

Data Structures and Algorithms Project

Project Theme: 23. ADT Priority Queue – implementation on a singly linked list on an array

1. **Domain** of the Abstract Data Type

$\mathcal{PQ} = \{ pq \mid pq \text{ is a priority queue with elements } (e, p), e \in \text{TElem}, p \in \text{TPriority} \}$

2. **Interface** of the Abstract Data Type

- `init(pq, R)`
// Description: creates a new empty priority queue
// Pre: R is a relation over the priorities, R: $\text{TPriority} \times \text{TPriority}$
// Post: $pq \in \mathcal{PQ}$, pq is an empty priority queue
- `destroy(pq)`
// Description: destroys a priority queue
// Pre: $pq \in \mathcal{PQ}$
// Post: pq was destroyed
- `push(pq, e, p)`
// Description: pushes (adds) a new element to the priority queue
// Pre: $pq \in \mathcal{PQ}$, $e \in \text{TElem}$, $p \in \text{TPriority}$
// Post: $pq' \in \mathcal{PQ}$, $pq' = pq \oplus (e, p)$
- `pop(pq, e, p)`
// Description: pops(removes) from the priority queue the element with the highest priority.
It returns both the element and its priority
// Pre: $pq \in \mathcal{PQ}$
// Post: $pq' \in \mathcal{PQ}$, $pq' = pq \ominus (e, p)$
 $e \in \text{TElem}$, $p \in \text{TPriority}$, e is the element with the highest priority from the priority queue, p is its priority
// Throws: an exception if the priority queue is empty
- `top(pq, e, p)`
// Description: returns from the priority queue the element with the highest priority and its priority
// Pre: $pq \in \mathcal{PQ}$
// Post: $pq' \in \mathcal{PQ}$, $pq' = pq \ominus (e, p)$
 $e \in \text{TElem}$, $p \in \text{TPriority}$, e is the element with the highest priority from the priority queue, p is its priority
// Throws: an exception if the priority queue is empty

- `isEmpty(pq)`
// Description: checks if the priority queue is empty
// Pre: $pq \in \mathcal{PQ}$
// Post: $isEmpty \leftarrow \begin{matrix} \text{true, if } pq \text{ has no elements} \\ \text{false, otherwise} \end{matrix}$
- `increasePriority(pq, e, p)`
// Description: increases the priority of a given elements
// Pre: $pq \in \mathcal{PQ}$, $e \in TElem$, $p \in Tpriority$
// Post: $pq \in \mathcal{PQ}$, $e \in TElem$, $p' \in Tpriority$, $p' > p$, p' is the new priority of e

3. Representation of the ADT Priority Queue on a singly linked list on an array

<u>Pair:</u> elem: TElem p: TPriority	<u>Node:</u> e: Pair next: Integer	<u>SLLA:</u> elems: Node[] cap: Integer size: Integer head: Integer firstEmpty: Integer R: Relation
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4. Implementations of the operations for ADT Priority Queue on a SLLA

subalgorithm initPriorityQueue(pq, R, capacity) is:

```

// Description: creates a new empty priority queue
// Pre: R is a relation over the priorities, R: TPriority x Tpriority
// Post:  $pq \in \mathcal{PQ}$ ,  $pq$  is an empty priority queue
// Complexity:  $O(n)$ 
{
    pq.size  $\leftarrow$  0
    pq.cap  $\leftarrow$  capacity
    pq.elems =  $\uparrow$ initNode(Node, pq.cap)
    for i  $\leftarrow$  0, i < pq.cap, i++
        new_elems[i].next  $\leftarrow$  i + 1
    new_elems[pq.cap].next  $\leftarrow$  -1
    pq.head = -1
    pq.firstEmpty = 0
    pq.priority = R
}

