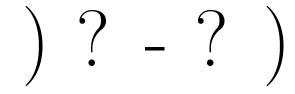


# Chulalongkorn University



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template from KACTL 2024-11-27

### template c IntPerm multinomial

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Template (1)	
template.cpp	
#pragma once	27 line

```
#include <bits/stdc++.h>
#define sz(x) (int)(x).size()
\#define all(x) (x).begin(), (x).end()
using namespace std;
using 11 = long long;
using db = long double;
using vi = vector<int>;
using vl = vector<11>;
using vd = vector<db>;
using pii = pair<int, int>;
using pll = pair<11, 11>;
using pdd = pair<db, db>;
const int INF = 0x3ffffffff;
// const int MOD=1000000007;
const int MOD = 998244353;
const 11 LINF = 0x1fffffffffffffffff;
const db DINF = numeric_limits<db>::infinity();
const db EPS = 1e-9;
const db PI = acos(db(-1));
```

L	<pre>cin.tie(nullptr)-&gt;sync_with_stdio(false);</pre>	
ı	}	
1	c.sh	2 lines
_	g++ -std=gnu++2a -Wall \$1 -o a.out	

## Mathematics (2)

./a.out

### 2.1 Goldbatch's Conjecture

- Even number can be written in sum of two primes (Up to 1e12)
- Range of  $N^{th}$  prime and  $N + 1^{th}$  prime will be less than or equal to 300 (Up to 1e12)

### 2.2 Divisibility

Number of divisors of N is given by  $\prod_{i=1}^{k} (a_i + 1)$  where  $N = \prod_{i=1}^{k} p_i^{a_i}$  and  $p_i$  are prime factors of N.

# Combinatorial (3)

### 3.1 Permutations

#### 3.1.1 Factorial

n	1 2 3	4	5 6	7	8	9	10	
n!	1 2 6	24 1	20 720	5040	40320	362880	3628800	-
n	11	12	13	14	15	16	17	
n!	1						3.6e14	
n	20	25	30	40	50 - 10	00 - 150	0 171	
$\overline{n!}$	2e18	2e25	3e32 8	647.3	e64 9e	157   6e26	$52 > \text{DBL}_1$	MAX

#### IntPerm.h

**Description:** Permutation -> integer conversion. (Not order preserving.) Integer -> permutation can use a lookup table. **Time:**  $\mathcal{O}(n)$ 

int permToInt(vi &v){
 int use = 0, i = 0, r = 0;
 for (int x : v) r = r \* ++i + \_\_builtin\_popcount(use & -(1 << x)),
 use |= 1 << x; // (note: minus, not ~!)
 return r;
}</pre>

### **3.1.2** Cycles

Let  $g_S(n)$  be the number of *n*-permutations whose cycle lengths all belong to the set S. Then

$$\sum_{n=0}^{\infty} g_S(n) \frac{x^n}{n!} = \exp\left(\sum_{n \in S} \frac{x^n}{n}\right)$$

### 3.1.3 Derangements

Permutations of a set such that none of the elements appear in their original position.

$$D(n) = (n-1)(D(n-1) + D(n-2)) = nD(n-1) + (-1)^n = \left\lfloor \frac{n!}{e} \right\rfloor$$

#### 3.1.4 Burnside's lemma

Given a group G of symmetries and a set X, the number of elements of X up to symmetry equals

$$\frac{1}{|G|} \sum_{g \in G} |X^g|,$$

where  $X^g$  are the elements fixed by g(g.x = x).

If f(n) counts "configurations" (of some sort) of length n, we can ignore rotational symmetry using  $G = \mathbb{Z}_n$  to get

$$g(n) = \frac{1}{n} \sum_{k=0}^{n-1} f(\gcd(n,k)) = \frac{1}{n} \sum_{k|n} f(k)\phi(n/k).$$

### 3.2 Partitions and subsets

### 3.2.1 Partition function

Number of ways of writing n as a sum of positive integers, disregarding the order of the summands.

$$p(0) = 1, \ p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k+1} p(n - k(3k - 1)/2)$$

### 3.2.2 Lucas' Theorem

Let n, m be non-negative integers and p a prime. Write  $n = n_k p^k + ... + n_1 p + n_0$  and  $m = m_k p^k + ... + m_1 p + m_0$ . Then  $\binom{n}{m} \equiv \prod_{i=0}^k \binom{n_i}{m_i} \pmod{p}$ .

### 3.2.3 Binomials

multinomial.h

return c;

Description: Computes  $\binom{k_1+\cdots+k_n}{k_1,k_2,\ldots,k_n} = \frac{(\sum k_i)!}{k_1!k_2!\ldots k_n!}$ .

11 multinomial (vi& v) {
 ll c = 1, m = v.empty() ? 1 : v[0];
 rep(i,1,sz(v)) rep(j,0,v[i])
 c = c \* ++m / (j+1);

### 3.3 General purpose numbers

### 3.3.1 Bernoulli numbers

EGF of Bernoulli numbers is  $B(t) = \frac{t}{e^t - 1}$  (FFT-able).  $B[0,...] = [1, -\frac{1}{2}, \frac{1}{6}, 0, -\frac{1}{30}, 0, \frac{1}{42},...]$ 

Sums of powers:

$$\sum_{i=1}^{n} n^{m} = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_{k} \cdot (n+1)^{m+1-k}$$

Euler-Maclaurin formula for infinite sums:

$$\sum_{i=m}^{\infty} f(i) = \int_{m}^{\infty} f(x)dx - \sum_{k=1}^{\infty} \frac{B_{k}}{k!} f^{(k-1)}(m)$$

$$\approx \int_{m}^{\infty} f(x)dx + \frac{f(m)}{2} - \frac{f'(m)}{12} + \frac{f'''(m)}{720} + O(f^{(5)}(m))$$

### 3.3.2 Stirling numbers of the first kind

Number of permutations on n items with k cycles.

$$c(n,k) = c(n-1,k-1) + (n-1)c(n-1,k), \ c(0,0) = 1$$
$$\sum_{k=0}^{n} c(n,k)x^{k} = x(x+1)\dots(x+n-1)$$

c(8, k) = 8, 0, 5040, 13068, 13132, 6769, 1960, 322, 28, 1 $c(n, 2) = 0, 0, 1, 3, 11, 50, 274, 1764, 13068, 109584, \dots$ 

#### 3.3.3 Eulerian numbers

Number of permutations  $\pi \in S_n$  in which exactly k elements are greater than the previous element. k j:s s.t.  $\pi(j) > \pi(j+1)$ , k+1 j:s s.t.  $\pi(j) \geq j$ , k j:s s.t.  $\pi(j) > j$ .

