

Graph Data Structures

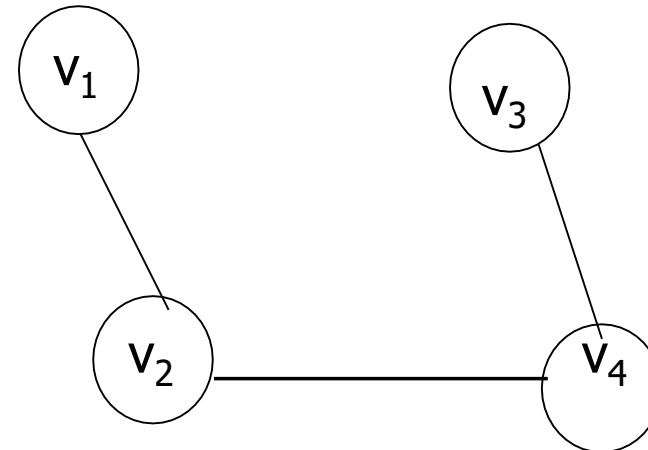
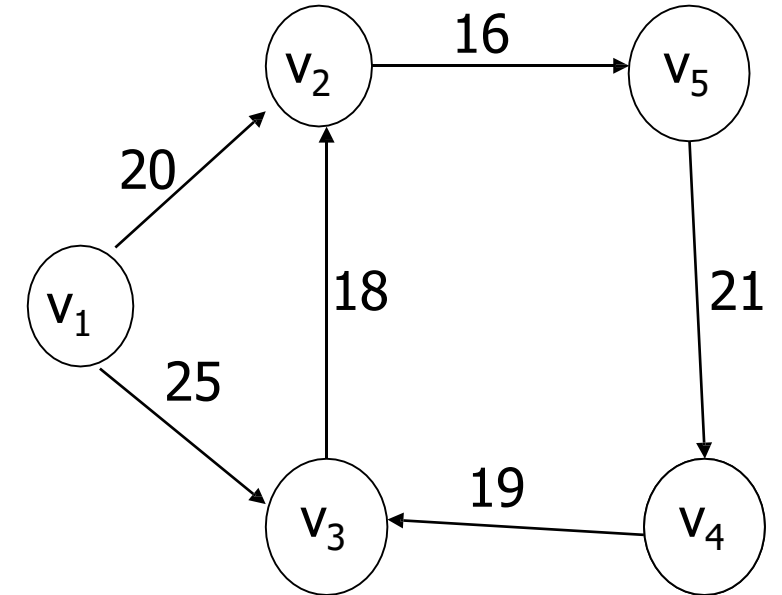
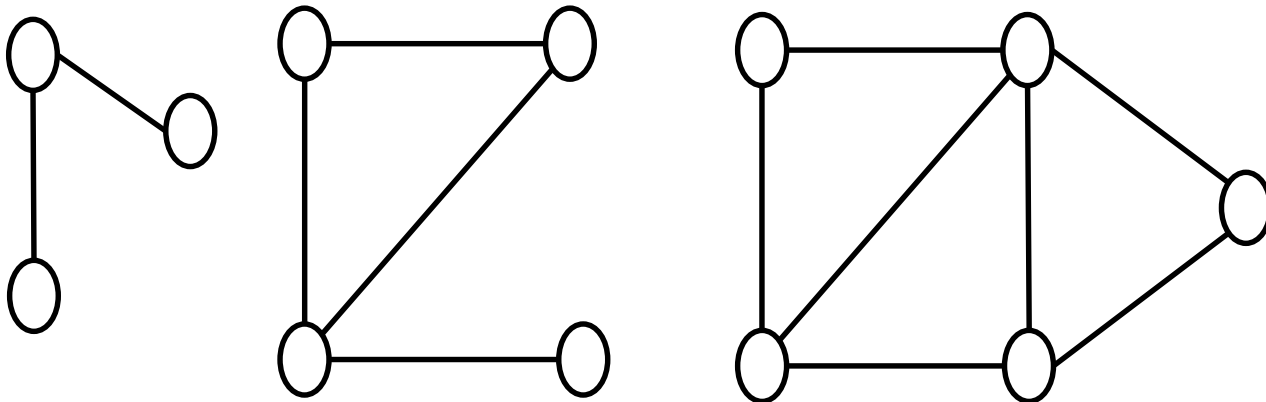
PROF. NAVRATI SAXENA

Graphs: Definition

- A mathematical non-linear data structure
- Capable of representing many kinds of physical structures
- A graph $G(V, E)$ is defined as a set of vertices (V) and a set of edges (E)
 - Vertices: Collection of nodes, represented as points or circles
 - Edges: Connect a pair of vertices and can have weights attached

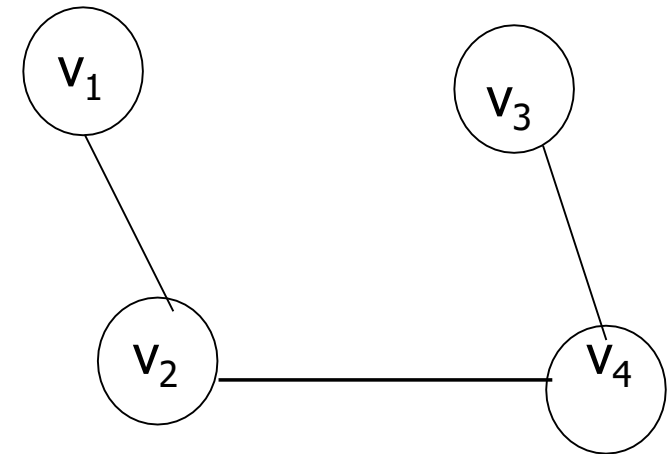
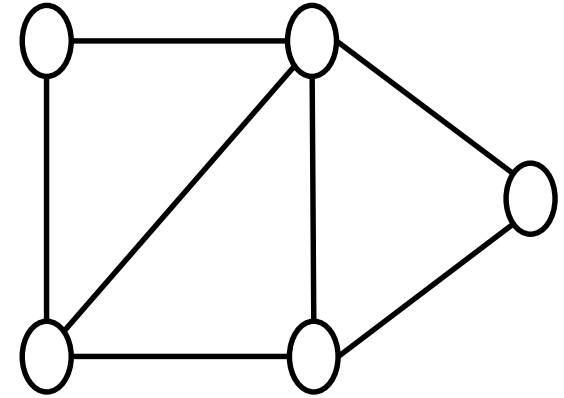
Graphs: Examples

- Most un-restricted form of data organization
- Final destination for problem solving
- Many variances (Many Types of Graphs)
 - Directed and un-directed graphs
 - Connected and un-connected graphs
 - Weighted graphs



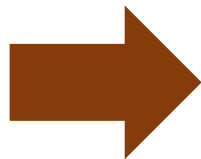
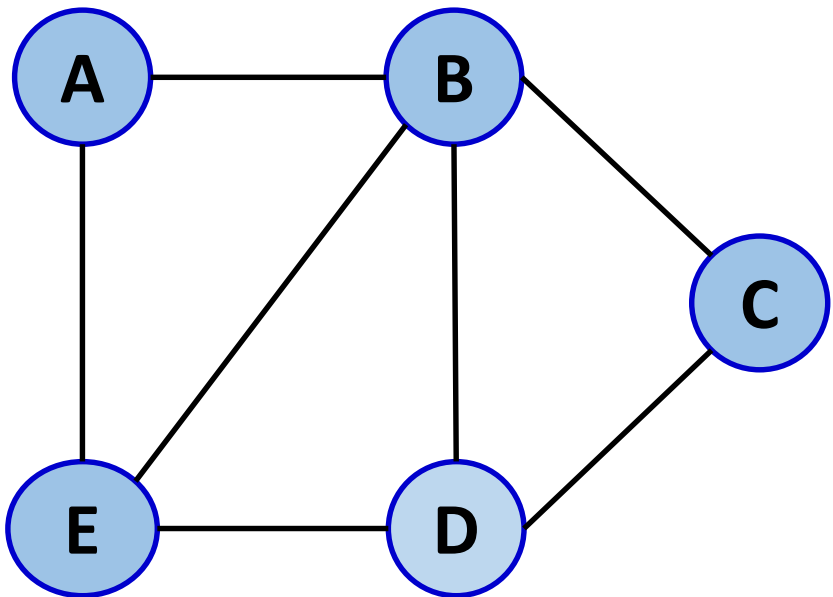
Types of Graph

- Undirected graph
- Directed graph
- Weighted graph
- Connected graph
- Non-connected graph



Adjacency Matrix Representation

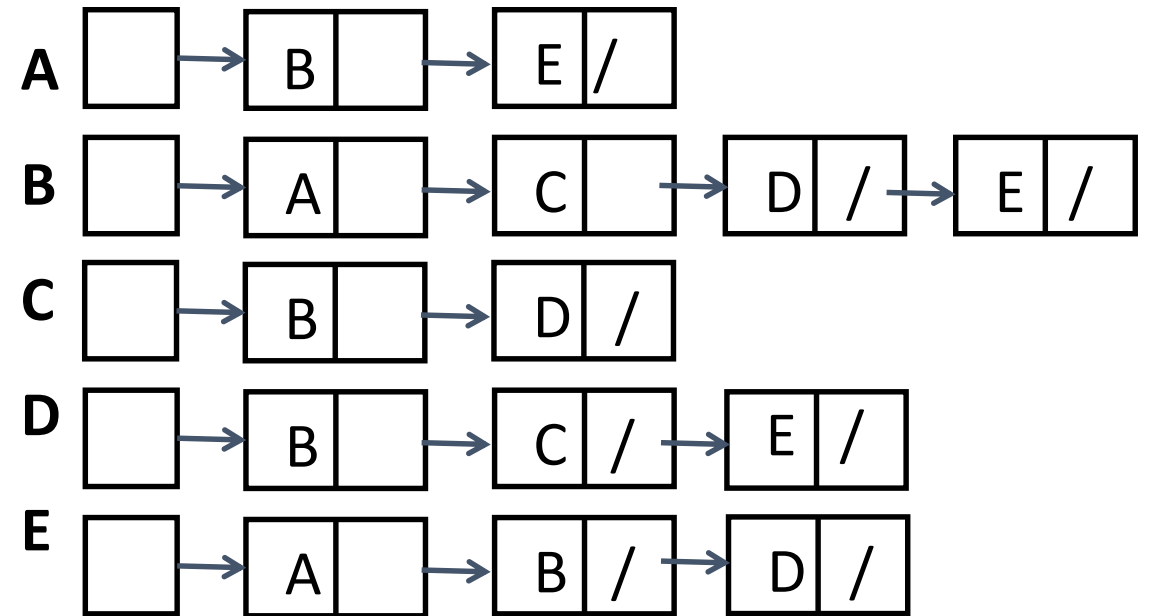
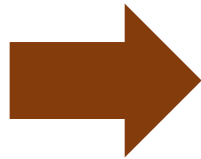
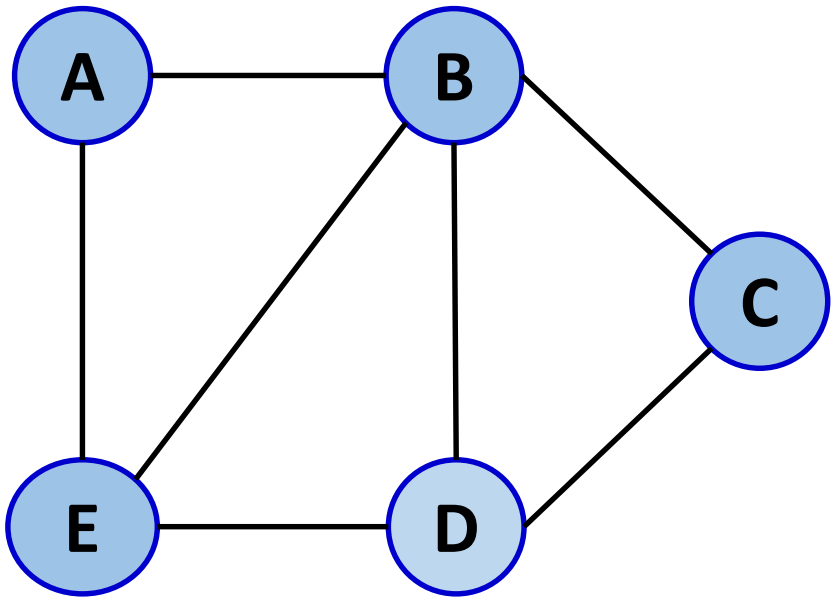
1. Create a Matrix (e.g. use a 2-dimensional array)
2. Any index i represents a node
3. Any entry (i, j) in the matrix represents connectivity between two nodes i, j
 1. $\text{entry}(i, j) = 1 \Rightarrow \text{an edge exists}$
 2. $\text{entry}(i, j) = 0 \Rightarrow \text{no edge exists}$



