Sparse automatic differentiation

From theory to practice

Guillaume Dalle* - LVMT, École des Ponts (gdalle.github.io)

Inria Paris, 07.03.2025

^{*}joint work with Adrian Hill & Alexis Montoison

Agenda

- 1. Introduction
- 2. Automatic differentiation
- 3. Exploiting sparsity
- 4. Pattern detection and coloring
- 5. Implementation
- 6. Conclusion

Introduction

Newton's method

Root-finding

Solve F(x) = 0 by iterating

$$x_{t+1} = x_t - \underbrace{\left[\partial F(x_t)\right]^{-1} F(x_t)}_{\text{jacobian}}$$

Optimization

Solve $\min f(x)$ by iterating

$$x_{t+1} = x_t - \underbrace{\left[\nabla^2 f(x_t)\right]^{-1}}_{\text{hessian}} \nabla f(x_t)$$

Linear system involving a derivative matrix A.

Implicit differentiation

Differentiate $x \to y(x)$ knowing **optimality conditions** c(x, y(x)) = 0.

Applications: fixed-point iterations, optimization problems.

Implicit function theorem²

$$\partial_1 c(x, y(x)) + \partial_2 c(x, y(x)) \cdot \partial y(x) = 0$$

$$\partial y(x) = -\underbrace{\left[\partial_2 c(x,y(x))\right]^{-1}}_{\text{jacobian}} \partial_1 c(x,y(x))$$

Linear system involving a derivative matrix A.

²Blondel et al. (2022)

Conventional wisdom

- · Jacobian and Hessian matrices are too large to compute or store
- We can only access lazy maps $u \mapsto Au$ (JVPs, VJPs, HVPs³)
- Linear systems $A^{-1}v$ must be solved with iterative methods
- Downsides: each iteration is expensive, convergence is tricky

³Dagréou et al. (2024)

The benefits of sparsity

- Jacobian and Hessian matrices have mostly zero coefficients
- We can compute and store A explicitly
- Linear systems $A^{-1}v$ can be solved with iterative or direct methods
- Upsides: faster iterations, or even exact solves

Automatic differentiation

Pocket AD

```
import Base: +, * # overload standard operators
struct Dual
    val::Float64
    der::Float64
end
+(x::Dual, y::Dual) = Dual(x.val + y.val, x.der + y.der)
*(x::Dual, y::Dual) = Dual(x.val * y.val, x.der*y.val + x.val*y.der)
+(x, y :: Dual) = Dual(x, 0) + y
\star(x, y :: Dual) = Dual(x, 0) \star y
```

Pocket AD

Does it work?

```
julia > f(x) = 1 + 2 * x + 3 * x * x;
julia> f(4)
57
julia> f(Dual(4, 1)) # exact derivative
Dual(57.0, 26.0)
julia> (f(4 + 1e-5) - f(4)) / 1e-5 # approximate derivative
26.000029998840542
```

What is AD?

input	output
program to compute the function $x \longmapsto f(x)$	program to compute the differential
	$x \longmapsto \partial f(x)$
	which is a linear map $u\mapsto \partial f(x)[u]$

Two ingredients only:

- 1. hardcode basic derivatives (+, \times , exp, \log , ...)
- 2. handle compositions $f = g \circ h$

Composition

For a function $f = g \circ h$, the chain rule gives

standard
$$\partial f(x) = \partial g(h(x)) \circ \partial h(x)$$

adjoint $\partial f(x)^* = \partial h(x)^* \circ \partial g(h(x))^*$

These linear maps apply as follows:

forward
$$\partial f(x): u \xrightarrow{\partial h(x)} v \xrightarrow{\partial g(h(x))} w$$

reverse $\partial f(x)^*: u \xleftarrow{\partial h(x)^*} v \xleftarrow{\partial g(h(x))^*} w$

