Real-Time Rendering Doubly Connected Edge List (DCEL) WS 2019/20

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Literature

- Computational Geometry, Algorithms and Applications (3rd edition), Mark de Berg, Otfried Cheong, Marc van Kreveld, Mark Overmars, Springer 2008
 Course follows partially this book
 Lecture design and slide set by M. v. Kreveld
- Inspired by further lectures
 - Bernd Fröhlich and Stephan Beck, Bauhaus Universität Weimar, Real Time Rendering 2019.
 - Pjotr Indyk, MIT
 - Thomas Ottmann, Freiburg
 - Bernd Gärtner, Michael Hoffmann, ETH
 - Prof. Stefan Schirra, Madgeburg

Motivation

- Networks, rivers, and some other map layers can be represented as collections of line segments.
- Layers representing labeled regions have a more complicated structure; e.g. a thematic map of vegetation
- For this kind of layers, storing a planar subdivision as a collection of line segments is not such a good idea
- Operations like reporting the boundary of a region, or accessing adjacent regions would be rather complicated

Motivation

- We need to incorporate topological information
 - Which segments bound a given region
 - Visit all edges around a given vertex
 - Which regions are adjacent, and so on
- A simple data structure called Doubly Connected Edge List (DCEL) is commonly used to represent this kind of planar subdivision

SchleswigHolstein

MecklenburgVorpommern

Hamburg

Bremen

Niedersachsen

NordrheinWestfalen

SachsenAnhalt

Sachsen

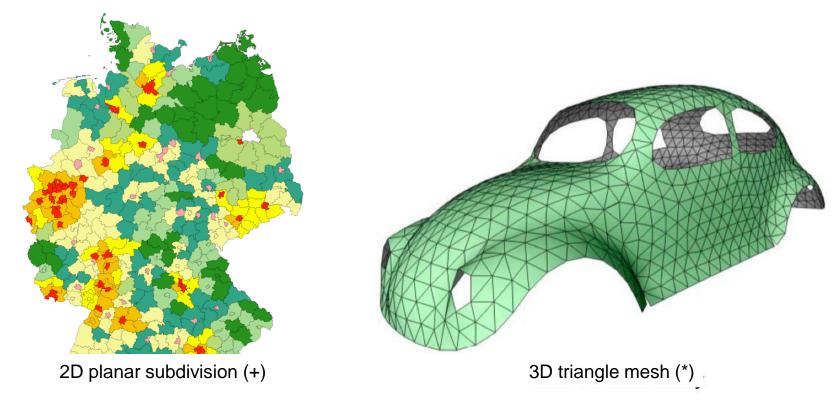
Thüringen

Bayern

BadenWürttemberg

^{*} from http://turismo.org/mapa-de-alemania/

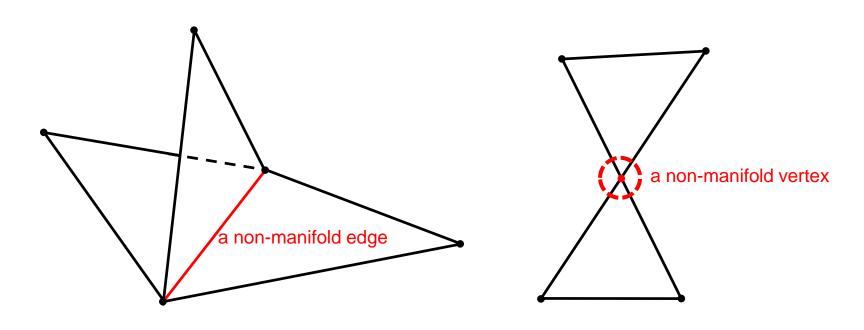
- Subdivisions can be
 - in 2D plane (planar)
 - in 3D (mesh), also known as Boundary Representations (B-Rep)



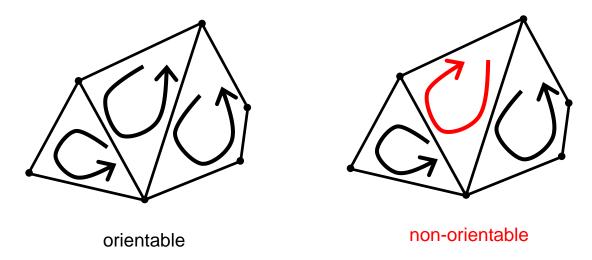
⁺ from M. Dörrbecker

^{*} from http://www.pointclouds.org/blog/nvcs/martin/index.php

- A 3D mesh is manifold if:
 - every edge separates 2 faces
 - no three or more faces share an edge
 - faces are not connected by single vertex

