TÉCNICO LISBOA

Peer-to-Peer Systems and Overlay Networks

1º Semester – Academic Year 2013/2014

Kuganesan Srijeyanthan, 79531 - kuganesan.srijeyanthan@ist.utl.pt
Gayana Ranganatha Chandrasekara Pilana Withanage, 79529 - gayana.withanage@ist.utl.pt
Group 3 - Taguspark

Introduction

Peer to Peer distributed file system is an emerging method of utilizing commodity hardware with high degree of fault tolerance and scalability. Our group has researched, implemented some sample programs and tested with local virtual nodes in order to finalize the communication protocol for this application development. We have chosen Pastry [1, 2] as a main DHT service on top of virtual file system.

Protocol Selection and Design

Why Pastry

Refer the below table in order to compare some features of different DHT protocols and it will help to justify our selection of protocol PASTRY.

Features \ Protocol	Chord	Pastry	Kademlia
Routing	Log(n)	Log(n) but can be improved with neighbor and leaf nodes	Log(n)
Locality	No	Yes	No
Node Join Complexity	$Log^2(n)$	Log(n)	$Log^2(n)$
Open Source Implementation	Yes	Yes	Yes
File Systems related prior applications	No	Yes	No
Security Implementation	No	Yes	Yes

The following diagram clearly explains how file systems are going to be dispersed in to several nodes which are belongs to the Pastry ring. Let's assume UserA is going to send p2p_smaple.txt file via peer to peer network , once the file is created inside the fuse file system all the contents will be send via Pastry chunk by chunk, which we could be able set from configuration file.