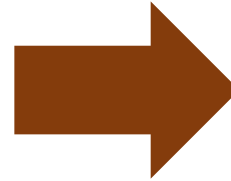
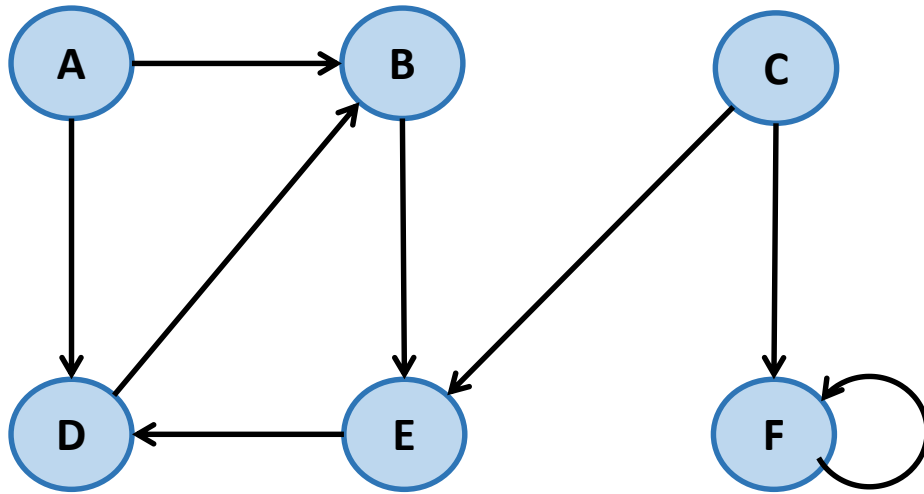
	A	B	C	D	E
A	0	1	0	0	1
B	1	0	1	1	1
C	0	1	0	1	0
D	0	1	1	0	1
E	1	1	0	1	0

Adjacency List Representation

1. Create an array of Lists (e.g. use a linked-list)
2. The headers of the lists represents the nodes
3. Every list represents the nodes, connected from the header node

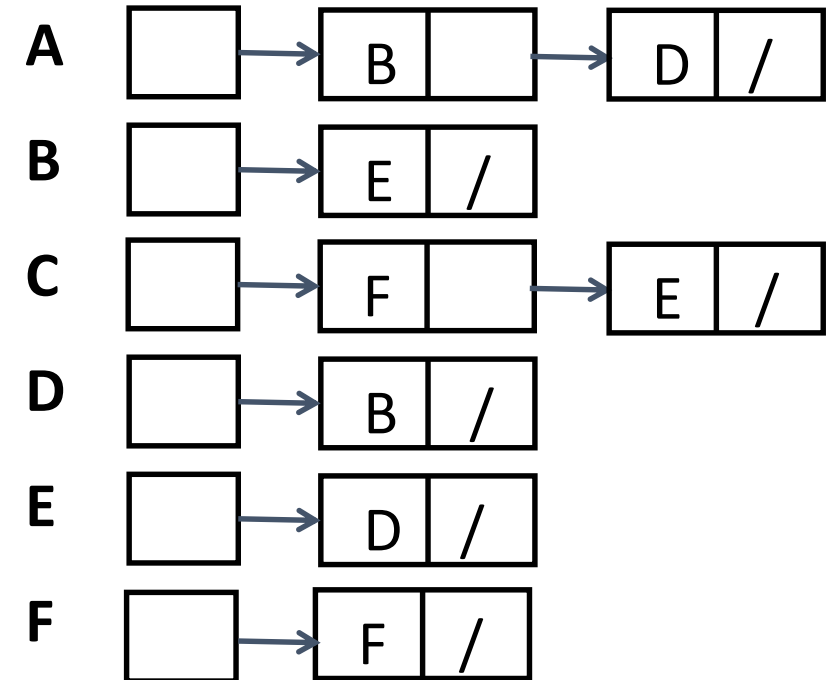
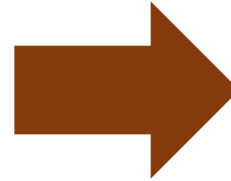
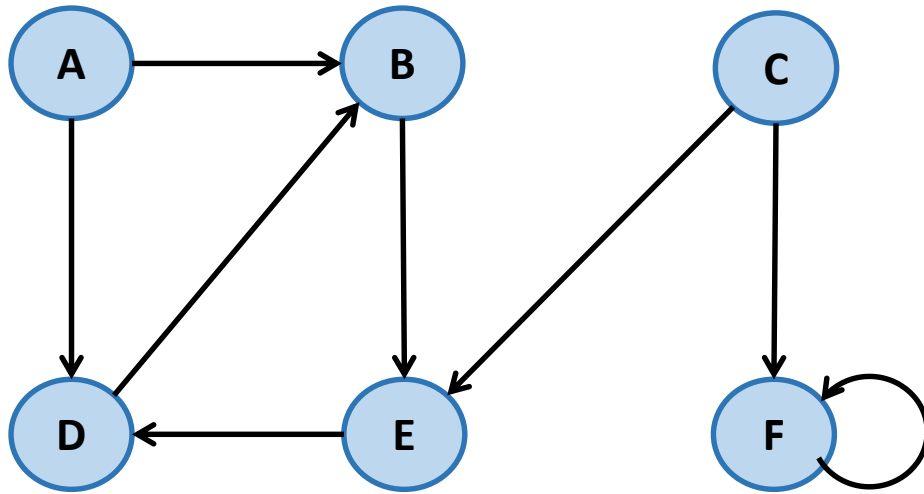


Adjacency Matrix for Directed Graph



	A	B	C	D	E	F
A	0	1	0	1	0	0
B	0	0	0	0	1	0
C	0	0	0	0	1	1
D	0	1	0	0	0	0
E	0	0	0	1	0	0
F	0	0	0	0	0	1

Adjacency List for Directed Graph



Major Graph Operations

- Insert
- Delete
- Search (Traversal)
 - Breadth-First Search (BFS)
 - Depth-First Search (DFS)

Breadth-First Search (BFS Traversal)

- **Idea**

- Start from the starting node (this will be given to you)
- Generate a traversal tree while traversing the graph
- A node is **traversed** when its all successor nodes are generated