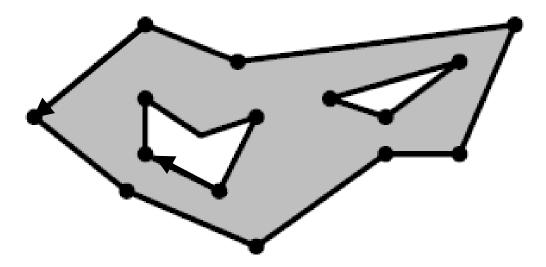


- A (2D/3D) mesh is orientable if:
 - orientation for all faces are consistent
 - inside or outside is defined by vertex ordering



 We focus on orientable and manifold planar subdivisions

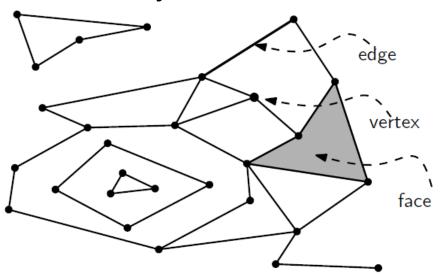
Mesh polygons may contain holes

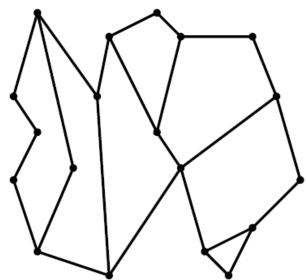


Typically CCW is used for inside and CW for outside

Planar Subdivisions

- A planar subdivision is a structure induced by a set of line segments in the plane
 - Line segments only intersect at common endpoints
 - Consists of vertices, edges and faces
- In computer graphics we often talk about polygon meshes instead of a planar subdivision – which is not exactly the same but close …





Vertices

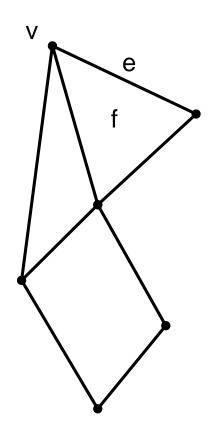
 Endpoints of line segments (nodes of the graph)

Edge

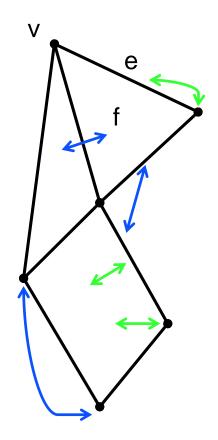
 Interiors of line segments (open interval, edges of a graph). It does not include endpoints.

Faces

 Interiors of connected twodimensional regions that do not contain any point of any line segment.
 Open polygonal region whose boundary is formed by edges and vertices from the subdivision, but that boundary is not part of the face (open)



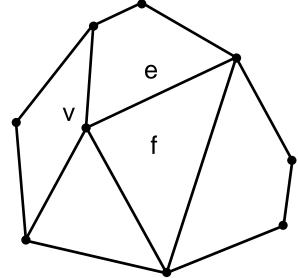
- Objects of the same dimensionality are adjacent or not
- Objects of different dimensionality are incident or not



Planar Subdivisions

We need a data structure to respond these useful queries:

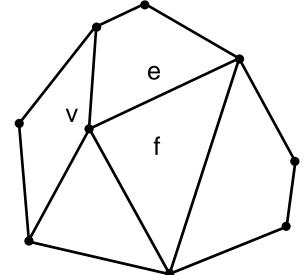
- Which faces use this vertex?
- Which edges use this vertex?
- Which faces share this edge?
- Which edges border this face?
- Which faces are adjacent to this face?