DHT implementation of P2P distributed file system using Pastry

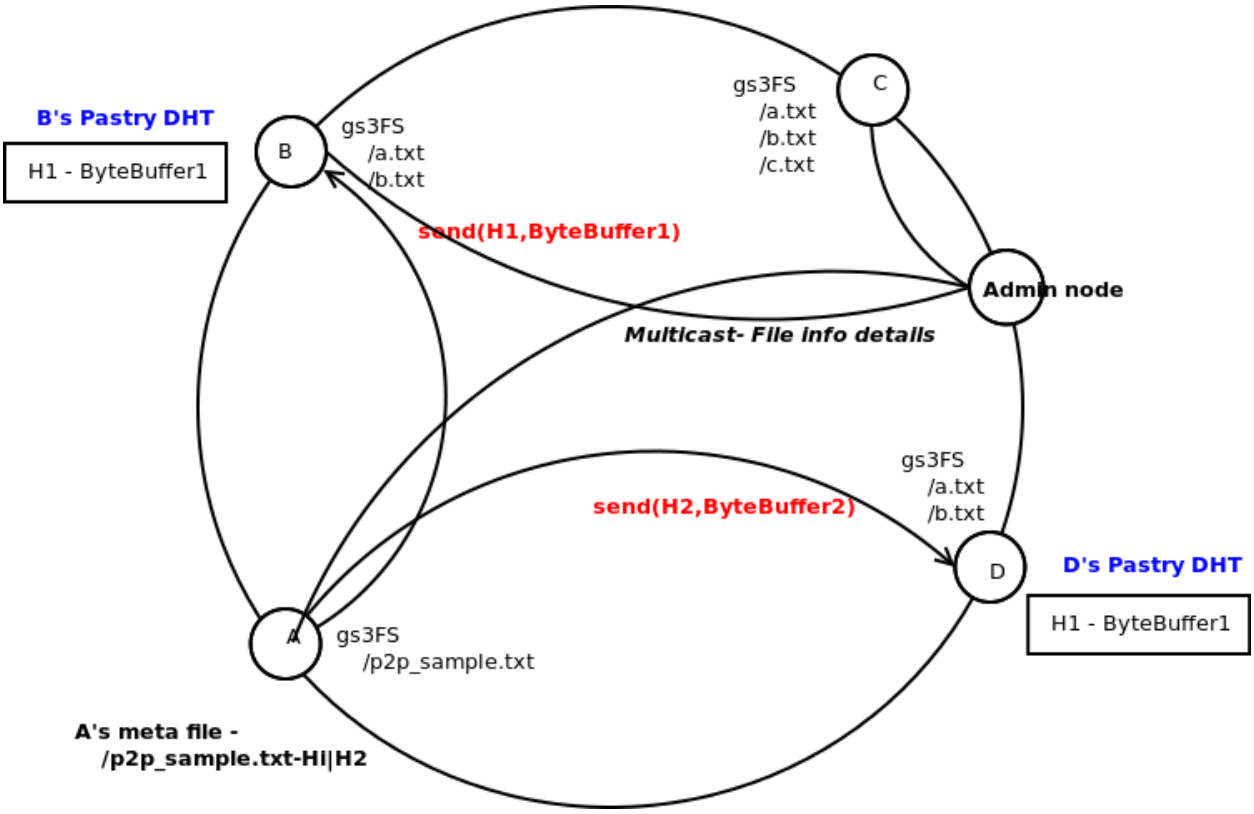


Fig.1 DHT architecture of P2P distributed file system

Architecture of the system

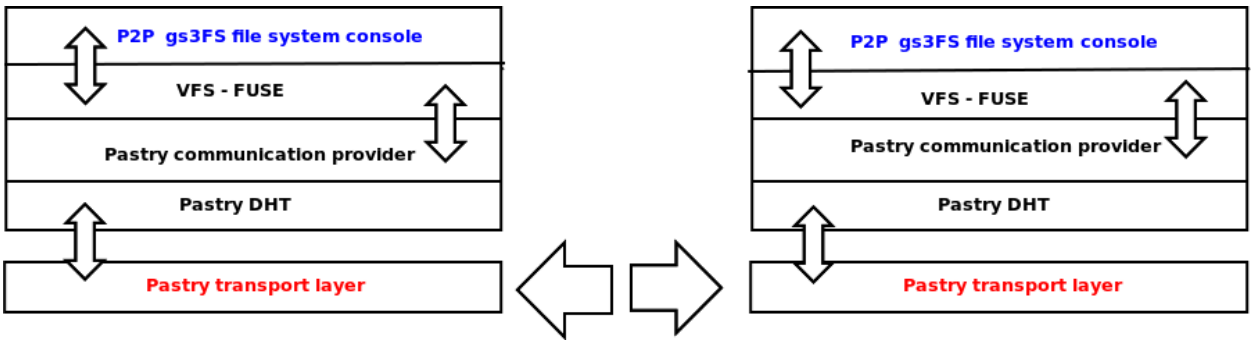


Fig.2 System architecture how message flows in between two nodes

Fig.2 shows the node architecture of P2P DFS implementation, all the operations are transparent to user and will be handled by Pastry communication provider module which we are going to implement. It will handle node entry, file management, and file replication according to Pastry protocol concepts.

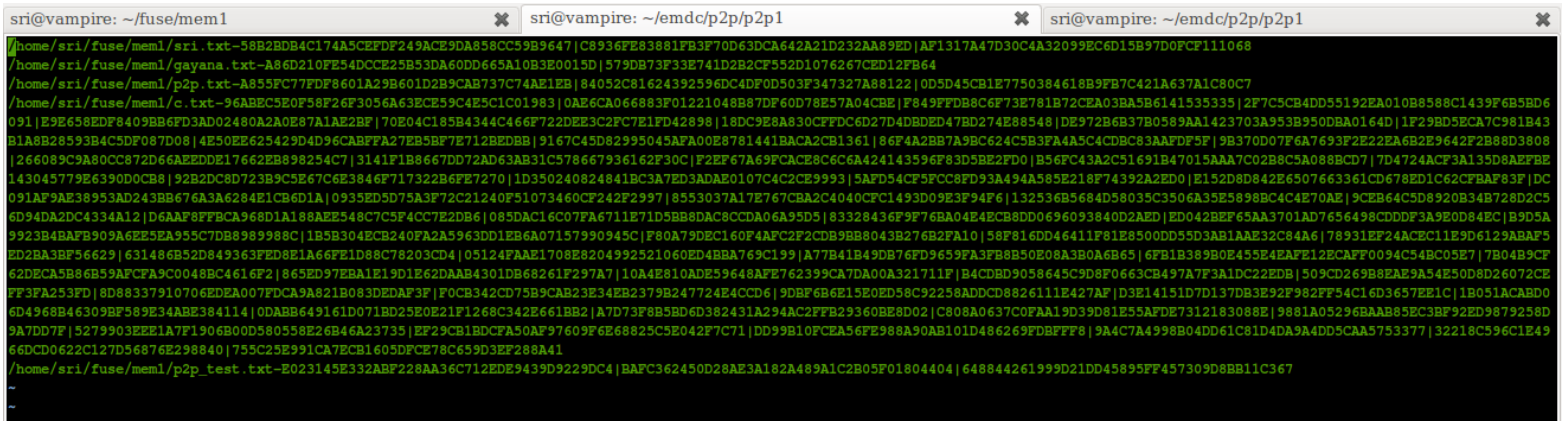
File replication management

File replication will be handled by the Pastry implementation according to the configuration value we provide. In our implementation we used the replication node count as 2 which will make 3 replicas where 2 in the leaf set and other in the node itself.