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n,n-1) = 1$$

$$E(n,k) = \sum_{j=0}^{k} (-1)^{j} \binom{n+1}{j} (k+1-j)^{n}$$

### 3.3.4 Stirling numbers of the second kind

Partitions of n distinct elements into exactly k groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k)$$

$$S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-j} \binom{k}{j} j^{n}$$

### 3.3.5 Bell numbers

Total number of partitions of n distinct elements. B(n) = $1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, \dots$  For p prime,

$$B(p^m + n) \equiv mB(n) + B(n+1) \pmod{p}$$

### 3.3.6 Labeled unrooted trees

# on n vertices:  $n^{n-2}$ # on k existing trees of size  $n_i$ :  $n_1 n_2 \cdots n_k n^{k-2}$ # with degrees  $d_i$ :  $(n-2)!/((d_1-1)!\cdots(d_n-1)!)$ 

#### 3.3.7 Catalan numbers

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \binom{2n}{n} - \binom{2n}{n+1} = \frac{(2n)!}{(n+1)!n!}$$

$$C_0 = 1, \ C_{n+1} = \frac{2(2n+1)}{n+2} C_n, \ C_{n+1} = \sum_{n=1}^{\infty} C_i C_{n-n}$$

 $C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, \dots$ 

- sub-diagonal monotone paths in an  $n \times n$  grid.
- strings with n pairs of parenthesis, correctly nested.
- binary trees with with n+1 leaves (0 or 2 children).
- ordered trees with n+1 vertices.
- ways a convex polygon with n+2 sides can be cut into triangles by connecting vertices with straight lines.
- permutations of [n] with no 3-term increasing subseq.

## Numerical (4)

### 4.1 Newton's Method

if 
$$F(Q) = 0$$
, then  $Q_{2n} \equiv Q_n - \frac{F(Q_n)}{F'(Q_n)} \pmod{x^{2n}}$   

$$Q = P^{-1} : Q_{2n} \equiv Q_n \cdot (2 - P \cdot Q_n^2) \pmod{x^{2n}}$$

$$Q = \ln P = \int \frac{P'}{P} dx$$

$$Q = e^p : Q_{2n} \equiv Q_n (1 + P - \ln Q_n) \pmod{x^{2n}}$$

$$Q = \sqrt{P} : Q_{2n} \equiv \frac{1}{2} (Q_n + P \cdot Q_n^{-1}) \pmod{x^{2n}}$$

$$Q = P^k = \alpha^k x^{kt} e^{k \ln T} : P = \alpha \cdot x^t \cdot T \cdot T(0) = 1$$

# Group (5)

### 5.1 Monoid

monoid/MonoidBase.hpp Description: Monoid Base class

e75b74, 6 lines

```
template < class T, T (*combine) (T, T), T (*identity) () >
struct MonoidBase
    using value_type = T;
    static constexpr T op(const T &x,const T &y) {return combine
    static constexpr T unit() {return identity();}
```

### Action

action/MonoidActionBase.hpp Description: Monoid Action Base class.

template < class MInfo, class MTaq, typename MInfo::value\_type (\*combine) (typename MInfo::value\_type,typename MTag:: value\_type)> struct MonoidActionBase{ using InfoMonoid = MInfo;

```
using Info = typename InfoMonoid::value_type;
    using Tag = typename TagMonoid::value type;
    static constexpr Info op(const Info &a,const Tag &b) {
        return combine (a, b);
};
action/DefaultAction.hpp
Description: Default Action class.
                                                      e45000. 10 lines
template < class Monoid>
struct DefaultAction{
    using InfoMonoid = Monoid;
    using TagMonoid = Monoid;
    using Info = typename Monoid::value_type;
    using Tag = typename Monoid::value_type;
    static constexpr Info op(const Info &a,const Tag &b) {
        return Monoid::op(a,b);
};
```

## Data Structures (6)

using TagMonoid = MTag;

```
OrderedSet.hpp
Description: Ordered Set
"../template/Header.hpp", <bits/extc++.h>
```

1a7f5f, 14 lines

```
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
    tree_order_statistics_node_update>;
// can be change to less equal
void usage() {
 ordered set<int> st, st 2;
 st.insert(2);
 st.insert(1);
 cout << st.order_of_key(2);</pre>
 cout << *st.find_by_order(1);</pre>
 st.join(st_2); // merge
```

### FenwickTree.hpp

Description: Fenwick / Binary Indexed Tree

43767a, 41 lines

```
template<class T>
struct Fenwick{
    int n,logn;
    vector<T> t:
    Fenwick(){}
    Fenwick(int _n) {init(vector<T>(_n,T{}));}
    template<class U>
    Fenwick(const vector < U > &a) {init(a);}
    template<class U>
    void init(const vector<U> &a) {
        n=(int)a.size();
        logn=31- builtin clz(n);
        t.assign(n+1,T\{\});
        for (int i=1; i<=n; i++) {</pre>
             t[i]=t[i]+a[i-1];
             int j=i+(i\&-i);
             if (j<=n)t[j]=t[j]+t[i];</pre>
    void update(int x,const T &v) {
        for (int i=x+1; i<=n; i+=i&-i) t[i] =t[i]+v;</pre>
```

```
Chula[) ? - ? )]
```

```
void update(int 1, int r, const T &v) {
    update (1, v), update (r+1, -v);
T query(int x) {
    T res{};
    for(int i=x+1;i>0;i-=i&-i)res=res+t[i];
    return res;
T query(int 1,int r){
    return query (r) -query (1-1);
int find(const T &k){
    int x=0;
    T cur{};
    for (int i=1<<logn; i>0; i>>=1)
        if (x+i<=n&&cur+t[x+i]<=k)x+=i,cur=cur+t[x];</pre>
```

### SmallSegmentTree.hpp

};

Description: Zero-indexed max-tree. Bounds are inclusive to the left and exclusive to the right. Can be changed by modifying T, f and unit.

