Limited use cryptographic tokens in securing ephemeral cloud servers

Gautam Kumar, Prof Brent Lagesse Computing and Software Systems University of Washington Bothell {gautamk,lagesse}@uw.edu

Abstract—Many enterprises and consumers today are dependent on services which deploy their computational resources on Infrastructure as a Service (IaaS) cloud providers such as Amazon's AWS, Google's Cloud platform and Windows Azure. Cloud deployments at scale can have hundreds of virtual servers running and the same time sharing sensitive configuration information such as DB passwords and API Keys. This paper aims to propose a risk limiting architecture to safeguard against a potential information leak of sensitive configuration data using a Central Trusted Authority with hash chains as an authentication mechanism for ephemeral servers in the cloud to provide moving target defence against attackers.

Introduction

The essence of securing cloud systems is using multiple layers [16] of security to increase an attacker's cost for taking over the system. One of the possible layer of security is using moving target defences [6].

In this paper we propose an implementation of moving target defence using ephemeral servers and a central trusted authority which acts on behalf an ephemeral server and proxies requests to sensitive resources such as database servers, caching servers and REST end points. Hash chains are used as an authentication mechanism by the central trusted authority. We take advantage of the limited use property of hash chains to secure authenticate ephemeral servers for a limited period of time.

BACKGROUND

Hash Chains [8]

Leslie Lamport [10] first proposed the use of hash chains in his paper on a method for secure password authentication over an insecure medium.

"A hash chain is a sequence of values derived via consecutive applications of a cryptographic hash

function to an initial input. Due to the properties of the hash function, it is relatively easy to calculate successive values in the chain but given a particular value, it is infeasible to determine the previous value"

A hash chain in essence is merely the successive computation of a Cryptographic hash function on a given value.

As an example, Let x be the initial password a hash chain of length 2 would be $H^2(x) = H(H(x))$. A hash chain of n values is denoted as $H^n(x)$ and the i^{th} value in the chain would be computed as $x_i = H(x_{i-1})$.

For a given value in the chain x_i its computationally infeasible to determine the previous value in the chain x_{i-1} .

MOTIVATION

Many organisations are moving their server infrastructure away from physical hardware on-site to virtual hardware in the cloud. Such a migration brings a wide range of problems and potential solutions. TODO: ??

POTENTIAL THREATS

According to OWASP's Top 10 security threats, "Sensitive data exposure" is the 6^{th} most critical type of security threat in web applications as of 2013 [23].

Sensitive data exposure simply refers to unintended exposure of sensitive information such as passwords, social security numbers, date of birth and so on. In the context of a cloud systems sensitive information may also include credentials to access a database, email server, REST API keys and so on. These credentials are usually stored as part of

a configuration file which cloud servers can use to authenticate themselves with third party services within or outside the private cloud network.

According to a report by Risk Based Security [19] [20] the number of data leaks has dramatically increased from 2012 to 2013, to the tune of \$812 million. Though sensitive data exposure in the context of cloud configurations would only constitute a small part of these leaks, leaking of such credentials can potentially lead to massive data leaks or other potential vulnerabilities being exposed to potential attackers.

Sensitive data exposure can potentially be a consequence of other threats such as cross site scripting (XSS) [13], Injection is the most critical threat while XSS is the 3^{rd} most critical threat as classified by OWASP in 2013 [23].

PROPOSED SOLUTION ARCHITECTURE

The proposed architecture to defend against the threat of sensitive data exposure is to use a Central Trusted Authority (CTA) responsible for storing sensitive information. The CTA would act as a proxy and would make requests on behalf of client facing servers, refer Fig. 1.

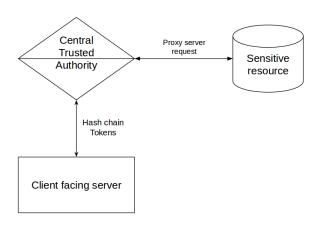


Fig. 1. Architectural overview of the system. This figure describes the three primary modules involved. The client facing server, The central trusted authority and a sensitive resource.

The client facing servers would use hash chains to authenticate with the CTA. Hash chains cryptographically limit the number of times a key can be used. Limited use was intentionally selected to promote the creation of a moving target for attackers.

Assumptions

Client facing servers are the servers which are exposed outside the private cloud network environment. These client facing servers could potentially be load balancing servers, compute servers.