Metadata file

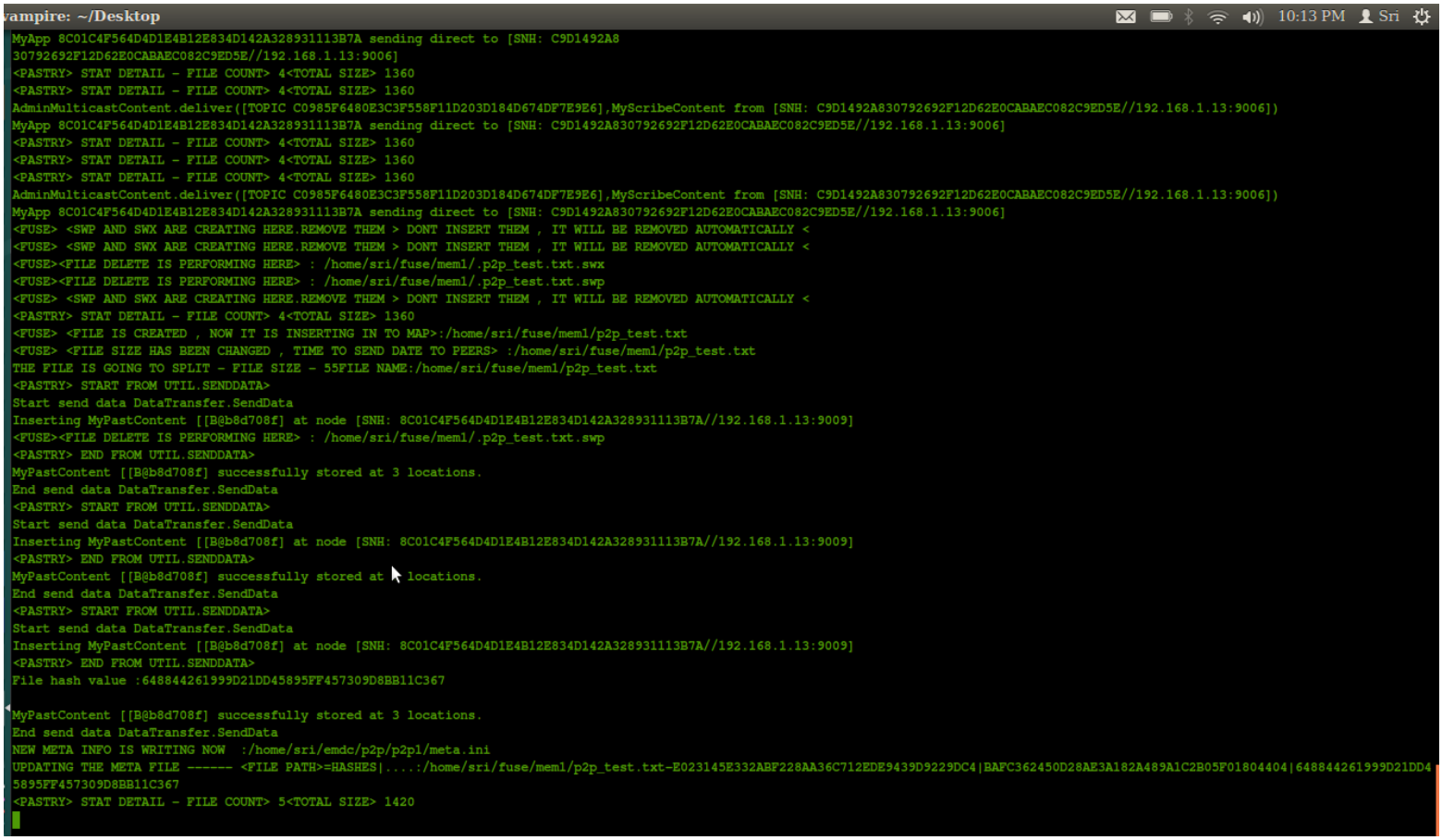
Each node has its own meta data file stored in to its persistent storage. Metafile will generate before the peer enter into the Pastry ring if the user is a new comer.
If a previous user joins again to the ring, application will check the availability of the Metafile and will request the files from peers and mount the files to the user's FUSE file system.
Metafile keep the records as <Key,Value> pairs where
Key = unique file path
Value = Pipe separated Hash value set of each file chunks of the desired size those user shared previously.

Below figure shows a sample Meta info file generated by our application.

