

## Time & space Complexity - 2

Q1  $\text{int } c = 0;$   $\rightarrow i = i+1 \rightarrow i = i*2 \rightarrow i*2$   
 $\text{for}(\text{int } i = 1; i < n; i = i) \{$   
 $\quad c++;$   
 $\}$

$i = 1, 2, 4, 8, 16 \dots n$

no. of iterations = no. of values of  $i$

$i = 1, 2^1, 2^2, 2^3, 2^4 \dots 2^x$   
 $\underbrace{\hspace{10em}}_{x+1 \text{ terms}}$   
 $2^x = n$

$$T.C = O(x+1) = O(x) = O(\log n)$$

$$2^x = n$$

$$\log_2 n = x$$

Q2  $\text{int } c = 0$

$\text{for}(\text{int } i = 1; i \leq n; i++) \{$   
 $\quad \text{for}(\text{int } j = 0; j < i; j++) \{$   
 $\quad \quad c++;$   
 $\quad \}$   
 $\}$

$i = 1, j = 0 \rightarrow 1$

$i = 2, j = 0, 1 \rightarrow 2$

$i = 4, j = 0, 1, 2, 3 \rightarrow 4$

$i = 8, j = 0, 1, 2, 3, 4, 5, 6, 7 \rightarrow 8$

$i = n, j = 0, 1, 2, 3 \dots, n-1 \rightarrow n$

(m-1)

$$n = 2^k$$

no. of iterations =  $1 + 2 + 4 + 8 + \dots + n$

$$= 1 + 2^1 + 2^2 + 2^3 + 2^4 + \dots + 2^x$$

(x+1) terms

$$S_n = \frac{a(r^n - 1)}{r - 1}$$

$$= \frac{1(2^{x+1} - 1)}{2 - 1}$$

$$= 2^{x+1} - 1$$

$$\Rightarrow T.C = O(2^{x+1} - 1) = O(2 \cdot 2^x) = O(2^{x+1})$$

$$T.C = O(n)$$

(m-2)

Total no. of iterations:

$$\Rightarrow 1 + 2 + 4 + 8 + 16 + \dots + \frac{n}{2} + n$$

$$\Rightarrow 2 + 2 + 4 + 8 + 16 + \dots + \frac{n}{2} + n - 1$$

$$\Rightarrow 2n - 1$$

$$T.C = O(2n - 1)$$

$$T.C = O(n)$$

$$ex = n = 16$$

$$1 + 1 + 2 + 4 + 8 + 16 - 1$$

$$2 + 2 + 4 + 8 + 16 - 1$$

$$4 + 4 + 8 + 16 - 1$$

$$8 + 8 + 16 - 1$$

$$16 + 16 - 1$$

$$32 - 1$$

$$2n - 1$$

Q3  $\int$  int c = 0;  
 for (int i = 1; i < n; i \*= 2) {  $\rightarrow$  T.C =  $O(\log n)$   
 for (int j = n; j >= 0; j--) {  $\rightarrow$  T.C =  $O(n)$   
 c++; }  
 3 3

j' ka koi lena - Denai  
 hi nhi hai

T.C:  $O(n \cdot \log n)$

i = 1  $\rightarrow$  n+1  
 i = 2  $\rightarrow$  n+1

!

i = n  $\rightarrow$  n+1  
 $\rightarrow O(\log n)$

Q4  $\int$  int c = 0;  
 for (int i = 1; i < n; i \*= 2) {  
 for (int j = 0; j < i; j++) {  
 c++; }  
 }  
 3

T.C:  $O(n)$

i = 1  $\rightarrow$  j = 0, 1  $\rightarrow$  1

i = 2, j = 0, 1  $\rightarrow$  2

i = 3, j = 0, 1, 2, 3  $\rightarrow$  4



Q5

```

int c = 0;
for (int i = 1; i * i < n; i += 2) {
    for (int j = 0; j < i; j++) {
        c++;
    }
}

```

$i^2 < n \rightarrow i < \sqrt{n}$

$i = 1, 2, 4, 8, 16, \dots, 2^x$

$\downarrow$   
 $\sqrt{n}$

$$\rightarrow 2^x \cdot 2^x = n$$

$$\Rightarrow 2^x = \sqrt{n}$$

Total no. of iteration =  $1 + 2 + 3 + 4 + \dots + 2^x$

$$= 2^{x+1} - 1$$

$$\rightarrow T.C = O(2^x) = O(\sqrt{n})$$

$$\boxed{T.C := O(\sqrt{n})}$$

Q6

```

int c = 0;
for (int i = 1; i * i < n; i += i) {
    for (int j = n; j > i; j--) {
        c++;
    }
}