The client facing servers are assumed to be ephemeral. This is common in many cloud deployments [21] and contributes to the moving target nature of the security architecture. Companies such as Netflix expect this behaviour with their chaos engineering architecture [1]. This allows for higher reliability of their server infrastructure.

Client facing servers are expected to shut down after their hash chain expires. This contributes to the ephemeral nature and also to moving target defence of the overall system.

CTA Architecture and Implementation

The Central Trusted Authority consists of three primary components, A hash chain verifier, A storage backend and a request proxy. The CTA generally performs three roles within the system which are

- · Create new hash chain
- · Verify hash chains
- Proxy requests

Creating new hash chain: Hash chains are created by iteratively hashing a secret key K, n number of times. After the hash chain is created the CTA stores $H^{100}(K)$ in the storage backend and returns K and n to the client facing server. The secret key K is not stored by the CTA.

The client facing server can now use the the secret key n number of times.

```
Data: Hash Chain secret K and Hash chain length N
Result: Hash Chain H^N(K)

1 i \leftarrow 1;
2 H^1(K) \leftarrow H(K);
3 while i <= N do
4 | i \leftarrow i+1;
5 | H^i(K) \leftarrow H(H^{i-1}(K));
6 end
7 return H^i(K) where i equals N;
```

Algorithm 1: Generating a Hash Chain

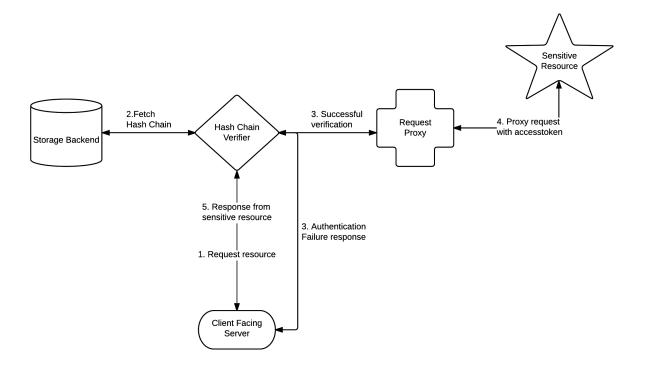


Fig. 2. Architecture of the Central Trusted Authority. The CTA consists of a storage backend, a hash chain verifier and a request proxy.

Hash chain verification: Hash chains are used to authenticate client facing server requests which require access to a sensitive resource. The hash chain verification is detailed in algorithm 2.

```
Data: Authentication key H^{i-1}(K) from the client
```

Result: Response from sensitive resource

- 1 Let $H_{client}^{i} = H(H^{i-1}(K))$;
- 2 Let $AuthenticationData = fetch(H_{client}^{i})$;
- 3 if AuthenticationData exists in storage backend then
- replace H^i_{client} with $H^{i-1}(K)$ in storage backend;
- fetch and return response from sensitive resource;
- 6 else
- return authentication failure;
- 8 end

Algorithm 2: Verification of Hash Chain authentication

REQUEST PROXY

The proposed solution architecture of using a Central Trusted Authority (CTA) to proxy requests on behalf of all the clients, places the CTA as a single point of failure. We can overcome this limitation by improving the reliability and trustability of the CTA.

Reliability

A lot of research [5] [22] [24] [9] has been conducted into improving the reliability of cloud systems through vertical, horizontal scaling and automatic provisioning. Much of that work can be leveraged for the purpose of improving the reliability of the proposed architecture. For example, Beltrán [2] proposes an architecture of utilizing multi-tired load balancers to improve the reliability of a service.

Large scale database deployments typically deploy databases behind a reverse proxy for load balancing and Geo distribution [7]. Thus a reverse proxy such as Vitess [7] is ideally suited to act as the CTA in such architectures.

Security and Trust

Trust and security of the CTA is another primary area of concern with a distributed cloud system. The CTA as a central point of access to all sensitive resources would be a prime target for malicious actors and thus cannot be completely trusted. Chen et al [4] propose a solution to this problem of a Malicious proxy using trusted hardware such as Trusted Platform Module (TPM) or the IBM 4758 cryptographic coprocessor [17].

Chen et al assume the proxy, the CTA in our architecture, is malicious but is incapable of modifying the underlying hardware. The CTA executable is also verified by a trusted third party to operate correctly as a proxy as described in [17].