Forward chain rule, illustrated

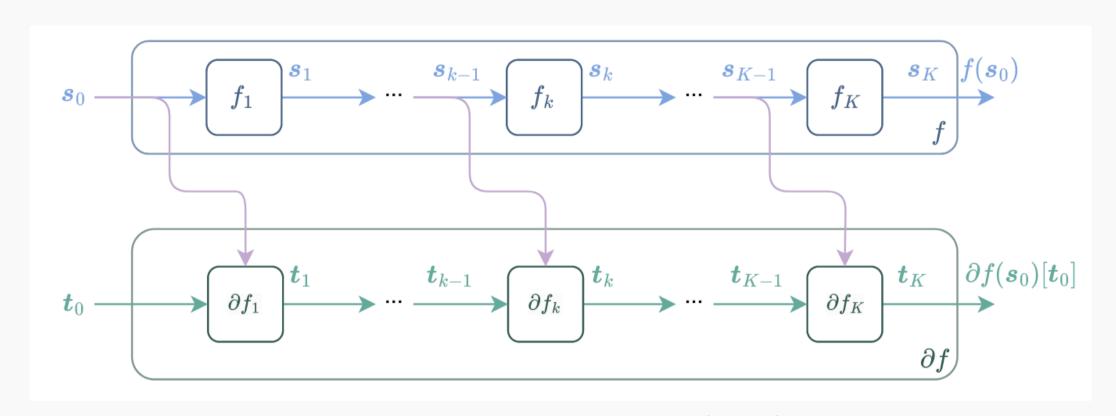


Figure 1: Blondel & Roulet (2024)

Reverse chain rule, illustrated

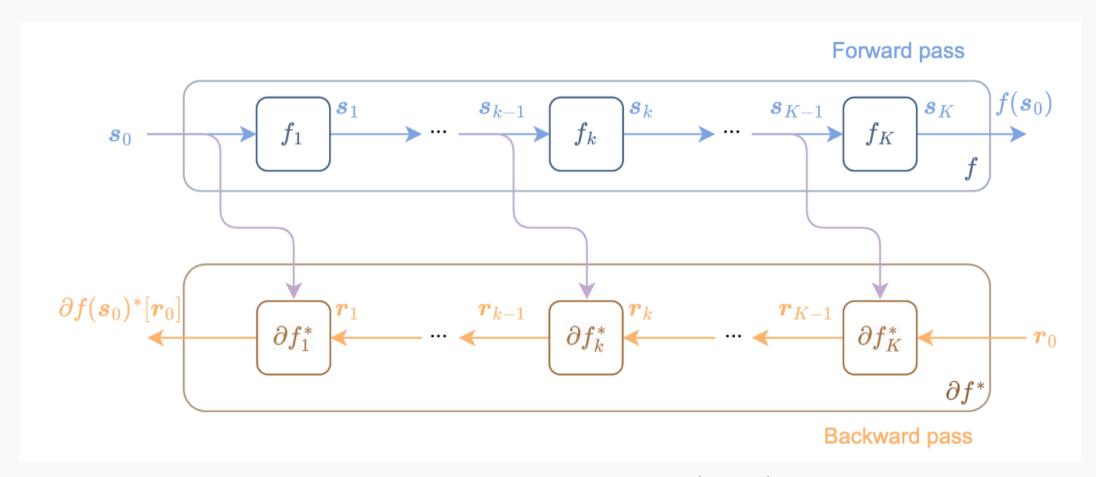


Figure 2: Blondel & Roulet (2024)

Two modes

Forward-mode AD computes Jacobian-Vector Products:

$$u \longmapsto \partial f(x)[u] = Ju$$

Reverse-mode AD computes Vector-Jacobian Products:

$$w \longmapsto \partial f(x)^*[w] = w^*J$$

No need to materialize intermediate Jacobian matrices!