```

subalgorithm push(pq, a, t) is:

```
// Description: pushes (adds) a new element to the priority queue
// Pre: pq ∈ PQ, a ∈ Airplane, t ∈ Double
// Post: pq' ∈ PQ, pq' = pq ⊕ (e, p)
// Complexities:
//   Best case: θ(1)
//   Worst case: θ(size)
//   Average case: θ(size) amortized =  $\frac{1 + 2 + \dots + \text{size}}{\text{size}} = \frac{\sum_{i=1}^{\text{size}} i}{\text{size}} = \frac{\text{size}(\text{size}+1)}{2 * \text{size}}$ 
//   (amortized, as we are not taking into consideration the resizing operation, considering the small
//   number of times it should be performed
{
    if size = capacity then
        resize(pq)
    pair ← init(Pair, e, p);
    if head = -1 then
        insertEmptyQueue(pq, pair);
    else
        prev_pos ← -1, current_pos ← head
        condition ← True
        while condition execute
            if current_pos = -1
                condition = False
            else
                airplane ← elems[current_pos].pair.airplane
                if Relation.compare(a, airplane) == True
                    prev_pos ← current_pos
                    current_pos ← elems[current_pos].next
                else condition ← False
        if prev_pos = -1
            insertFirst(pq, pair)
        else
            insertOnPosition(pq, pair, position)
    size = size + 1
}
```

subalgorithm resize(pq, factor) is:

```
// Description: multiplies the capacity of the priority queue by factor
// Pre: pq ∈ PQ, factor ∈ Integer
// Post: pq' ∈ PQ, pq.cap = pq.cap * factor
// Complexity: O(size)
{
    pq.cap ← pq.cap * factor
    new_elems[] ← init(Node, pq.cap)

    for i ← pq.size, i < pq.cap - 1, i++ execute
        new_elems[i].next ← i + 1
    new_elems[pq.cap].next ← -1
    pq.firstEmpty = pq.size

    for i ← 0, i < pq.size, i++ execute
        new_elems[i] ← pq.elems[i]
}
```

subalgorithm insertEmptyQueue(pq, pair) is:

```
// Description: adds a Pair to an empty Priority Queue
// Pre:  $pq \in \mathcal{PQ}$ ,  $pair \in \text{Pair}$ 
// Post:  $pq' \in \mathcal{PQ}$ ,  $pair \in pq$ ,  $pq.top = pq.elems[pq.head].pair$ 
// Complexity:  $O(1)$ 
{
    new_node  $\leftarrow$  init(Node, pair, -1);
    new_position  $\leftarrow$  pq.firstEmpty
    pq.firstEmpty  $\leftarrow$  pq.elems[new_position].next
    pq.elems[new_position]  $\leftarrow$  new_node
    pq.head  $\leftarrow$  new_position
}
```

subalgorithm insertOnPosition(pq, pair, prev_position) is:

```
// Description: adds a Pair to a Priority Queue
// Pre:  $pq \in \mathcal{PQ}$ ,  $pair \in \text{Pair}$ 
// Post:  $pq' \in \mathcal{PQ}$ ,  $pair \in pq$ ,  $pq.top = pq.elems[pq.head].pair$ 
// Complexity:  $O(1)$ 
{
    new_node  $\leftarrow$  init(Node, pair, -1);
    new_position  $\leftarrow$  pq.firstEmpty
    pq.firstEmpty  $\leftarrow$  pq.elems[new_position].next
    pq.elems[new_position]  $\leftarrow$  new_node
    pq.elems[prev_position].next  $\leftarrow$  new_position
}
```

subalgorithm insertFirst(pq, pair) is:

```
// Description: adds a Pair to the top of the Priority Queue
// Pre:  $pq \in \mathcal{PQ}$ ,  $pair \in \text{Pair}$ 
// Post:  $pq' \in \mathcal{PQ}$ ,  $pair \in pq$ ,  $pq.top = pq.elems[pq.head].pair$ 
// Complexity:  $O(1)$ 
{
    new_node  $\leftarrow$  init(Node, pair, pq.head);
    new_position  $\leftarrow$  pq.firstEmpty
    pq.firstEmpty  $\leftarrow$  pq.elems[new_position].next
    pq.elems[new_position]  $\leftarrow$  new_node
    pq.head  $\leftarrow$  new_position
}
```

