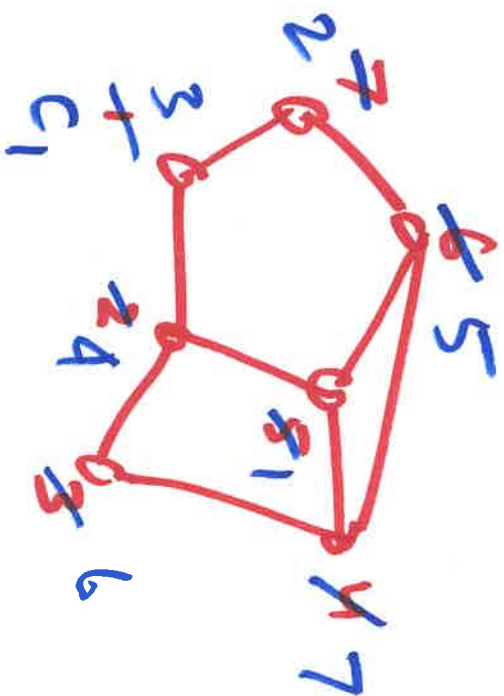


# Set 7: Greedy Algorithms



5  
1 2 3 4 5 6 7  
 $c_1 \neq c_1$   $c_2$   $\uparrow$   $q_1$   $q_2$   $q_3$   $c_6$

Bad  $\frac{2^n / n^n / n!}{\text{high}} \cdot \text{exp.}$  Polynomial  $\Rightarrow O(n^2)$  ~~high~~

Good low Polynomial  $\Rightarrow O(n)$  1, n

$n / n^2 / n^{10} / n^{100} \Rightarrow O(n^n) \sim O(2^n)$

$1 + 2 + \dots + (n-1) \Rightarrow O(n^2)$   
 Alg <  $\frac{\text{quality}}{\text{cost} \sim \text{Time/Space}}$  < high

1. 9 - 10
- 2 15 - 16
- 3 21 - 23
- 4 7 - 12
- 5 6 - 8
- 6 11 - 18
- 7 17 - 22
- 8 13 - 14
- 9 19 - 20

S

T

- |   |   |   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|---|---|
| 3 | 6 | 9 | 2 | 1 | 4 | 7 | 5 | 8 |
| 3 | 5 | 9 | 3 | 1 | 6 | 8 | 4 | 7 |

$\{5, 1, 6, 9, 3\}$   


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 $\{4, 8, 2, 7\}$

{ 4, 7, 8, 9 }

Max Ind Sub.