Theorem: cost of 1 JVP or VJP \propto cost of 1 function evaluation

Interpretations

- Forward mode: "pushforward" of an input perturbation
- Reverse mode: "pullback" of an output sensitivity

Reverse mode gives gradients for roughly the same cost as the function itself:

$$\nabla f(x) = \partial f(x)^*[1]$$

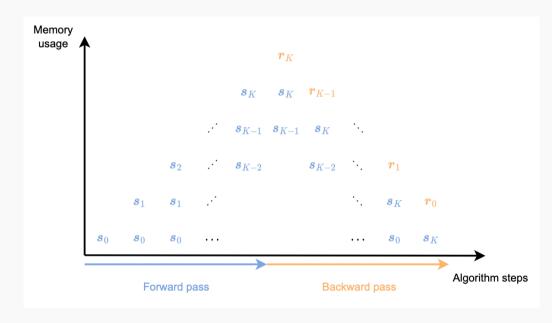


Figure 3: The devil is in the details (Blondel & Roulet, 2024)

Pocket AD, chain rule version

```
# Basic rules
using LinearAlgebra
A, b = rand(2, 3), rand(2)
residuals(x) = A * x - b
\partial(::typeof(residuals)) = x \rightarrow (u \rightarrow A * u) # \mathbb{R}^3 \rightarrow \mathbb{R}^2
\partial^{\mathsf{T}}(:: \mathsf{typeof}(\mathsf{residuals})) = \mathsf{X} \to (\mathsf{V} \to \mathsf{adjoint}(\mathsf{A}) \star \mathsf{V}) \# \mathbb{R}^2 \to \mathbb{R}^3
sgnorm(r) = sum(abs2, r)
\delta(::typeof(sqnorm)) = r \rightarrow (v \rightarrow dot(2r, v)) \# \mathbb{R}^2 \rightarrow \mathbb{R}
\partial^{\mathsf{T}}(::\mathsf{typeof}(\mathsf{sqnorm})) = r \to (\mathsf{w} \to 2r .* \mathsf{w}) \# \mathbb{R} \to \mathbb{R}^2
```

Pocket AD, chain rule version

```
# Composition
function ∂(f::ComposedFunction)
     g, h = f.outer, f.inner
     return x \to \partial(g)(h(x)) \cdot \partial(h)(x)
end
function ∂<sup>T</sup>(f::ComposedFunction)
     g, h = f.outer, f.inner
     return x \to \partial^{T}(h)(x) \cdot \partial^{T}(g)(h(x))
end
```

Pocket AD, chain rule version

```
julia > import ForwardDiff as FD, Zygote
julia> f = sqnorm ∘ residuals;
julia> x, \Delta x = rand(3), [1, 0, 0];
julia> \partial(f)(x)(\Delta x) # partial derivative
                                                       julia> FD.derivative(t \rightarrow f(x + t * \Deltax), 0)
0.8691056836969242
                                                       0.8691056836969242
julia> \partial^{\mathsf{T}}(\mathsf{f})(\mathsf{x})(1) # gradient
                                                       julia> Zygote.gradient(f, x)[1]
3-element Vector{Float64}:
                                                       3-element Vector{Float64}:
0.8691056836969242
                                                        0.8691056836969242
                                                        0.9973491983376236
0.9973491983376236
 0.5768822265195823
                                                        0.5768822265195823
```

Exploiting sparsity

From maps to matrices

To compute the Jacobian matrix J of a composition $f: \mathbb{R}^m \longrightarrow \mathbb{R}^n$:

- product of intermediate Jacobian matrices
- reconstruction from several JVPs or VJPs

	forward mode	reverse mode
idea	1 JVP gives 1 column	1 VJP gives 1 row
formula	$J_{\cdot,j} = \partial f(x) \left[e_j \right]$	$J_{i,\cdot} = \partial f(x)^*[e_i]$
cost	n JVPs (input dimension)	m JVPs (output dimension)

Using fewer products

When the Jacobian is sparse, we can compute it faster4.