Time:  $\mathcal{O}(\log N)$ 0f4bdb, 19 lines

```
struct Tree {
  typedef int T;
  static constexpr T unit = INT_MIN;
  T f (T a, T b) { return max(a, b); } // (any associative fn)
  vector<T> s; int n;
  Tree(int n = 0, T def = unit) : s(2*n, def), n(n) {}
  void update(int pos, T val) {
    for (s[pos += n] = val; pos /= 2;)
     s[pos] = f(s[pos * 2], s[pos * 2 + 1]);
  T query (int b, int e) { // query [b, e]
   T ra = unit, rb = unit;
    for (b += n, e += n; b < e; b /= 2, e /= 2) {
     if (b % 2) ra = f(ra, s[b++]);
     if (e % 2) rb = f(s[--e], rb);
    return f(ra, rb);
};
```

## Segment Tree.hpp

Description: Segment Tree

```
c51dec. 85 lines
template < class Monoid>
struct SegmentTree{
    using T = typename Monoid::value_type;
    int n;
    vector<T> t;
    SegmentTree(){}
    SegmentTree(int n, function<T(int)> create) {init(n, create);}
    SegmentTree(int n,T v=Monoid::unit()) {init(n,[&](int){
         return v; });}
    template < class U>
    SegmentTree(const vector<U> &a) {init((int)a.size(), [&](int
         i) {return T(a[i]);});}
    void init(int _n,function<T(int)> create){
        t.assign(4<<(31-__builtin_clz(n)), Monoid::unit());
        function<void(int,int,int)> build=[&](int l,int r,int i
            if (l==r) return void(t[i]=create(l));
            int m = (1+r)/2;
            build(l, m, i*2);
            build(m+1, r, i*2+1);
```

### SmallSegmentTree SegmentTreeBeats

```
pull(i);
        };
        build(0,n-1,1);
    void pull(int i){
        t[i]=Monoid::op(t[i*2],t[i*2+1]);
    void modify(int 1,int r,int i,int x,const T &v) {
        if (x<1 | | r<x) return;
        if(l==r)return void(t[i]=v);
        int m = (1+r)/2;
        modify(1, m, i*2, x, v);
        modify (m+1, r, i*2+1, x, v);
        pull(i);
    void modify(int x,const T &v) {
        modify (0, n-1, 1, x, v);
    template<class U>
    void update(int 1,int r,int i,int x,const U &v) {
        if (x<1 | | r<x) return;
        if(l==r)return void(t[i]=Monoid::op(t[i],v));
        int m = (1+r)/2;
        update (1, m, i*2, x, v);
        update (m+1,r,i*2+1,x,v);
        pull(i);
    template < class U>
    void update(int x,const U &v) {
        update (0, n-1, 1, x, v);
    T query(int 1, int r, int i, int x, int y) {
        if (y<1||r<x) return Monoid::unit();</pre>
        if (x<=1&&r<=y) return t[i];
        int m = (1+r)/2;
        return Monoid::op(query(1, m, i*2, x, y), query(m+1, r, i*2+1,
              x, y));
    T query(int x,int y) {
        return query (0, n-1, 1, x, y);
    template<class F>
    int findfirst(int 1, int r, int i, int x, int y, const F &f) {
        if (y<1||r<x||!f(t[i])) return n;</pre>
        if(l==r)return 1;
        int m = (1+r)/2;
        int res=findfirst(l,m,i*2,x,v,f);
        if (res==n) res=findfirst (m+1, r, i * 2+1, x, y, f);
         return res;
    template < class F >
    int findfirst(int x,int y,const F &f) {
         return findfirst(0,n-1,1,x,y,f);
    template < class F>
    int findlast(int 1,int r,int i,int x,int y,const F &f) {
        if (y<1 | | r<x | | ! f (t[i])) return -1;</pre>
        if(l==r)return 1;
        int m = (1+r)/2;
        int res=findlast (m+1, r, i*2+1, x, y, f);
        if (res==-1) res=findlast(1, m, i*2, x, y, f);
        return res:
    template < class F >
    int findlast(int x,int y,const F &f) {
         return findlast(0,n-1,1,x,y,f);
};
```

```
Segment TreeBeats.hpp
Description: Segment Tree Beats
"../template/Header.hpp"
                                                        efa2ef, 134 lines
const int N = 2e5 + 5;
const int K = 1 << 19;</pre>
int n, q;
ll a[N];
struct SegTree{
    struct Node {
        11 sum, add;
        11 mn, mn2, fn;
        11 mx, mx2, fx;
        Node(){
             sum=add=fn=fx=0, mn=mn2=LINF, mx=mx2=-LINF;
        Node(ll v){
             sum=mn=mx=v,add=0,mn2=LINF,mx2=-LINF,fn=fx=1;
        friend Node operator+(const Node &1,const Node &r) {
             Node res;
             res.sum=1.sum+r.sum;
             res.add=0;
             if(1.mx>r.mx){
                 res.mx=1.mx,res.fx=1.fx;
                 res.mx2=max(1.mx2,r.mx);
             }else if(r.mx>1.mx) {
                 res.mx=r.mx,res.fx=r.fx;
                 res.mx2=max(r.mx2,1.mx);
             }else{
                 res.mx=1.mx,res.fx=1.fx+r.fx;
                 res.mx2=max(1.mx2,r.mx2);
             if(1.mn<r.mn){
                 res.mn=1.mn, res.fn=1.fn;
                 res.mn2=min(1.mn2,r.mn);
             }else if(r.mn<1.mn){</pre>
                 res.mn=r.mn,res.fn=r.fn;
                 res.mn2=min(r.mn2,1.mn);
                 res.mn=1.mn, res.fn=1.fn+r.fn;
                 res.mn2=min(1.mn2,r.mn2);
             return res;
        void apply(int 1,int r,ll v){
             sum+=(r-1+1)*v;
             mx+=v, mx2+=v;
             mn+=v, mn2+=v;
             add+=v;
        void chmin(ll v) {
             if (v>=mx) return;
             sum+=(v-mx)*fx;
             if (mn==mx) mn=v;
             if (mn2==mx) mn2=v;
        void chmax(ll v){
             if (v<=mn) return;</pre>
             sum+=(v-mn)*fn;
             if (mx==mn) mx=v;
             if (mx2==mn) mx2=v;
             mn=v:
    }t[K];
    void pull(int i) {
        t[i]=t[i*2]+t[i*2+1];
```

