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# Algorithm for long integer multiplication

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## TABLE OF CONTENTS

### Contents

<b>1</b>	<b>Explanation of the algorithm</b>	<b>3</b>
<b>2</b>	<b>Pseudocode</b>	<b>4</b>
<b>3</b>	<b>Analysis of time complexity</b>	<b>5</b>

# 1 Explanation of the algorithm

This algorithm is really very simple. The main points are listed below

- Let us suppose we have to multiply  $a$  and  $b$  where both  $a$  and  $b$  are integers.
- Let the smaller one of them be  $a$  and so  $a$  is multiplier and  $b$  is multiplicand.
- First copy  $a$  and  $b$  into two other variables, say  $m$  and  $n$ .
- Divide  $m$  and  $n$  into its digits considering the units place, tens place, hundreds place and so on.  
for example :-  $324 = 300 + 20 + 4$
- Now, multiply both the digits of  $m$  and  $n$  taking into account the corresponding factors of 10 and then add up all the products to get the result.

## 2 Pseudocode

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Let  $a$  be the multiplier and  $b$  be the multiplier.

Let  $factora$  and  $factorb$  be the factor of multiplication for  $a$  and  $b$  respectively

Let  $rema$  be the last digit in  $a$  and  $remb$  be the last digit in  $b$

$factora=1, factorb=1, result=0$

$m \leftarrow a$

$n \leftarrow b$

**while** 1 **do**

$rema \leftarrow m \% 10$

**while** 1 **do**

$remb \leftarrow n \% 10$

$result \leftarrow result + (rema * remb * factora * factorb)$

$n \leftarrow n / 10$

**if**  $n == 0$  **then**

$n \leftarrow b$

            break

**end if**

**end while**

$m \leftarrow m / 10$

$factora \leftarrow factora * 10$

$factorb \leftarrow factorb * 10$

**if**  $m == 0$  **then**

$m \leftarrow b$

        break

**end if**

**end while**

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### 3 Analysis of time complexity

This algorithm has two loops which are nested into each other.

The worst case that will be faced by the algorithm is when both the loops run upto  $n$  times where  $n$  is the number of digits in the two numbers.

Some other operations are also executed in the algorithm like modulus operator,initialisation etc. but they all take constant time and so, the worst number of instructions executed by the algorithm is proportional to  $n^2 + k$  where  $k$  is a constant.

And so, the time complexity of the algorithm is  $O(n^2)$  where  $O$  stands for Big-O notation.