function pop(pq) is:

```
// Description: pops(removes) and returns the pair with the highest priority from the priority queue
// Pre:  $pq \in \mathcal{PQ}$ 
// Post:  $pq' \in \mathcal{PQ}$ ,  $pq' = pq(e, p) \ominus e \in \text{TElem}$ ,  $p$  Tpriority,  $e$  is the element with the highest priority from the
//       priority  $\in \in$  queue,  $p$  is its priority
// Throws: an exception if the priority queue is empty
// Complexity:  $O(1)$ 
{
    if pq.size = 0
        throw exception "Priority queue is empty!"
    int emptied_position = pq.head
    pair  $\leftarrow$  pq.elems[pq.head].pair
    pq.head = pq.elems[pq.head].next
    pq.elems[emptied_position].next  $\leftarrow$  pq.firstEmpty
    pq.firstEmpty  $\leftarrow$  emptied_position
    pop  $\leftarrow$  pair
}
```

function top(pq) is:

```
// Description: returns from the priority queue the element with the highest priority and its priority
// Pre: pq ∈ PQ
// Post: pq ∈ PQ, Pair contains the element with the highest priority from the priority queue and its priority
// Throws: an exception if the priority queue is empty
// Complexity: O(1)
{
    if pq.size = 0
        throw exception "Priority queue is empty!"
    int emptied_position = pq.head
    pair ← pq.elems[pq.head].pair
    top ← pair
}
```

function isEmpty(pq) is:

```
// Description: checks if the priority queue is empty
// Pre: pq ∈ PQ
// Post: isEmpty ← true, if pq has no elements
//           false, otherwise
// Complexity: O(1)
{
    if pq.size = 0
        isEmpty ← true
    isEmpty ← false
}
```

5. Tests for the Priority Queue ADT

void Test::test()

```
{
    PriorityQueue pq{new TimePriority()};

    Airplane a1{ 1, "WizzAir", "Bucharest", 15.00 };
    Airplane a2{ 2, "WizzAir", "CLuj-Napoca", 14.30 };
    Airplane a3{ 4, "Lufthansa", "Munich", 14.45 };
    Airplane a4{ 3, "Blue Air", "Bucharest", 17.00 };

    // A1 IS ADDED FIRST IN THE PQ -> TEST INSERT IN EMPTY QUEUE
    pq.push(a1, a1.getHour());
    assert(pq.getSize() == 1);
    assert(pq.top().getAirplane().getID() == a1.getID());

    // A2 IS ADDED SECOND BUT WILL BE THE FIRST IN THE PQ -> TEST INSERT FIRST
    pq.push(a2, a2.getHour());
    assert(pq.getSize() == 2);
    assert(pq.top().getAirplane().getID() == a2.getID());

    // A3 WILL BE THE SECOND IN THE PQ -> TEST INSERT ON POSITION
    pq.push(a3, a3.getHour());
    assert(pq.getSize() == 3);

    // A4 WILL BE THE LAST IN THE PQ -> TEST INSERT LAST
    pq.push(a4, a4.getHour());
    assert(pq.getSize() == 4);
}
```

```
// VERIFY IF A3 IS THE SECOND AND A4 IS THE FOURTH -> TEST TOP AND POP
pq.pop();
assert(pq.top().getAirplane().getID() == a3.getID());
pq.pop(); pq.pop();
assert(pq.top().getAirplane().getID() == a4.getID());

// EXTRACT THE LAST ELEMENT -> TEST isEmpty
assert(pq.isEmpty() == false);
pq.pop();
assert(pq.isEmpty() == true);
}

void Test::testExceptions()
{
    PriorityQueue pq{ new TimePriority() };

    // PRIORITY QUEUE IS EMPTY, CALLING POP OR TOP ON IT THROWS EXCEPTIONS
    try {
        pq.pop();
    }
    catch (std::exception & ex)
    {
        std::string exception = ex.what();
        assert(exception == "Priority queue is empty!");
    }

    try {
        pq.top();
    }
    catch (std::exception & ex)
    {
        std::string exception = ex.what();
        assert(exception == "Priority queue is empty!");
    }
}
```