### LazySegmentTree DynamicSegmentTree

```
void push(int 1,int r,int i) {
        int m = (1+r)/2;
        t[i*2].apply(l,m,t[i].add);
        t[i*2+1].apply(m+1,r,t[i].add);
        t[i*2].chmin(t[i].mx);
        t[i*2+1].chmin(t[i].mx);
        t[i*2].chmax(t[i].mn);
        t[i*2+1].chmax(t[i].mn);
        t[i].add=0;
    void build(int 1,int r,int i) {
        if(l==r) return void(t[i]=Node(a[l]));
        int m = (1+r)/2;
        build(1, m, i * 2);
        build(m+1, r, i*2+1);
        pull(i);
    void build() {
        build(1, n, 1);
    void range_add(int l,int r,int i,int x,int y,ll v){
        if (v<1 | | r<x) return;</pre>
        if (x<=l&&r<=y) return t[i].apply(1,r,v);</pre>
        int m = (1+r)/2;
        push(l,r,i);
        range_add(1, m, i*2, x, y, v);
        range_add(m+1, r, i*2+1, x, y, v);
        pull(i);
    void range_add(int x,int y,ll v){
         range_add(1, n, 1, x, y, v);
    void range_chmin(int 1,int r,int i,int x,int y,11 v) {
        if (y<1 | | r<x | | t [i] .mx<=v) return;
        if (x<=1&&r<=y&&t[i].mx2<v) return t[i].chmin(v);</pre>
        int m = (1+r)/2;
        push(l,r,i);
        range chmin(1, m, i*2, x, v, v);
         range_chmin(m+1, r, i*2+1, x, y, v);
        pull(i);
    void range_chmin(int x,int y,ll v){
        range_chmin(1, n, 1, x, y, v);
    void range_chmax(int 1,int r,int i,int x,int y,11 v) {
        if (y<1 | | r<x | |t[i].mn>=v) return;
        if (x<=1&&r<=y&&t[i].mn2>v) return t[i].chmax(v);
        int m = (1+r)/2;
        push(l,r,i);
        range_chmax(1, m, i*2, x, y, v);
        range_chmax(m+1, r, i*2+1, x, y, v);
        pull(i);
    void range_chmax(int x,int y,ll v) {
        range_chmax(1,n,1,x,y,v);
    11 query(int 1,int r,int i,int x,int y) {
        if (y<1 | |r<x) return 0;
        if (x<=l&&r<=y) return t[i].sum;</pre>
        int m = (1+r)/2;
        push(l,r,i);
        return query(1, m, i * 2, x, y) + query(m+1, r, i * 2+1, x, y);
    11 query(int x,int y) {
         return query(1,n,1,x,y);
};
```

```
LazySegmentTree.hpp
Description: Segment Tree with Lazy Propagation
                                                      91ab0c, 103 lines
template < class MonoidAction>
struct LazySegmentTree{
    using InfoMonoid = typename MonoidAction::InfoMonoid;
    using TagMonoid = typename MonoidAction::TagMonoid;
    using Info = typename MonoidAction::Info;
    using Tag = typename MonoidAction::Tag;
    int n:
    vector<Info> t;
    vector<Tag> lz;
    LazySegmentTree(){}
    LazySegmentTree(int n, function<Info(int)> create) {init(n,
    LazySegmentTree(int n, Info v=InfoMonoid::unit()) {init(n
         , [&] (int) {return v;});}
    template<class T>
    LazySegmentTree(const vector<T> &a) {init((int)a.size(), [&](
         int i) {return Info(a[i]);});}
    void init(int _n,function<Info(int)> create){
        int m=4 << (31- builtin clz(n));
        t.assign(m,InfoMonoid::unit());
        lz.assign(m, TagMonoid::unit());
        function<void(int,int,int)> build=[&](int 1,int r,int i
            if(l==r)return void(t[i]=create(l));
            int m = (1+r)/2;
            build(1, m, i * 2);
            build(m+1, r, i*2+1);
            pull(i);
        build(0, n-1, 1);
    void pull(int i){
        t[i]=InfoMonoid::op(t[i*2],t[i*2+1]);
    void apply(int i,const Tag &v) {
        t[i]=MonoidAction::op(t[i],v);
        lz[i]=TagMonoid::op(lz[i],v);
    void push(int i) {
        apply(i*2, lz[i]);
        apply(i*2+1,lz[i]);
        lz[i]=TagMonoid::unit();
    void modify(int l,int r,int i,int x,const Info &v){
        if (x<1 | | r<x) return;</pre>
        if(l==r)return void(t[i]=v);
        int m = (1+r)/2;
        push(i);
        modify(1, m, i*2, x, v);
        modify (m+1,r,i*2+1,x,v);
        pull(i);
    void modify(int x,const Info &v){
        modify (0, n-1, 1, x, v);
    void update(int 1,int r,int i,int x,int y,const Tag &v) {
        if (y<1 | | r<x) return;
        if (x<=l&&r<=y) return apply(i,v);</pre>
        int m = (1+r)/2;
        push(i);
        update(1, m, i*2, x, y, v);
        update (m+1,r,i*2+1,x,y,v);
        pull(i);
    void update(int x,int v,const Tag &v) {
```

```
update (0, n-1, 1, x, y, v);
    Info query(int 1,int r,int i,int x,int y) {
        if (y<1||r<x) return InfoMonoid::unit();</pre>
        if (x<=1&&r<=y) return t[i];</pre>
        int m = (1+r)/2;
        push(i);
        return InfoMonoid::op(query(1, m, i * 2, x, y), query(m+1, r, i
              *2+1, x, y));
    Info query(int x,int y) {
        return query (0, n-1, 1, x, y);
    template < class F >
    int findfirst(int 1,int r,int i,int x,int y,const F &f) {
        if(y<1||r<x||!f(t[i]))return n;
        if(l==r)return 1;
        int m = (1+r)/2;
        push(i);
        int res=findfirst(1, m, i*2, x, y, f);
        if (res==n) res=findfirst (m+1, r, i*2+1, x, y, f);
        return res;
    template < class F >
    int findfirst(int x,int y,const F &f) {
        return findfirst(0, n-1, 1, x, y, f);
    template<class F>
    int findlast(int 1,int r,int i,int x,int y,const F &f) {
        if (y<1 | | r<x | | ! f (t [i])) return -1;</pre>
        if(l==r)return 1;
        int m = (1+r)/2;
        int res=findlast (m+1, r, i*2+1, x, y, f);
        if (res==-1) res=findlast (1, m, i*2, x, y, f);
        return res;
    template<class F>
    int findlast(int x,int v,const F &f) {
        return findlast(0,n-1,1,x,y,f);
};
DynamicSegmentTree.hpp
Description: Dynamic Segment Tree
                                                        e84eeb, 106 lines
template < class MonoidAction >
struct DynamicSegmentTree{
    using InfoMonoid = typename MonoidAction::InfoMonoid;
    using TagMonoid = typename MonoidAction::TagMonoid;
    using Info = typename MonoidAction::Info;
    using Tag = typename MonoidAction::Tag;
    struct Node:
    using Ptr = Node*;
    struct Node{
        Info val;
        Tag lz;
        Ptr 1,r;
        Node (Info v): val(v), lz(TagMonoid::unit()), l(nullptr), r(
             nullptr) {}
        Node(Info v, Tag t):val(v),lz(t),l(nullptr),r(nullptr){}
    11 lb, ub;
    Ptr rt;
    function<Info(11,11)> create;
    DynamicSegmentTree() {init(0,0);}
    DynamicSegmentTree(ll n) {init(0, n-1);}
```