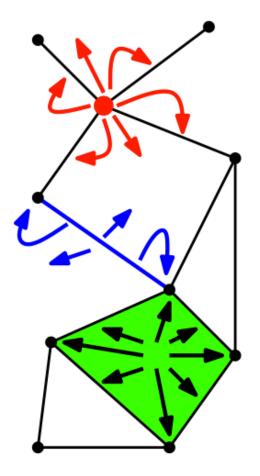
Planar Subdivisions

We need a data structure to respond these useful queries:

- Which faces use this vertex?
- Which edges use this vertex?
- Which faces share this edge?
- Which edges border this face?
- Which faces are adjacent to this face?
- We need a data structure that:
 - stores topology, geometry and attributes
 - incident and adjacent objects can be reached easily
 - enables traversal
 - with O(v+e+f) storage

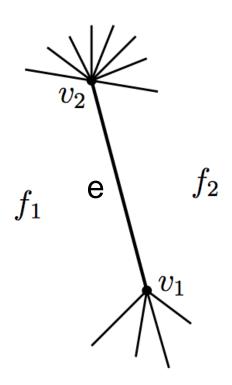


- A subdivision representation has:
 - a vertex class,
 - an edge class
 - a face class
- It is a pointers-based structure where objects can reach incident (or adjacent) objects easily

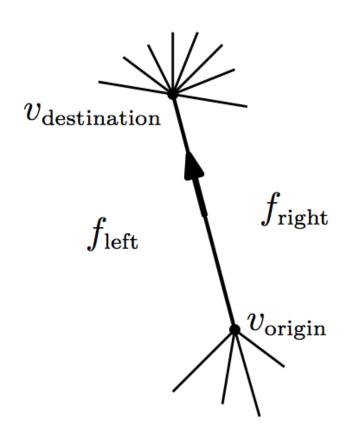


Use the edge as the central object

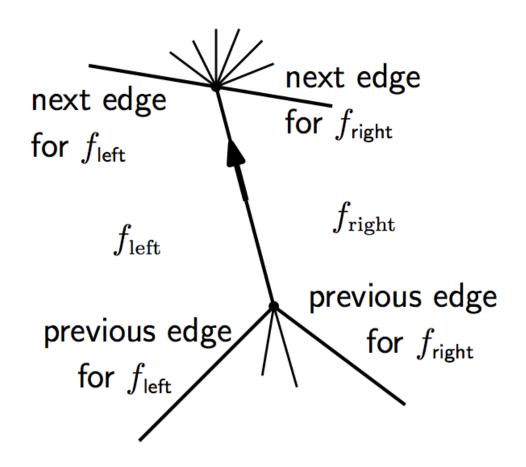
- For any edge:
 - exactly two vertices are incident
 - one or two faces are incident
 - zero or more other edges are adjacent



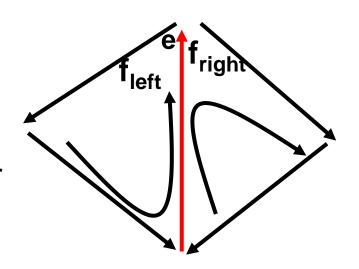
- Use the edge as the central object
 - \Rightarrow give it a direction!
- Any edge has:
 - an origin
 - a destination
 - a left face, and a right face



Four edges are of special interest:



- Each edge can be shared by 2 faces (f_{left} and f_{right})
- It would be nice if we could traverse a boundary cycle in CCW order by continuously following the next edge for f_{left} or f_{right}



... but, the orientation needs to be consistent...

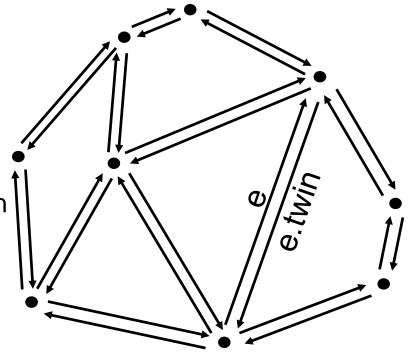
We apply a trick:split every edge into two half-edges