- **Implementation:** Use of a coloring scheme with a queue



Orange: Encountered but not traversed

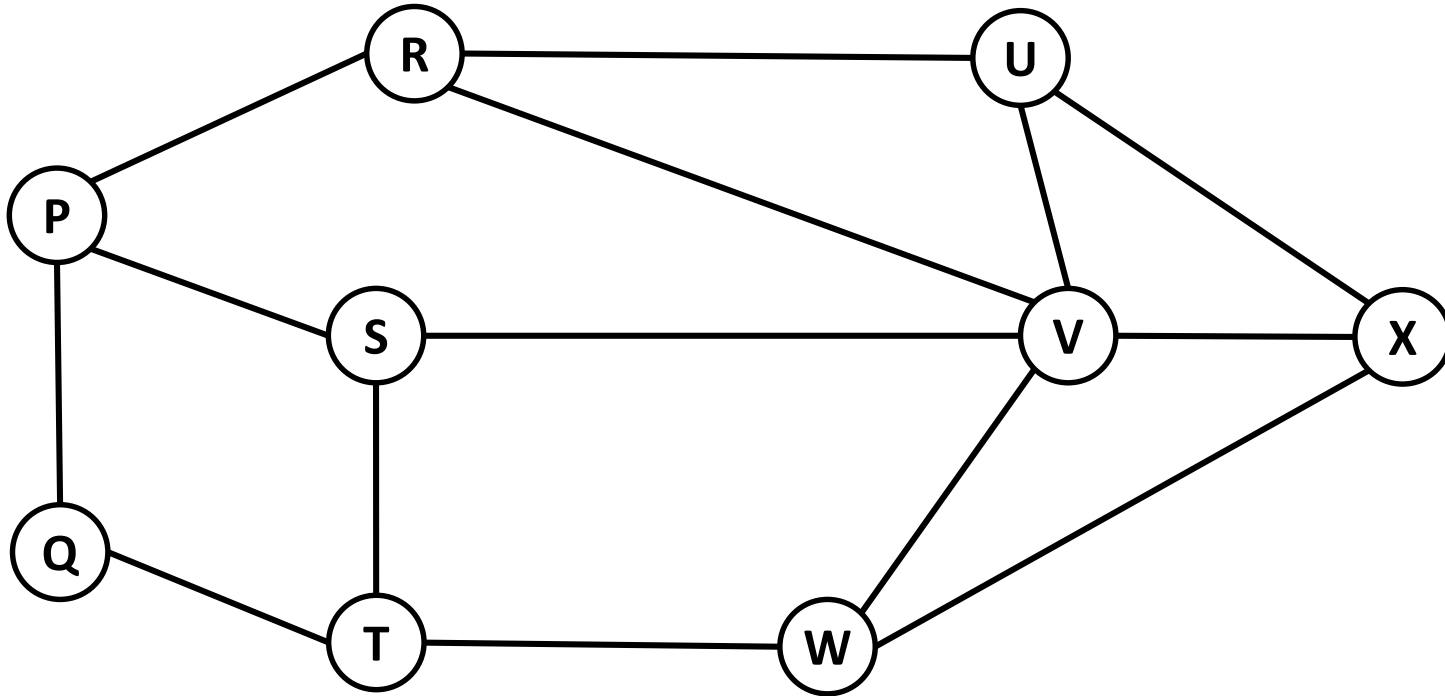


Green: traversed



White: not accessed/encountered/generated

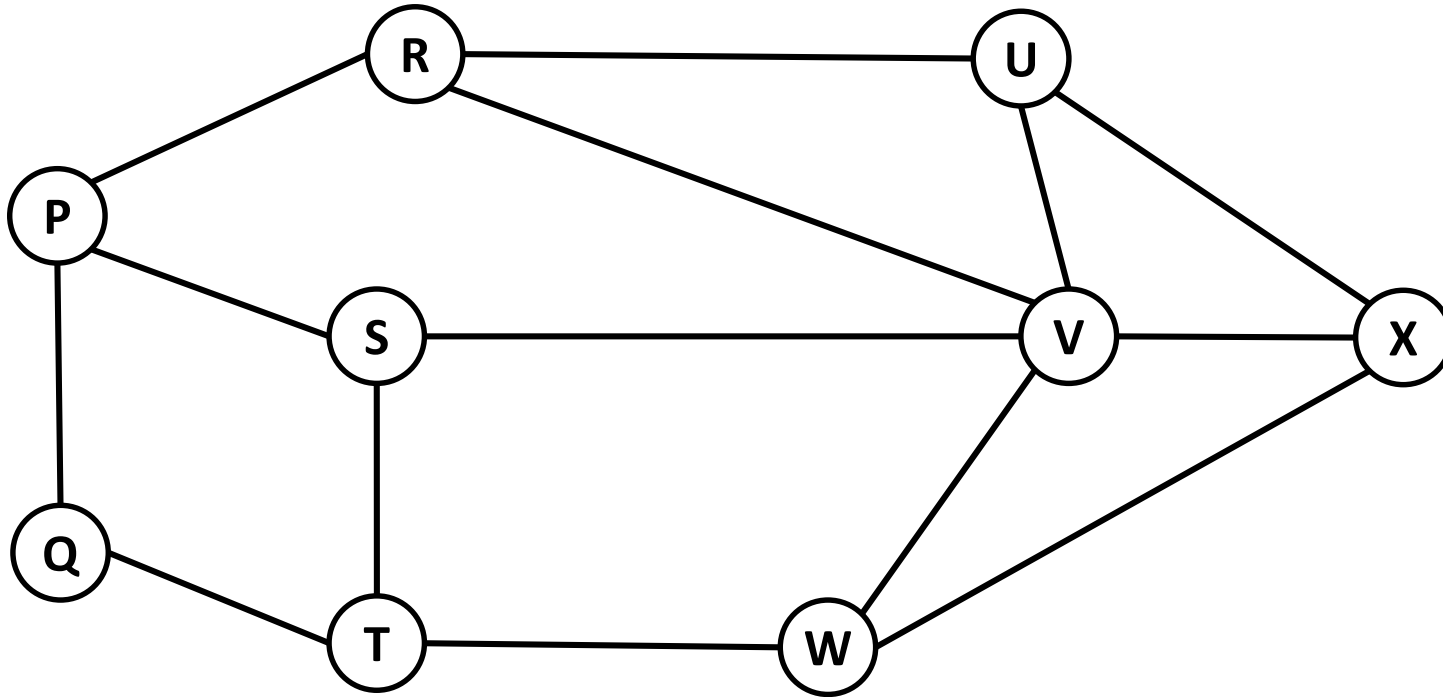
BFS Example – Initialization



Queue



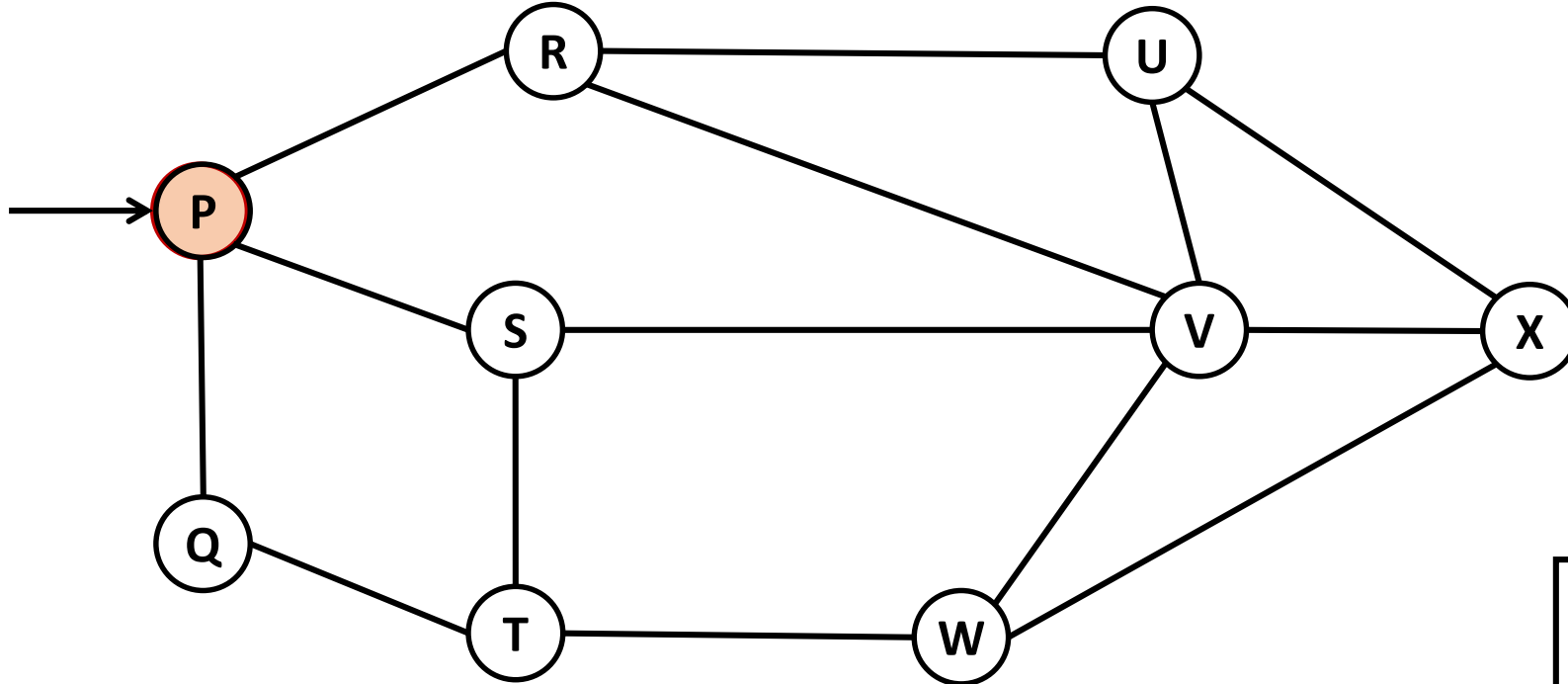
BFS Example: Step-1



Queue



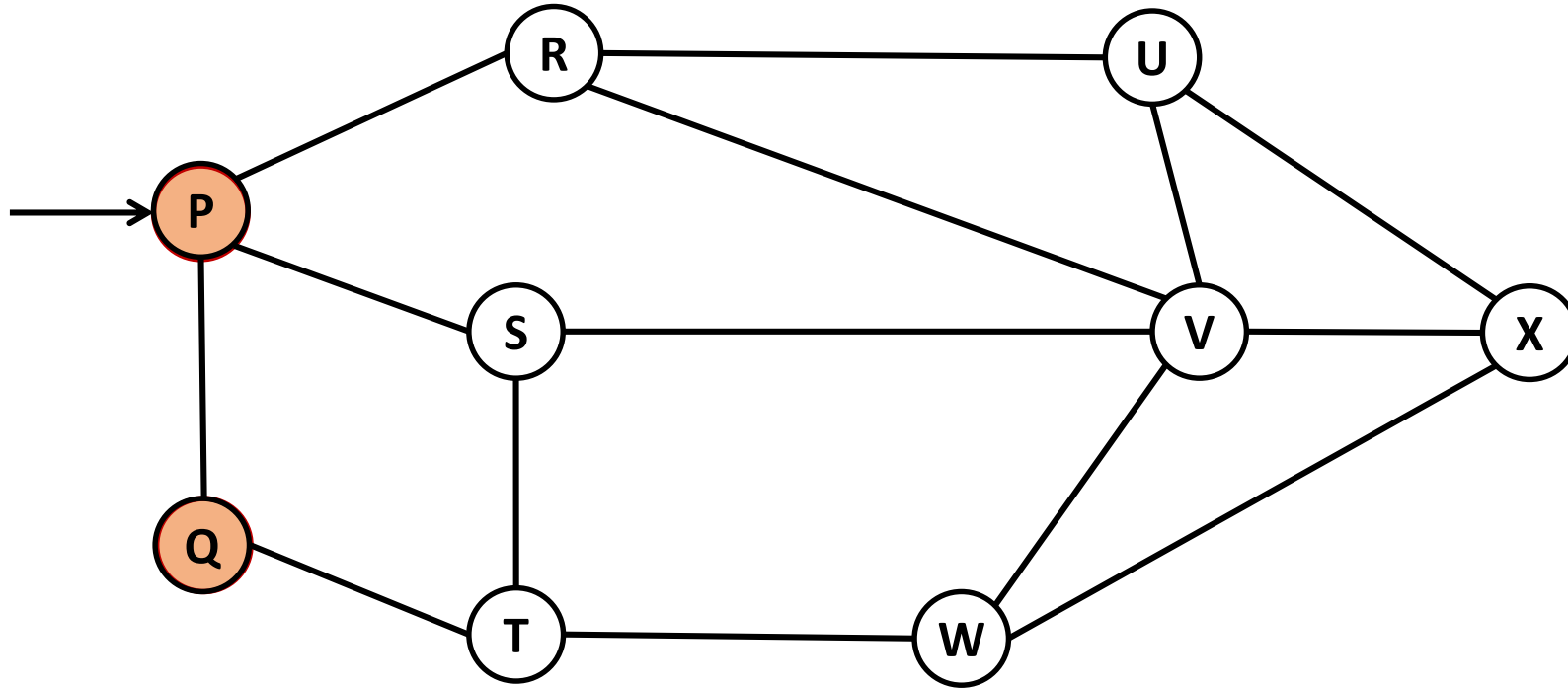
BFS Example: Step-2



Queue

P						
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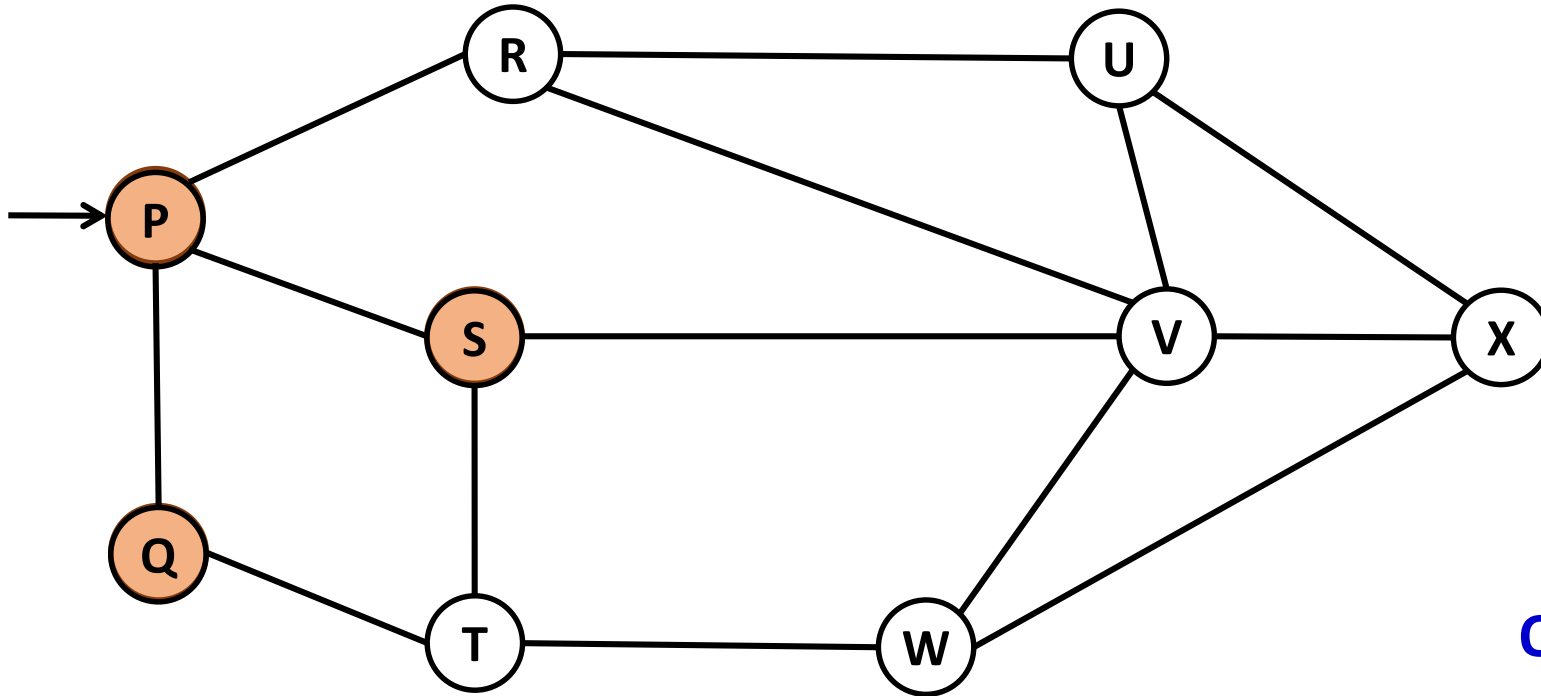
BFS Example: Step-3



