
 $2^{\text{nd}}$  stair. (क्योंकि एक या दो ~~steps~~ ~~steps~~ से  
चढ़ सकता है वहाँ से ॥

leetcode: Min Cost Climbing Stairs

- ① Can start from  $0^{\text{th}}$  or  $1^{\text{st}}$  stair
- ② have to pay the cost first, only  
then we can move 1 or 2 step ahead.
- ③ Return  $\min^{\text{m}}$  cost to reach ~~top~~  
last top stage



$$f(n) \Rightarrow \text{cost}(n) + \min(f(n-1) + f(n-2))$$

Recursion

```
int solve(vector<int> & cost, int n){
```

```
    if (n==0)
        return cost[0];
```

```
    if (n==1)
        return cost[1];
```

```
    int ans = cost[n] + min(solve(cost, n-1) + solve(cost, n-2));
```

```
    return ans; }
```

```
int minCost (vector<int> & cost){
```

```
    int ans = min(solve(cost, n-1), solve(cost, n-2));
```

```
}
```

Bottom up (dp)

```
int solve3 (vector<int> & cost, int n) {
    vector<int> dp(n+1); // dp[0]
```



$dp[0] = cost[0];$

$dp[1] = cost[1];$

for ( $i=2; i < n; i++$ ) {

$dp[i] = cost[i] + \min(dp[n-1], dp[n-2]);$   
}

return  $\min(dp[n-1], dp[n-2]);$

}

Space optimization - 9th वीजारा !!

★  $E^2$  value; 3rd previous last two values 42  
depend करती है ॥