Every shared half-edge:

has exactly one half-edge as twin

is directed opposite to its twin

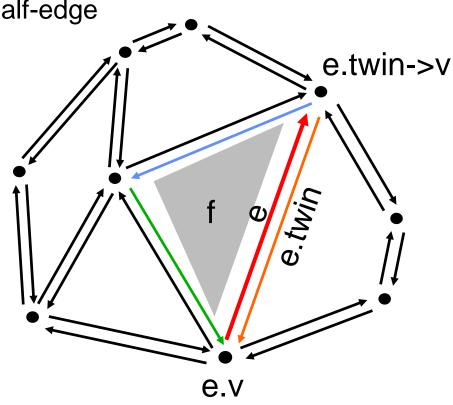
is incident to only one face



- Half-edge objects store pointers to:
 - incident starting vertex
 - adjacent next and previous half-edge
 - adjacent twin half-edge
 - incident left face

```
class h_edge {
   vertex* v;
   h_edge* twin;
   face* f;

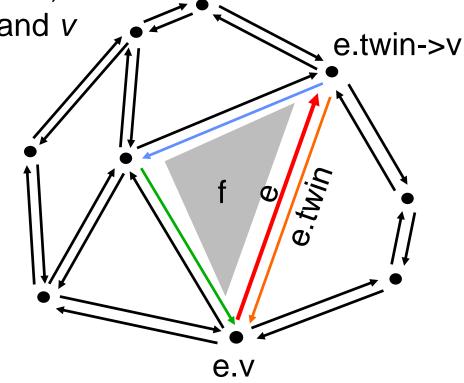
   h_edge* next;
   h_edge* prev;
};
```



Half-edge are oriented: origin → destination. If e has v as its origin and w as its destination, e. Twin has w as its origin and v as its destination.

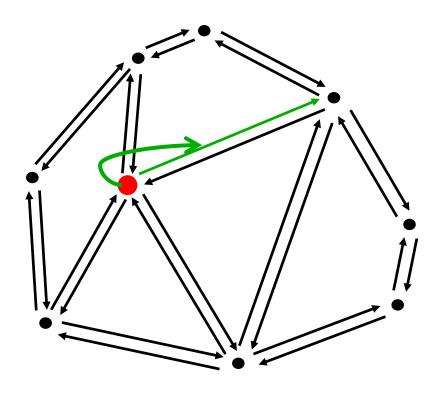
```
class h_edge {
   vertex* v;
   h_edge* twin;
   face* f;

   h_edge* next;
   h_edge* prev;
};
```



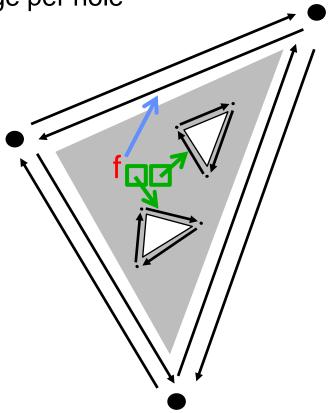
- Vertex objects store:
 - coordinates
 - pointer to one of the outgoing half edges

```
class vertex {
    point2(3)D p;
    h_edge* e;
};
```



- **Face** objects store pointers to:
 - outer component (one half-edge of outer cycle)
 - Inner components: one half edge per hole

```
class face {
   h_edge* outer;
   list<h_edge*> inner;
}
```



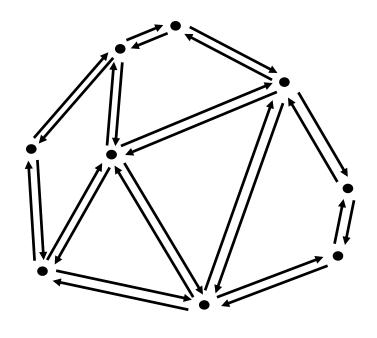
 A DCEL data structure is a collection of vertex, half-edge and face objects that are connected correctly by pointers.

```
class vertex {
      point2(3)D p;
      h_edge* e;
}
class h_edge {
   vertex* v;
   h_edge* twin;
   face* f;
   h_edge* next;
   h_edge* prev;
```