If columns $j_1, ..., j_k$ of J are structurally orthogonal (their nonzeros never overlap), we deduce them all from a single JVP:

$$J_{j_1} + \ldots + J_{j_k} = \partial f(x) \left[e_{j_1} + \ldots + e_{j_k} \right]$$

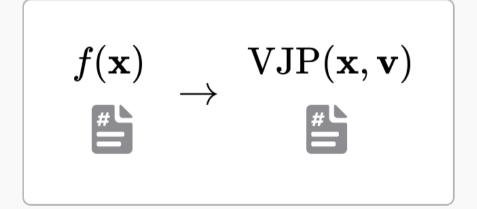
Once we have grouped columns, sparse AD has two steps:

- 3. one JVP for each group $c = \{j_1, ..., j_k\}$
- 4. decompression into individual columns $j_1,...,j_k$

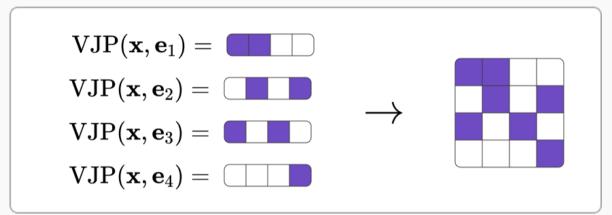
⁴Curtis et al. (1974)

The gist in one slide

(a) AD code transformation



(b) Standard AD Jacobian computation



(c) ASD Jacobian computation

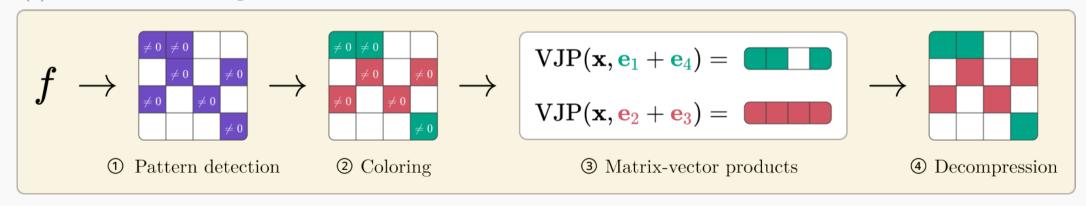


Figure 4: Hill & Dalle (2025)

Two preliminary steps

When grouping columns, we want to

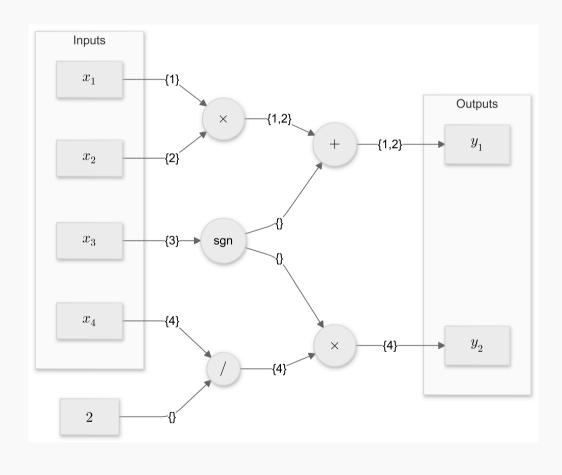
- guarantee structural orthogonality (correctness)
- form the smallest number of groups (efficiency)

preparation	execution	
1. pattern detection	3. matrix-vector products	
2. coloring	4. decompression	

The preparation phase can be amortized across several inputs.

Pattern detection and coloring

Tracing dependencies in the computation graph



Computation graph for

$$y_1 = x_1 x_2 + \operatorname{sign}(x_3)$$

$$y_2 = \operatorname{sign}(x_3) \times \left(\frac{x_4}{2}\right)$$

Its Jacobian will have 3 nonzero coefficients.

Pocket pattern detection

```
import Base: +, *, /, sign
struct Tracer
  indices::Set{Int}
end
Tracer() = Tracer(Set{Int}())
+(a::Tracer, b::Tracer) = Tracer(a.indices ∪ b.indices)
*(a::Tracer, b::Tracer) = Tracer(a.indices ∪ b.indices)
/(a::Tracer, b) = Tracer(a.indices)
sign(a::Tracer) = Tracer() # zero derivatives
```

Pocket pattern detection

Does it work? julia> f(x) = [x[1] * x[2] * sign(x[3]), sign(x[3]) * x[4] / 2];julia > x = Tracer.(Set.([1, 2, 3, 4]))4-element Vector{Tracer}: Tracer(Set([1])) Tracer(Set([2])) Tracer(Set([3])) Tracer(Set([4])) julia> f(x) 2-element Vector{Tracer}: Tracer(Set([2, 1])) Tracer(Set([4]))

Partitions of a matrix

Orthogonal for all (i, j) s.t. $A_{ij} \neq 0$,

• column j is alone in group c(j) with a nonzero in row i

Symmetrically orthogonal for all (i, j) s.t. $A_{ij} \neq 0$,

- either column j is alone in group c(j) with a nonzero in row i
- or column i is alone in group c(i) with a nonzero in row j

Each partition can be reformulated as a specific coloring problem⁵.