Queue

P	Q					
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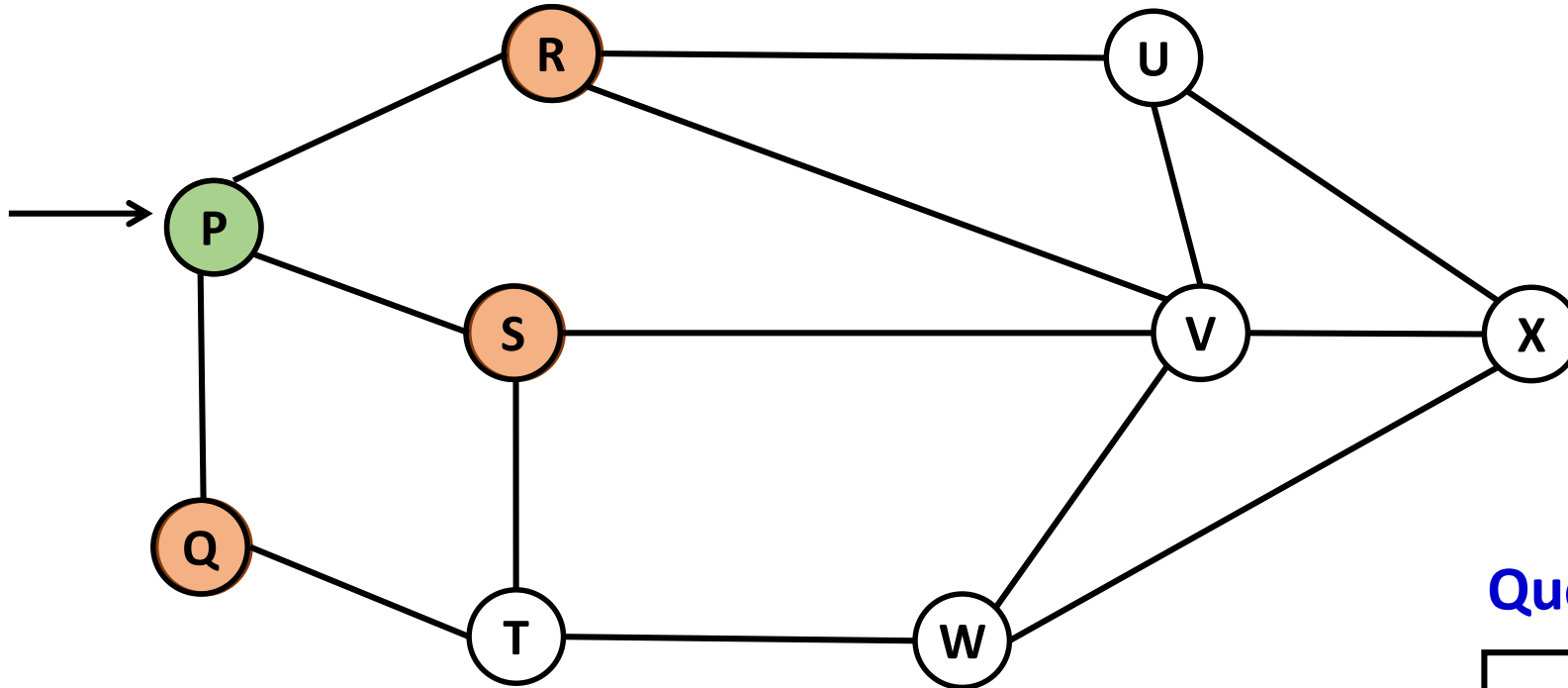
BFS Example: Step-4



Queue

P	Q	S				
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BFS Example: Step-5



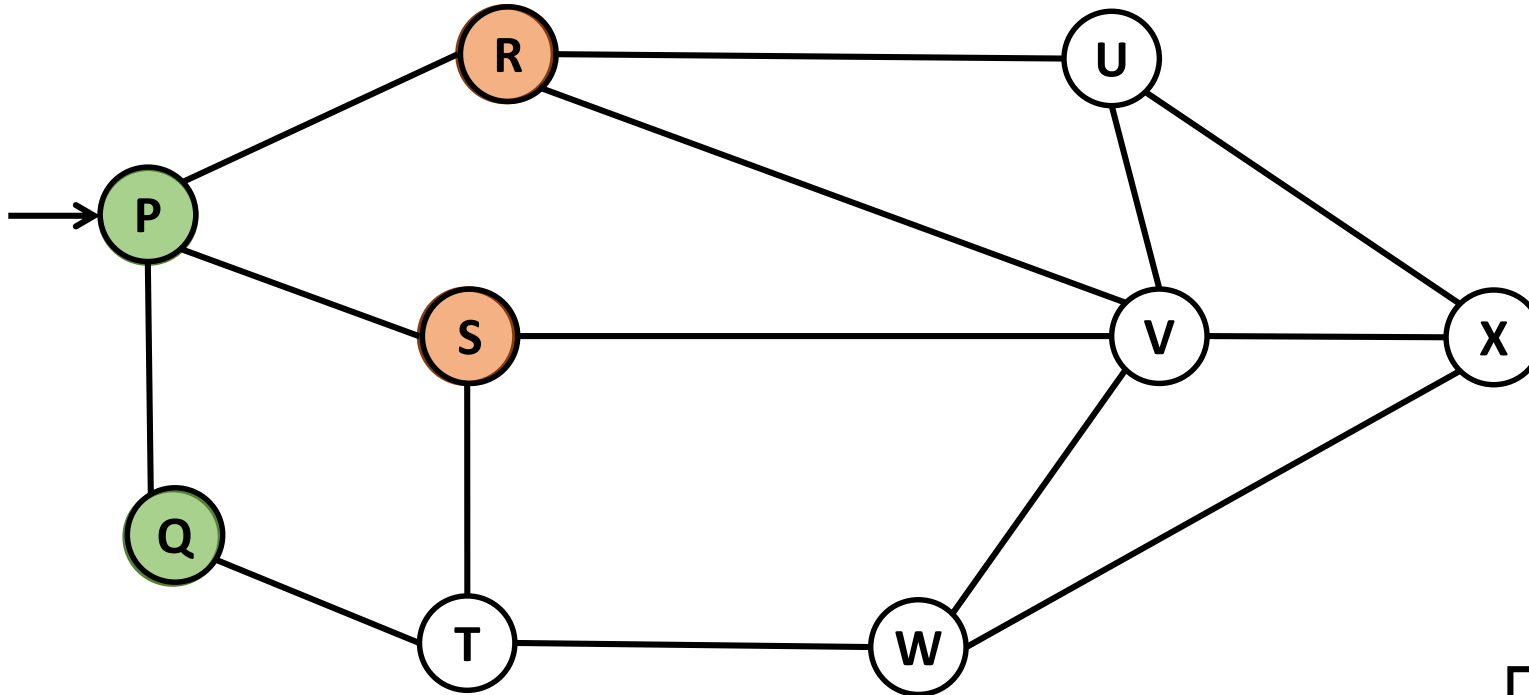
Traversed List

P

Queue

Q	S	R	T			
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BFS Example: Step-6