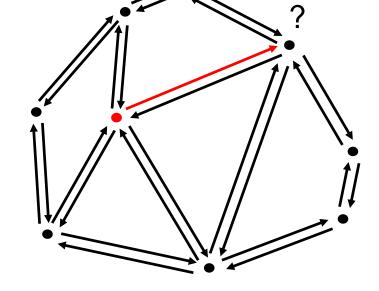
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class face {
   h_edge* outer;
   list<h_edge*> inner;
```

- Recap some questions:
 - Which is the destination vertex for a half-edge?
 - Which faces use this vertex?
 - Which edges use this vertex?
 - Which faces border this edge?
 - Which edges border this face?





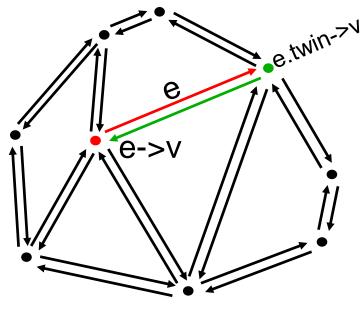
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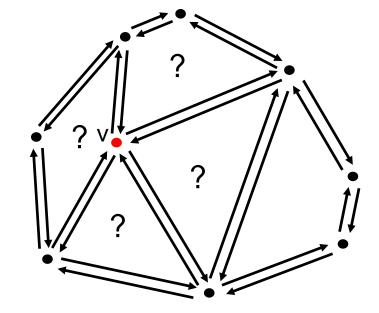
Which are all adjacent faces of a face?

- Recap some questions:
 - Which is the destination vertex for a half-edge?
 Use pointer to twin half-edge
 - Which faces use this vertex?
 - Which edges use this vertex?
 - Which faces border this edge?
 - Which edges border this face?





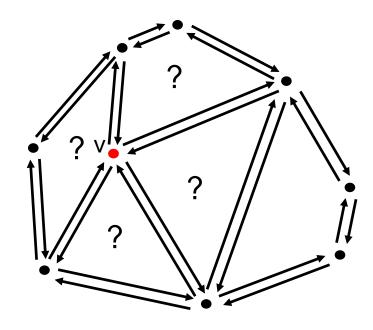
- Recap some questions:
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Which are all adjacent faces of a face?

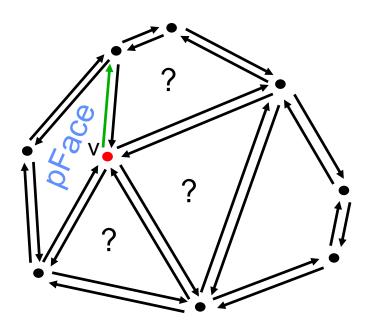
- Which faces use this vertex?
 - Iterate over the outgoing and incoming half-edges of vertex v

```
h_edge* pEdge = v.e;
do
{
   face* pFace = pEdge->f;
   if (pFace != NULL)
       report(pFace);
   pEdge = pEdge->twin->next;
} while (pEdge != v.e);
```



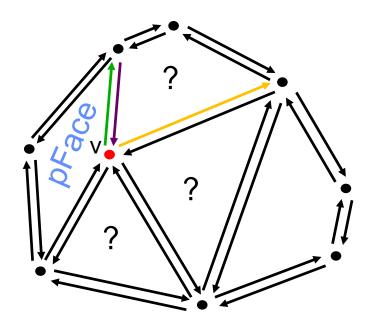
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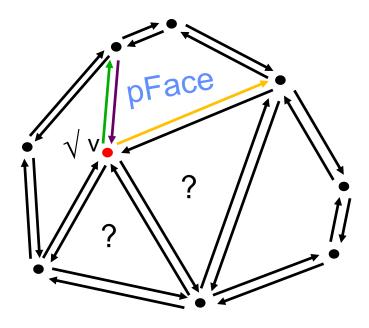
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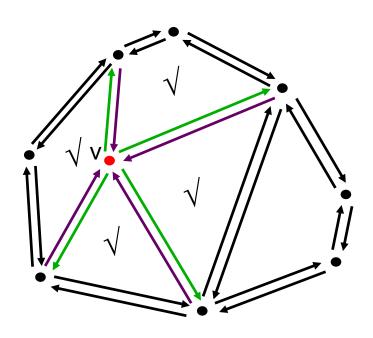
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   if (pFace != NULL)
       report(pFace);
   pEdge = pEdge->twin->next;
} while (pEdge != v.e);
```



