

Graph Data Structures

Minimum Spanning Tree

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Problems Specific with Graphs

Minimum Spanning Tree

- Path Problems

1. Simple Paths

- 2. Shortest Path Problem**

- Single source shortest paths
- All-pair shortest paths

3. Find Cycles

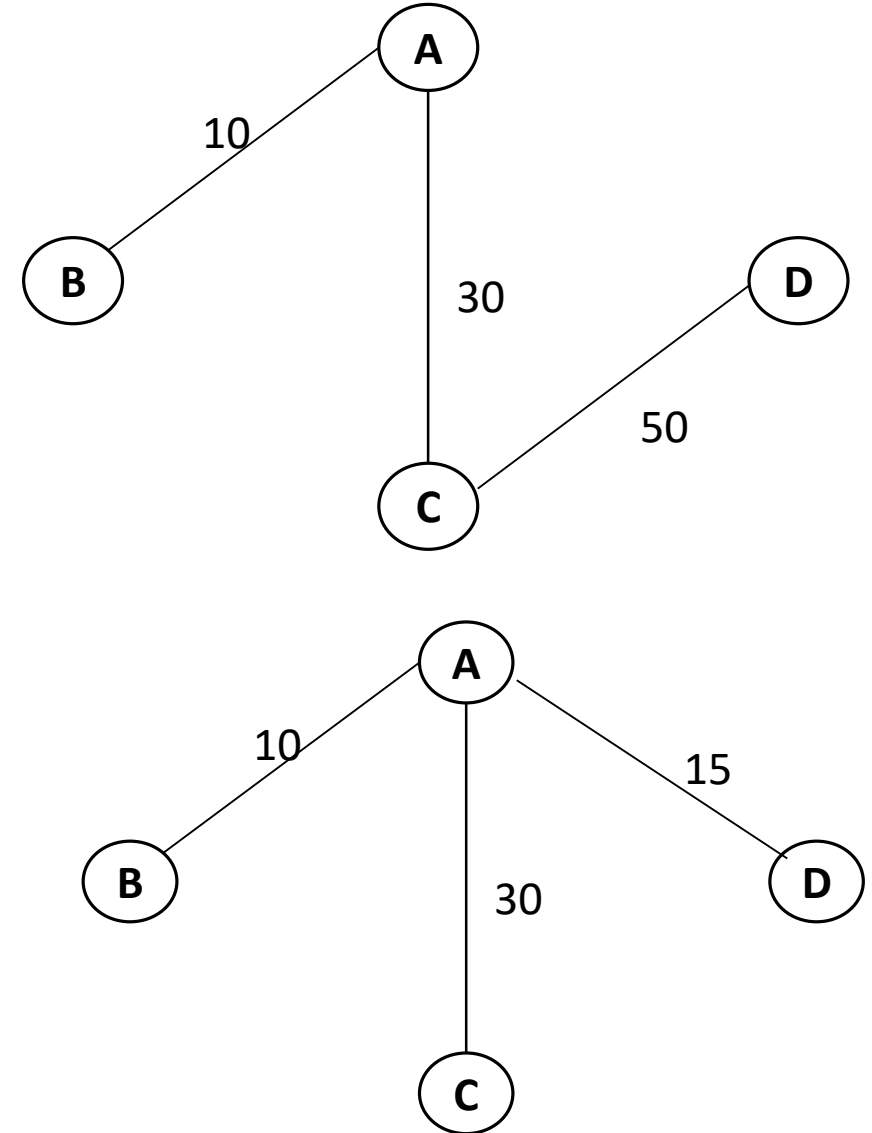
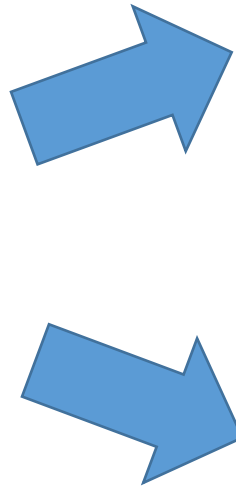
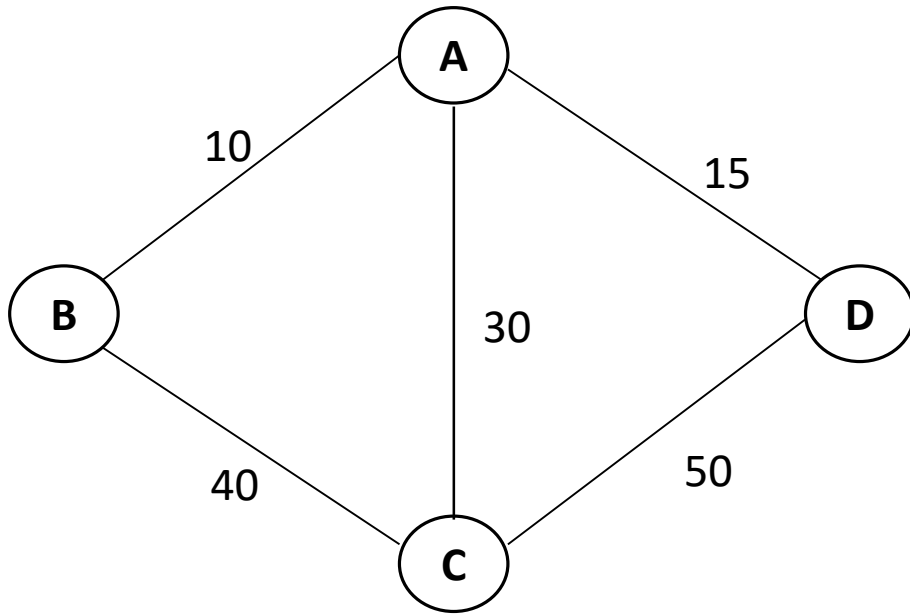
4. Euler Path and Circuit Problem

5. Hamiltonian Path and Circuit Problem (or TSP)

Spanning Tree

- For a weighted undirected graph, Spanning Tree is a sub-graph that connects all the vertices together using the minimum number of edges required
- The graph should be connected: There should be no cycles
- For n number of *vertices*, $(n-1)$ edges are needed
- Hence for four vertices, we need three edges without any cycles