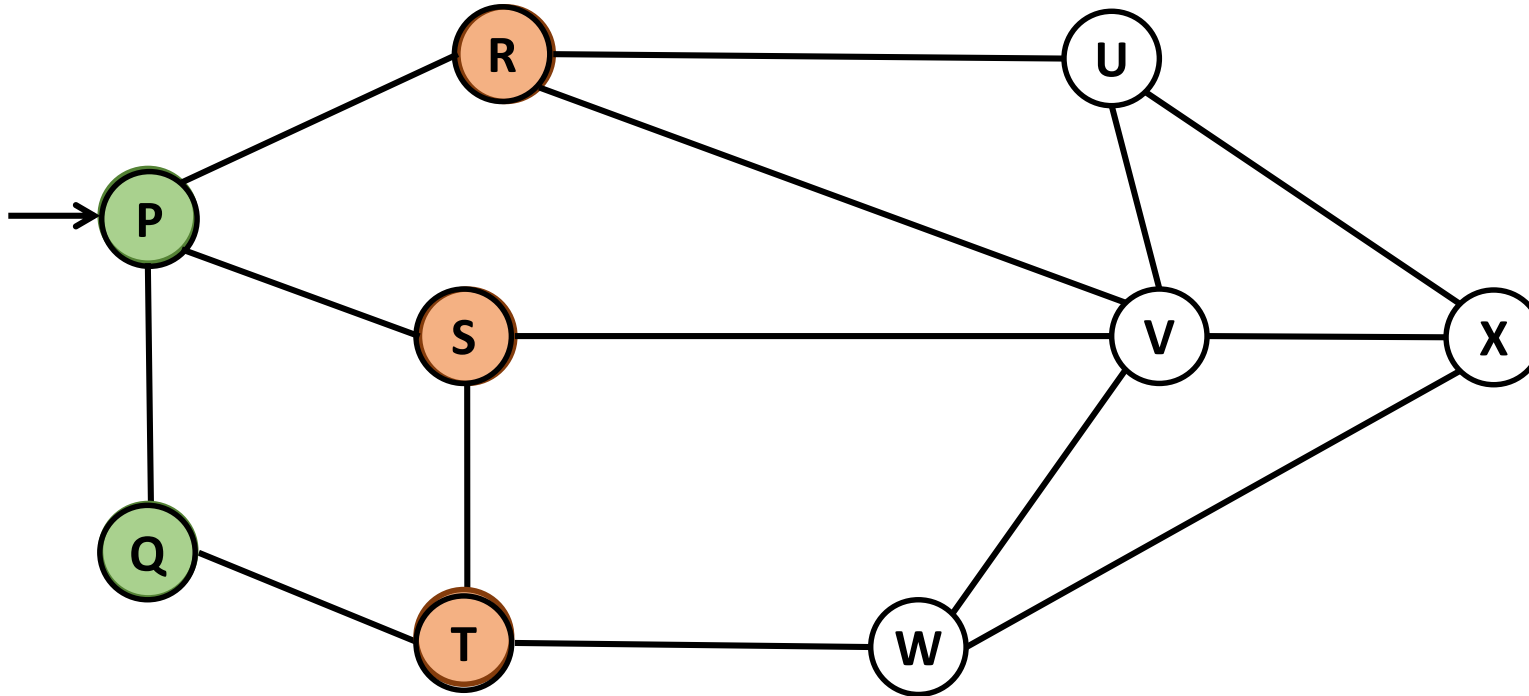
Queue

R	S	T				
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Traversed List

PQ

BFS Example: Step-7



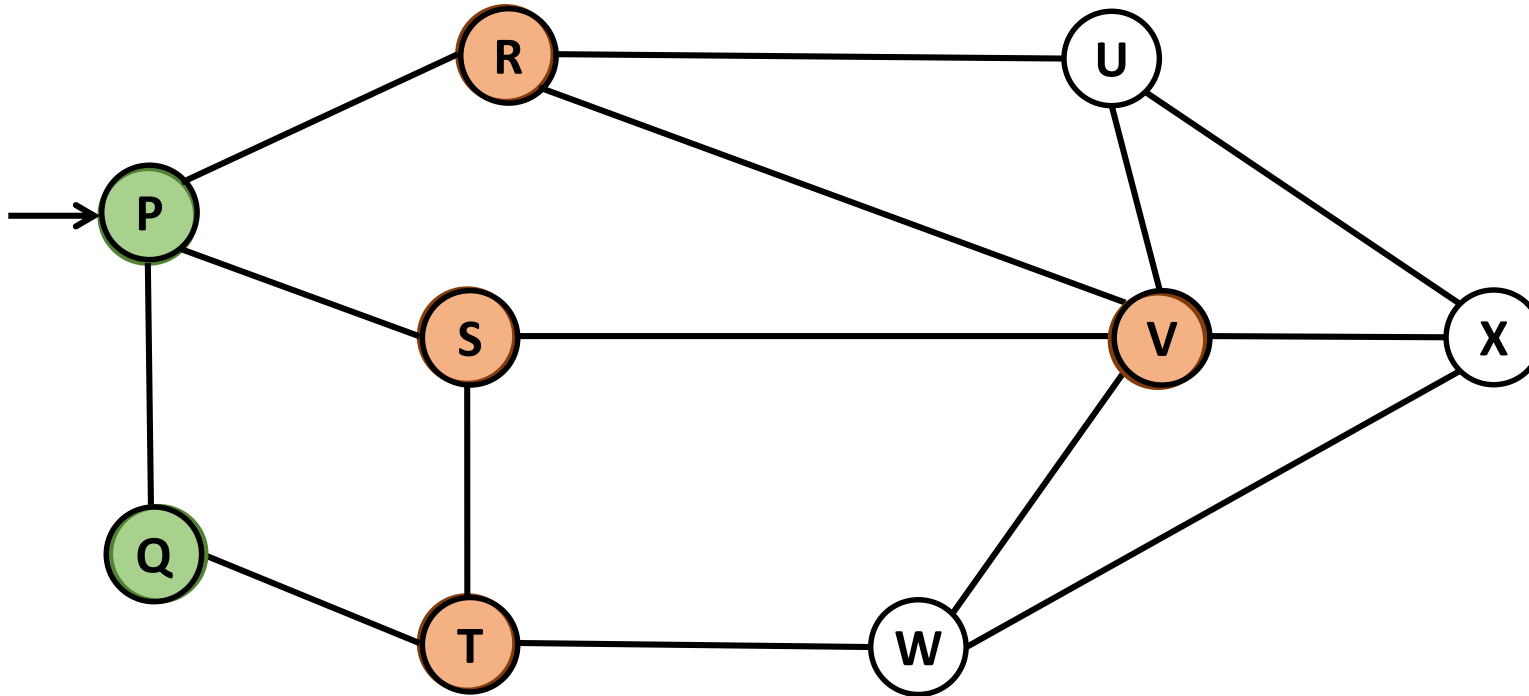
Queue

R	S	T				
---	---	---	--	--	--	--

Traversed List

PQ

BFS Example: Step-8



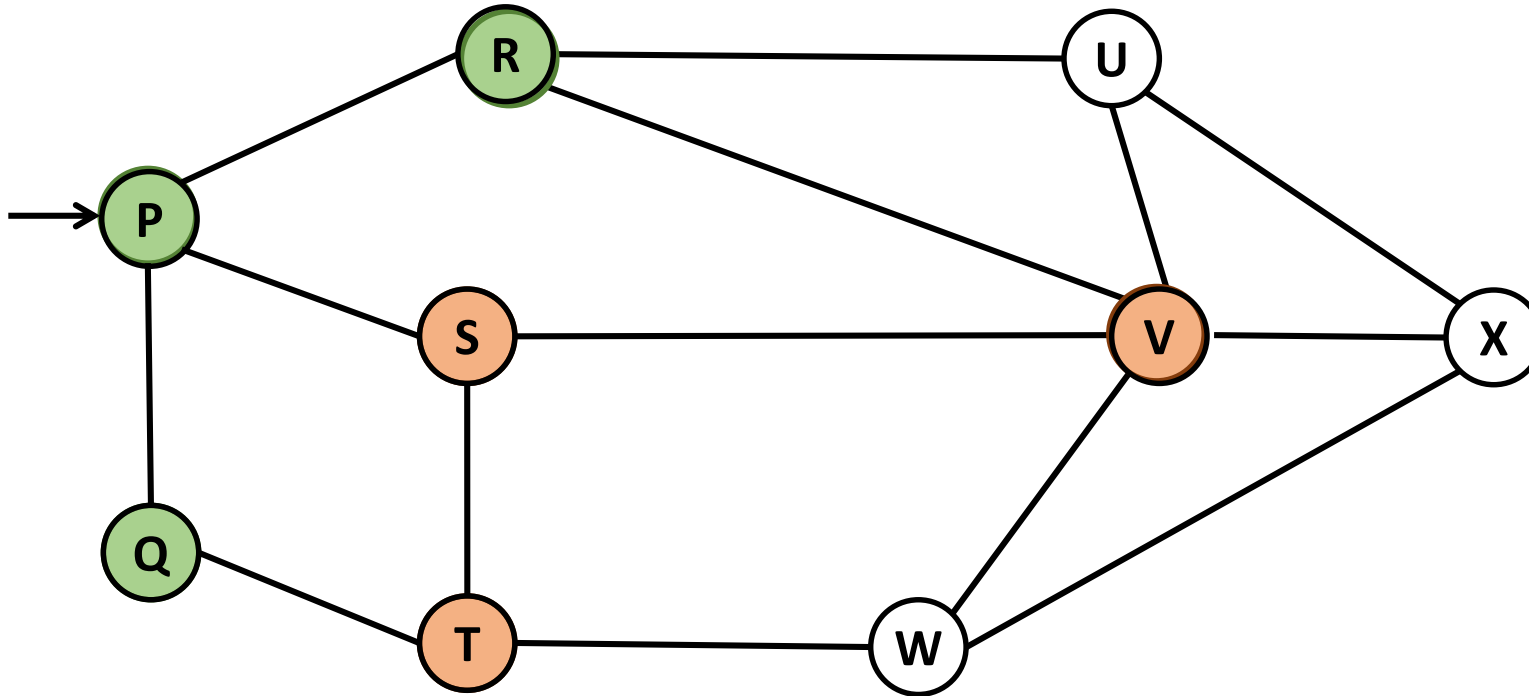
Traversed List

PQ

Queue

R	S	T	V			
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BFS Example: Step-9



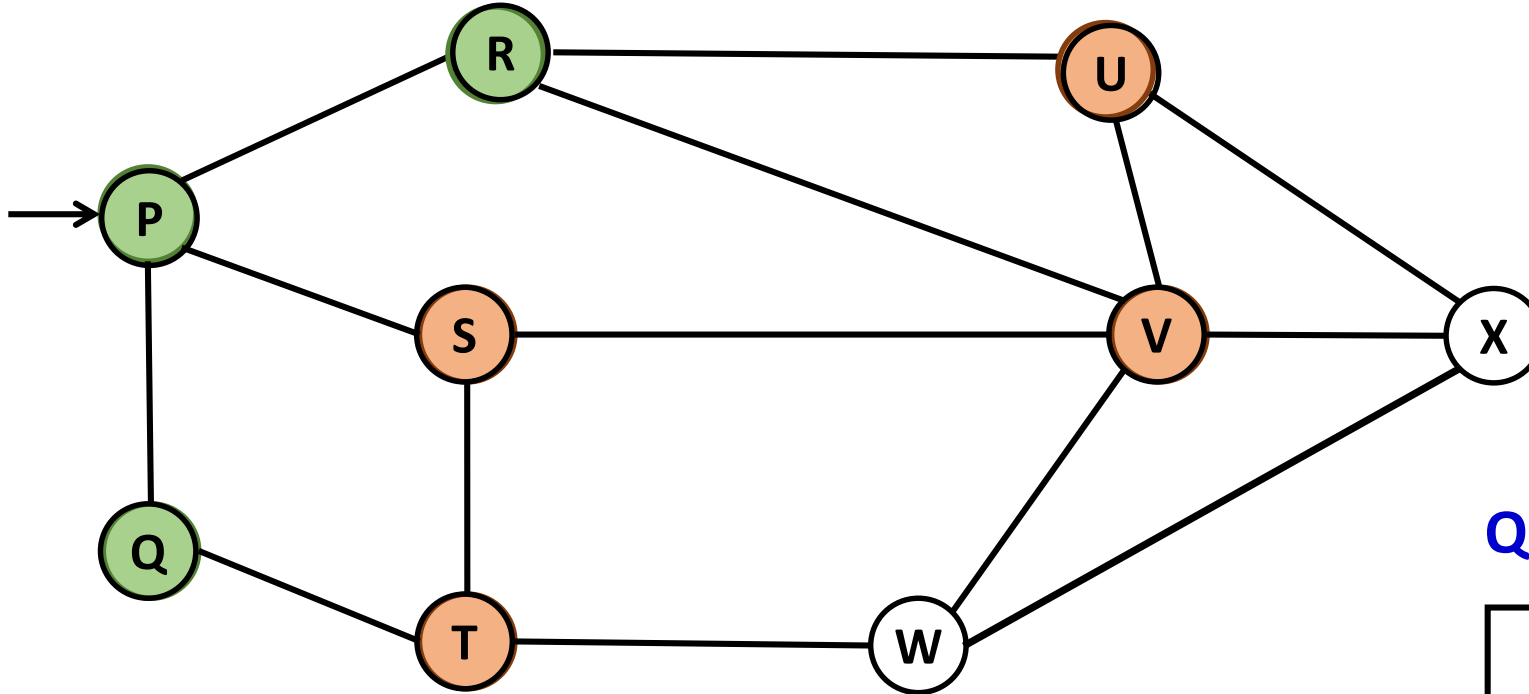
Traversed List

PQR

Queue

S	T	V				
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BFS Example: Step-10



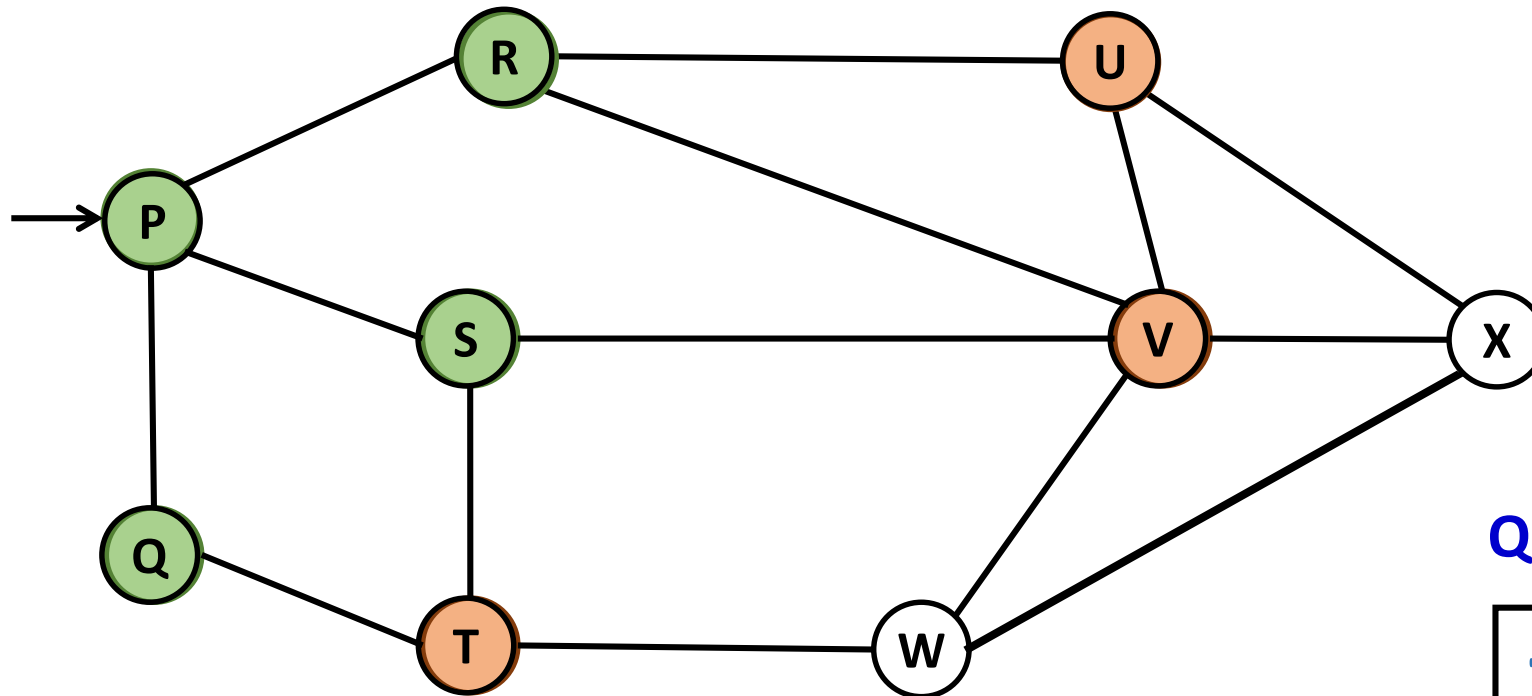
Traversed List

PQR

Queue

S	T	V	U			
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BFS Example: Step-11