DynamicSegmentTree(11 lb, 11 ub) {init(lb, ub);}

5

### DSU BinaryTrie LiChaoTree

```
DynamicSegmentTree(ll lb, ll ub, function < Info(ll, ll) > create
     ) {init(lb,ub,create);}
void init(ll _lb,ll _ub,function<Info(ll,ll)> _create=[](ll
      1,11 r) {return InfoMonoid::unit();}) {
    lb=_lb, ub=_ub;
    create=_create;
    rt=new Node(create(lb,ub));
Info val(Ptr t){
    return t?t->val:InfoMonoid::unit();
void pull(Ptr &t){
    t->val=InfoMonoid::op(val(t->1),val(t->r));
void apply(Ptr &t,const Tag &v,ll 1,ll r){
    if(!t)t=new Node(create(1,r));
    t->val=MonoidAction::op(t->val,v);
    t->1z=TagMonoid::op(t->1z,v);
void push(Ptr &t, ll l, ll m, ll r) {
    apply (t->1, t->1z, 1, m);
    apply (t->r,t->lz,m+1,r);
    t->1z=TagMonoid::unit();
void modify(ll 1,ll r,Ptr &t,ll x,const Info &v) {
    if (x<1 | r<x) return;
    if (l==r) return void(t->val=v);
   11 m=1+(r-1)/2;
    push(t,1,m,r);
    modify(1, m, t->1, x, v);
    modify (m+1, r, t->r, x, v);
    pull(t);
void modify(ll x,const Info &v){
    modify(lb,ub,rt,x,v);
void update(ll 1,ll r,Ptr &t,ll x,ll y,const Tag &v) {
    if (y<1 | | r<x) return;</pre>
    if (x<=1&&r<=y) return apply(t,v,l,r);</pre>
    11 m=1+(r-1)/2;
    push(t,l,m,r);
    update(1, m, t \rightarrow 1, x, v, v);
    update (m+1,r,t->r,x,y,v);
    pull(t);
void update(ll x,ll y,const Tag &v) {
    update(lb,ub,rt,x,y,v);
Info query(11 1,11 r,Ptr &t,11 x,11 y){
    if(v<1||r<x)return InfoMonoid::unit();</pre>
    if (x<=1&&r<=y) return t->val;
   11 m=1+(r-1)/2;
    push(t,l,m,r);
    return InfoMonoid::op(query(1, m, t->1, x, y), query(m+1, r, t
         ->r,x,y));
Info query(ll x,ll y){
    return query(lb,ub,rt,x,y);
template < class F>
11 findfirst(ll 1,11 r,Ptr t,11 x,11 y,const F &f) {
    if (y<1 | | r<x | | ! f (t->val) ) return -1;
    if(l==r)return 1;
   11 m=1+(r-1)/2;
    push(t,l,m,r):
    11 res=findfirst(1,m,t->1,x,y,f);
    if (res==-1) res=findfirst (m+1, r, t->r, x, y, f);
    return res;
```

```
template<class F>
    11 findfirst(ll x, ll y, const F &f) {
        return findfirst(lb,ub,rt,x,y,f);
    template<class F>
    11 findlast(ll 1,ll r,Ptr t,ll x,ll y,const F &f) {
        if (v<1||r<x||!t||!f(t->val))return -1;
        if(l==r)return 1;
        11 m=1+(r-1)/2;
        push(t,l,m,r);
        11 res=findlast(m+1,r,t->r,x,y,f);
        if (res==-1) res=findlast(1, m, t->1, x, y, f);
        return res:
    template<class F>
    11 findlast(ll x,ll y,const F &f) {
        return findlast(lb,ub,rt,x,y,f);
};
DSU.hpp
Description: Disjoint Set Union.
                                                      0b3cb8, 26 lines
struct DSU{
    vector<int> p,sz;
    DSU() {}
   DSU(int n) {init(n);}
   void init(int n) {
        p.resize(n);
        iota(p.begin(),p.end(),0);
        sz.assign(n,1);
    int find(int u){
        return p[u] == u?u:p[u] = find(p[u]);
    bool same(int u,int v){
        return find(u) == find(v);
    bool merge(int u,int v) {
        u=find(u), v=find(v);
        if(u==v)return false;
        sz[u]+=sz[v];
        p[v]=u;
        return true;
    int size(int u) {
        return sz[find(u)];
};
Binary Trie. hpp
Description: Binary Trie
                                                      ae5b7a, 66 lines
template<int BIT, class T = uint32_t, class S = int>
struct BinaryTrie{
    struct Node{
        array<int,2> ch;
        S cnt:
        Node():ch\{-1,-1\}, cnt(0){}
    vector<Node> t;
    BinaryTrie():t{Node()}{}
    int new_node(){
        t.emplace back(Node());
        return t.size()-1;
   S size(){
        return t[0].cnt;
```

```
bool empty() {
        return size()==0;
    S get cnt(int i) {
        return i!=-1?t[i].cnt:S(0);
    void insert(T x,S k=1){
        int u=0;
        t[u].cnt+=k;
        for(int i=BIT-1;i>=0;i--){
            int v=x>>i&1;
            if(t[u].ch[v]==-1)t[u].ch[v]=new_node();
            u=t[u].ch[v];
            t[u].cnt+=k;
    void erase(T x,S k=1){
        int u=0;
        assert(t[u].cnt>=k);
        t[u].cnt-=k;
        for(int i=BIT-1; i>=0; i--) {
            int v=x>>i&1;
            u=t[u].ch[v];
            assert (u!=-1&&t[u].cnt>=k);
            t[u].cnt-=k;
    T kth(S k, T x=0) {
        assert(k<size());
        int u=0;
        T res=0;
        for (int i=BIT-1; i>=0; i--) {
            int v=x>>i&1;
            if (k < get_cnt (t[u].ch[v])) {
                u=t[u].ch[v];
            }else{
                res|=T(1)<<i;
                k-=get_cnt(t[u].ch[v]);
                u=t[u].ch[v^1];
        return res;
    T min(T x){
        return kth(0,x);
    T max(T x){
        return kth(size()-1,x);
};
LiChaoTree.hpp
Description: Li-Chao Tree (minimize)
                                                      4ab713, 52 lines
template<class T>
struct LiChaoTree{
    static const T INF=numeric_limits<T>::max()/2;
    struct Line{
        T m,c;
        Line(T _m,T _c):m(_m),c(_c){}
        inline T eval(T x)const{return m*x+c;}
    vector<T> xs;
    vector<Line> t;
    LiChaoTree(){}
    LiChaoTree (const vector<T> &x):xs(x) {init(x);}
    LiChaoTree(int n):xs(n){
        vector<T> x(n);
        iota(x.begin(), x.end(), 0);
```

