

Abstract

Some exercises from http://brendanfong.com/programmingcats_files/ps1.pdf

Exercise 2

a) $Ob(\mathcal{C}) = 1, 2$; The sets of morphisms: $\mathcal{C}(1, 1), \mathcal{C}(1, 2), \mathcal{C}(2, 1), \mathcal{C}(2, 2)$. The composition rule $f \circ id_1 = id_2 \circ f : 1 \rightarrow 2$. The identity morphisms: id_1 and id_2 .

b) Unit: $f \circ id_1 = f$ and $id_2 \circ f = f$. Associative: $f \circ id_1 = id_2 \circ f$.

Exercise 3

If there are no identity morphisms id_c and id_d , then we can't affirm an isomorphism with $g = f^{-1}$, given that we can't prove $g \circ f = id_c$ and $f \circ g = id_d$?

Exercise 4

a) Following associative, but not unity law:

$$1 \begin{array}{c} \xrightarrow{f} \\ \xleftarrow{g} \end{array} 2$$

Here if $h = f \circ g$, we have $(f \circ g) \circ h = f \circ (g \circ h) = f$, but since there aren't any identity morphisms unity is not followed.

b) Following unit, but not associative law: Can build free categories where each object has an identity morphism?

Exercise 5

a) The carrier set is \mathbb{N} , operation or function is $+$ and the identity element 0. Unity is followed as any $n \in \mathbb{N} = n$, when $e * n$ or $n * e$, given that $n + 0 = n$. Associativity follows given that addition is commutative.

b) Here $List_{\{0,1\}}$ is the carrier set, concatenation the operation, and empty string \square the identity. Unity is followed given that concatenating an empty string to an initial string s will output the s . Associativity isn't followed? Order of concatenation matters?

c) This is because a monoid is a category that only has one homset. Consider the \mathbb{N} carrier set in a). All morphisms in this monoid are of the type $\mathbb{N} \rightarrow \mathbb{N}$, with the corresponding monoid operation and identity element. This is the case for any monoid independent of the triple (M, e, \diamond) .

Exercise 6

a) $id_{12} = 1$ given that for $x * 12 = 12$ only $x = 1$ suffices.

b) If $x : a \rightarrow b$ and $y : b \rightarrow c$ then $y \circ x : a \rightarrow c$ implies that if b is divisible by a and c is divisible by b , then c is divisible by a , which follows given that division is transitive.

c) If \mathcal{P} would be given by all \mathbb{N} , then the previous arguments wouldn't work because division by 0 is undefined.