⁵Gebremedhin et al. (2005)

Graph representations of a matrix

Column intersection $(j_1,j_2) \in \mathcal{E} \iff \exists i, A_{ij_1} \neq 0 \text{ and } A_{ij_2} \neq 0$ Bipartite $(i,j) \in \mathcal{E} \iff A_{ij} \neq 0$ (2 vertex sets \mathcal{I} and \mathcal{J}) Adjacency (sym.) $(i,j) \in \mathcal{E} \iff i \neq j \ \& \ A_{ij} \neq 0$

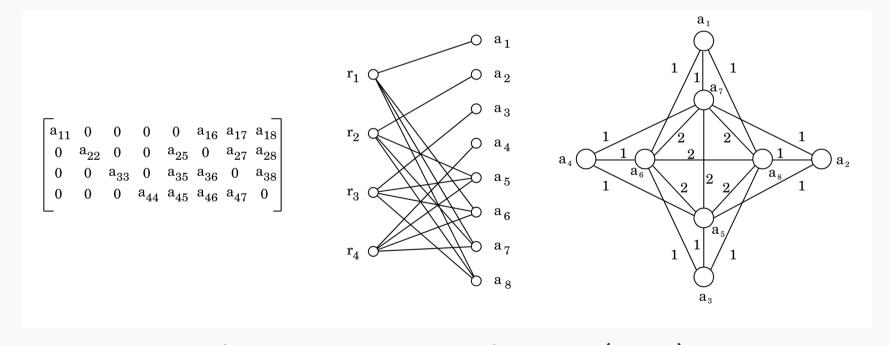


Figure 5: Gebremedhin et al. (2005)

Jacobian coloring

Coloring of intersection graph / distance-2 coloring of bipartite graph

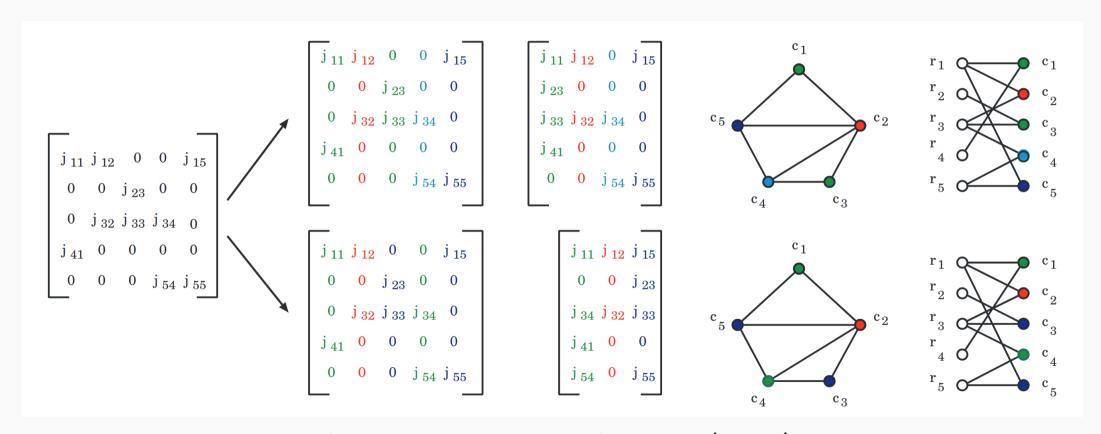


Figure 6: Gebremedhin et al. (2005)

