Decompose the given relation R(A, B, C, D) using 3NF.
FDs = { C- > A, C -> D, C -> C, AB -> C }
R(A, B, C, D)
1. Consider the following relation and functional dependencies



C -> A

C -> D

C -> C

AB -> C

(2)

C -> AD

AB -> C

(3)

The Fc is:

AB -> CD

To convert into 3NF, put the left-hand-side (LHS) and the right-hand-side (RHS) of each FD in Fc together in one relation. Doing so, we get ABCD.

### Based on the decomposed relations (above), discuss to show that they are in 3NF

- (1) Does this satisfy lossless-join decomposition?
  - a. Yes, as there is only one functional dependency.
- (2) Does this satisfy dependency preservation?
  - a. Yes, the dependencies can be checked without joining tables.

Consider the following relation and functional dependencies

# Decompose the given relation R using BCNF

(1)

C -> A

C -> D

C -> C

AB -> C

(2)

C -> ACD

AB -> ABC

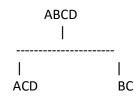
(3)

C -> ACD

AB -> ABCD

C -> ACD is not a trivial FD.

$$C + (ABCD - ACD) = BC$$



R(ABCDE) becomes ACD // BC

2. Consider the following relation and functional dependencies

# Decompose the given relation R using 3NF

(1)

A -> ABC

C -> D

A -> C

D -> D

(2)

A -> ABC

C -> D

(3)

The Fc is:

A -> BC

C -> D

(4)

ABC // CD

R1(ABC), R2(CD)

#### Based on the decomposed relations (above), discuss to show that they are in 3NF

- (3) Does this satisfy lossless-join decomposition?
  - a. Yes, as R1  $\cap$  R2 = {C} and C -> D (C is super key of R2)
- (4) Does this satisfy dependency preservation?
  - a. Yes, the dependencies can be checked without joining tables.

Consider the following relation and functional dependencies

# Decompose the given relation R using BCNF

(1)

A -> ABC

C -> D

A -> C

D -> D

(2)

A -> ABC

C -> CD

A -> AC

D -> D

(3)

A -> ABCD

C -> CD

D -> D

(4)

A is a super key, C->CD is a non-trivial FD, and D->D is a trivial FD

$$C + (ABCD - CD) = ABC$$

R(ABCDE) becomes CD // ABC