CHAPTER 5: PAIRWISE SEQUENCE ALIGNMENT AND DATABASE SEARCHING

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Alignment algorithms (preview)

- Needleman-Wunsch (1970) and variations:
 - for aligning two sequences
 - uses dynamic programming to "consider" all possible alignments (10⁶⁰⁰ for two sequences of length 1000!)
- FASTA: uses a heuristic method for efficient searches (though not guaranteed to find the optimal solution)
 - Creates dictionary of k-tuples for the query sequence which is checked against sequences in the database
 - A local alignment algorithm is used to complete the alignment
- BLAST (Basic Local Alignment Search Tool): also fast and uses a heuristic
 - Finds short matches (which do not have to be perfect)
 - Then uses local alignment to complete the alignment

Needleman and Wunsch Dynamic Programming method

- Dynamic programming method
 - The problem can be divided into smaller parts
- Consider the following alignment of the sequences

$$x = THISLINE$$
and $y = ISALIGNED$

$$X_1$$
 X_2 X_3 X_4 X_5 X_6 X_7 X_8 X_9 X_{10} X_{11} \mathbf{T} \mathbf{H} \mathbf{I} \mathbf{S} $\mathbf{-}$ \mathbf{L} \mathbf{I} $\mathbf{-}$ \mathbf{N} \mathbf{E} $\mathbf{-}$ $\mathbf{-}$ \mathbf{I} \mathbf{S} \mathbf{A} \mathbf{L} \mathbf{I} \mathbf{G} \mathbf{N} \mathbf{E} \mathbf{D} Y_1 Y_2 Y_3 Y_4 Y_5 Y_6 Y_7 Y_8 Y_9 Y_{10} Y_{11}

Dynamic Programming

$$X_1$$
 X_2 X_3 X_4 X_5 X_6 X_7 X_8 X_9 X_{10} X_{11}
 \mathbf{T} \mathbf{H} \mathbf{I} \mathbf{S} $\mathbf{-}$ \mathbf{L} \mathbf{I} $\mathbf{-}$ \mathbf{N} \mathbf{E} $\mathbf{-}$
 $\mathbf{-}$ \mathbf{I} \mathbf{S} \mathbf{A} \mathbf{L} \mathbf{I} \mathbf{G} \mathbf{N} \mathbf{E} \mathbf{D}
 Y_1 Y_2 Y_3 Y_4 Y_5 Y_6 Y_7 Y_8 Y_9 Y_{10} Y_{11}

In general, the alignment is

$$\begin{bmatrix} X_1 \cdots X_u \\ Y_1 \cdots Y_u \end{bmatrix} \begin{bmatrix} X_{u+1} \cdots X_v \\ Y_{u+1} \cdots Y_v \end{bmatrix} \begin{bmatrix} X_{v+1} \cdots X_L \\ Y_{v+1} \cdots Y_L \end{bmatrix}$$

where X_u and Y_v correspond to alignment positions.

The score of the optimal global alignment from positions $1 \rightarrow L$ is the sum of the partial alignment scores from positions $1 \rightarrow u$, $u+1 \rightarrow v$, and $v+1 \rightarrow L$

Dynamic Programming Algorithm

Let $S_{i,j}$ be the score of the optimal alignment of all characters up to x_i of sequence x and y_i of sequence y.

Then there are three possibilities for the end of this alignment

$$\begin{array}{c} \cdots \ \mathsf{X}_{i-1} \\ \cdots \ \mathsf{Y}_{j-1} \\ \mathsf{S}_{i-1,i-1} \end{array} \mathsf{X}_{i}$$

$$\begin{bmatrix} \cdots & \mathbf{x}_{i} \\ \cdots & \mathbf{y}_{j-1} \end{bmatrix} \mathbf{y}_{j}$$

$$\mathbf{S}_{i,j-1}$$

$$\begin{array}{c} \cdots & X_{i-1} \\ \cdots & Y_j \end{array}$$

The optimal alignment score $S_{i,i}$ is the maximum of the following

$$S_{i-1, j-1} + s(x_i, y_j)$$
 $S_{i, j-1} + g$

$$S_{i-1, j} + g$$

where $s(x_i, y_j)$ is the substitution (or match/mismatch score) for character x_i aligned with y_i , and g is a constant gap penalty

Example:

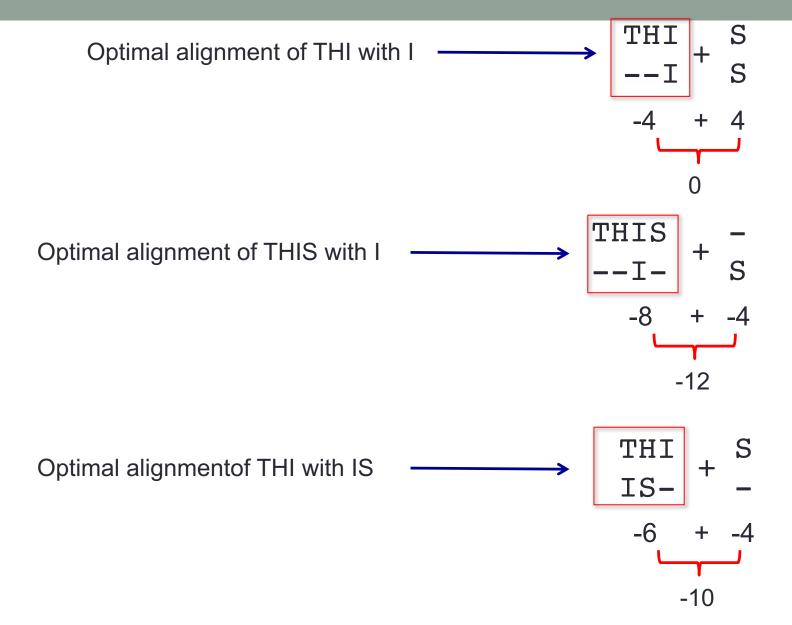
$$x = THISLINE and y = ISALIGNED$$
THISONED

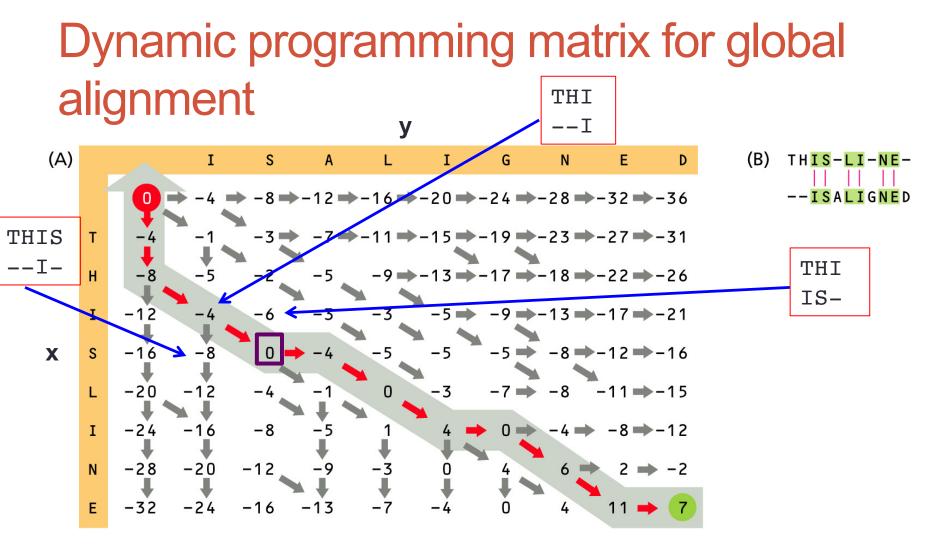
TO HISONED

TO HISONED

Assume a constant (linear) gap penalty g = 4, and $s(x_i, y_j)$ calculated using the BLOSUM-62 scoring matrix.