There are three primary attacks that a malicious actor may perform on the CTA. First, the attacker may try to expose the sensitive information stored, such as API keys and DB passwords, from the CTA using vulnerabilities in the CTA executable. Second, the attacker could potentially modify the requests / responses which are being proxied by the CTA. Third the attacker could potentially launch a reboot attack to inject a malicious executable after attestation to carry out one of the above attacks.

Prior work [12], [15] on preventing reboot-attacks can be leveraged to impede the attacker's ability to inject a malicious executable. The proposed architecture also allows for the CTA to be placed behind secure corporate firewalls to further limit the risk of a malicious take over by an attacker. Further work is needed to fully secure the CTA against a malicious proxy.

PERFORMANCE TESTING

In our testing hash chain initialization showed a linear increase in time complexity as shown in figure 4. This is in line with our expectations for hash chain implementations. A hash chain of length 100,000 takes 150 - 200 ms to initialize based on the storage backend used.

LIMITATIONS

The proposed architecture and implementation are ideally suited to cloud IaaS providers such as Amazon AWS where ephemeral servers are easily managed. Deploying the proposed system to a cloud provider without the ability to quickly provide large

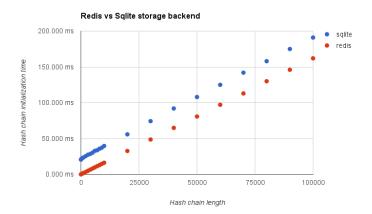


Fig. 4. Linear time complexity of the hash chain initialization. Execution platform Intel i5-5200U 2.2Ghz, 12GB RAM, SHA-512 implemented in python.

number of ephemeral servers could result is significant performance degradation.

Hash chains inherently posses certain vulnerabilities such as a Hash chain cycle and Hash chain length oracle attacks which can potentially reduce the effectiveness of the proposed architecture [11].

FUTURE WORK

Formulating an ideal balance between server lifespan and hash chain length to optimize computational resources in a cloud system, this can be derived from current work into load balancing in cloud systems [18]. Such a formulation can be used by dev ops engineers in choosing an ideal hash chain size based on the performance requirements of a cloud system.

Integrating Timed Release Cryptography [3] as a hash chain renewal mechanism to potentially increase the lifespan of a server after a certain cool off period.

A Merkle hash tree implementation middle-ware can potentially provide hierarchical authentication authorization [25] capabilities to the CTA.

RELATED WORK

Confidant [14] is a library maintained by Lyft, a transportation network company based out of San Francisco. Confidant provides an implementation of the Central trusted authority server with encryption at rest, authentication and authorization handled by

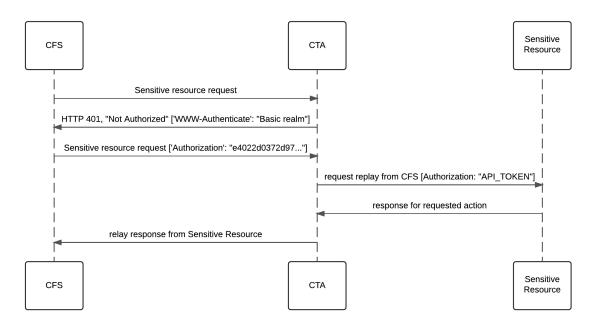


Fig. 3. Sequence diagram describing the authentication and proxying capabilities of the CTA. CFS refers to Client Facing Server while CTA refers to Central Trusted Authority

AWS's Key Management Service, KMS and a storage backend of DynamoDB. This severely hampers the ability of a potential user to deploy a Confidant instance on a different cloud IaaS provider beside Amazon's AWS.

Confidant also serves as an inspiration for the CTA component of the proposed architecture.

SUMMARY

In this paper we propose an architecture to provide moving target defence and minimize the damage that attackers can cause by centralizing the storage of sensitive configuration information and taking advantage of the limited use property of hash chains to authenticate potentially vulnerable client facing servers.

REFERENCES

- A. Basiri, N. Behnam, R. de Rooij, L. Hochstein, L. Kosewski, J. Reynolds, and C. Rosenthal. Chaos engineering. 33(3):35– 41.
- [2] Marta Beltrn. Automatic provisioning of multi-tier applications in cloud computing environments. 71(6):2221–2250.
- [3] Konstantinos Chalkias and George Stephanides. Timed release cryptography from bilinear pairings using hash chains. In Communications and Multimedia Security, pages 130–140. Springer.