### DynamicLiChaoTree SplayTreeBase

```
init(x);
    void init(const vector<T> &x){
        xs=x:
        sort(xs.begin(),xs.end());
        xs.erase(unique(xs.begin(), xs.end()), xs.end());
        t.assign(4<<(31-__builtin_clz(xs.size())),Line(0,INF));</pre>
    void insert(int 1,int r,int i,Line v) {
        int m = (1+r)/2;
        if(v.eval(xs[m]) <t[i].eval(xs[m])) swap(t[i],v);</pre>
        if(v.eval(xs[1]) <t[i].eval(xs[1]))insert(1, m, i*2, v);</pre>
        if(v.eval(xs[r]) < t[i].eval(xs[r]))insert(m+1,r,i*2+1,v)
    inline void insert(T m,T c){
        insert (0, (int) xs.size()-1,1, Line(m,c));
    void insert_range(int l, int r, int i, T x, T y, Line v) {
        if(y<xs[1]||xs[r]<x) return;
        if (x<=xs[1]&&xs[r]<=y) return insert(1,r,i,v);</pre>
        int m = (1+r)/2;
        insert_range(1, m, i * 2, x, y, v);
        insert_range(m+1, r, i*2+1, x, y, v);
    inline void insert_range(T m, T c, T x, T y) {
        insert_range(0,(int)xs.size()-1,1,x,y,Line(m,c));
    T query(int 1,int r,int i,T x) {
        if (l==r) return t[i].eval(x);
        int m = (1+r)/2;
        if (x \le xs[m]) return min (t[i].eval(x), query(1, m, i*2, x));
        return min(t[i].eval(x), query(m+1, r, i*2+1, x));
    inline T query(T x){
        return query(0,(int)xs.size()-1,1,x);
};
```

## DynamicLiChaoTree.hpp

Description: Dynamic Li-Chao Tree (minimize).

void splay(Ptr t) {

push(t);

if(!t)return;

while(!is\_root(t)){

Ptr x=t->p;

if(is root(x)){

```
b8af36, 50 lines
template<class T>
struct DvnamicLiChaoTree{
    static const T INF=numeric_limits<T>::max()/2;
    struct Line{
       T m,c;
        Line(T _m,T _c):m(_m),c(_c){}
        inline T eval(T x)const{return m*x+c;}
   struct Node:
   using Ptr = Node*;
    struct Node {
        Line v:
       Ptr l,r;
        Node():v(0,INF),l(nullptr),r(nullptr){}
        Node(Line _v):v(_v),l(nullptr),r(nullptr){}
   11 lb, ub;
   Ptr root:
   DynamicLiChaoTree(ll _lb,ll _ub):lb(_lb),ub(_ub),root(
         nullptr) {}
    void insert(T 1,T r,Ptr &t,Line v){
        if(!t)return void(t=new Node(v));
        T m=1+(r-1)/2;
        if(v.eval(m) < t->v.eval(m)) swap(t->v,v);
        if(v.eval(1) <t->v.eval(1)) insert(1, m, t->1, v);
        if (v.eval(r) <t->v.eval(r)) insert(m+1, r, t->r, v);
```

```
inline void insert(T m,T c){
        insert(lb,ub,root,Line(m,c));
    void insert_range(T 1, T r, Ptr &t, T x, T y, Line v) {
        if (y<1 | |r<x) return;
        if(!t)t=new Node();
        if (x<=1&&r<=y) return insert(1,r,t,v);</pre>
        T m=1+(r-1)/2;
        insert_range(l, m, t->l, x, y, v);
        insert_range (m+1, r, t->r, x, y, v);
    inline void insert_range(T m,T c,T x,T y){
        insert_range(lb, ub, root, x, y, Line(m, c));
    T query(T 1,T r,Ptr t,T x){
        if(!t)return INF;
        T m=1+(r-1)/2;
        if (x<=m) return min (t->v.eval(x), query(1, m, t->1, x));
        return min(t->v.eval(x), query(m+1, r, t->r, x));
    inline T query(T x){
        return query(lb,ub,root,x);
};
SplayTreeBase.hpp
Description: Splay Tree. splay (u) will make node u be the root of the tree
in amortized O(log n) time.
                                                       cc90a9, 113 lines
template < class Node >
struct SplayTreeBase{
    using Ptr = Node*;
    bool is root(Ptr t){
        return ! (t->p) | | (t->p->l!=t&&t->p->r!=t);
    } // The parent of the root stores the path parant in link
    int size(Ptr t){
        return t?t->size:0;
    virtual void push(Ptr t){};
    virtual void pull(Ptr t){};
    int pos(Ptr t){
        if(t->p){
             if (t->p->1==t) return -1;
             if(t->p->r==t)return 1;
        return 0;
    void rotate(Ptr t) {
        Ptr x=t->p, y=x->p;
        if (pos(t) ==-1) {
             if ((x->1=t->r))t->r->p=x;
             t->r=x, x->p=t;
        }else{
             if ((x->r=t->1))t->1->p=x;
             t->1=x, x->p=t;
        pull(x),pull(t);
        if((t->p=y)){
             if (y->l==x)y->l=t;
             if (y->r==x) y->r=t;
```

```
push(x), push(t);
            rotate(t);
        }else{
            Ptr y=x->p;
            push(y),push(x),push(t);
            if(pos(x) == pos(t)) rotate(x), rotate(t);
            else rotate(t), rotate(t);
Ptr get_first(Ptr t){
    while (t->1) push (t), t=t->1;
    splay(t);
    return t;
Ptr get_last(Ptr t){
    while (t->r) push (t), t=t->r;
    splay(t);
    return t;
Ptr merge(Ptr 1,Ptr r) {
    splay(1), splay(r);
    if(!1)return r;
    if(!r)return 1;
    l=get_last(l);
    1->r=r;
    r->p=1;
    pull(1);
    return 1;
pair<Ptr,Ptr> split(Ptr t,int k){
    if(!t)return {nullptr,nullptr};
    if (k==0) return {nullptr,t};
    if(k==size(t))return {t,nullptr};
    push(t);
    if(k<=size(t->1)){
        auto x=split(t->1,k);
        t->1=x.second;
        t->p=nullptr;
        if(x.second) x.second->p=t;
        pull(t);
        return {x.first,t};
        auto x=split(t->r,k-size(t->1)-1);
        t->r=x.first;
        t->p=nullptr;
        if(x.first)x.first->p=t;
        pull(t);
        return {t,x.second};
void insert(Ptr &t,int k,Ptr v) {
    splav(t);
    auto x=split(t,k);
    t=merge(merge(x.first, v), x.second);
void erase(Ptr &t,int k) {
    splav(t);
    auto x=split(t,k);
    auto y=split(x.second,1);
    // delete y.first;
    t=merge(x.first,y.second);
template<class T>
Ptr build(const vector<T> &v) {
    if(v.empty())return nullptr;
    function<Ptr(int,int)> build=[&](int 1,int r){
        if(l==r)return new Node(v[1]);
        int m = (1+r)/2;
```

```
return build(0, v.size()-1);
};
LazyReversibleBBST.hpp
Description: Lazy Reversible BBST Base.
                                                     904708, 81 lines
template < class Tree, class Node, class MonoidAction >
struct LazyReversibleBBST:Tree{
    using Tree::merge;
    using Tree::split;
    using typename Tree::Ptr;
    using InfoMonoid = typename MonoidAction::InfoMonoid;
    using TagMonoid = typename MonoidAction::TagMonoid;
    using Info = typename MonoidAction::Info;
    using Tag = typename MonoidAction::Tag;
    LazyReversibleBBST() = default;
    Info sum(Ptr t){
        return t?t->sum:InfoMonoid::unit();
    void pull(Ptr t) {
        if(!t)return;
        push(t);
        t->size=1;
        t->sum=t->val;
        t->revsum=t->val:
        if(t->1){
            t->size+=t->l->size;
            t->sum=InfoMonoid::op(t->l->sum,t->sum);
            t->revsum=InfoMonoid::op(t->revsum,t->l->revsum);
       if(t->r){
            t->size+=t->r->size;
            t->sum=InfoMonoid::op(t->sum,t->r->sum);
            t->revsum=InfoMonoid::op(t->r->revsum,t->revsum);
    void push(Ptr t){
        if(!t)return;
        if(t->rev){
            toggle(t->1);
            toggle(t->r);
            t->rev=false;
        if(t->lz!=TagMonoid::unit()){
            propagate (t->1,t->1z);
            propagate(t->r,t->lz);
            t->1z=TagMonoid::unit();
    void toggle (Ptr t) {
        if(!t)return;
        swap (t->1,t->r);
        swap (t->sum, t->revsum);
        t->rev^=true;
    void propagate(Ptr t,const Tag &v){
        if(!t)return;
        t->val=MonoidAction::op(t->val,v);
        t->sum=MonoidAction::op(t->sum,v);
        t->revsum=MonoidAction::op(t->revsum, v);
        t->1z=TagMonoid::op(t->1z,v);
    void apply(Ptr &t,int 1,int r,const Tag &v){
        if (1>r) return;
```

