

CS251–ASSIGNMENT 3

DOCUMENT ON L^AT_EX

Algorithm for long integer multiplication

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Algorithm for Long Integer Multiplication

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1 Explanation of the algorithm

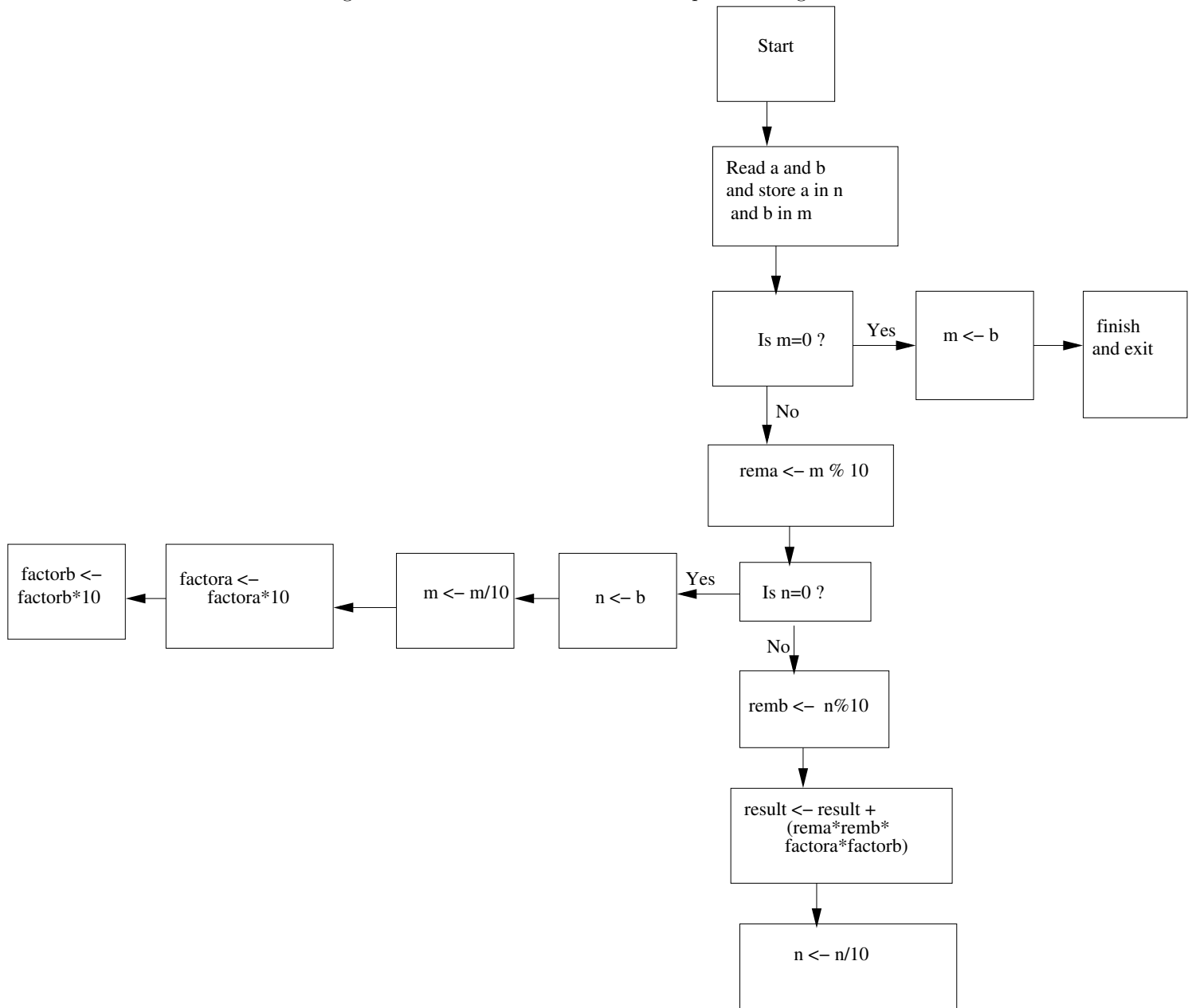
This algorithm is really very simple. The main points are listed below

- Let us suppose we have to multiply a and b where both a and b are integers.
- Let the smaller one of them be a and so a is multiplier and b is multiplicand.
- First copy a and b into two other variables, say m and n .
- Divide m and n into its digits considering the units place, tens place, hundreds place and so on.
for example :- $324 = 300 + 20 + 4$
- Now, multiply both the digits of m and n taking into account the corresponding factors of 10 and then add up all the products to get the result.

2 Pseudocode

Let a be the multiplier and b be the multiplier. Let $factora$ and $factorb$ be the factor of multiplication for a and b respectively. Let $rema$ be the last digit in a and $remb$ be the last digit in b . $factora=1, factorb=1, result=0$
 $m \leftarrow a$ $n \leftarrow b$ 1 $rema \leftarrow m \% 10$ 1 $remb \leftarrow n \% 10$ $result \leftarrow result + (rema * remb * factora * factorb)$ $n \leftarrow n / 10$
 $n == 0$ $n \leftarrow b$ break $m \leftarrow m / 10$ $factora \leftarrow factora * 10$ $factorb \leftarrow factorb * 10$ $m == 0$ $m \leftarrow b$ break

Figure 1: Flowchart to illustrate the steps of the algorithm



3 Flowchart

Figure 2: Example of multiplication of two integers

$$\begin{array}{r} 283 \\ \times 42 \\ \hline 566 \quad (= 283 \times 2) \\ 11320 \quad (= 283 \times 40) \\ \hline 11886 \\ \hline \end{array}$$

4 Example to illustrate the algorithm

5 Analysis of time complexity

This algorithm has two loops which are nested into each other.

The worst case that will be faced by the algorithm is when both the loops run upto n times where n is the number of digits in the two numbers.

Some other operations are also executed in the algorithm like modulus operator,initialisation etc. but they all take constant time and so, the worst number of instructions executed by the algorithm is proportional to $n^2 + k$ where k is a constant.

And so, the time complexity of the algorithm is $O(n^2)$ where O stands for Big-O notation.