 $S_{4,2}$ is the optimal score for aligning THIS with IS and is the maximum of

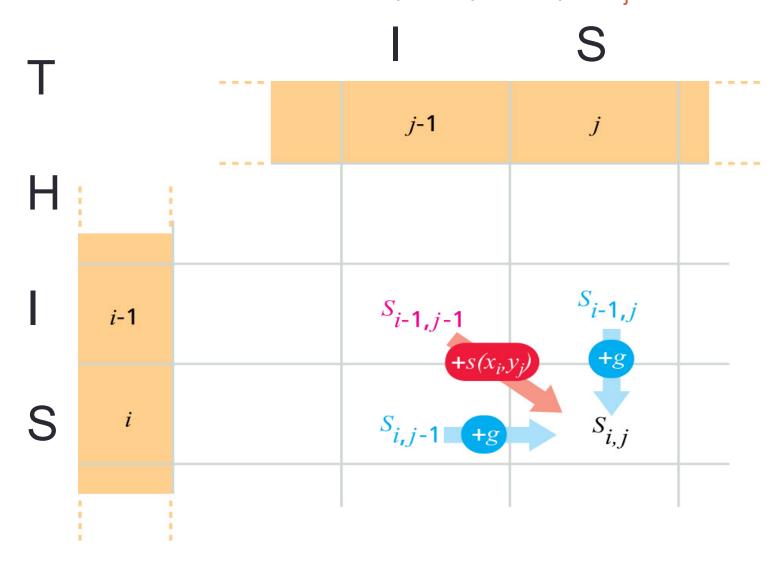




 $S_{i,j}$ is the score of the optimal alignment of all characters up to x_i of sequence \mathbf{x} and y_i of sequence \mathbf{y} .

The score of the optimal alignment of THIS with IS is $S_{4,2} = 0$.

Dynamic programming matrix illustrating possible optimal alignments for subsequences $x_1, \ldots x_i$ and $y_1, \ldots y_i$



Algorithm for finding the optimal global alignment

- For two sequences, $x_1, \ldots x_m$ and $y_1, \ldots y_n$
- Construct an (m+1) by (n+1) matrix (starting with element $S_{0,0}$ at the top left)
- For each i, the element $S_{i,0}$ corresponds to the alignment

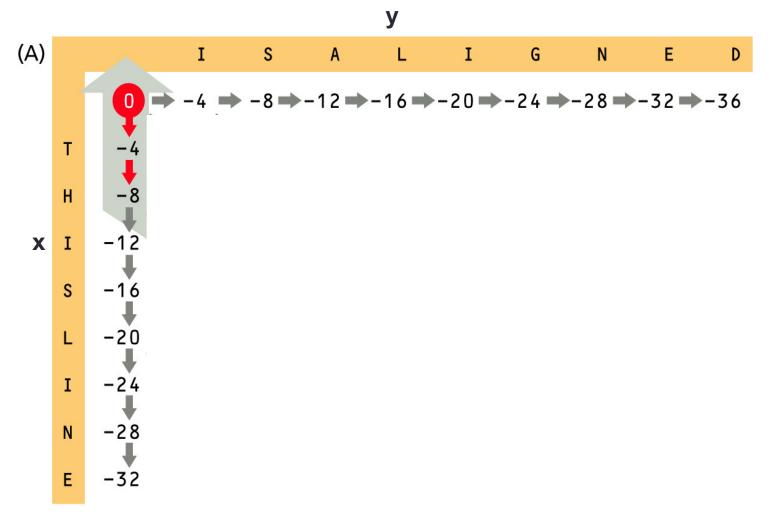
$$X_1 \quad X_2 \quad ... \quad X_1 \\ - \quad - \quad ... \quad -$$

For each j, the element $S_{0,j}$ corresponds to the alignment

$$- - ... - y_1 y_2 ... y_i$$

Fill in the matrix for all $S_{i,0}$ and $S_{0,i}$ appropriately

Assignment of $S_{i,0}$ and $S_{0,j}$ when a linear gap penalty of 4 is used



Algorithm for finding the optimal global alignment (con't)

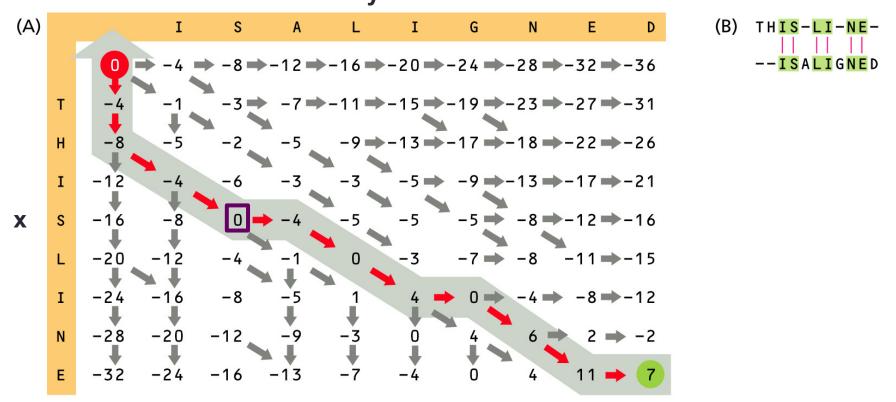
• For i = 1, ..., m and j = 1, ..., n

$$\operatorname{set} S_{i,j} = \max \begin{cases}
S_{i-1,j-1} + s(x_i, x_j) \\
S_{i-1,j} + g
\end{cases}$$

$$S_{i,j-1} + g$$

and draw an arrow to $S_{i,j}$ from the cell where the score was derived $(S_{i-1,j-1}, S_{i-1,j}, or S_{i,j-1})$