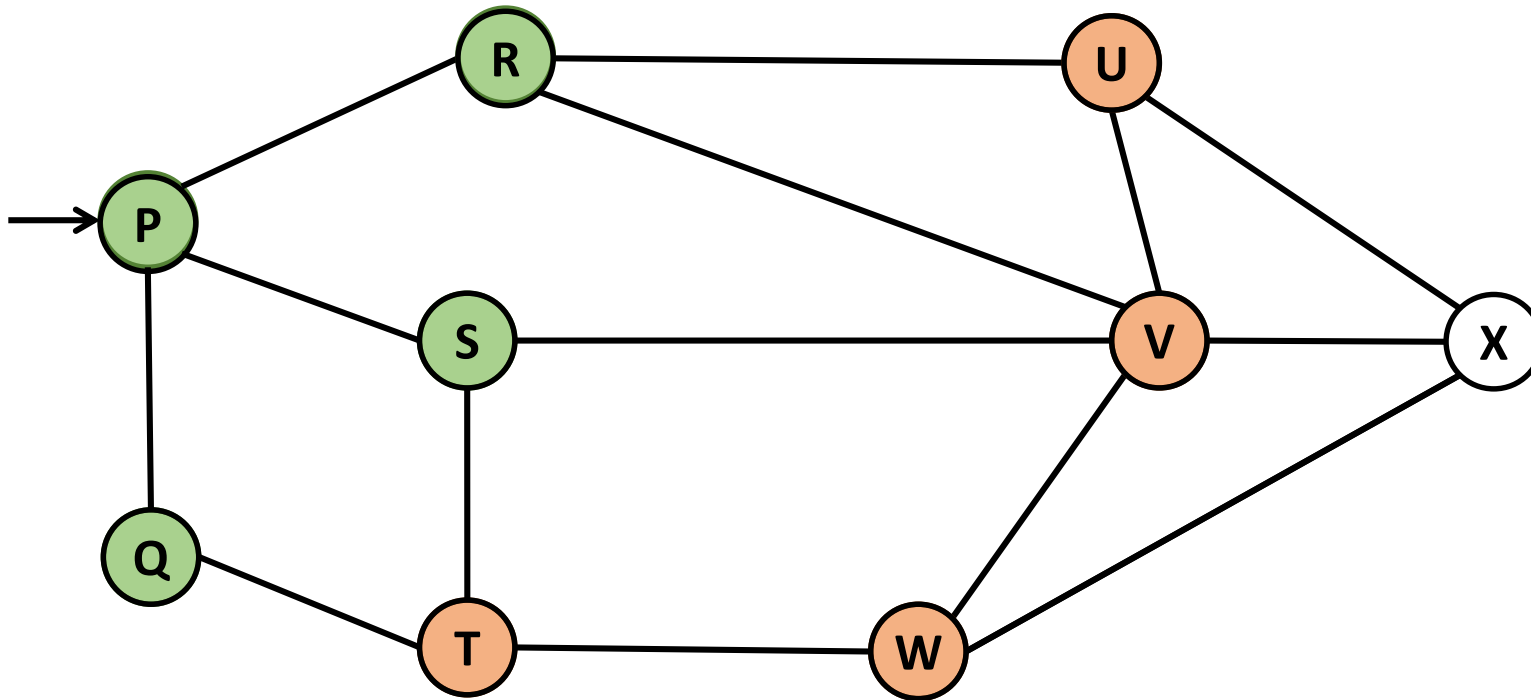
Traversed List

PQRS

Queue

T	V	U				
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BFS Example: Step-12



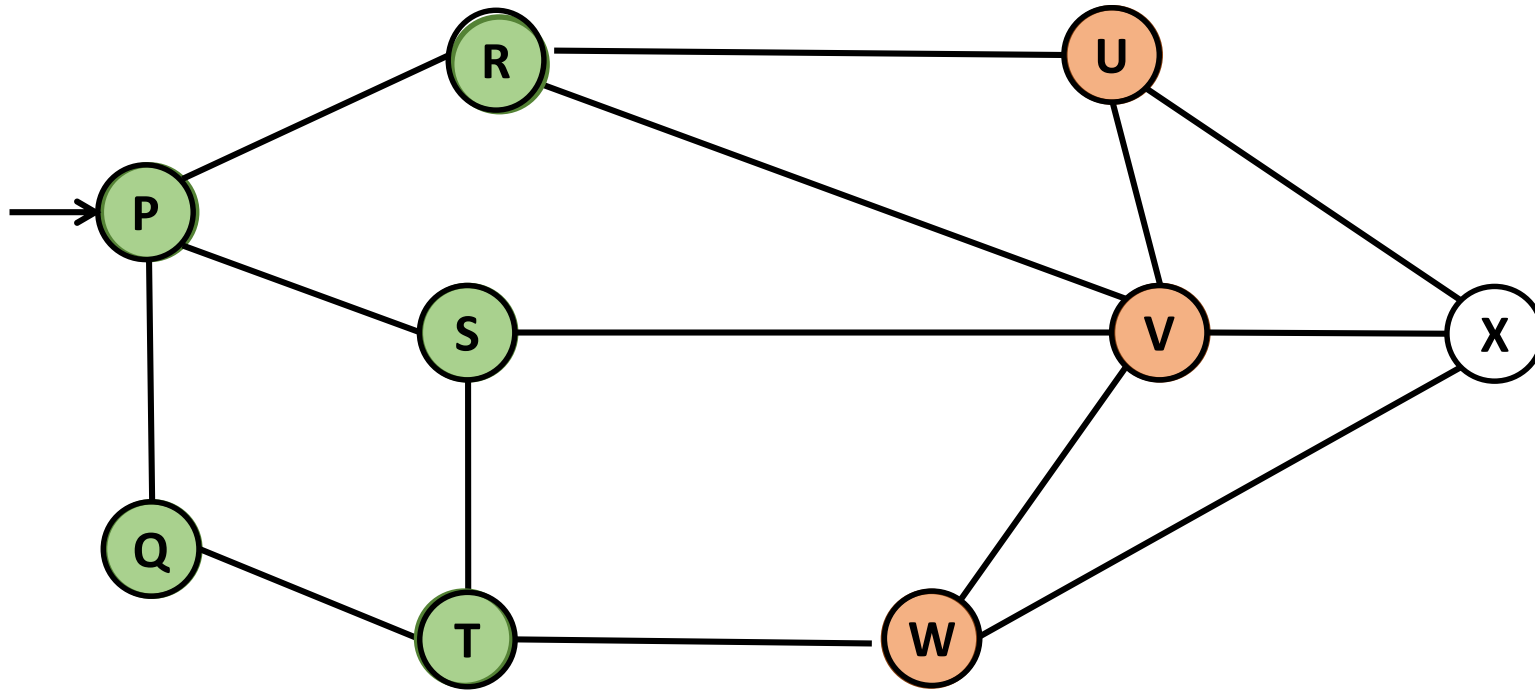
Queue

T	V	U	W			
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Traversed List

PQRS

BFS Example: Step-13



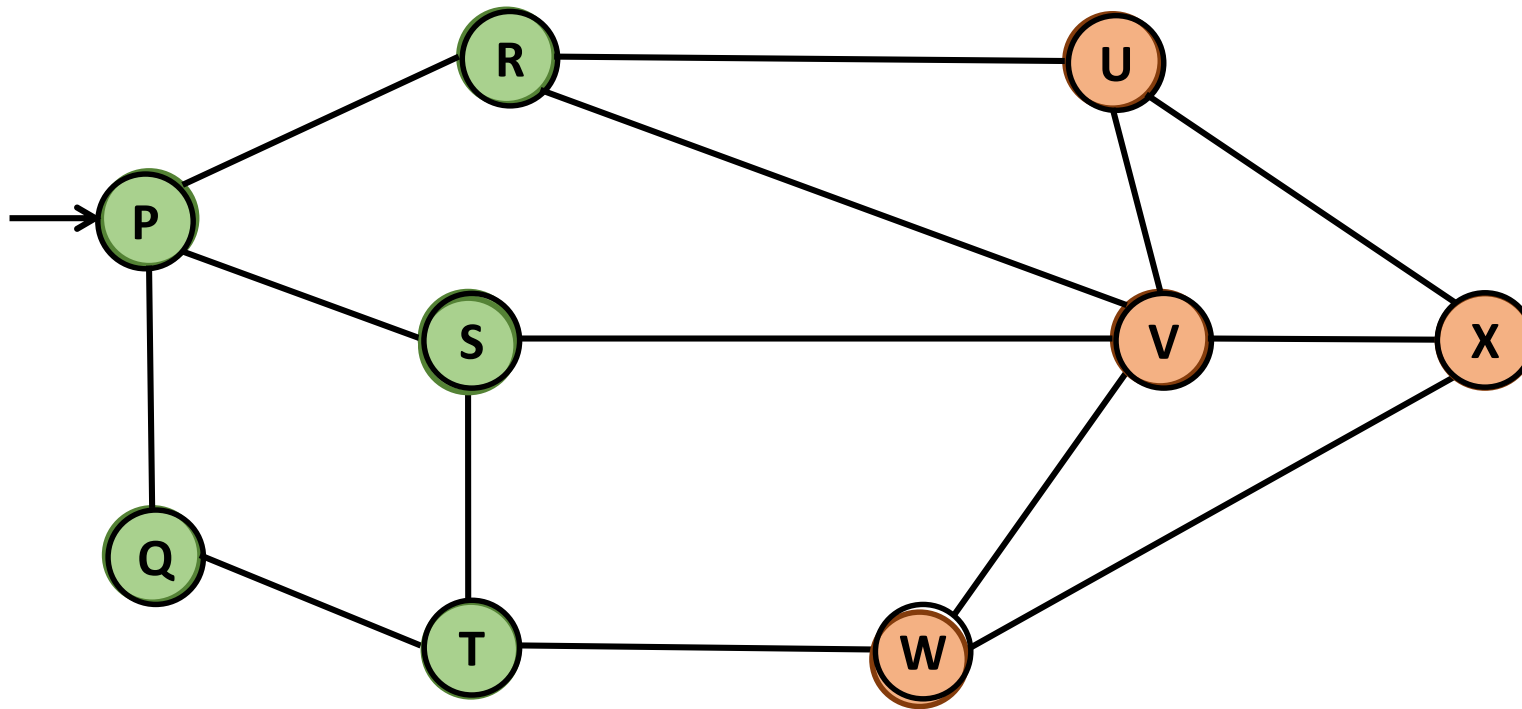
Traversed List

PQRST

Queue

V	U	W				
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BFS Example: Step-14



