

a)

```

void f1(int n)
{
    int i=2;
    while(i < n){
        /* do something that takes O(1) time */  $\Theta(1)$ 
        i = i*i;  $\Theta(1)$ 
    }
}

```

i
 i^2
 i^3
 i^4
 \vdots
 $i^{k-1} < n$
 $i^k \geq n$ terminate
 $k = \log_i n$

$$T(n) = 2 \sum_{i=1}^{\log_2 n} \Theta(1) = \Theta(\log n)$$

b)

```

void f2(int n)
{
    for(int i=1; i <= n; i++){
        if( (i % (int)sqrt(n)) == 0){
            for(int k=0; k < pow(i,3); k++) {
                /* do something that takes O(1) time */  $\Theta(1)$ 
            }
        }
    }
}

```

same as example in class

$$T(n) = \sum_{i=1}^n (\Theta(1) + O(\sum_{k=0}^{i^3} \Theta(1))) = \sum_{i=1}^n \Theta(1) + \underbrace{\sum_{i=1}^n \sum_{k=0}^{i^3} \Theta(1)}_{\sum_{i=1}^n \Theta(i^3)} = \Theta(n) + \sum_{k=1}^{\sqrt{n}} \Theta(n^3) = \Theta(n) + \Theta(n^3 \sqrt{n})$$

$\sum_{i=1}^n \Theta(i^3)$
 \hookrightarrow upper bound n^3

$$T(n) = \Theta(n^3 \sqrt{n})$$

c)



```

for(int i=1; i <= n; i++){
    for(int k=1; k <= n; k++){
        if( A[k] == i){
            for(int m=1; m <= n; m=m+m){
                // do something that takes O(1) time
                // Assume the contents of the A[] array are not cha
            }
        }
    }
}

```

 $\left. \begin{array}{l} \Theta(1) \\ \Theta(n/2) \end{array} \right\}$

$$\begin{array}{r}
 m \\
 m \\
 2m \\
 4m \\
 8m \\
 \vdots \\
 2 \cdot Km \geq n \\
 k = \frac{n}{2m} = \frac{n}{2} \cdot \frac{1}{m}
 \end{array}$$

$$T(n) = \sum_{i=0}^{n-1} \sum_{k=0}^{n-1} (\Theta(1) + O(n/2)) = \sum_{i=0}^{n-1} \sum_{k=0}^{n-1} \Theta(1) + \sum_{i=0}^{n-1} \sum_{k=0}^{n-1} O(n/2) = \sum_{i=0}^{n-1} \Theta(n) + \sum_{i=0}^{n-1} \Theta(n^2/2) = \Theta(n^2) + \Theta(n^3/2) = \Theta(n^3/2)$$

d)



```

int f (int n)
{
    int *a = new int [10];
    int size = 10;
    for (int i = 0; i < n; i ++)  $\Theta(n)$ 
    {
        if (i == size)  $\Theta(1)$ 
        {
            int newsize = 3*size/2;
            int *b = new int [newsize];
            for (int j = 0; j < size; j ++) b[j] = a[j];  $\Theta(size)$ 
            delete [] a;
            a = b;
            size = newsize;
        }
        a[i] = i*i;
    }
}

```

$$\begin{array}{l}
 \Theta(1) \left\{ \begin{array}{l} \text{int newsize} = 3 \cdot \text{size} / 2; \\ \text{int } *b = \text{new int} [\text{newsize}]; \\ \text{for (int j = 0; j < size; j ++)} \text{ b[j] = a[j]; } \\ \text{delete [] a;} \\ \text{a = b;} \\ \text{size = newsize;} \end{array} \right. \Theta(\text{size})
 \end{array}$$

$$\begin{array}{l}
 \frac{3}{2}k \cdot \text{size} \leq n \\
 k = \frac{2n}{3 \cdot \text{size}} \approx n
 \end{array}$$

$$T(n) = \Theta(1) + \sum_{i=0}^{n-1} \sum_{j=0}^{n-1} \Theta(1) = \Theta(n^2)$$