

# Find out the vertices contributing to vertex connectivity

Submitted by:

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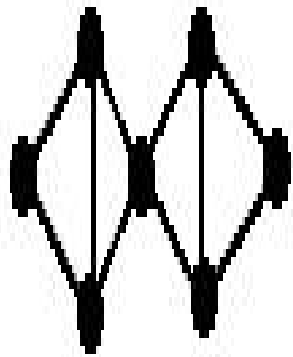
IIT2014121

Nazish Tabassum

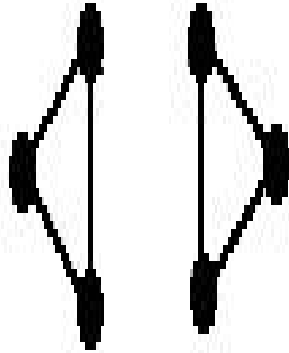
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Arqum Ahmed

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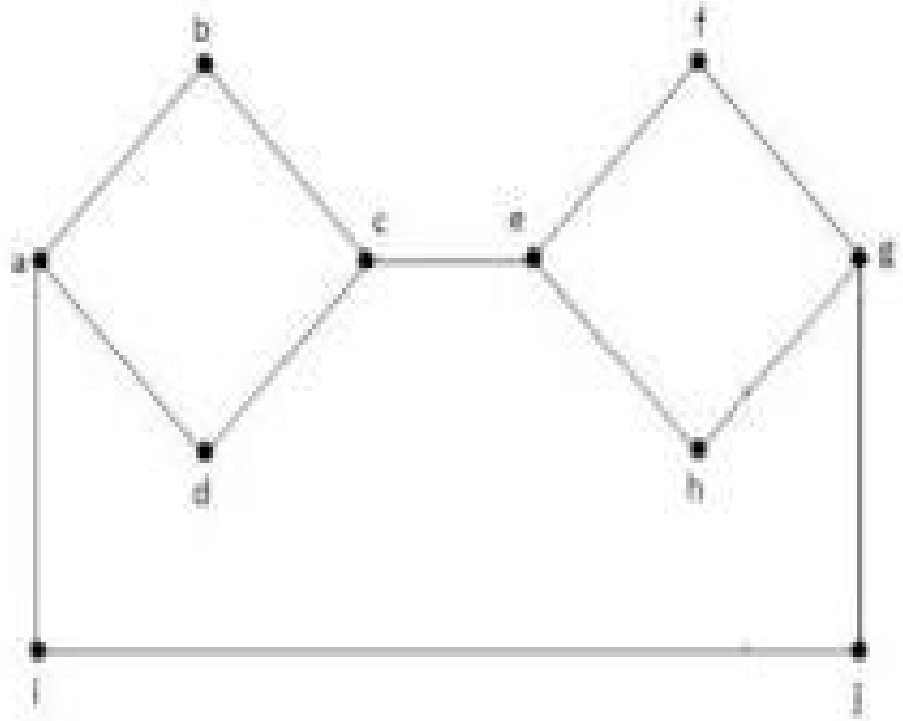


Vertices deleted: 0



Vertices deleted: 1

Vertex connectivity is 1



Vertex Connectivity is 2

# Algorithm 1

## A. Brute Force Approach:

Input : A graph  $G$  with cardinality equal to  $n$ .

Output : A set of vertices  $S$  whose size is equal to vertex connectivity of graph  $G$  and it contains all vertices enlisted in vertex connectivity.