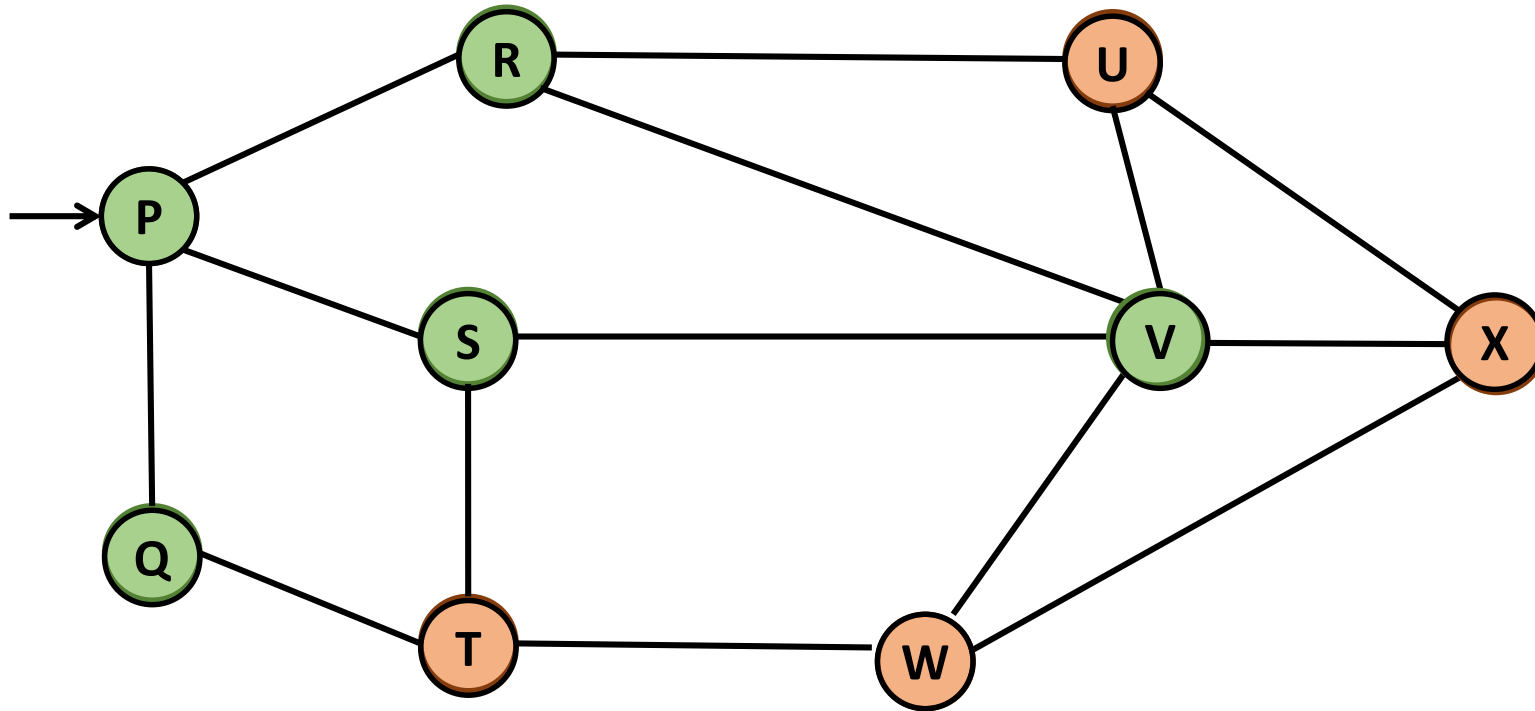
Traversed List

PQRST

Queue

V	U	W	X			
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BFS Example: Step-15



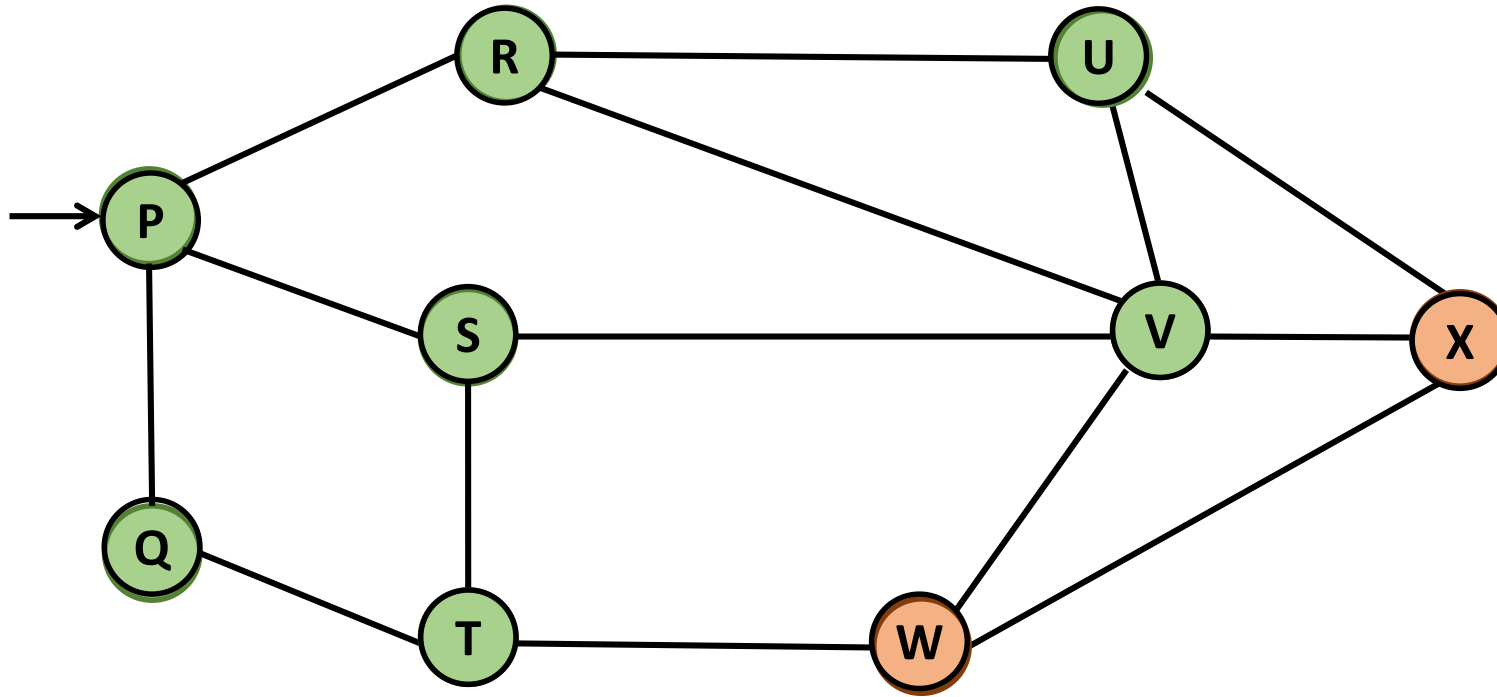
Traversed List

PQRSTV

Queue

U	W	X				
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BFS Example: Step-16