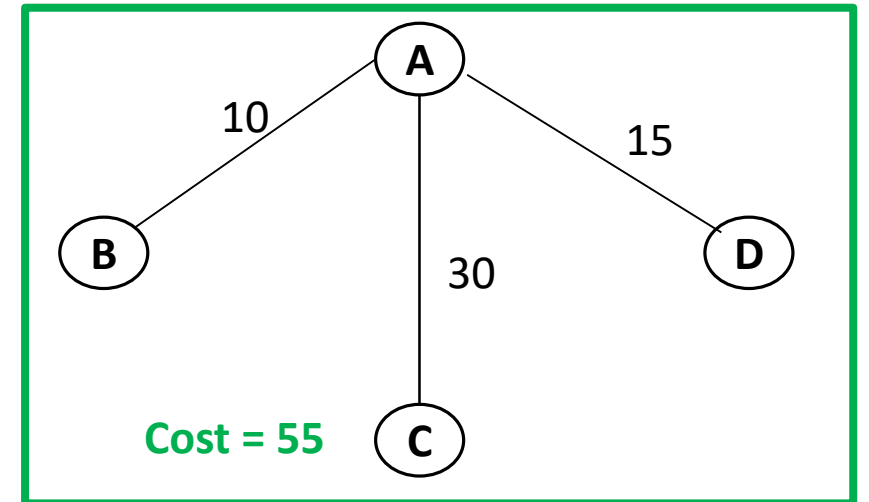
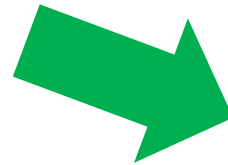
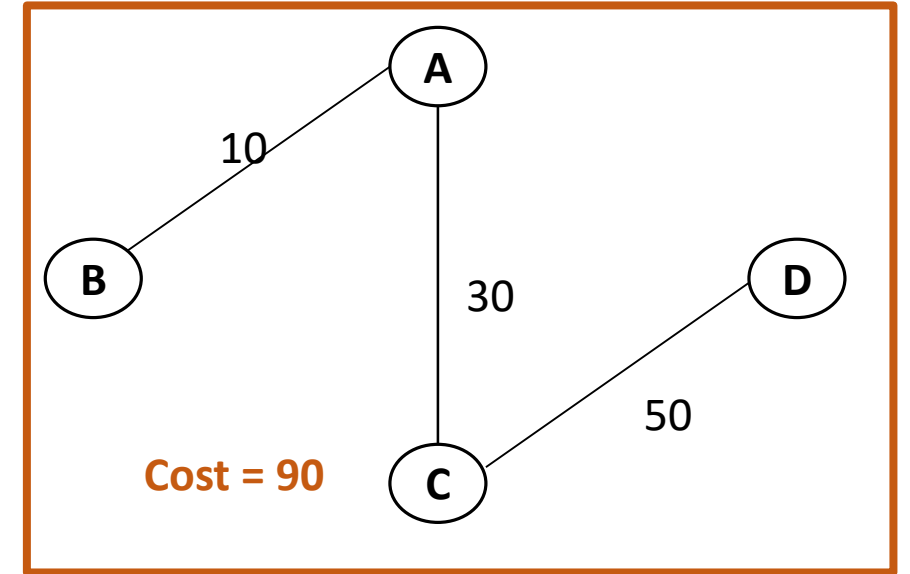
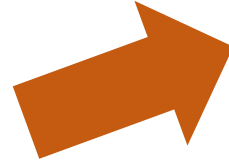
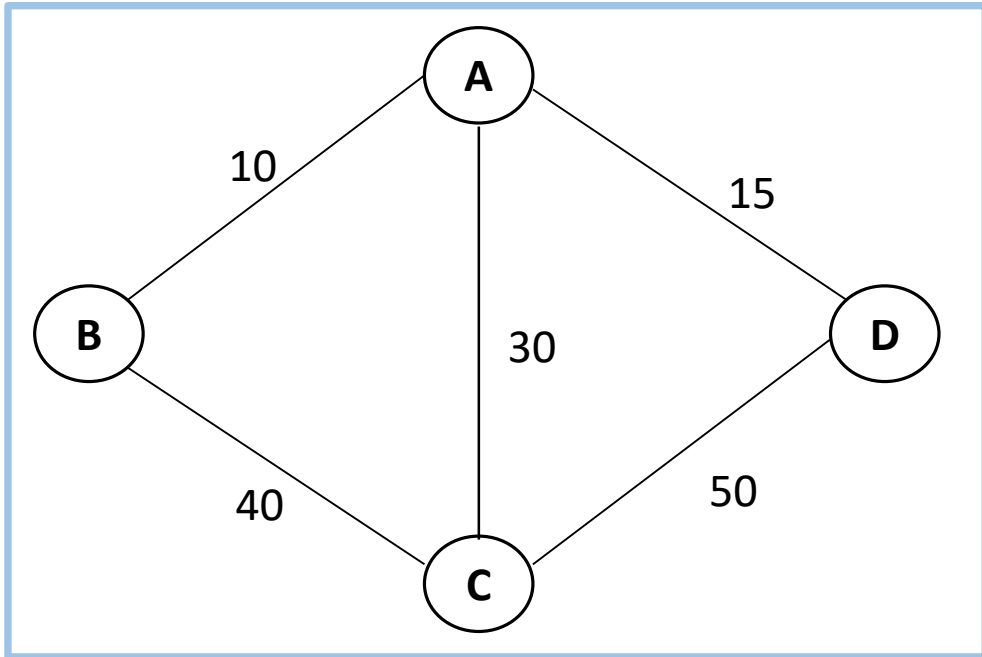
Many Spanning Trees for a Single Graph



Minimum Spanning Tree (MST)

- **Cost of any spanning tree:** Add weights of all the edges
- Many spanning trees for a single graph
- Calculate cost of all spanning trees and pick the minimum cost spanning tree
 - A tree with minimum weights is MST
- **Time complexity:** *Exponential*
- Two Algorithms
 - Prim's Algorithm
 - Kruskal's Algorithm

Minimum Spanning Trees



Prim's Algorithm: At a Glance

- Maintains two sets
- **MST Set**: vertices that are included in the spanning tree
- **Set 2**: vertices that are not yet included in the spanning tree

Prim's MST Algorithm

- **Key Idea:** Find the local optimum in the hopes of finding a global optimum.
- Start from one vertex and keep adding edges with the lowest weight until all vertices are reached.