FindVerticesInVertexConnectivity( $G$ )

- 1    RemoveSelfLoop
- 2    RemoveParallelEdges
- 3     $S = \text{emptySet}$
- 4    If graph is disconnected or  $n = 1$
- 5        return  $S$
- 6    If graph is complete
- 7        Select any  $n-1$  vertices and add it to  $S$
- 8        Return  $S$
- 9    For  $i = 1$  to  $(n-2)$
- 10        Generate every combination of vertices of size  $i$
- 11        remove those vertices and add it to set  $S$
- 12        If graph is disconnected
- 13            return  $S$
- 14        empty set  $S$

# Algorithm 2

```
1  RemoveSelfLoop
2  RemoveParallelEdges
3  S = emptySet
4  If graph is disconnected or  $n = 1$ 
5    return S
6  If graph is complete
7    Select any  $n-1$  vertices and add it to S
8    Return S
9  Replace each edge  $(x, y) \in E$  with arcs  $(x, y)$  and
   $(y, x)$ , and call the resulting digraph D.
10 For each pair  $v, w$  in  $G$  which are non-adjacent
11   For each vertex  $u$  other than  $v$  and  $w$  in  $G$ ,
     replace  $u$  with two new vertices  $u_1$  and  $u_2$ , and
     then add the new arc  $(u_1, u_2)$ . Connect all the
     arcs that were coming to  $u$  in  $G$  to  $u_1$ , and
     similarly, connect all the arcs that were going
     out of  $u$  in  $G$  to  $u_2$  in  $D$ .
12   Assign  $v$  as the source vertex and  $w$  as the sink
     vertex.
13   Assign the capacity of each arc to 1, and call
     the resulting network  $H$ .
14    $T = \text{ComputeMinCutUsingMaxFlow}(H, v, w)$ 
15   if  $(|T| < |S|)$ 
16      $S = T$ 
17   Restore network  $H$  back to  $D$ .
18 return S
```

ComputeMinCutUsingMaxFlow( $H, v, w$ )

- 1 Run Ford-Fulkerson algorithm and consider the final residual graph
- 2 Find the set of vertices that are reachable from source in the residual graph.
- 3 All edges which are from a reachable vertex to non-reachable vertex are minimum cut edges.
- 4 Find the corresponding vertices to edges all the edges that form minimum cut edges.
- 5 Add all those vertices to Set  $T$
- 6 Return  $T$

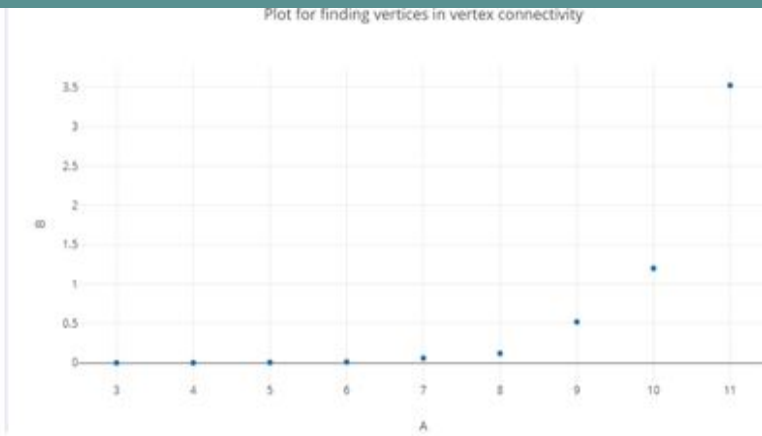


Fig. 1. Algorithm1 graph

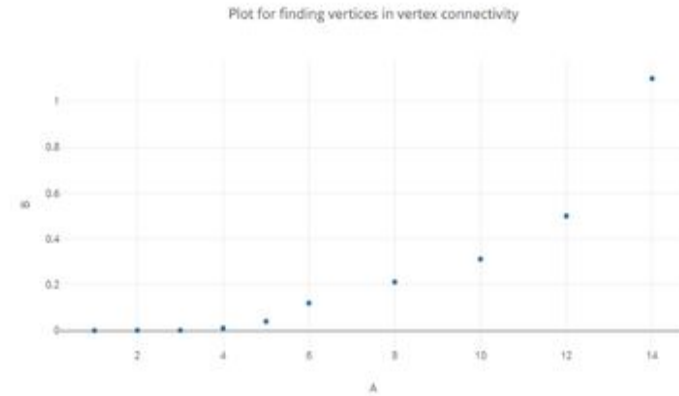


Fig. 2. Algorithm2 graph