return merge(build(1,m),build(m+1,r));

```
auto x=split(t,1);
        auto y=split(x.second, r-1+1);
        propagate(y.first,v);
        t=merge(x.first, merge(y.first, y.second));
    Info query(Ptr &t,int 1,int r){
        if(l>r)return InfoMonoid::unit();
        auto x=split(t,1);
        auto y=split(x.second, r-l+1);
        Info res=sum(y.first);
        t=merge(x.first, merge(y.first, y.second));
        return res;
    void reverse(Ptr &t,int 1,int r) {
        if (1>r) return;
        auto x=split(t,1);
        auto y=split(x.second, r-1+1);
        toggle(y.first);
        t=merge(x.first,merge(y.first,y.second));
};
LazyReversibleSplayTree.hpp
Description: Lazy Reversible Splay Tree.
"SplayTreeBase.hpp", "LazyReversibleBBST.hpp"
                                                     b8455b, 23 lines
template < class MonoidAction>
struct LazyReversibleSplayTreeNode{
    using Ptr = LazyReversibleSplayTreeNode*;
    using InfoMonoid = typename MonoidAction::InfoMonoid;
    using TagMonoid = typename MonoidAction::TagMonoid;
    using Info = typename MonoidAction::Info;
    using Tag = typename MonoidAction::Tag;
    using value_type = Info;
   Ptr 1, r, p;
    Info val, sum, revsum;
    Tag lz;
    int size;
    LazyReversibleSplayTreeNode(const Info &_val=InfoMonoid::
        unit(),const Tag &_lz=TagMonoid::unit())
        :1(),r(),p(),val(_val),sum(_val),revsum(_val),lz(_lz),
             size(1), rev(false) {}
};
template < class MonoidAction >
struct LazyReversibleSplayTree
    : LazyReversibleBBST<SplayTreeBase<
         LazyReversibleSplayTreeNode<MonoidAction>>,
     LazyReversibleSplayTreeNode<MonoidAction>, MonoidAction>{
    using Node = LazyReversibleSplayTreeNode<MonoidAction>;
LinkCutTreeBase.hpp
Description: Link Cut Tree Base.
Usage: evert(u): make u be the root of the tree.
link(u,v): attach u to v.
cut(u,v): remove edge between u and v.
get_root(u): get the root of the tree containing u.
lca(u,v): get the lowest common ancestor of u and v.
fold(u,v): get the value of the path from u to v. _{\rm b432c3,\,59\;lines}
template<class Splay>
struct LinkCutTreeBase:Splay{
    using Node = typename Splay::Node;
    using Ptr = Node*;
    using T = typename Node::value_type;
   Ptr expose(Ptr t){
        Ptr pc=nullptr; // preferred child
        for (Ptr cur=t; cur; cur=cur->p) {
```

```
this->splay(cur);
            cur->r=pc;
            this->pull(cur);
            pc=cur;
        this->splay(t);
        return pc;
    void evert(Ptr t) { // make t be the root of the tree
        expose(t);
        this->toggle(t);
        this->push(t);
    void link(Ptr u,Ptr v) { // attach u to v
        evert(u);
        expose (v);
        u->p=v;
    void cut(Ptr u,Ptr v){ // cut edge between u and v
        evert(u);
        expose (v);
        assert (u->p==v);
        v->l=u->p=nullptr;
        this->pull(v);
    Ptr get_root(Ptr t){
        expose(t);
        while (t->1) this->push (t), t=t->1;
        this->splay(t);
        return t;
    Ptr lca(Ptr u,Ptr v) {
        if (get_root(u)!=get_root(v))return nullptr;
        expose(u);
        return expose(v);
    void set_val(Ptr t,const T &val){
        this->evert(t);
        t->val=val;
        this->pull(t);
    T get val(Ptr t){
        this->evert(t);
        return t->val;
    T fold (Ptr u, Ptr v) {
        evert(u);
        expose (v);
        return v->sum;
};
```