6. Problem Statement + Justification

Simulate the flights display in the arrival's terminal of an airport with two gates. You should be able to see at any time the last arrived airplanes for each of the two gates and also the next flight scheduled to arrive. Marking the next scheduled flight as arrived should put it on the display of the right gate. An airplane needs a total time of 30 minutes to disembark passengers and make the preparations for the next flight or leave the gate. If an airplane arrives before any gate can be emptied, it should be displayed that it is redirected to another terminal. You should also be able to add at any time a new airplane scheduled later for the day by adding its ID, company, origin and arrival time.

I have chosen the above problem statement because priority queues are of great use in managing events in simulations like an airport scenario. In this case the priority of the element is represented by the scheduled time of landing. All the main typical operations for a priority queue can be used in such a problem.

The Next Arrival display will always contain the data of the top of the priority queue, while the gate displays will show the data of the last popped items. Adding a new flight is done by entering info like ID, company and origin, and the last argument, arrival time represents the priority of the element: earlier flights are located nearer the top. The flight will later be displayed only at the right time according to its arrival time. Exceptions thrown by pop and top functions are handled by displaying the message "No incoming airplane!" instead of the next flight information.

7. Pseudocode for the solution of the problem

subalgorithm initController(ctr, PriorityQueue pq) is:

```
// Description: initializes controller with a given priority queue of objects of type Airplanes
// Pre: pq ∈ PQ
// Post: ctr ∈ PQ
// Complexity: O(1)
{
    ctr.airplanes ← pq;
    ctr.gate_1 ← initPair()
    ctr.gate_2 ← initPair()
}
```

subalgorithm set_landed(ctr) is:

```
// Description: Redirects the arriving airplane to the right gate, but only if there is one available; pops an item
// from the priority queue, the one set as next arrival before
// Post: if possible, gate_1 or gate_2 are assigned a new Airplane and last_gate is updated
// Complexity: O(1)
{
    if ctr.is_available_1 then
        ctr.gate_1 ← ctr.airplanes.pop()
        ctr.last_gate ← 1
    else if ctr.is_available_2 then
        ctr.gate_2 ← ctr.airplanes.pop()
        ctr.last_gate ← 2
    else ctr.airplanes.pop()
}
```

function get_last_arrival(ctr) is:

```
// Description: returns information about the last popped item, i.e., the last airplane arrived
// Complexity: O(1)
{
    if ctr.last_gate = 1 then
        function ← ctr.gate_1.airplane
    else
        function ← ctr.gate_2.airplane
}
```

function is_available_1(ctr) is:

```
// Description: returns true if the airplane at gate 1 arrived for more than 30 minutes and false otherwise
// Complexity: O(1)
{
    next ← ctr.airplanes.top()
    if time between next.priority and ctr.gate_1.priority is less than 30 minutes then
        is_available_1 ← True
    else
        is_available_1 ← False
}
```

function is_available_2(ctr) is:

```
// Description: returns true if the airplane at gate 2 arrived for more than 30 minutes and false otherwise
// Complexity: O(1)
{
    next ← ctr.airplanes.top()
    if time between airplanes is less than 30 minutes then
        is_available_2 ← True
    else
        is_available_2 ← False
}
```

subalgorithm add_airplane(ctr, airplane) is:

```
// Description: pushes a new pair to the priority queue
// Pre: airplane Airplane
// Post: pq' ∈ PQ, airplane and priority ∈ pq
// Complexity of push operation
{
    ctr.airplanes.push(airplane, airplane.arrivalTime);
}
```