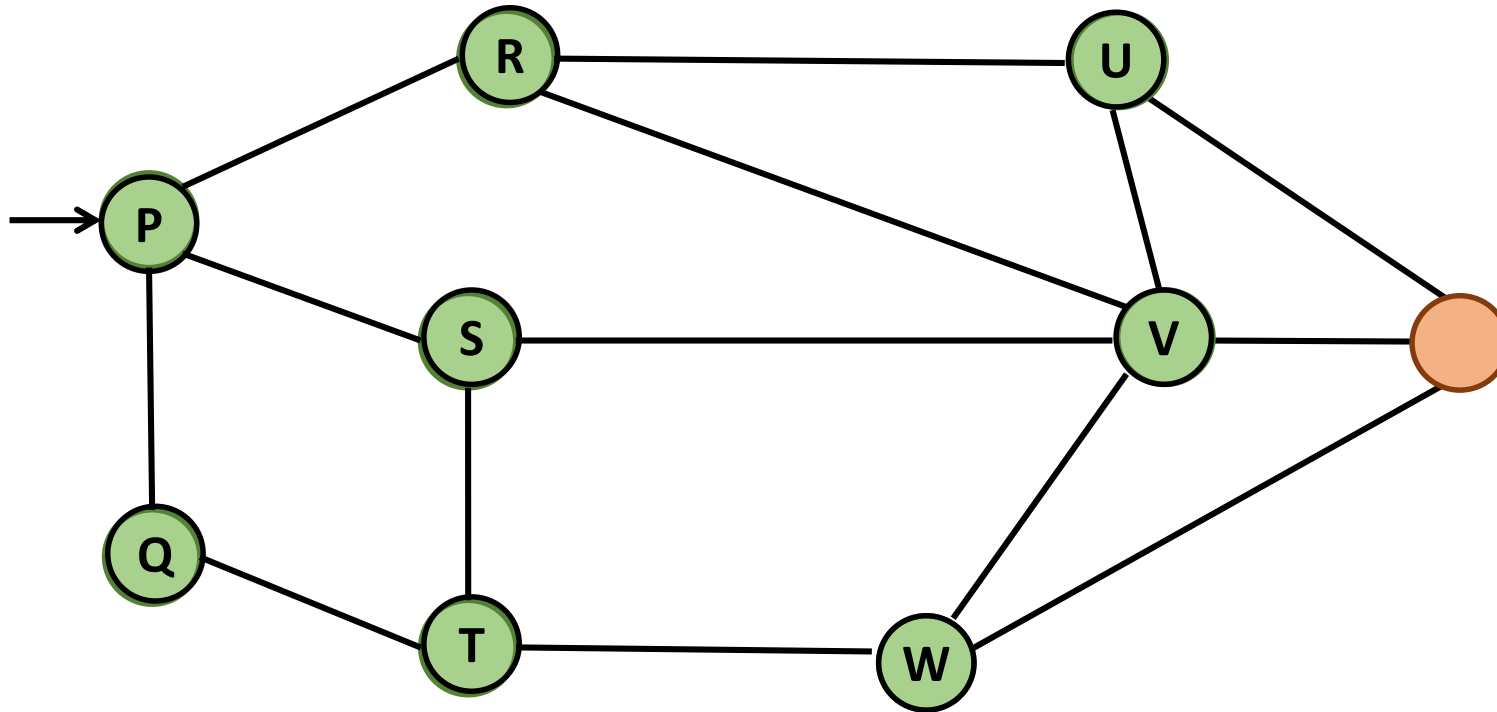
Traversed List

PQRSTVU

Queue



BFS Example: Step-17



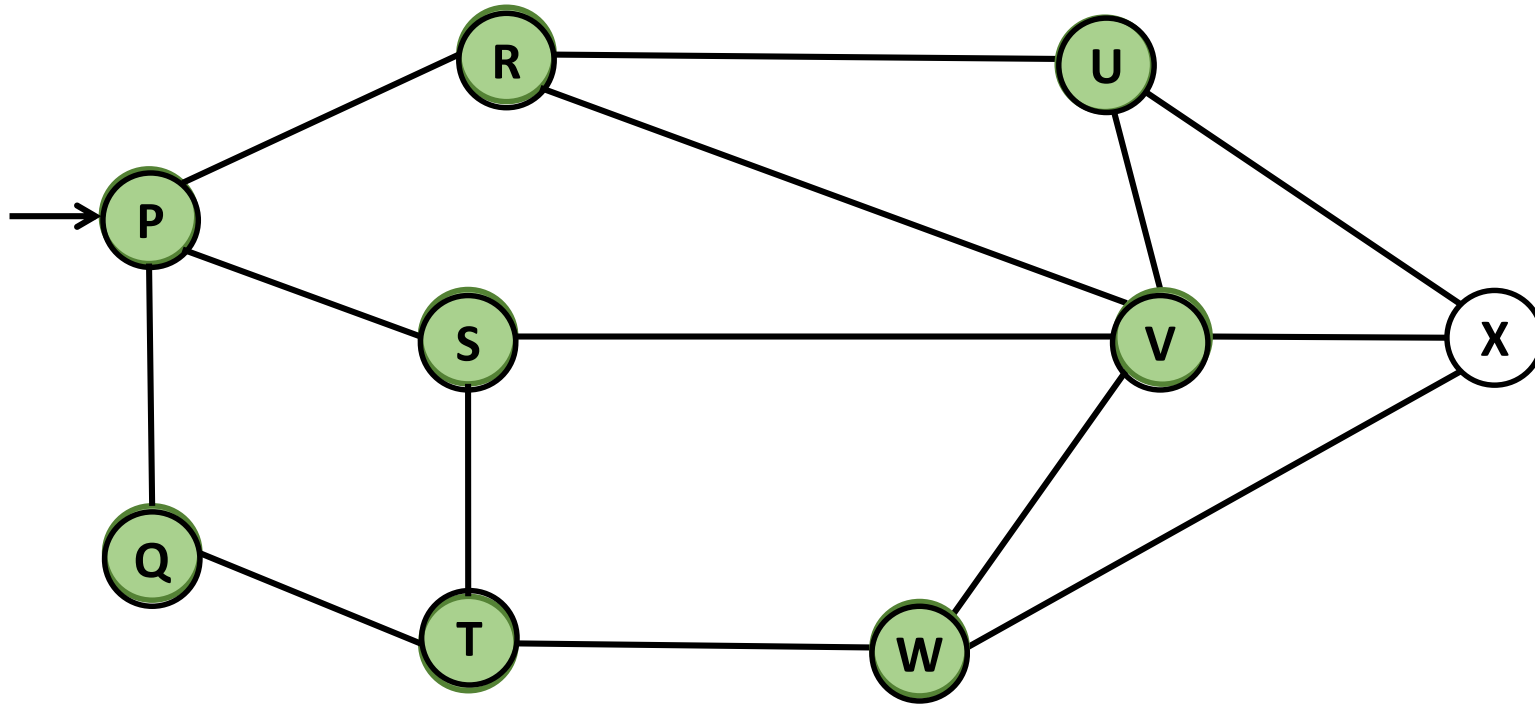
Traversed List

PQRSTUVWXYZ

Queue

X						
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BFS Example: Step-18



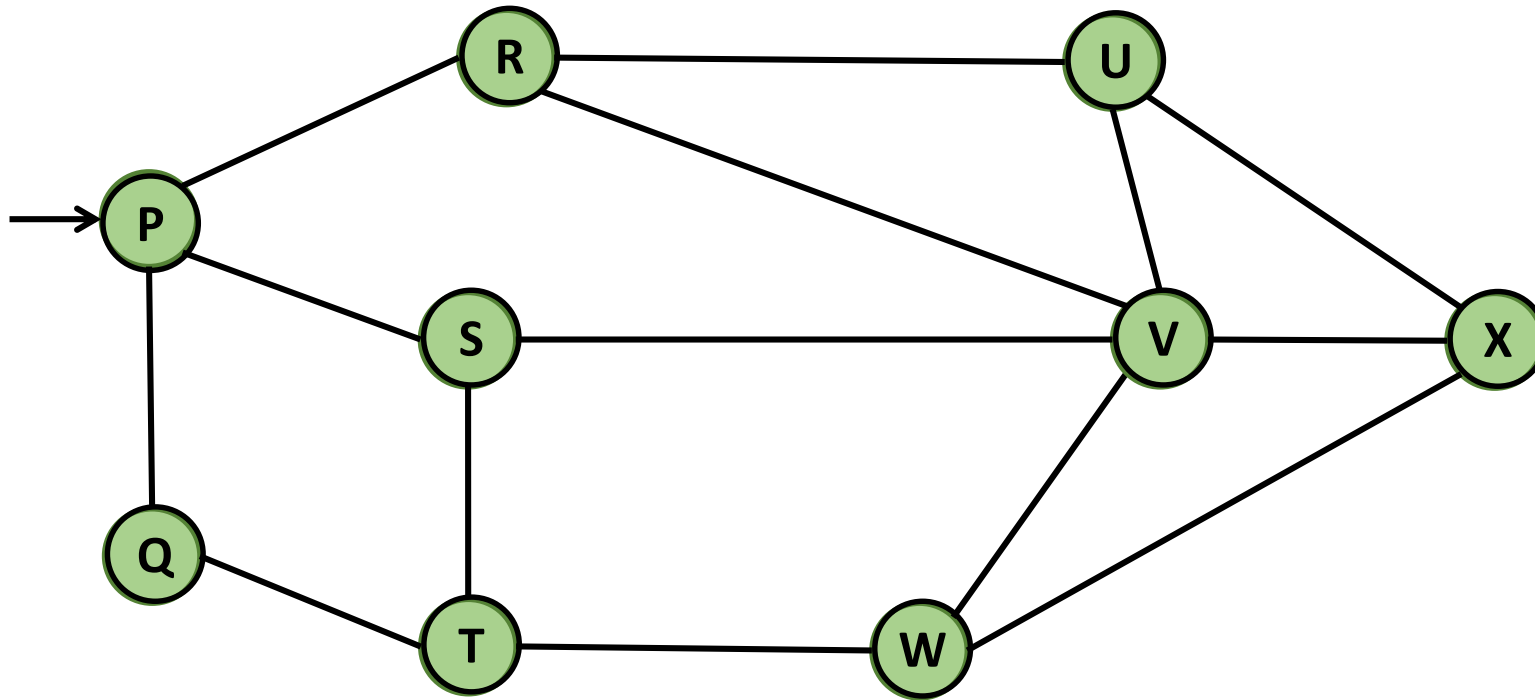
Queue



Traversed List

PQRSTUVWXYZ

BFS Implementation: Step-23



Queue



Traversed List

PQRSTVUWX

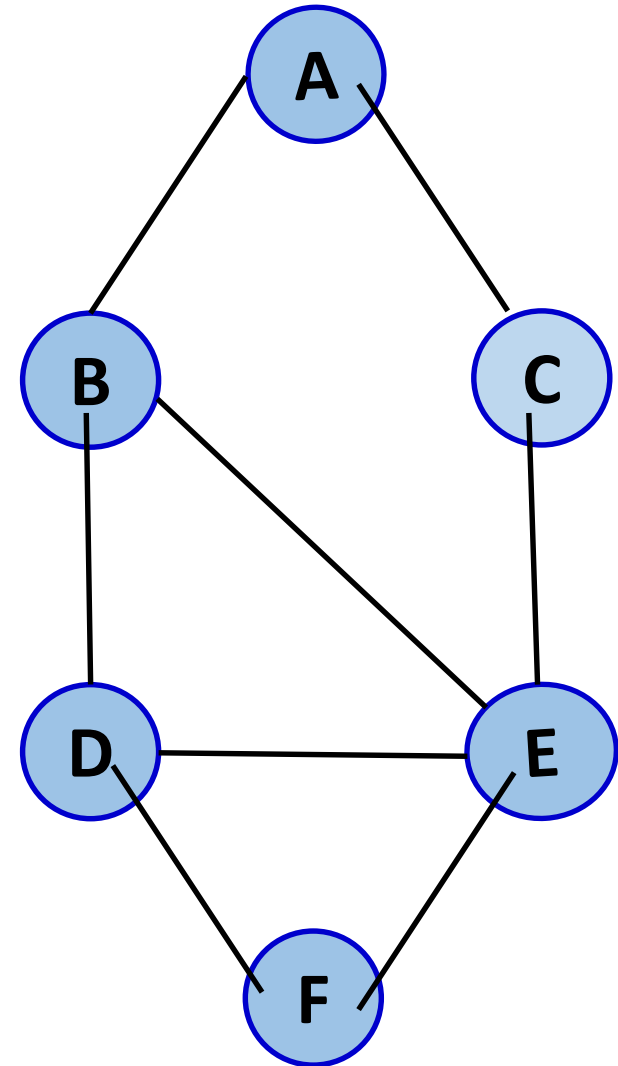
Graph: Depth-First Traversal

- **Key Idea**

1. Same as with tree traversal using DFS
2. Use a stack
3. Use three colors White, Orange, and Green, similarly

Depth-First Traversal

- **Idea:** If possible, go (in depth) forward else back track
- **Problem:** Since we have cycles, each node may be visited infinite times
- **Solution:** use a Boolean visited array



Problems Specific with Graphs

Minimum Spanning Tree

- Path Problems

1. Simple Paths

- 2. Shortest Path Problem**

- Single source shortest paths
- All-pair shortest paths

3. Find Cycles

4. Euler Path and Circuit Problem

5. Hamiltonian Path and Circuit Problem (or TSP)

Single Source Shortest Path Algorithms

Two famous algorithms for single source shortest paths

- **Dijkstra's Algorithm**
- Bellman-Ford's Algorithm
- Both use “edge relaxation” approach
- **Dijkstra's algorithm** (greedy approach) is faster
 - Can not be used with graphs having negative weight edges
- Bellman-Ford's algorithm (dynamic programming approach) can work with graphs having negative weight edges
 - But not with “negative cycles”

Greedy Algorithms

- Always makes local optimal choice at each stage with the intent of finding a global optimum.
- Might not always produce an optimal solution,
- Nonetheless yields local optimal solutions that approximate a global optimal solution in a reasonable amount of time.

Five components of Greedy Algorithms

1. **Candidate set:** From which a solution is created
2. **Selection function:** Chooses the best candidate to be added to the solution
3. **Feasibility function:** Determines if a candidate can be used to contribute to a solution
4. **Objective function:** Assigns a value to a solution, or a partial solution, and
5. **Solution function:** Indicates when we have discovered a complete solution

Thank you!

BFS: Algorithm

35

BFS(G, s)

```
1   for each vertex  $u \in V[G]$  do
2        $color[u] \leftarrow WHITE$ 
3        $dist[u] \leftarrow \infty$ 
4        $\pi[u] \leftarrow NULL$  //traversed set
5    $color[s] \leftarrow ORANGE$ 
6    $dist[s] \leftarrow 0$ 
7    $\pi[s] \leftarrow NIL$ 
8    $Q \leftarrow \Phi$ 
9   ENQUEUE ( $Q, s$ )
10  while  $Q \neq \Phi$  do
11       $u \leftarrow DEQUEUE(Q)$ 
12      for each  $v \in Adj[u]$  do
13          if  $color[v] = WHITE$ 
14               $color[v] \leftarrow ORANGE$ 
```

```
void Graph::DFSUtil(int v, bool visited[])
{
    visited[v] = true; // Mark the current node as visited and print it
    print v;
    list<int>::iterator i; // Recur for all the vertices adjacent to this vertex
    for(i = adj[v].begin(); i != adj[v].end(); ++i)
        if(!visited[*i])
            DFSUtil(*i, visited);
}

void Graph::DFS() // The function to do DFS traversal. It uses recursive DFSUtil()
{
    bool *visited = new bool[V]; // Mark all the vertices as not visited
    for (int i = 0; i < V; i++)
        visited[i] = false;
    for (int i = 0; i < V; i++) /* Call the recursive helper function to print
                                DFS traversal starting from all vertices one by one*/
        if (visited[i] == false)
            DFSUtil(i, visited);
}
```

Depth-First Traversal

DFS(G)

```

1 for each vertex  $u \in V[G]$  do
2    $color[u] \leftarrow WHITE$ 
3    $\pi[u] \leftarrow NIL$ 
4    $time \leftarrow 0$ 
5 for each vertex  $u \in V[G]$  do
6   if  $color[u] = WHITE$  then
7     DFS-VISIT( $u$ )
  
```

DFS-VISIT(u)

```

1   $color[u] \leftarrow ORANGE$ 
2   $time \leftarrow time + 1$ 
3   $d[u] \leftarrow time$       //discovery time
4  for each  $v \in Adj[u]$  do
5    if  $color[v] = WHITE$ 
6      then  $\pi[v] \leftarrow u$ 
7      DFS-VISIT( $v$ )
8   $color[u] \leftarrow GREEN$ 
9   $f[u] \leftarrow time \leftarrow time + 1$  //finish time
  
```


Thank you!