Steps:

1. Initialize the minimum spanning tree with a vertex chosen at random.
2. Find all the edges that connect the tree to new vertices, select the minimum and add it to the tree
3. Keep repeating step 2 until we get a minimum spanning tree

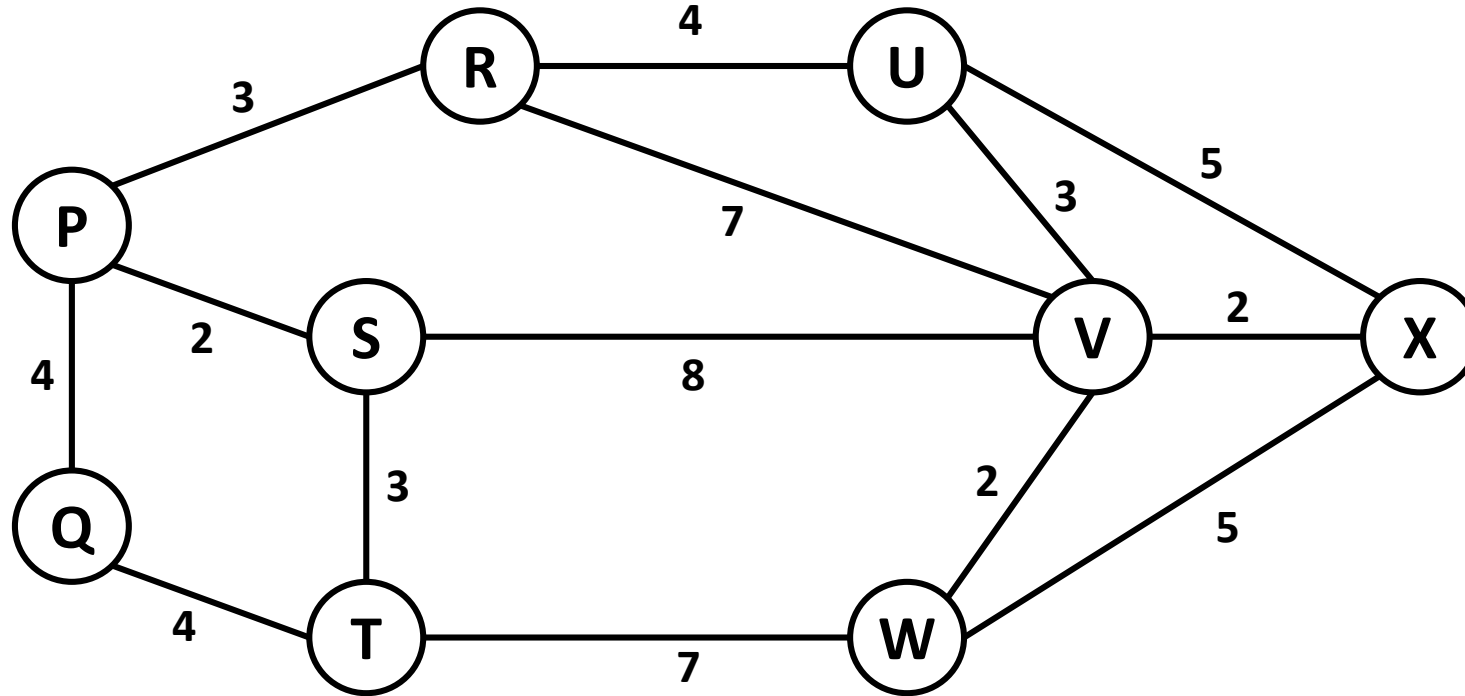
1. Create an array of size **V** and initialize it with NIL
2. Create a Priority Queue (min Heap) of size **V**. Let the **min Heap be Q**
3. Insert all vertices into **Q**. Assign **cost** of starting vertex to **0** and the **cost** of other vertices to infinite.

Prim's Algorithm: Pseudo-code

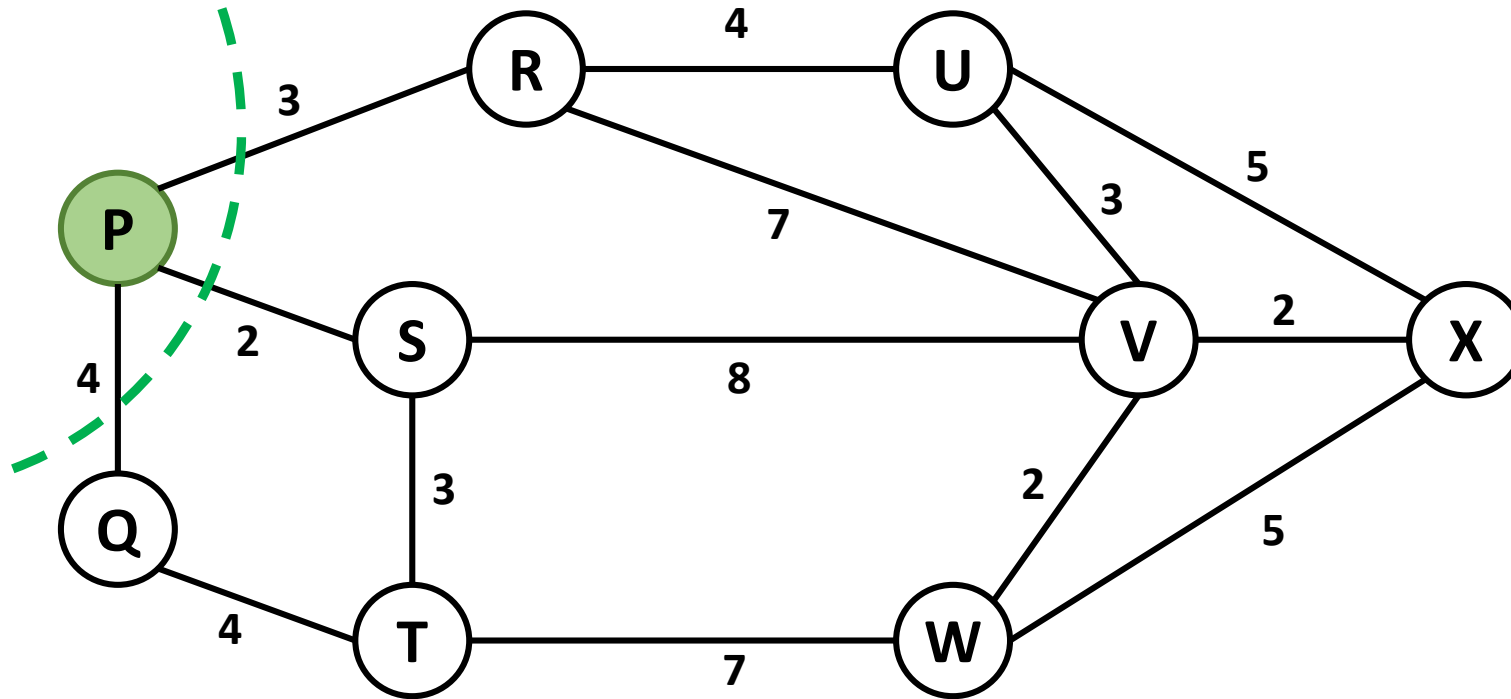
MST-PRIM($G, w, source$)