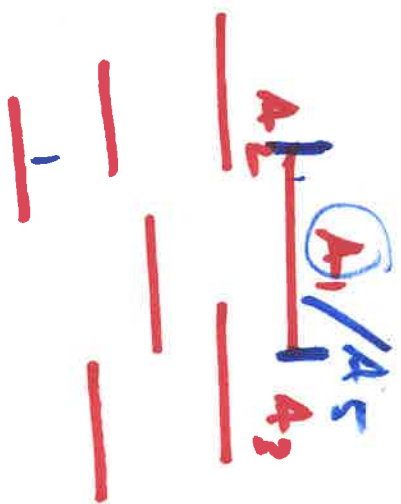
{ 1, 7, 8, 9 }

1 7 8 9

7 8 9

1

1



$$A = \{A_1, A_2, \dots, A_n\}$$

$$I_1, I_2 \dots I_n$$

$$A' \subseteq A$$

↖ Largest Ind. Set  
↖ Largest Set of Compatible Activities

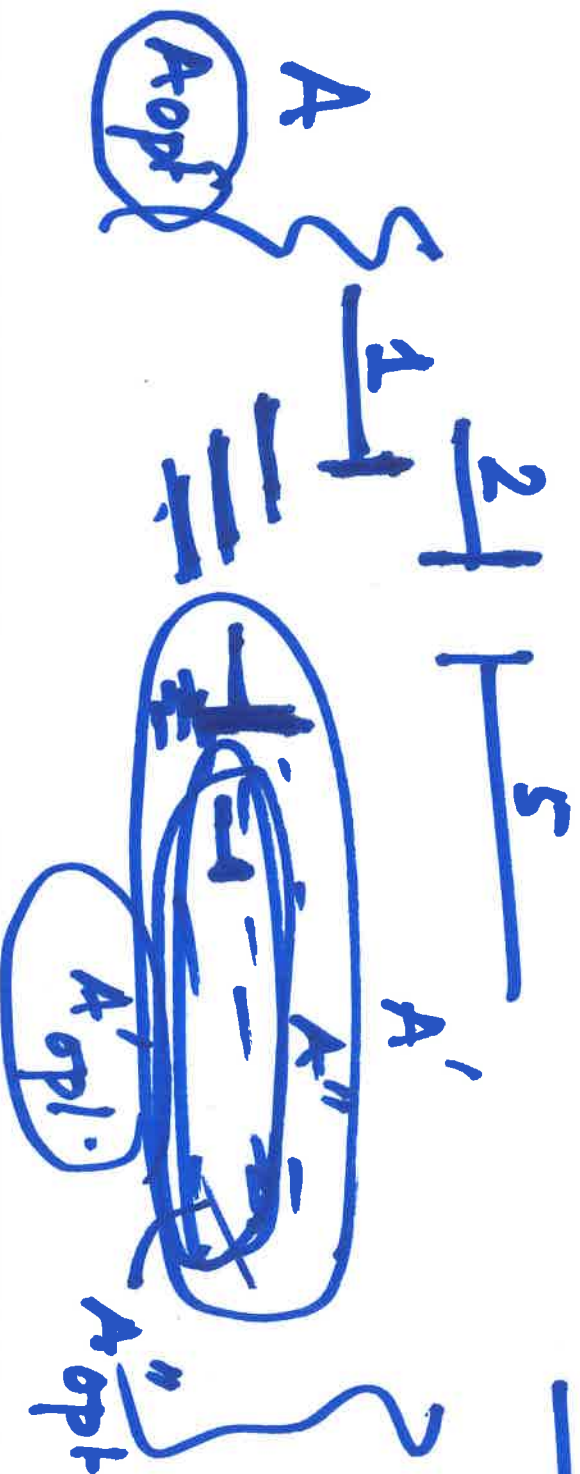
$$A' = \{$$

$a_1$

$a_2$

OPT  $\Rightarrow \{2, 5, 7, 9\}$

$\{ \textcircled{1}, 5, 7, 9 \}$   $\Leftarrow$  Valid OPT



$$\frac{LIS}{OPT} \rightarrow A' \subseteq A$$





$$1.c_1 + 1.c_2 + \dots + 1.c_9$$

Σ

$$6c_1 + 2c_2 + c_3$$

$$5c_1 + 4c_2$$

$$c_1 \leq c_2 \leq c_3 \leq \dots$$

$$c_1 < c_2 < c_3 \dots$$

$$c_1 + c_2 + c_3 + \dots + c_9$$

$$c_1 + c_1 + c_1 + c_1 + c_1$$

$$+ c_2 + c_2 + c_3$$

$$9c_1 + 0.c_2 + 0.c_3$$



$S_1: \underline{6C_1 + 2C_2 + C_3} \rightarrow \text{Max } \underline{\text{End Sol}} \underline{\text{Alg 1}}$   
based Soln Proc

$5C_1 + 4C_2 \rightarrow \text{Min \# Proc.}$

based Soln Alg 2

Chomoko New base

Soln

Least Cost Soln.