Global alignment: Dynamic programming matrix using linear gap penalty of 4



The score for the optimal alignment of THIS with IS is highlighted

Dynamic programming matrix for global alignment

y

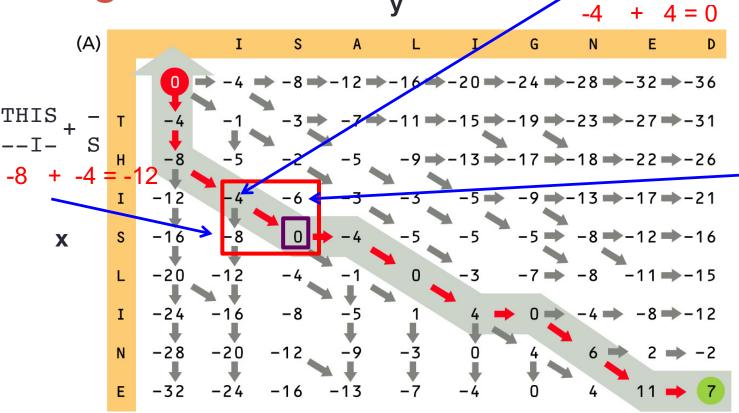
(A)

I S A L I G N E D (B) THIS-I

(B) THIS-I

(B) THIS-I

(C) A -4 -8 -12 -16 -20 -24 -28 -32 -36

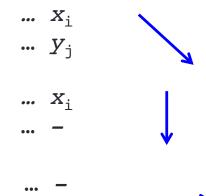


 $S_{i,j}$ is the score of the optimal alignment of all characters up to x_i of sequence x and y_i of sequence y.

The score of the optimal alignment of THIS with IS is $S_{4,2} = 0$.

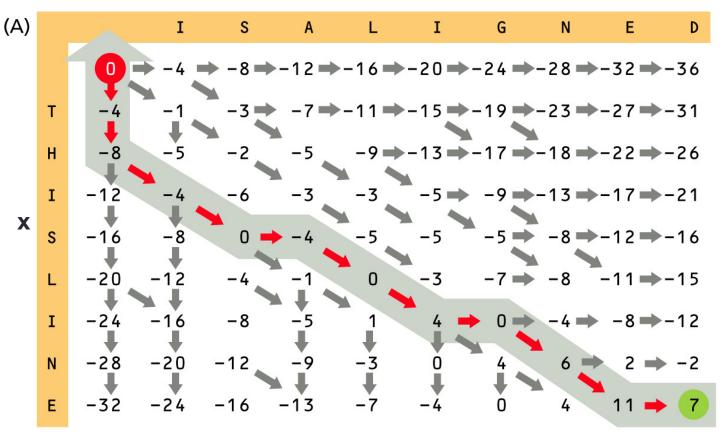
Algorithm for finding the optimal global alignment (con't)

- After the dynamic programming matrix is filled in, $S_{m,n}$ contains the score of the optimal global alignment
- A process called *traceback* is then used to recover the alignment itself, based on the path used to generate the optimal alignment score (based on the arrows), starting with $S_{m,n}$ and working backwards to $S_{0,0}$
- If the score $S_{i,j}$ was derived
 - diagonally from $S_{i-1,j-1}$, then align x_i with y_i
 - vertically from $S_{i-1,i}$, then align a gap with y_i
 - horizontally from $S_{i,j-1}$, then align x_i with a gap