Hessian coloring

Star coloring of adjacency graph

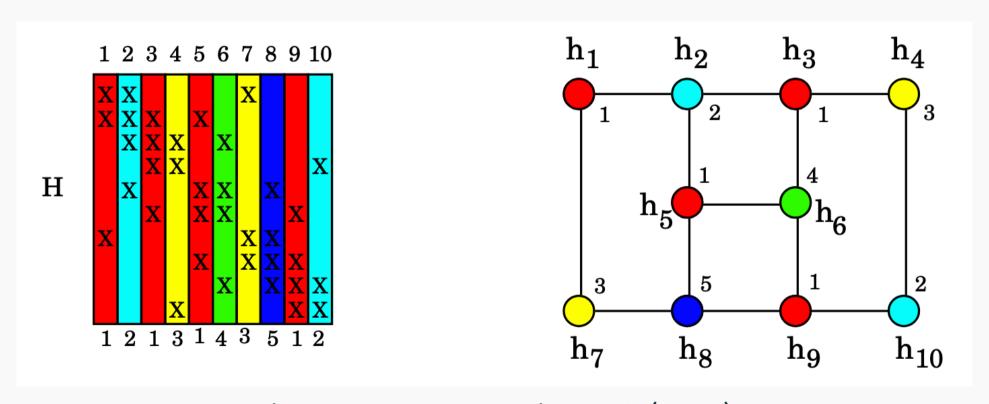


Figure 7: Gebremedhin et al. (2009)

Hessian coloring

Why a "star" coloring⁶? Consider

$$A = \begin{pmatrix} A_{kk} & A_{ki} & \cdot & \cdot \\ A_{ik} & A_{ii} & A_{ij} & \cdot \\ \cdot & A_{ji} & A_{jj} & A_{jl} \\ \cdot & \cdot & A_{lj} & A_{ll} \end{pmatrix}$$

If coloring c yields a symmetrically orthogonal partition:

- $c(i) \neq c(j)$
- $c(i) \neq c(k)$
- $c(j) \neq c(l)$

Any path on 4 vertices (i, j, k, l) must use at least 3 colors \iff any 2-colored subgraph is a collection of disjoint stars (it contains no path longer than 3).

⁶Coleman & Moré (1984)

Jacobian bicoloring

Bidirectional coloring of bipartite graph, with neutral color

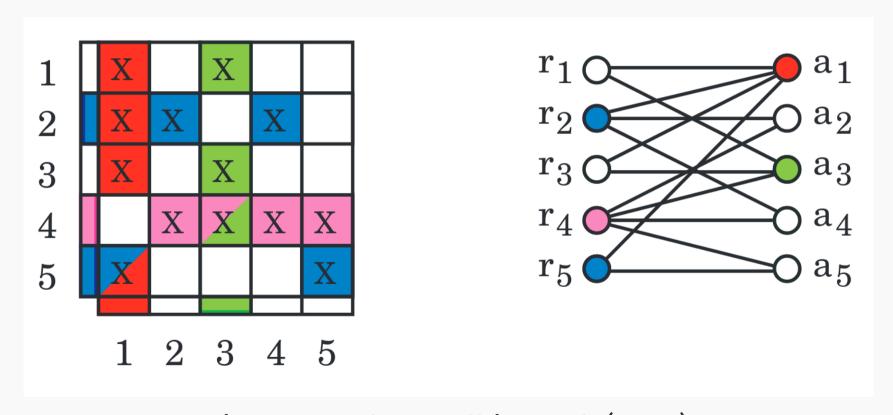


Figure 8: Gebremedhin et al. (2005)

Bicoloring from symmetric coloring [new]

To color the rows and columns of J, color the columns of $H=\left(egin{smallmatrix} 0 & J \ J & 0 \end{smallmatrix}\right)$

It sounds simple, but:

- Some colors may be redundant
- Detecting these is tightly linked to the two-colored structures
- Efficient decompression requires lots of preprocessing

