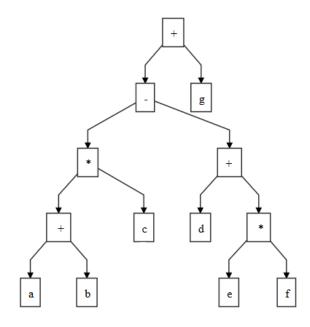
DSA - Seminar 7

- Project stage for projects with trees (binary tree, binary search tree)
- 1. Build the binary tree for an arithmetic expression that contains the operators +, -, *, /. Use the postfix notation of the expression.

Ex:
$$(a + b)^* c - (d + e * f) + g =>$$

Postfix notation: ab+c*def*+-g+

The corresponding binary tree is:

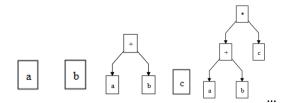


If we traverse the tree in postorder, we will get the postfix notation.

Algorithm:

- 1. Use an auxiliary stack that contains the address of nodes from the tree
- 2. Start building the tree from the bottom up.
- 3. Parse the postfix expression
- 4. If we find an operand -> push it to the stack
- 5. If we find an operator->
 - a. Pop an element from the stack left child
 - b. Pop an element from the stack right child
 - c. Create a node containing as information the operator and the left and right child
 - d. Push this new node to the stack
- 6. The root of th tree will be the last element from the stack.

Stack:



Assume we have a binary tree with dynamically allocated nodes.

Node:

e: TElem <u>BT</u>:

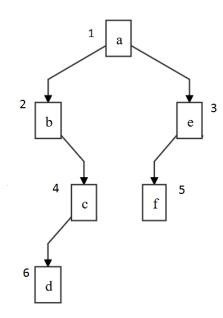
left, right: ↑Node root: ↑Node

The stack will contain elements of type \tauNode and we will only use the interface of the stack

- Init
- Push
- Pop
- Top

```
Subalgorithm buildTree (postE, tree) is:
    init(s)
    for every e in postE execute:
                                                                    If, instead of node, we want
         if e is an operand then:
                                                                    to use binary trees: Stack
              alocate (newNode)
                                                                    will contain elements of type
              [newNode].e \leftarrow e
                                                                    AB and we will call
              [newNode].left ← NIL
                                                                    operations from AB's
              [newNode].right ← NIL
              push (s, newNode)
                                                                    interface.
         else
                                                                    initLeaf(bt, e)
              pop(s, p1)
                                                                    push(s, bt)
              pop(s, p2)
              alocate (newNode)
              [newNode].e \leftarrow e
              [newNode].left \leftarrow p2
              [newNode].right \leftarrow p1
                                                                    initTree(bt, p2, e, p1)
              push (s, newNode)
                                                                    push(s, bt)
         end-if
    end-for
                                                                    pop(s, tree)
    pop(s, p)
    tree.root ← p
end-subalgorithm
```

2. Generate the table with information from a binary tree. Node numbering is done according to levels.



	1	2	3
	Info	Index Left	Index Right
1	а	2	3
2	b	0	4
3	е	5	0
4	С	6	0
5	f	0	0
6	d	0	0

- Divide the solution into 2 functions: addNumbers and buildTable
- We use a queue for storing the nodes (we need level-order traversal)
- Assume that each Node has a field *nr:Integer* (we are going to store the number of a node here).

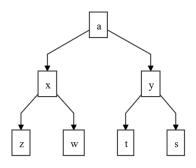
```
Subalgorithm addNumbers (tree, k)
//pre: tree is a binary tree
//post: nr field from every node is set to the correct value, k is an integer
number, it represents the number of nodes from the tree.
    k ← 0
    init(q)
    if tree.root # NIL then
        push(q, tree.root)
        k ← 1
         [tree.root].nr \leftarrow k
    end-if
    while (¬ isEmpty(q)) execute
        pop (q, p)
        if ([p].left # NIL) then
             k \leftarrow k + 1
             [[p].left].nr \leftarrow k
             push(q, [p].left)
         end-if
         if ([p].right # NIL) then
             k \leftarrow k + 1
             [[p].right].nr \leftarrow k
             push(q, [p].right)
         end-if
    end-while
end-subalgorithm
```

```
subalgorithm buildTable(p, T) is:
//pre: p is a pointer to a node, T is a matrix that holds the information
from the tree (column 1 node info, column 2 index of left, column 3 index of
right
    if (p \neq NIL) then
        T[[p].nr, 1] \leftarrow [p].e
        if ([p].left # NIL) then
             T[[p].nr, 2] \leftarrow [[p].left].nr
        else
             T[[p].nr, 2] \leftarrow 0
        end-if
        if ([p].right # NIL) then
             T[[p].nr, 3] \leftarrow [[p].right].nr
        else
             T[[p].nr, 3] \leftarrow 0
        end-if
        buildTable([p].left, T)
        buildTable([p].right, T)
    end-if
end-subalgorithm
subalgorithm table(tree, T, k) is:
    addNumbers(tree, k)
    @define T as a matrix with k lines and 3 columns
    buildTable(tree.root, T)
end-subalgorithm
```

- How could I solve the problem if I do not want to add the nr field to the Node structure?
 - I can write one single function (addNumbers + buildTable together) and can add to the queue pairs of the form <node, nr>. When we pop a node from the queue we set the data for the corresponding row from T.
 - o If I still want to use two function, the *addNumbers* can create a Map with elements <node, nr> or <tree, nr> and pass this map as a parameter to the *buildTable* function.
- What happens if I do not want recursive traversal for building the table?
 - Use a stack of a queue for the traversal.
- 3. Given a binary tree that represents the ancestors of a person up to the nth generation, where the left subtree represents the maternal line and the right subtree represents the paternal line:
 - a. Display all the females from the tree (root can be either male of female)

b. Display all ancestors of degree k (root has degree 0)

$$\circ$$
 K = 2 – z, w, t, s



a. Traverse the tree using a queue (or stack) and print only the left subtrees

```
Subalgorithm females (tree) is:
    init(q)
    if tree.root ≠ NIL then
        push (q, tree.root)
        print [tree.root].e
    end-if
    while ¬isEmpty(q) execute
        pop(q, p)
        if ([p].left # NIL) then
            print [[p].left].e
            push(q, [p].left)
        end-if
        if ([p].right # NIL) then
            push(q, [p].right)
        end-if
    end-while
end-subalgorithm
```

b. Recursive version – using the tree's interface (we do not care/do not use the representation of the tree)

```
Subalgorithm level(tree, k, v) is
// v is a vector in which we will add the elements from level k, assume
it has an insert operation that adds a new element.
    if ¬isEmpty(tree) then
        if k = 0 then
            insert(v, root(tree))
        else
            if ¬ isEmpty(left (tree)) then
                level(left (tree), k-1, v)
            end-if
            if ¬ isEmpty(right (tree)) then
                level(left (tree), k-1, v)
            end-if
        end-if
    end-if
end-subalgorithm
subalgorithm ancestors (tree, k, v) is:
    init (v) // initialize an empty vector
    level (tree, k, v)
    for i ← 1, dim(v) execute
        print element(v, i)
    end-for
end-subalgorithm
```

- O How can we solve the problem with a non-recursive function?
 - Put <node/tree, level> pairs in the queue
 - Or use two queues
 - Or count on the fact the tree should be complete (in real life you might have missing nodes/subtrees)