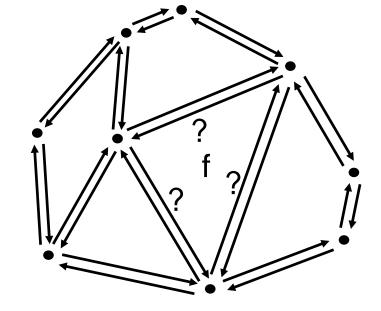
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      report(pFace);
   pEdge = pEdge->twin->next;
} while (pEdge != v.e);
```



After 4 iterations v.e is reached and all faces around v were reported ($\sqrt{\ }$)

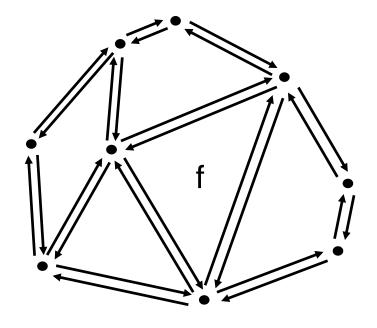
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 - Which faces use this vertex?
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 - Which edges border this face?



Which are all adjacent faces of a face?

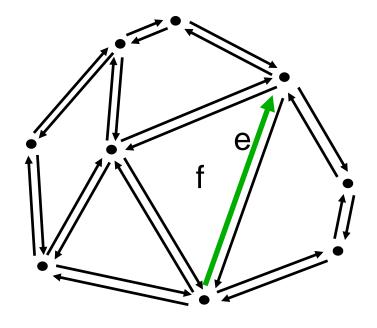
- Which edges border this face?
 - Iterating over the half-edges of a given face f

```
h_edge* e = f.outer;
do {
    Report(e);
    e = e->next;
} while (e!= f.outer);
```



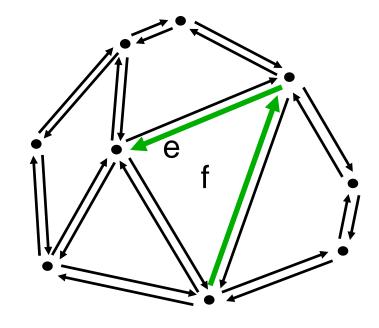
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 - Iterating over the half-edges of a given face f

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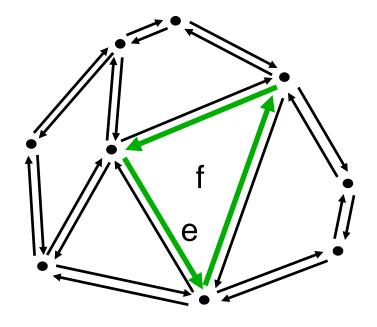
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 - Iterating over the half-edges of a given face f

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h_edge* e = f.outer;
do {
    Report(e);
    e = e->next;
} while (e!= f.outer);
```



- Which edges border this face?
 - Iterating over the half-edges of a given face f

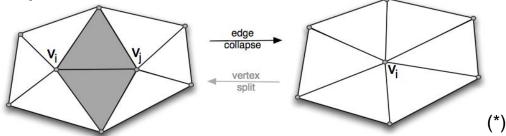
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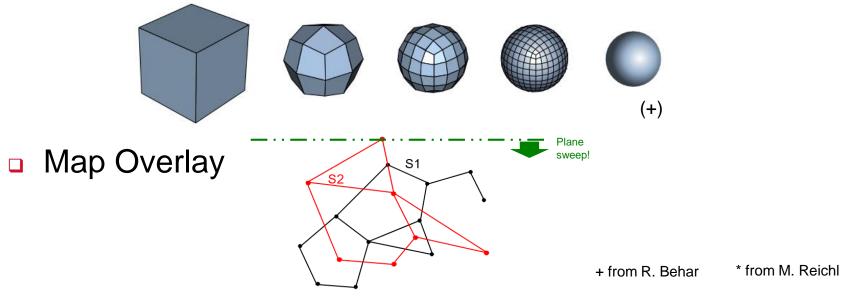
DCEL Applications

Mesh processing (e.g. calculate normal per face, per vertex)

Mesh simplification



Subdivision surfaces (Catmull-Clark)



Bauhaus-Universität Weimar

Summary

- Doubly-Connected Edge List
 - Common data structure for planar subdivisions
 - Boundary representation (B-rep) of a planar subdivision
 - Forms a graph of connected vertices, half-edges and faces
 - Half-edge is the central component for traversal
 - Supports adjacency queries
 - Storage is O(e+v+f) (if no holes are present)
 - Many applications

End