The sharp bits

Pattern detection

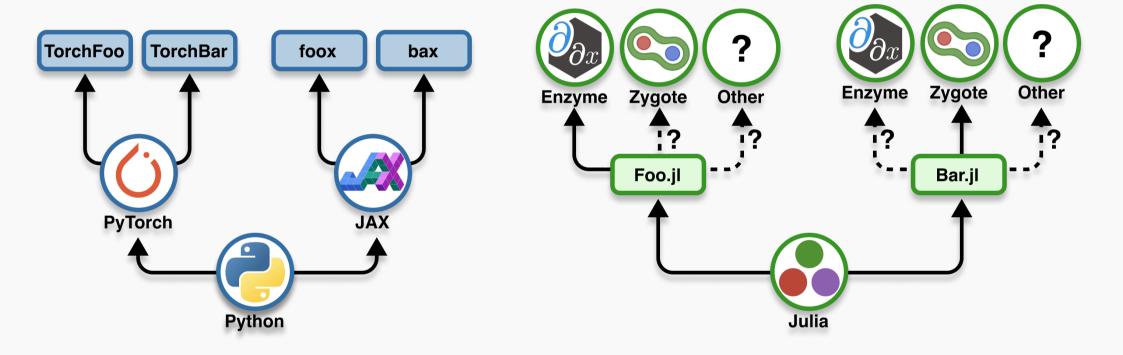
- Linear versus nonlinear interactions
- Local versus global sparsity

Coloring

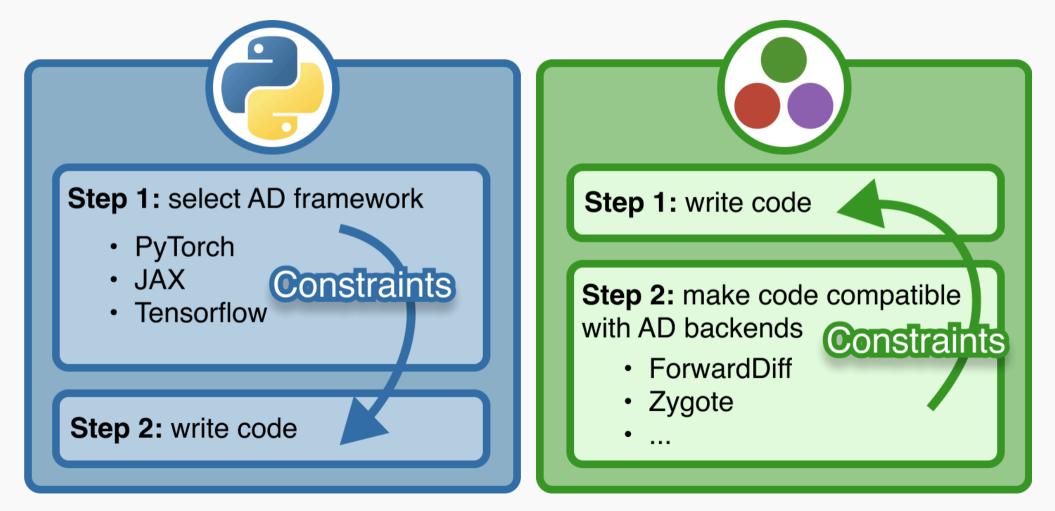
- Only heuristic algorithms
- Vertex ordering matters a lot

Implementation

AD in Python & Julia



AD in Python & Julia



Interfaces for experimenting [new]

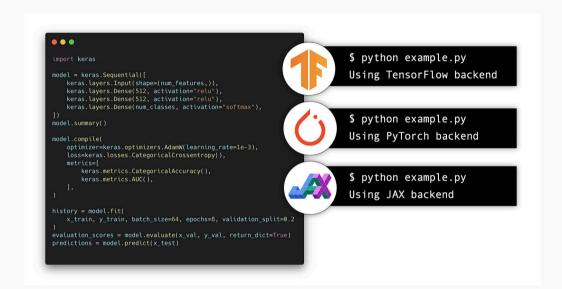


Figure 11: In Python, Keras supports Tensorflow, PyTorch and JAX.



Figure 12: In Julia, 14 AD backends inside Differentiationterface.jl

Once we have a common syntax, we can do more!