$$C_1 + C_3 : 2C_2$$

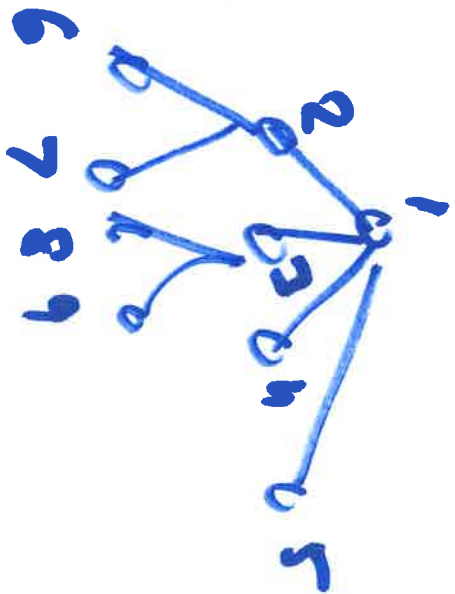
$$\boxed{C_1 + C_3 < 2C_2}$$

$$C_1 + C_3 > 2C_2$$

$$C_1 \leq C_2 \leq C_3$$

4  
5

At



$$\Rightarrow TC_1 = 6C_1 + 2C_2 + C_3 \quad (A_1)$$

$$TC_2 = 5C_1 + 4C_2 \quad (A_2)$$

Least Cost Soln

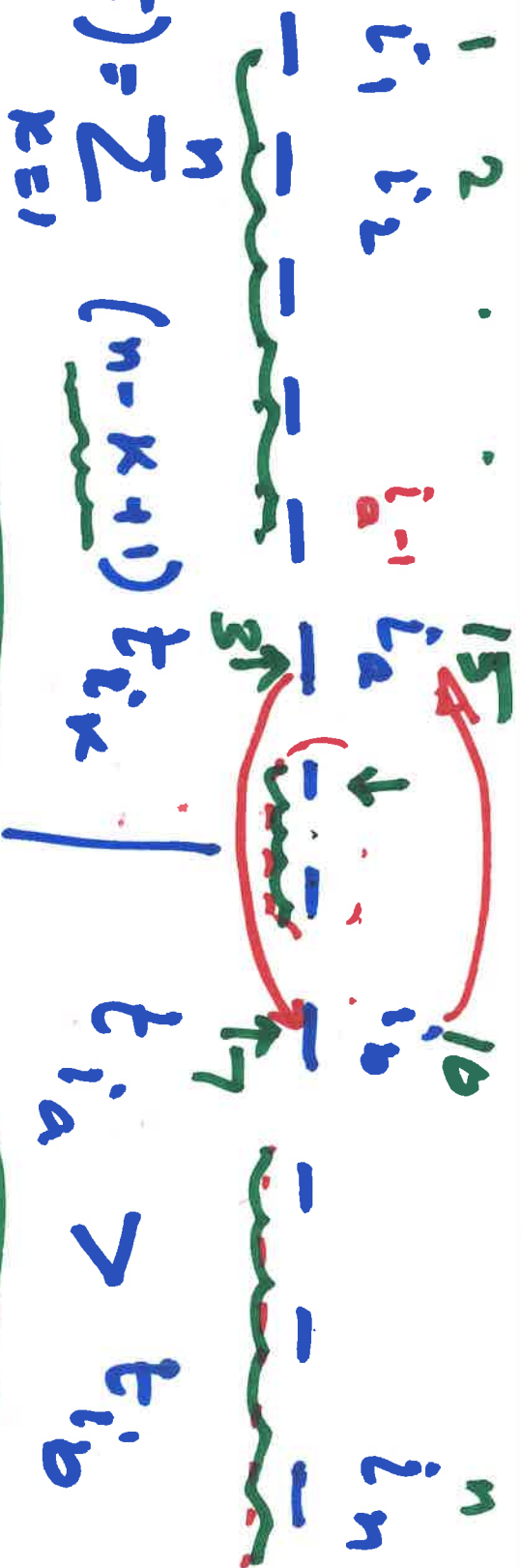
$$6C_1 + 2C_2 + C_3 : 5C_1 + 4C_2$$

$$C_1 + C_3 : 2C_2$$

$$C_1 \leq C_2 \leq C_3$$

$$\frac{C_1 + C_3 > 2C_2}{C_1 + C_3 < 2C_2}$$

$$C_1 + C_3 < 2C_2$$



$$T(I) = \sum_{k=1}^n \underbrace{(n-k+1)}_{\text{weight}} t_{i_k}$$

$$t_{i_a} > t_{i_b}$$

$$\begin{aligned}
 T(I) &= \cancel{x} + \underbrace{(n-a+1)}_{\text{weight}} t_{i_a} + \cancel{1} + \dots + \underbrace{(n-b+1)}_{\text{weight}} t_{i_b} + \cancel{2} \\
 T(I') &= \cancel{x} + \underbrace{(n-a+1)}_{\text{weight}} t_{i_b} + \cancel{1} + \dots + \underbrace{(n-b+1)}_{\text{weight}} t_{i_a} + \cancel{2}
 \end{aligned}$$

$T(I) : T(I')$

$$\underbrace{T(I) - T(I')} = > 0$$

$$\begin{aligned}
 & (n-a+1)t_{i_a} - (n-b+1)t_{i_a} \\
 & + (n-b+1)t_{i_b} - (n-a+1)t_{i_b}
 \end{aligned}$$

$$\begin{aligned}
 & = (\cancel{n-a} + \cancel{n} - \cancel{n} + b - \cancel{n})t_{i_a} \\
 & + (\cancel{n} - b + \cancel{n} - \cancel{n} + a - \cancel{n})t_{i_b}
 \end{aligned}$$

$$= (b-a)t_{i_a} + (a-b)t_{i_b}$$

$$\begin{aligned}
 & a \div b \mid t_{i_a} \div t_{i_b} \\
 & a < b \quad t_{i_a} > t_{i_b}
 \end{aligned}$$

$$= (b-a)t_{i_a} + (b-a)(t_{i_b})$$

$$= (b-a)(t_{i_a} - t_{i_b})$$