1. *for each $u \in V[G]$ do // for all vertices*
2. $cost[u] \leftarrow \infty$
3. $\pi[u] \leftarrow NULL$
4. $cost[source] \leftarrow 0$
5. $Q \leftarrow V[G]$ // set 2
6. *while $Q \neq \emptyset$ do*
7. $u \leftarrow EXTRACT-MIN(Q)$
8. *for each $v \in Adj[u]$ do*
9. *if $v \in Q$ and $w(u, v) < cost[v]$*
10. $\pi[v] \leftarrow u$
11. $cost[v] \leftarrow w(u, v)$

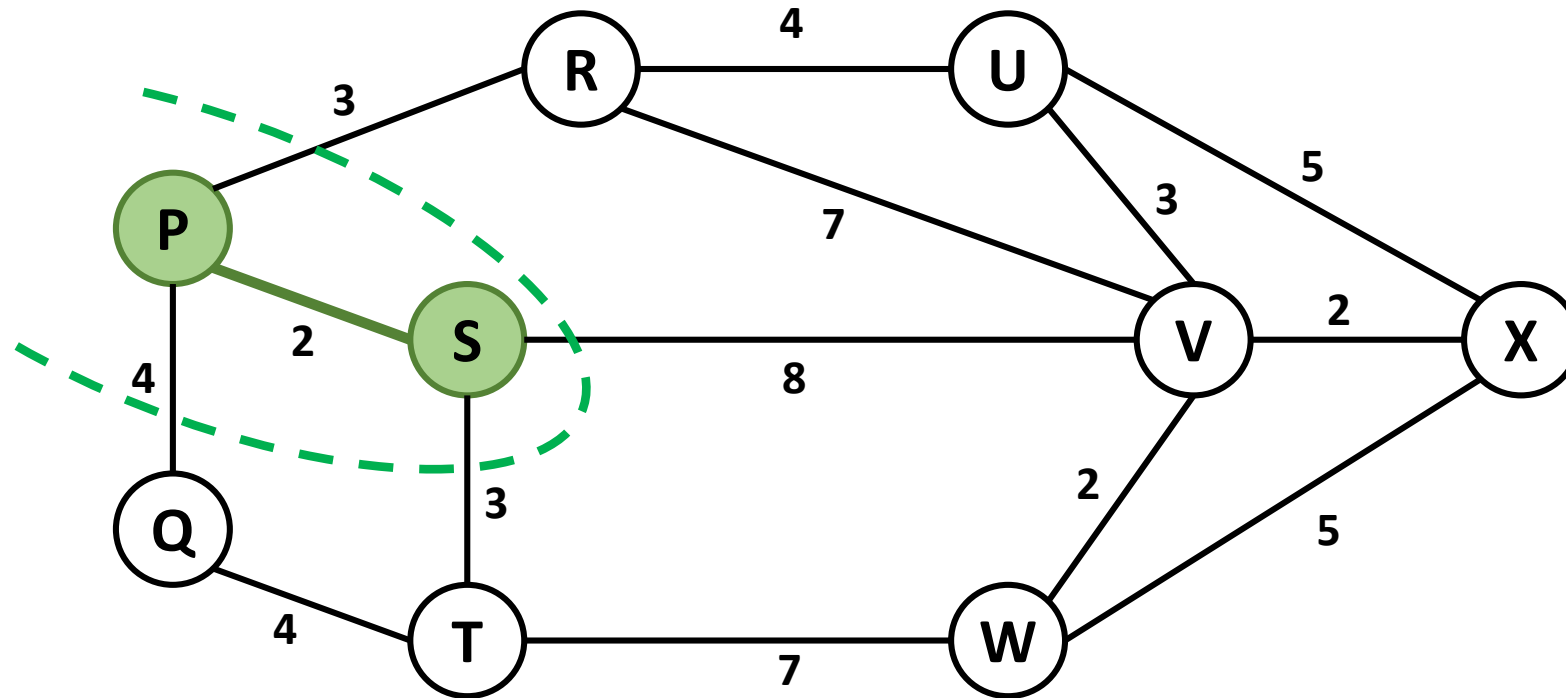
Prim's Algorithm: An Example



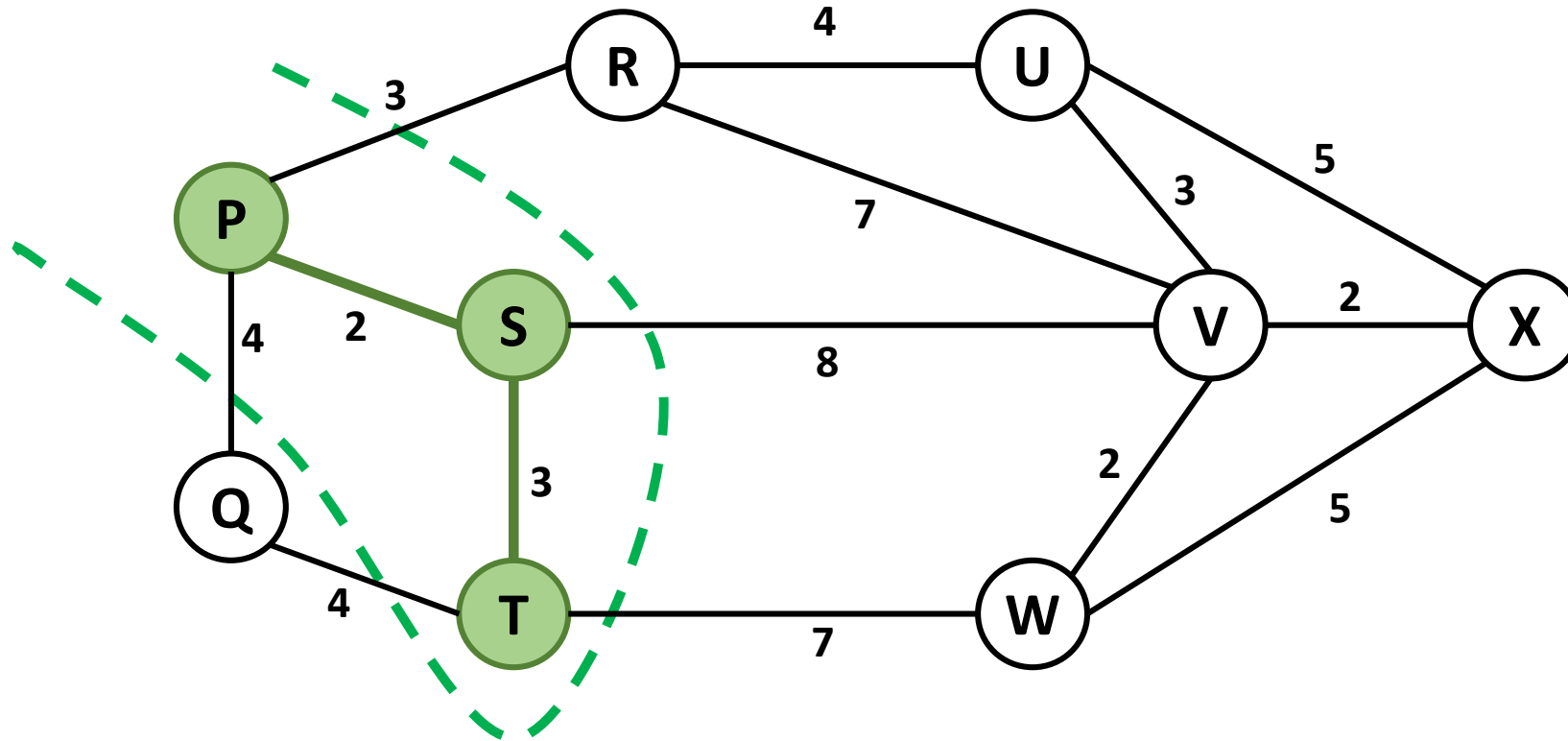
Prim's Algorithm – An Example: Step 1



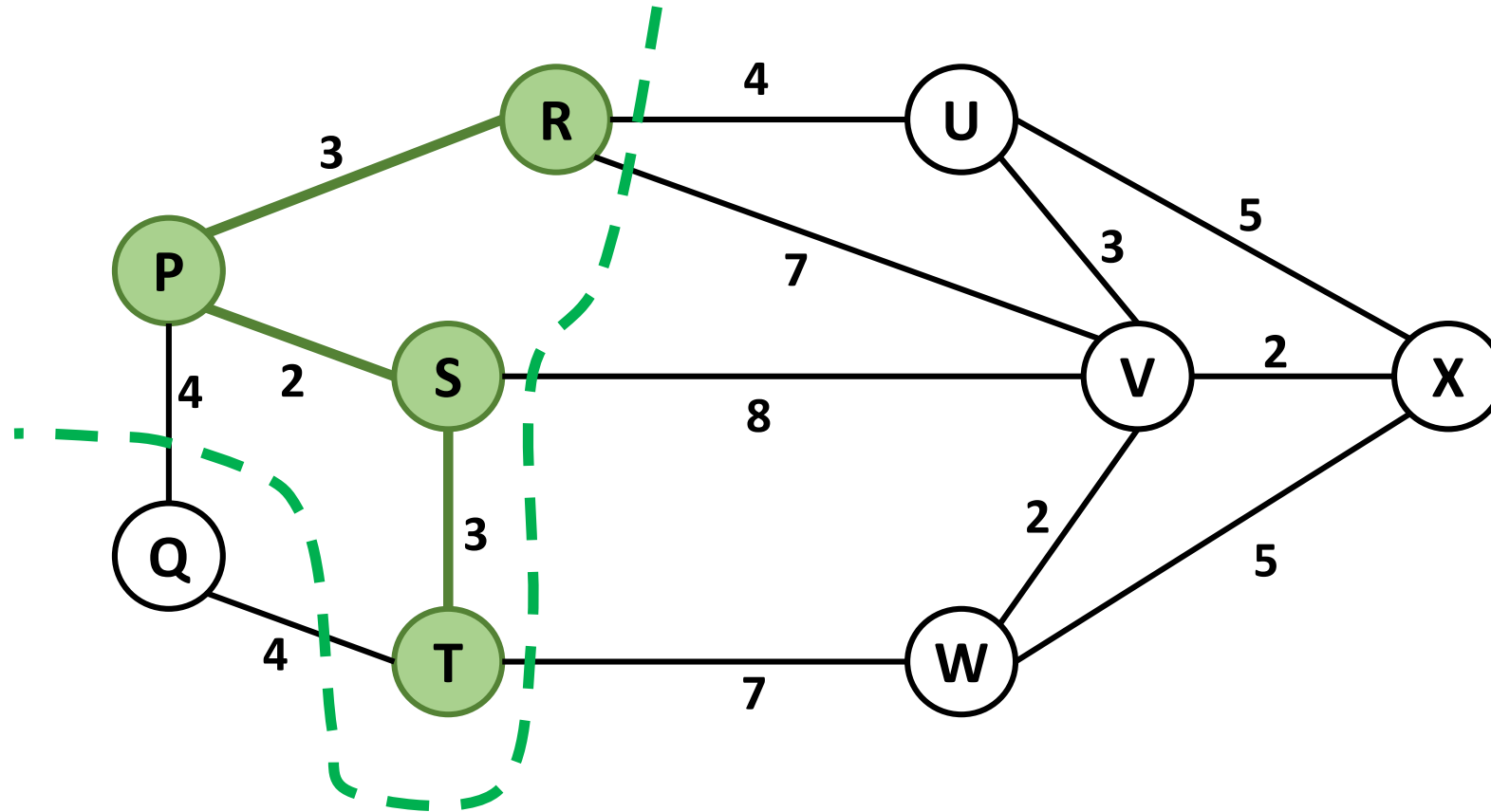
Prim's Algorithm – An Example: Step 2



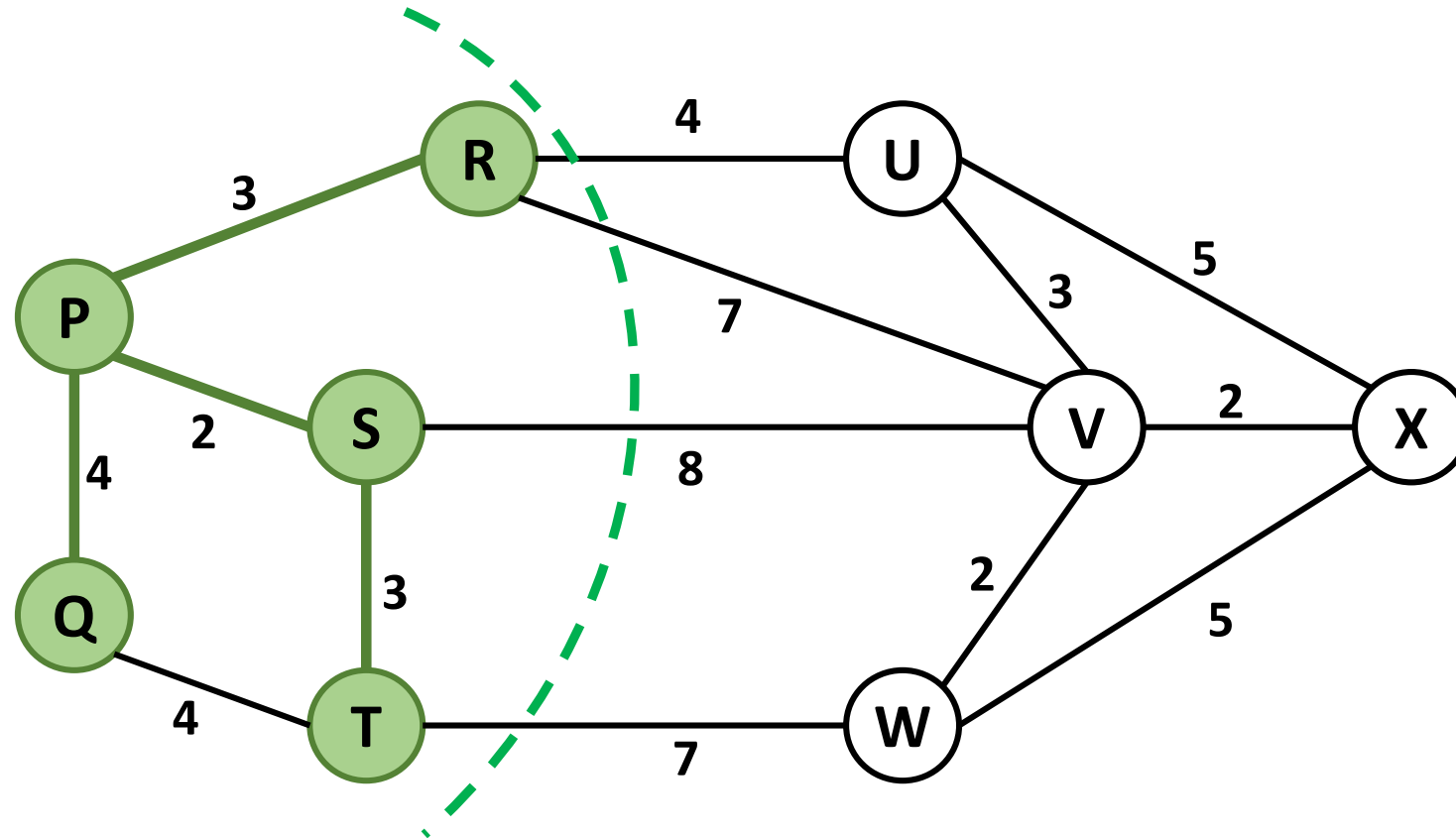
Prim's Algorithm – An Example: Step 3



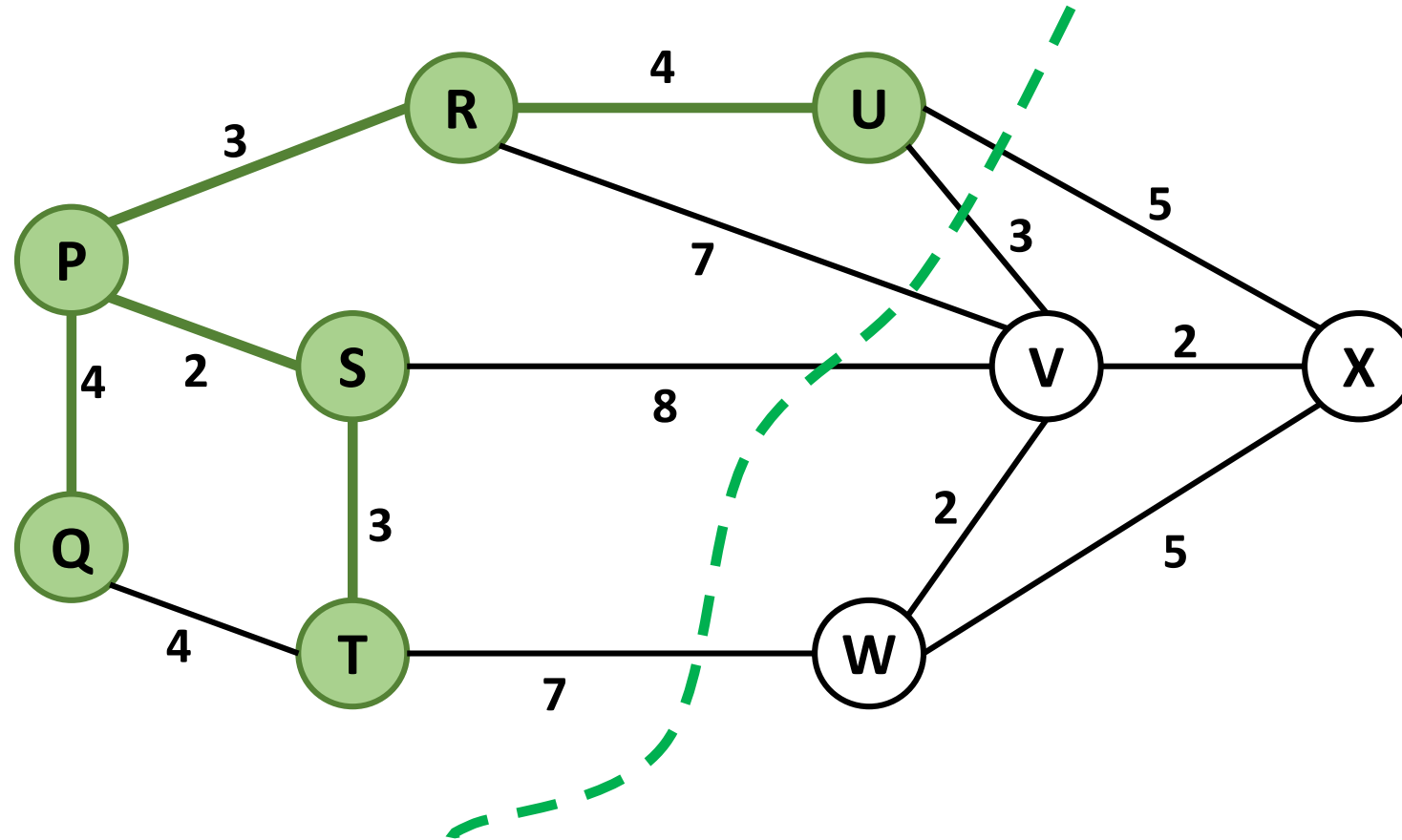