```

$i < \sqrt{n}$        $i += i$

$i = 1, 2, 4, 8, \dots, 2^x$

$\Rightarrow 2^x = \sqrt{n}$

$$2^x = \sqrt{n}$$

$$x = \log_2 \sqrt{n}$$

$i = 1, j = n, n-1, \dots, 3, 2 = n-1$

$i = 2, j = n, n-1, \dots, 4, 3 = n-2$

$i = 4, j = n, n-1, \dots, 6, 5 = n-4$

$i = \sqrt{n}, j = n - \sqrt{n}$

$$\text{Total iteration} = (n-1) + (n-2) + (n-4) + \dots + (n-\sqrt{n})$$

$$= (n-1) + (n-2) + (n-4) + (n-8) + \dots + (n-2^x)$$

$(x+1)$  terms

$$= \underbrace{(n+n+n+\dots)}_{(x+1)} - \underbrace{(1+2+4+8+\dots+2^x)}_{(x+1)}$$

$$= n(x+1) - [2^{x+1} - 1]$$

$$= n \cdot x + n - 2 \cdot 2^x + 1$$

$$T.C = O(n \cdot \log_2 \sqrt{n} + n - \sqrt{n})$$

$$\boxed{T.C = O(n \cdot \log \sqrt{n})}$$

$$\left\{ \begin{array}{l} \log_a b^m = m \log_a b \\ \log_2 \sqrt{n} = \log_2 n^{1/2} = \frac{1}{2} \log_2 n \end{array} \right\}$$

$$T.C = O(n \cdot \log n^{1/2}) = O\left(n \cdot \frac{1}{2} \log n\right)$$

$$\boxed{T.C = O(n \cdot \log n)} \quad \checkmark$$

$$\rightarrow T.C = \log(\log n)$$

Q6 `int c=0`  
`for (int i=2; i<n; i=i*i) {`  
 `c++;`  
`}`

$$i = 2, 4, 16, 256, \dots, 65536$$

khatm

$$i = 2^1, 2^2, 2^4, 2^8, \dots, 2^{16}$$

$$i = 2^1, 2^2, 2^4, 2^8, \dots, 2^{n/2}, 2^n$$

$$\Rightarrow 2^1, 2^2, 2^4, 2^8, 2^{16}, 2^{32}, 2^{64}, \dots$$

$$\underbrace{2^1, 2^{2^1}, 2^{2^2}, 2^{2^3}, 2^{2^4}, 2^{2^5}, 2^{2^6}, \dots, 2^{2^n}}_{(n+1) \text{ terms}}$$

$$2^{2^n} = n$$

$$T.C = O(n+1) \approx O(n)$$

$$2^{2^n} = n \Rightarrow \log_2 n = 2^n$$

$$2^n \log_2(2) = \log_2(n) \quad 2^n = k$$

$$T.C = O(\log(\log n)) \quad \log_2 k = n$$

$$\log_2(\log_2 n) = n$$



Q T.C —  $i \leq \sqrt{n}$   
 $\approx \text{int } c = 0;$   
 for (int  $i = 2$ ;  $i * i \leq n$ ;  $i = i + 1$ ) {  
 $c++;$   
 }

$$i = 2, 4, 16, 256, 65536$$

$$i = 2, 2^2, 2^4, 2^8, 2^{16}, \dots$$

$$i = 2^1, 2^{2^1}, 2^{2^2}, 2^{2^3}, 2^{2^4} \dots 2^{2^n}$$

$n+1$

$$T.C = O(n+1)$$

$$= O(n)$$

$$= O(\log(\log \sqrt{n}))$$

$$= O(\log(\log n^{1/2}))$$

$$= O\left(\log\left(\frac{1}{2} \log n\right)\right) \quad n = \log_2(\log_2 \sqrt{n})$$

$$= O\left(\log \frac{1}{2} + \log(\log n)\right)$$

$$= O(\log(\log n))$$

$$2^{2^n} = \sqrt{n}$$

$$\log_2 \sqrt{n} = 2^n$$

$$\boxed{\begin{matrix} K = 2^n \\ \log_2 K = n \end{matrix}}$$