

# Pthread Implementation of Gaussian Elimination

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## 1 Introduction

Let the system be

$$Ax = b \quad (1)$$

Here  $A$  is a dense  $n \times n$  matrix of coefficients such that  $A[i][j] = a_{i,j}$ ,  $b$  is an  $n \times 1$  vector and  $x$  is the desired solution vector.

A system of equations is solved in two stages. First, through a series of algebraic manipulations, the original system of equations is reduced to an upper triangular system of the form  $Ux = y$ , where  $U$  is a unit upper-triangular matrix. In the second stage of solving a system of linear equations, the upper-triangular system is solved for the variables in reverse order from  $x_{n-1}$  to  $x_0$  by a procedure known as back-substitution.

**Algorithm 1** *gaussian\_elimination*

1. procedure gaussian\_elimination( $A, b, y$ )
2. begin
3.     for  $k := 0$  to  $n-1$  do
4.         for  $j := k+1$  to  $n-1$  do
5.              $A[k][j] := A[k][j] / A[k][k]$ ; //division step
6.          $y[k] := b[k] / A[k][k]$ ;
7.          $A[k][k] := 1$ ;
8.         for  $i := k+1$  to  $n-1$  do
9.             for  $j := k+1$  to  $n-1$  do
10.                  $A[i][j] := A[i][j] - A[i][k] \times A[k][j]$ ; //elimination step
11.                  $b[i] := b[i] - A[i][k] \times y[k]$ ;
12.                  $A[i][k] := 0$ ;
13.             endfor;
14.         endfor;
15. end gaussian\_elimination;

## 2 Pipeline Communication and Computation

The matrix is scattered to all processor so two consecutive row are in two consecutive processors.

During the  $k^{th}$  iteration processor  $P_k$  broadcast part of the  $k^{th}$  row of the matrix to processors  $P_{k+1}, \dots, P_{p-1}$ . Assume that the processors  $P_0 \dots P_{p-1}$  are connected in a linear array, and  $P_{k+1}$  is the first processor to receive the  $k^{th}$  row from processor  $P_k$ . Then the processor  $P_{k+1}$  must forward this data to  $P_{k+2}$ . However, after forwarding the  $k^{th}$  row to  $P_{k+2}$ , processor  $P_{k+1}$  need not wait to perform the elimination step until all the processors up to  $P_{p-1}$  have received the  $k^{th}$  row. Similarly,  $P_{k+2}$  can start its computation as soon as it has forwarded the  $k^{th}$  row to  $P_{k+3}$ , and so on. Meanwhile, after completing the computation for the  $k^{th}$  iteration,  $P_{k+1}$  can perform the division step, and start the broadcast of the  $(k+1)^{th}$  row by sending it to  $P_{k+2}$ . In this case of shared memory the receive, send and broadcast from MPI are replaced by pthread mutex procedures but without pipeline communication.

This algorithm is made for comparing the MPI, pthread and OpenMP implementation.

## 3 Results

I have compiled the Gaussian Elimination without load balancing and pivoting program with native gcc of the Fedora Core 1 with maximum optimization "-O3".

The executable were run on a dual pentium II at 500MHz with 256MB RAM and with LINUX Fedora Core 1.

The following result was made for a average of 10 runs for serial and parallel programs.