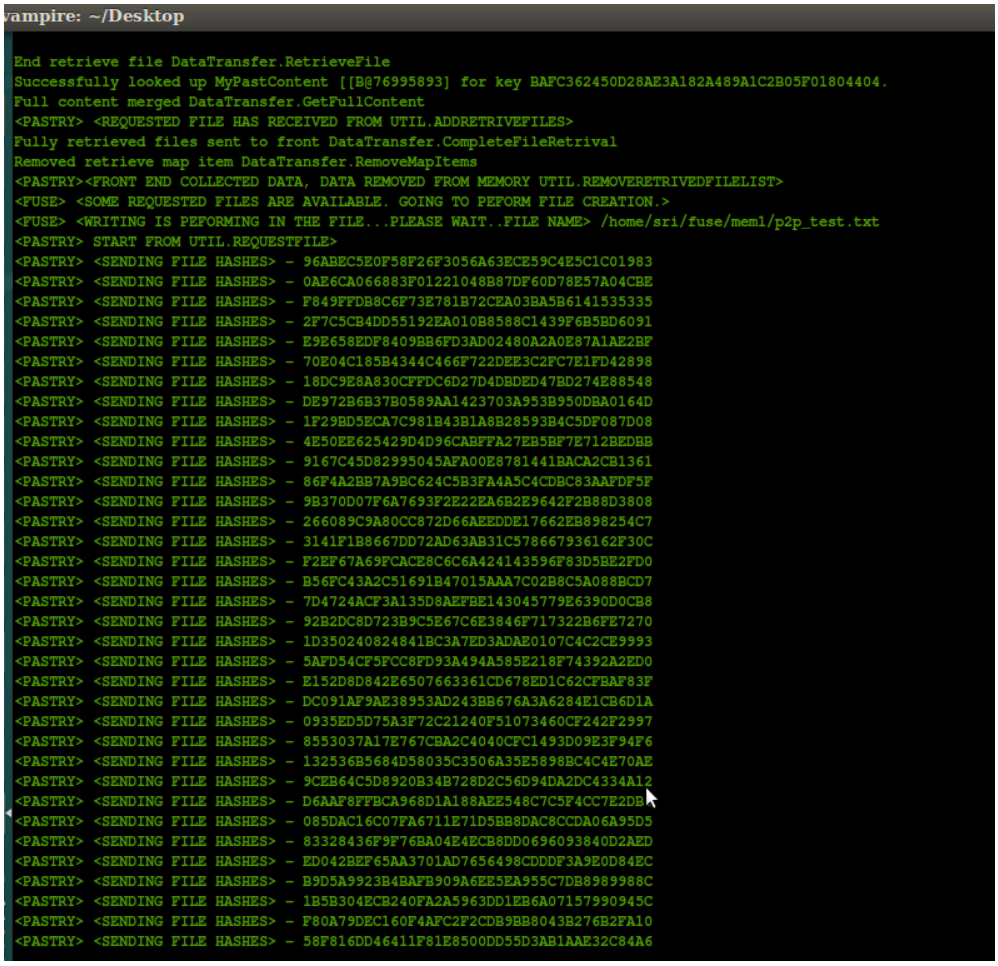


Steps of operation

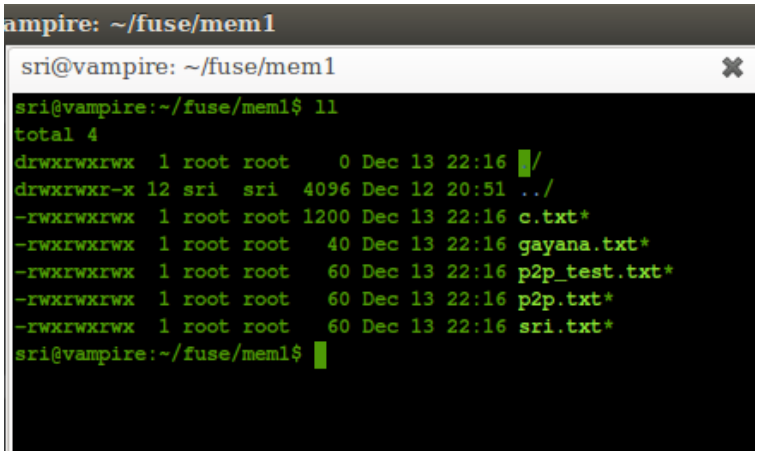
1. User executes the application and join the Pastry network. This will mount a FUSE file system for the user.
2. User can create files in the newly mounted file system. When user create and save the files, application will send the files, chunking them into pieces of pre-configured size to the Pastry DHT. Below image will show the console output provided when sending files.