Global alignment: Dynamic programming matrix using linear gap penalty of 4

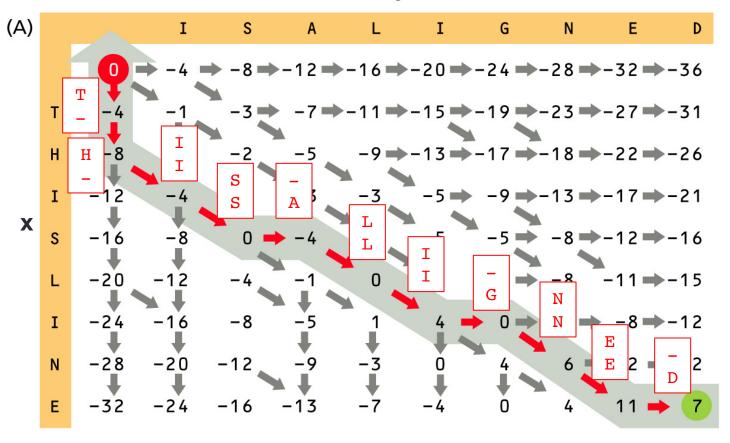
y



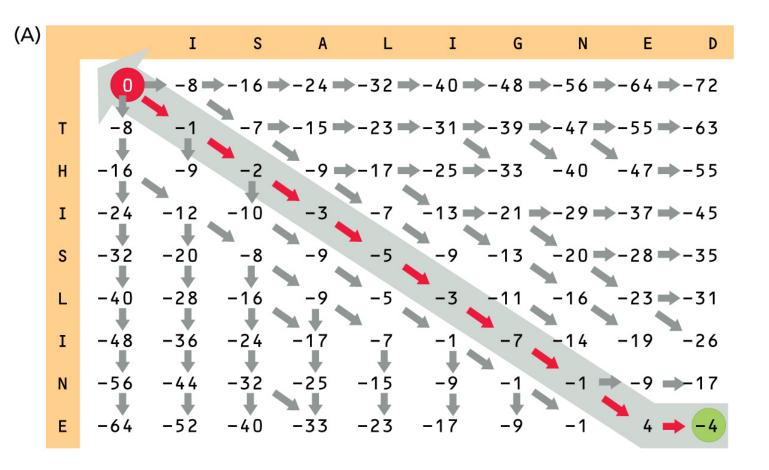


Global alignment: Dynamic programming matrix using linear gap penalty of 4

y



Global alignment: Dynamic Programming Matrix using linear gap penalty of 8

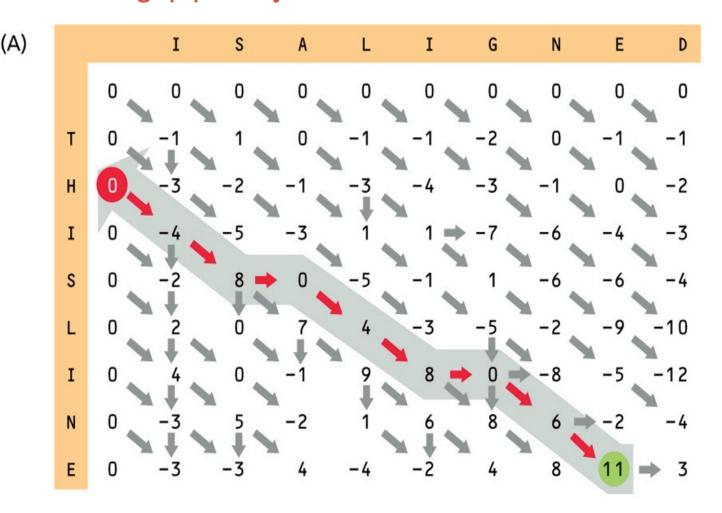


(B) THISLINE-|-| ISALIGNED

Modification of algorithm for semiglobal alignments

- A semiglobal alignment does not penalize for gaps at the beginning and end of an alignment
- Appropriate for global alignments of sequences that are not the same length
- Modifications to algorithm for semiglobal alignment:
 - Set $S_{i,0}$ and $S_{0,i}$ to 0 for all i and j
 - Start the traceback from the highest scoring element in the bottom row or last column

Semiglobal alignment: Dynamic Programming Matrix using a linear gap penalty of 8





Modification of algorithm for local alignments (1st proposed by Smith-Watterman)

- A local alignment aligns regions of two sequences, will not necessarily span the length of each sequence
- Appropriate for identifying functional domains of a protein
- Modifications to algorithm for local alignment:

$$S_{i-1,j-1} + s(x_i, x_j)$$

$$S_{i-1,j} + g$$

$$S_{i,j-1} + g$$

$$0$$

- assumes that the expected alignment score is negative for random sequences and is positive for similar sequences
- Traceback starts from highest scoring matrix element and ends at 0

Local alignment: Dynamic Programming Matrix using a linear gap penalty of 8

