Algorithms and Data Structures Homework 2, Problem 2

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f is of order at most g, written f(n) is O(g(n)) or $f(n) \in O(g(n))$ if and only if there exists a positive real k and nonnegative real number N such that $f(n) \leq k g(n)$ for all real numbers n > N.

Intuitively, the definition that f(n) is O(g(n)) says that f(n) grows slower than some fixed multiple of g(n) as n grows without bound. We can also use $f \in O(g)$ if and only if $\lim_{n \to \infty} \frac{f(n)}{g(n)} = c$ where c is finite.

The constant N and k in the definition above are called witnesses to the relationship f(n) is O(g(n)).

① $4n < n^9$ for all n72; witnesses are N=2 and k=1. Also

 $\lim_{n\to\infty} \frac{f(n)}{g(n)} = \lim_{n\to\infty} \frac{7n}{n^7} = 0$, finite. Thus $7n \in O(n^7)$ or $f(n) \in \mathcal{G}(0(g(n)))$

1
$$\lim_{n\to\infty} \frac{f(n)}{g(n)} = \frac{n^2}{n \cdot \log_2 n} = \frac{n}{\log_2 n}$$
 (Then apply L'Hapital's

rule) =
$$\lim_{n\to\infty} \left(\frac{1}{\frac{1}{n\ln 2}}\right) = \lim_{n\to\infty} \ln(n) = \infty$$
.

Thus,
$$n^2/\log_2 n \notin O(n)$$
 but $n \in O(n^2/\log_2 n)$ or $g(n) \in O(f(n))$.

$$\frac{3}{1 + \infty} \frac{f(n)}{g(n)} = \frac{27 \cdot (\log_2(n))}{(\log_2(u_n))^2} = \frac{27 \log_2(n)}{(\log_2(n)+2)^2} = \frac{17 \log_2(n)}{(\log_2(n))^2 + 4 \log_2(n+4)}$$

$$= 27 \cdot \lim_{n \to \infty} \left(\frac{\log_2(n)}{\log_2(n)^2 + 4\log_2(n) + 4} \right) = 27 \lim_{n \to \infty} \left(\frac{\log_2(n)}{1 + \log_2(n)} + \frac{4}{(\log_2(n))^2} \right)$$

$$= 27 \cdot \frac{\lim_{n \to \infty} \left(\frac{1}{\log_2(n)}\right)}{\lim_{n \to \infty} \left(\frac{1}{\log_2(n)}\right)} = 0.$$

$$\lim_{n \to \infty} \left(\frac{1}{\log_2(n)}\right) = 0.$$

$$= 27 \cdot \frac{0}{1} = 0$$
 Thus $\frac{1}{1}$



1 m (1+4 +4 m)=1 =1

 $f(n) \in O(g(n))$.

(4)
$$\lim_{n\to\infty} \frac{f(n)}{g(n)} = \lim_{n\to\infty} \lim_{n\to\infty} \frac{f(n)}{5h^2 + n^{0.3} + n^{0.$$