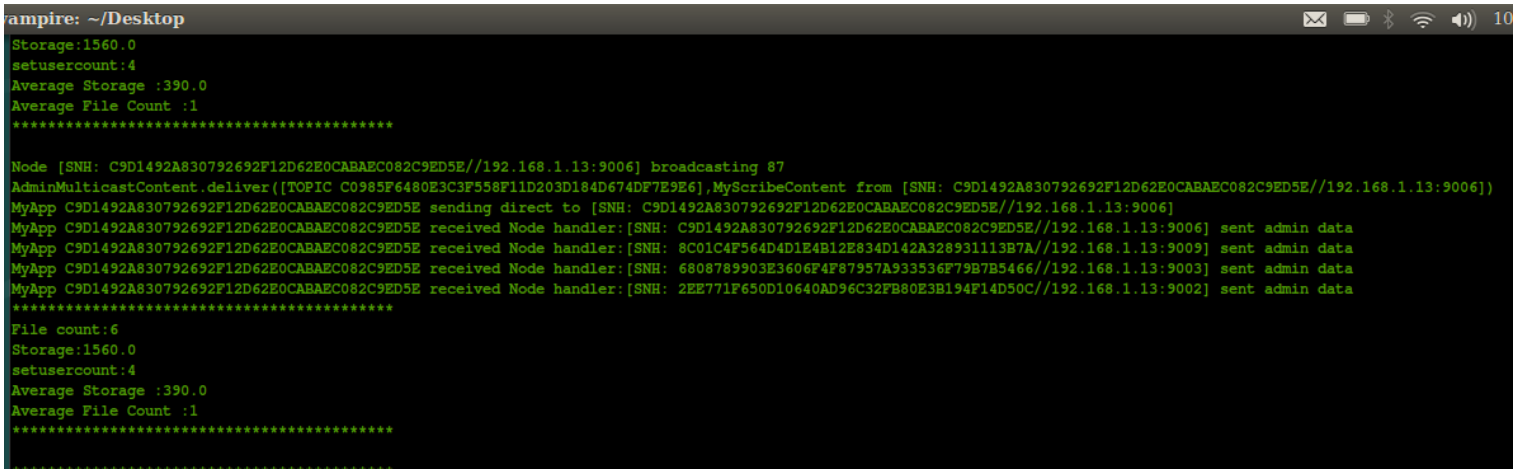
3. If the user update the file and save again new Hash values will be generated and above step 2 will be proceeded again.
4. When a previous user joined the network again, application will request the files in the metafile from the peers. Below figure shows the console output for the file request at start up.



5. After the retrieval of files from peers, user can see the files mounted in the FUSE file system. Below image shows how user see the file system.



6. The admin user can see the statistic information such as files shared in the network, active user count, total storage etc. Admin module will multicast messages time to time via the pastry network and collect above details. This will give approximate value for those information. Below image illustrates an admin view for information.



Test Observations

We generated 10 nodes with our application in a local newtwork and shared files in the network. We observed that the node joining consumed a considerable amount of time when the number of peers increased.

After joining the network the peers seemed to be communicating properly and in a stable manner. We removed upto 2 random nodes at a time but it did not make an impact to lose the files shared in the network. Once they rejoin they could retrieve all the files they shared earlier.

Note: To run the application we used below sample command line instruction

java -Djna.nosys=true -cp gs3FS.jar:lib/fusejan.jar:lib/fuse-jna-uber.jar:lib/Pastry_v1.jar:lib/xmlpull_1_1_3_4a.jar:lib/xpp3-1.1.3.4d_b2.jar

rice.uproject.implmeation.StartUp 9006 192.168.1.13 9002 /home/sri/fuse/mem1 /home/sri/emdc/p2p/p2p1

For admin console only use the below command line instruction.

java -Djna.nosys=true -cp gs3FS.jar:lib/fusejan.jar:lib/fuse-jna-uber.jar:lib/Pastry_v1.jar:lib/xmlpull_1_1_3_4a.jar:lib/xpp3-1.1.3.4d_b2.jar rice.uproject.implmeation.StartUp 9006 192.168.1.13 9002 /home/sri/fuse/mem1 /home/sri/emdc/p2p/p2p1 1

For the testing we have used the multicast rerun duration as 25 seconds. When the network is large this value should be increased in to a considerably higher value in order to control the Pastry network congestion.

References

[1] A. Rowstron and P. Druschel. Pastry: Scalable, distributed object location and routing for large-scale peer-to-peer systems. In Proc. IFIP/ACM Middleware, November 2001

[2] Antony Rowstron and Peter Druschel . Storage management and caching in PAST, a large-scale, persistent peer-to-peer storage utility