Previous implementations of sparse AD

- In low-level programming languages (C, Fortran)
- In closed-source languages (Matlab)
- In domain-specific languages (AMPL, CasADi)

Basically nothing in Python (either in JAX or PyTorch).

First drafts in Julia for scientific machine learning, but severely limited: single-backend, slow.

A modern sparse AD ecosystem [new]

Independent packages working together:

- Step 1: SparseConnectivityTracer.jl
- Steps 2 & 4: SparseMatrixColorings.jl
- Step 3: Differentiationterface.jl

	SCT.jl	SMC.jl	DI.jl
lines of code	4861	4242	16971
indirect dependents	420	437	426
downloads / month	4.2k	16k	20k

Compatible with generic code!

Impact

Users already include...

- Scientific computing: SciML (Julia's scipy)
 - Differential equations
 - Nonlinear solvers
 - Optimization
- Probabilistic programming: Turing.jl
- Symbolic regression: PySR

Live demo

This is the part where things go sideways.

Conclusion

Perspectives

- GPU-compatible pattern detection and coloring
- Adaptation in JAX with program transformations
- · New, unsuspected applications "just because we can"

Going further

On general AD:

- Baydin et al. (2018)
- Margossian (2019)
- Blondel & Roulet (2024)

On sparse AD:

- Gebremedhin et al. (2005)
- Griewank & Walther (2008)
- Hill & Dalle (2025)

Bibliography

- Baydin, A. G., Pearlmutter, B. A., Radul, A. A., & Siskind, J. M. (2018). Automatic Differentiation in Machine Learning: a Survey. *Journal of Machine Learning Research*, 18(153), 1–43. http://jmlr.org/papers/v18/17-468.html
- Blondel, M., & Roulet, V. (2024, July). The Elements of Differentiable Programming. arXiv. https://doi.org/10.48550/arXiv.2403.14606
- Blondel, M., Berthet, Q., Cuturi, M., Frostig, R., Hoyer, S., Llinares-Lopez, F., Pedregosa, F., & Vert, J.-P. (2022). Efficient and Modular Implicit
 Differentiation. Advances in Neural Information Processing Systems, 35, 5230–5242. https://proceedings.neurips.cc/paper_files/paper/2022/hash/228b9279ecf9bbafe582406850c57115-Abstract-Conference.html
- Coleman, T. F., & Moré, J. J. (1984). Estimation of sparse hessian matrices and graph coloring problems. *Mathematical Programming*, 28(3), 243–270. https://doi.org/10.1007/BF02612334
- Curtis, A. R., Powell, M. J. D., & Reid, J. K. (1974). On the Estimation of Sparse Jacobian Matrices. *IMA Journal of Applied Mathematics*, 13(1), 117–119. https://doi.org/10.1093/imamat/13.1.117
- Dagréou, M., Ablin, P., Vaiter, S., & Moreau, T. (2024, February). How to compute Hessian-vector products?. *The Third Blogpost Track at ICLR 2024*. https://openreview.net/forum?id=rTgjQtGP3O
- Gebremedhin, A. H., Tarafdar, A., Pothen, A., & Walther, A. (2009). Efficient Computation of Sparse Hessians Using Coloring and Automatic Differentiation. *INFORMS Journal on Computing*, 21(2), 209–223. https://doi.org/10.1287/ijoc.1080.0286
- Gebremedhin, A. H., Manne, F., & Pothen, A. (2005). What Color Is Your Jacobian? Graph Coloring for Computing Derivatives. SIAM Review, 47(4), 629–705. https://doi.org/10/cmwds4
- Griewank, A., & Walther, A. (2008). Evaluating derivatives: principles and techniques of algorithmic differentiation (2nd ed). Society for Industrial and Applied Mathematics. https://epubs.siam.org/doi/book/10.1137/1.9780898717761
- Hill, A., & Dalle, G. (2025, January). Sparser, Better, Faster, Stronger: Efficient Automatic Differentiation for Sparse Jacobians and Hessians. arXiv. https://doi.org/10.48550/arXiv.2501.17737