Prim's Algorithm – An Example: Step 4



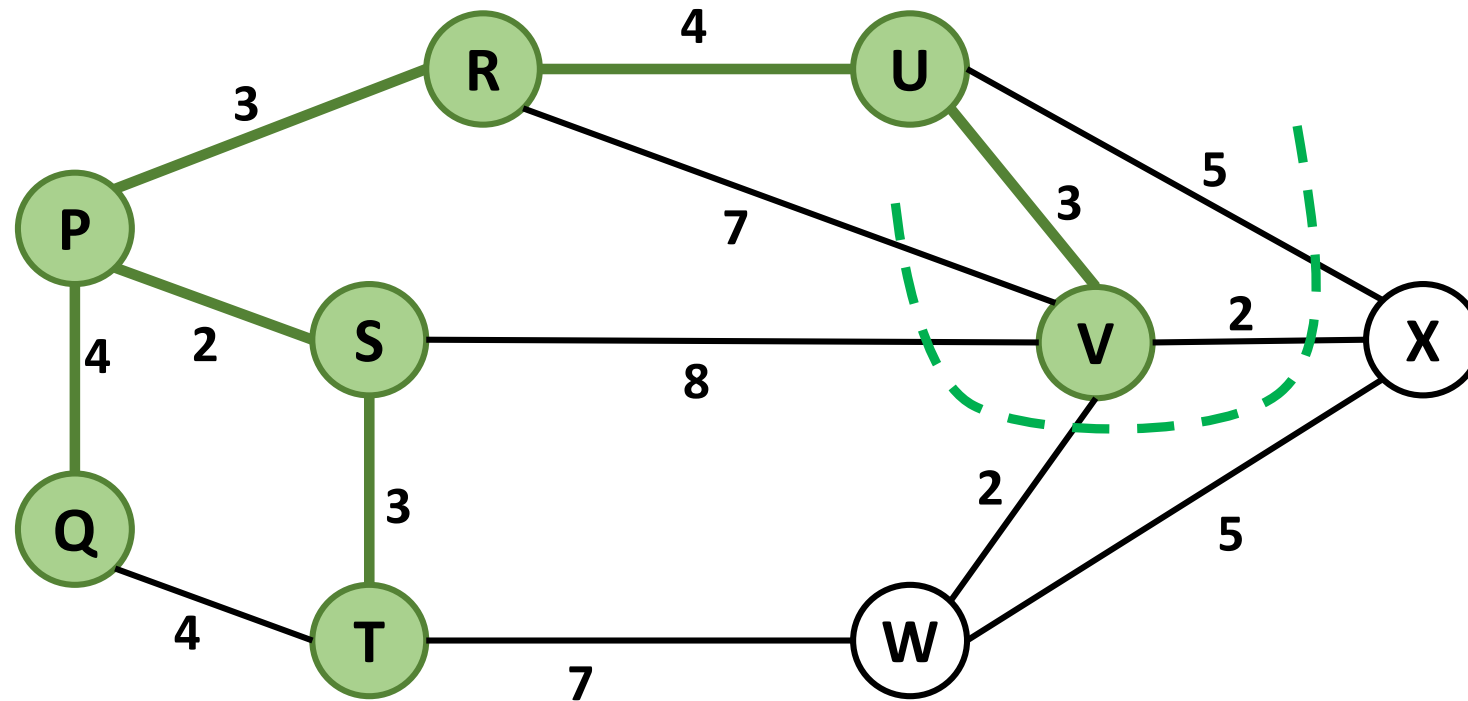
Prim's Algorithm – An Example: Step 5



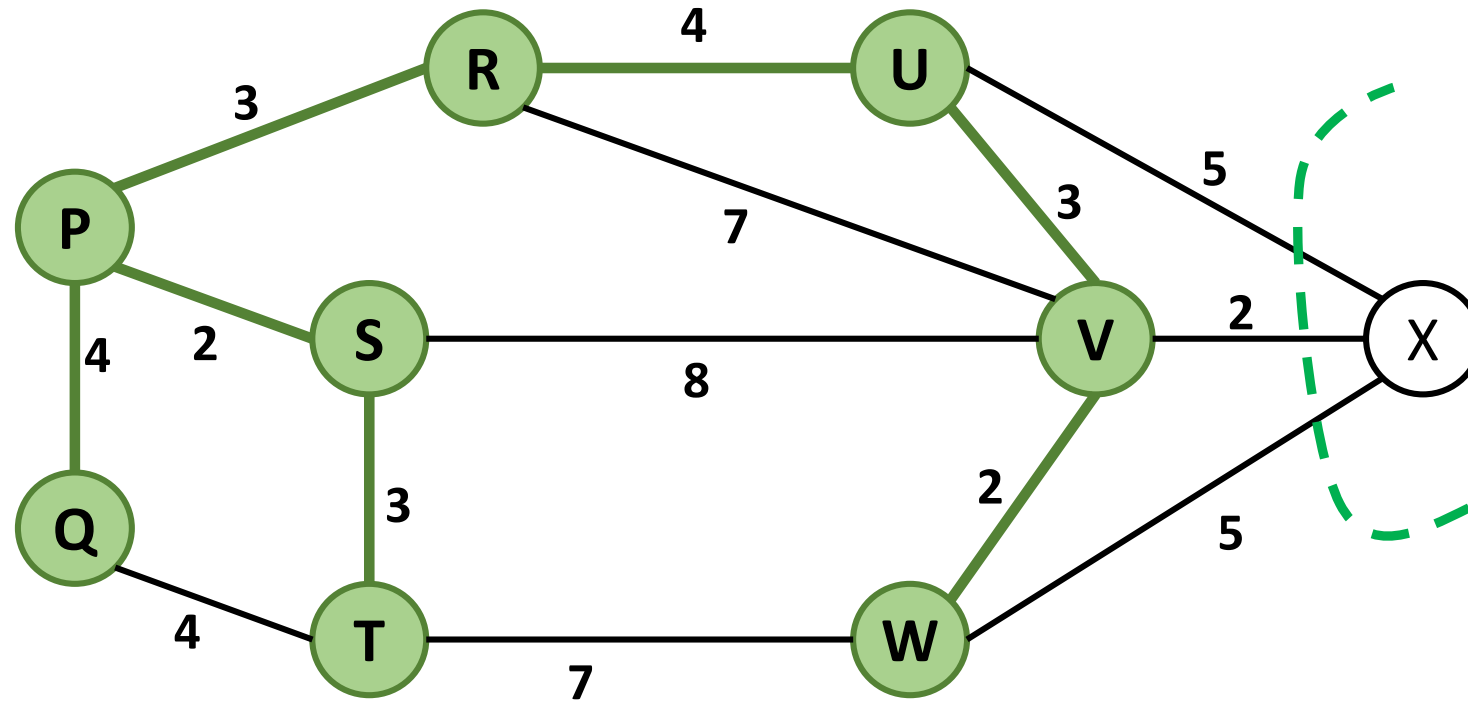
Prim's Algorithm – An Example: Step 6



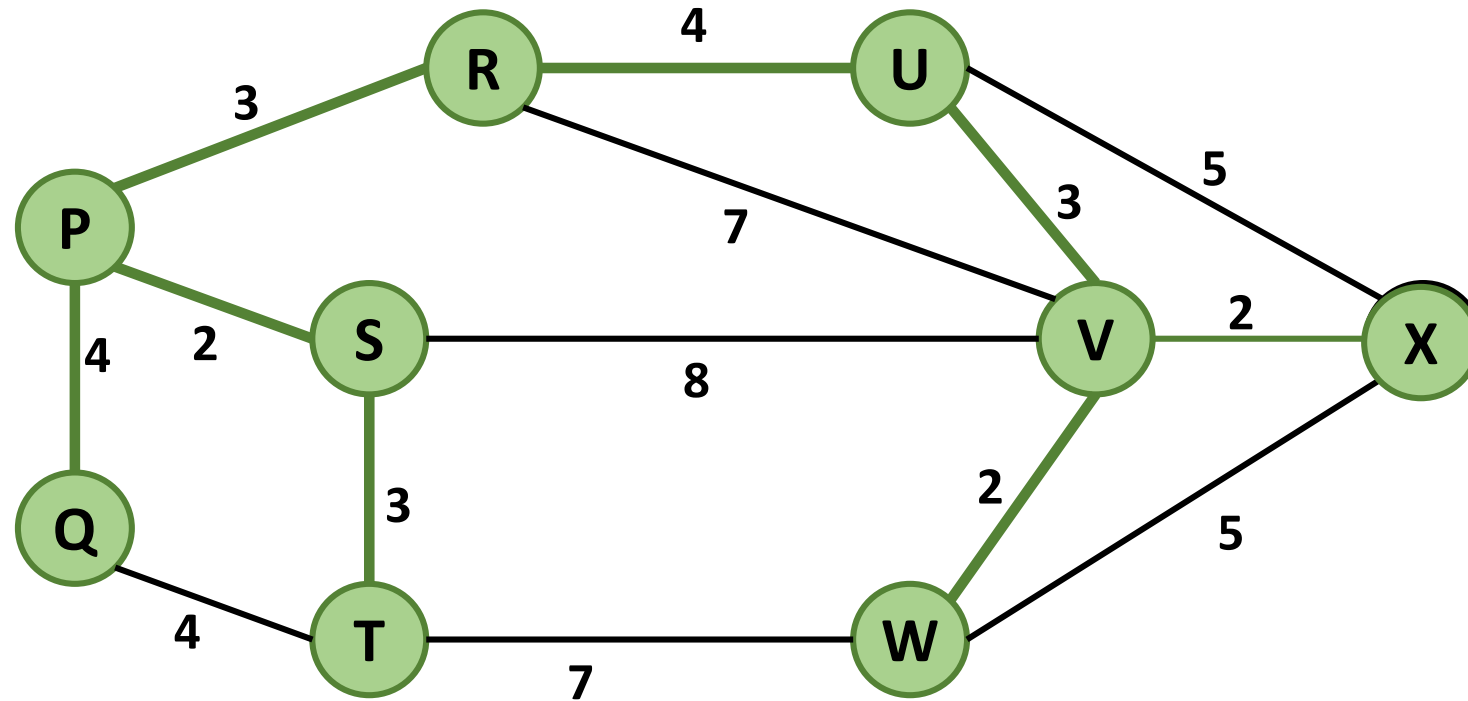
Prim's Algorithm – An Example: Step 7



Prim's Algorithm – An Example: Step 8



Prim's Algorithm – An Example: Step 9



Prim's Algorithm

- The algorithm was developed in 1930 by Czech mathematician Vojtěch Jarník
- Later rediscovered and republished by computer scientists Robert C. Prim in 1957 and Edsger W. Dijkstra in 1959.
- Therefore, it is also sometimes called the **Jarník's algorithm**, **Prim–Jarník algorithm**, **Prim–Dijkstra algorithm** or the **DJP algorithm**.

Minimum Spanning Trees

- A tree of the nodes of the graph with minimum total edge weight.
- Both Prim's and Kruskal's algorithms are examples of greedy algorithms
- **Applications**
 - Reducing copper to connect multiple nodes in a electrical/electronic circuit
 - Minimizing network length (cable cost) to connect multiple routers / computers etc.

Thank you!