- [4] Ruichuan Chen, Alexey Reznichenko, Paul Francis, and Johanes Gehrke. Towards statistical queries over distributed private user data. In Presented as part of the 9th USENIX Symposium on Networked Systems Design and Implementation (NSDI 12), pages 169–182.
- [5] S. Dutta, S. Gera, A. Verma, and B. Viswanathan. SmartScale: Automatic application scaling in enterprise clouds. In 2012 IEEE 5th International Conference on Cloud Computing (CLOUD), pages 221–228.
- [6] David Evans, Anh Nguyen-Tuong, and John Knight. Effectiveness of moving target defenses. In Sushil Jajodia, Anup K. Ghosh, Vipin Swarup, Cliff Wang, and X. Sean Wang, editors, *Moving Target Defense*, number 54 in Advances in Information Security, pages 29–48. Springer New York. DOI: 10.1007/978-1-4614-0977-9 2.
- [7] Todd Hoff. 7 years of YouTube scalability lessons in 30 minutes. http://highscalability.com/blog/2012/3/26/7-years-of-youtube-scalability-lessons-in-30-minutes.html.
- [8] Dwight Horne. Hash chain. In Henk C. A. van Tilborg and Sushil Jajodia, editors, *Encyclopedia of Cryptography and Security*, pages 542–543. Springer US. DOI: 10.1007/978-1-4419-5906-5_780.
- [9] Senthil SK Kumar and P. Balasubramanie. Dynamic scheduling for cloud reliability using transportation problem. 8(10):1615.
- [10] Leslie Lamport. Password authentication with insecure communication. 24(11):770–772.
- [11] D. Lee. Hash function vulnerability index and hash chain attacks. In 2007 3rd IEEE Workshop on Secure Network Protocols, pages 1–6.
- [12] Benot Libert and Damien Vergnaud. Tracing malicious proxies in proxy re-encryption. In Steven D. Galbraith and Kenneth G. Paterson, editors, *Pairing-Based Cryptography Pairing 2008*, number 5209 in Lecture Notes in Computer Science, pages

- 332–353. Springer Berlin Heidelberg. DOI: 10.1007/978-3-540-85538-5 22.
- [13] M. T. Louw and V. N. Venkatakrishnan. Blueprint: Robust prevention of cross-site scripting attacks for existing browsers. In 2009 30th IEEE Symposium on Security and Privacy, pages 331–346.
- [14] lyft. Confidant: Your secret keeper. A library to store and retrive senstive configuration within a central trusted authority encrypted at rest using Amazon KMS. https://github.com/lyft/confidant.
- [15] Jonathan M. McCune, Bryan J. Parno, Adrian Perrig, Michael K. Reiter, and Hiroshi Isozaki. Flicker: An execution infrastructure for TCB minimization. In ACM SIGOPS Operating Systems Review, volume 42, pages 315–328. ACM.
- [16] A. Panwar, R. Patidar, and V. Koshta. Layered security approach in cloud. In 3rd International Conference on Advances in Recent Technologies in Communication and Computing (ARTCom 2011), pages 214–218.
- [17] Bryan Parno, Jonathan M. McCune, and Adrian Perrig. Bootstrapping trust in commodity computers. In 2010 IEEE Symposium on Security and Privacy, pages 414–429. IEEE.
- [18] M. Randles, D. Lamb, and A. Taleb-Bendiab. A comparative study into distributed load balancing algorithms for cloud computing. In 2010 IEEE 24th International Conference on Advanced Information Networking and Applications Workshops (WAINA), pages 551–556.
- [19] Risk Based and Security. An executives guide to 2013 data breach trends.
- [20] Xiaokui Shu, Danfeng Yao, and Elisa Bertino. Privacypreserving detection of sensitive data exposure. 10(5):1092– 1103.
- [21] Luis M. Vaquero, Luis Rodero-Merino, and Rajkumar Buyya. Dynamically scaling applications in the cloud. 41(1):45–52.
- [22] Kashi Venkatesh Vishwanath and Nachiappan Nagappan. Characterizing cloud computing hardware reliability. In *Proceedings of the 1st ACM symposium on Cloud computing*, pages 193–204. ACM.
- [23] Dave Wichers. OWASP top-10 2013.
- [24] Zhang Xuejie, Wang Zhijian, and Xu Feng. Reliability evaluation of cloud computing systems using hybrid methods. 19(2):165–174.
- [25] X. Yi and W. Wang. The cloud access control based on dynamic feedback and merkle hash tree. In 2012 Fifth International Symposium on Computational Intelligence and Design (ISCID), volume 1, pages 217–221.