LazyLinkCutTree.hpp
Description: Lazy Link Cut Tree.

```
Chula[) ? - ? )]
```

using Ptr = Lct::Ptr;

vector<Ptr> ptr(n);

using Node = Lct:: Node;

auto link=[&](int u,int v){
Lct::link(ptr[u],ptr[v]);

int to t=0;

Usage: using Lct = LazyLinkCutTree<Action>;

for(int i=0;i<n;i++)ptr[i]=new Node(val[i]);</pre>

```
};
auto cut = [\&] (int u, int v) {
Lct: : cut(ptr[u],ptr[v]);
auto update=[&](int u,int v,const Action:: Tag &val){
Lct: : apply(ptr[u],ptr[v],val);
};
auto query=[&](int u,int v){
return Lct::fold(ptr[u],ptr[v]);
"LazyReversibleSplayTree.hpp", "LinkCutTreeBase.hpp"
                                                       ead3da, 12 lines
template < class MonoidAction >
struct LazyLinkCutTree:LinkCutTreeBase<LazyReversibleSplayTree</pre>
     MonoidAction>>{
    using base = LinkCutTreeBase<LazyReversibleSplayTree<</pre>
         MonoidAction>>;
    using Ptr = typename base::Ptr;
    using Tag = typename MonoidAction::Tag;
    void apply (Ptr u, Ptr v, const Tag &val) {
        this->evert(u);
        this->expose(v);
        this->propagate(v, val);
};
StaticTopTree.hpp
Description: Static Top Tree.
                                                      7e10be, 186 lines
template<class HLD>
struct StaticTopTree{
    using P = pair<int,int>;
    enum Type{Compress, Rake, AddEdge, AddVertex, Vertex};
    int n, root;
    HLD &hld;
    vector<int> lch,rch,par;
    vector<Type> type;
    StaticTopTree(HLD &_hld):hld(_hld){build();}
    void build(){
        n=hld.n:
        lch=rch=par=vector<int>(n,-1);
        type.assign(n,Compress);
        root=compress(hld.root).second;
    int add(int i,int l,int r,Type t){
        if (i==-1) {
            i=n++;
            lch.emplace_back(1);
            rch.emplace_back(r);
            par.emplace_back(-1);
            type.emplace_back(t);
        }else{
            lch[i]=l,rch[i]=r,type[i]=t;
        if(1!=-1)par[1]=i;
        if(r!=-1)par[r]=i;
        return i;
    pair<int, int> merge(vector<pair<int, int>> a, Type t) {
        if(a.size()==1)return \ a[0];
```

```
vector < pair < int, int >> l, r;
        for (auto [i, s]: a) to t+=s;
        for(auto [i,s]:a) {
             (tot>s?l:r).emplace back(i,s);
             t \circ t = s * 2;
        auto [i, si]=merge(l, t);
        auto [j, sj] = merge(r, t);
        return \{add(-1,i,j,t), si+sj\};
    P compress(int i) {
        vector<P> a{add_vertex(i)};
        auto work=[&](){
            auto [sj,j]=a.back();
            a.pop_back();
            auto [si,i]=a.back();
            a.back() = \{\max(si, sj) + 1, add(-1, i, j, Compress)\};
        while(hld.hv[i]!=-1){
            a.emplace_back(add_vertex(i=hld.hv[i]));
            while(true) {
                if(a.size() >= 3&& (a.end()[-3].first == a.end()
                      [-2].first||a.end()[-3].first<=a.back().
                     first)){
                    P tmp=a.back();
                    a.pop_back();
                     work();
                     a.emplace_back(tmp);
                 }else if(a.size()>=2&&a.end()[-2].first<=a.back
                     ().first){
                     work();
                }else break;
        while (a.size()>=2) work();
        return a[0];
        priority_queue<P, vector<P>, greater<P>> pq;
        for(int j:hld.g[i])if(j!=hld.par[i]&&j!=hld.hv[i])pq.
             emplace(add edge(j));
        while (pq.size()>=2) {
            auto [si,i]=pq.top();pq.pop();
            auto [sj,j]=pq.top();pq.pop();
            pq.emplace (max(si,sj)+1,add(-1,i,j,Rake));
        return pg.empty()?make_pair(0,-1):pg.top();
    P add_edge(int i) {
        auto [sj,j]=compress(i);
        return {sj+1,add(-1,j,-1,AddEdge)};
    P add vertex(int i) {
        auto [sj,j]=rake(i);
        return {sj+1,add(i,j,-1,j==-1?Vertex:AddVertex)};
};
struct\ TreeDP{
    struct Path f
        static Path unit();
    };
    struct Points
        static Point unit();
    static Path compress(Path l, Path r);
    static Point rake(Point 1, Point r);
```

```
static\ Point\ add\_edge(Path\ p);
    static Path add vertex(Point p, int u);
    static\ Path\ ver\overline{t}ex(int\ u);
template < class HLD, class TreeDP>
struct StaticTopTreeRerootingDP{
    using Path = typename TreeDP::Path;
    using Point = typename TreeDP::Point;
    StaticTopTree<HLD> stt;
    vector<Path> path,rpath;
    vector<Point> point;
    StaticTopTreeRerootingDP(HLD &hld):stt(hld) {
        int n=stt.n;
        path.resize(n);
        point.resize(n);
        rpath.resize(n);
        dfs(stt.root);
    void _update(int u) {
        if(stt.type[u] == stt.Vertex) {
            path[u]=rpath[u]=TreeDP::vertex(u);
        }else if(stt.type[u] == stt.Compress) {
            path[u]=TreeDP::compress(path[stt.lch[u]],path[stt.
            rpath[u] = TreeDP::compress(rpath[stt.rch[u]], rpath[
                 stt.lch[u]]);
        }else if(stt.type[u] == stt.Rake) {
            point[u] = TreeDP::rake(point[stt.lch[u]], point[stt.
                 rch[u]]);
        }else if(stt.type[u] == stt.AddEdge) {
            point[u]=TreeDP::add_edge(path[stt.lch[u]]);
        }else{
            path[u]=rpath[u]=TreeDP::add_vertex(point[stt.lch[u
    void dfs(int u){
        if (u==-1) return;
        dfs(stt.lch[u]);
        dfs(stt.rch[u]);
        _update(u);
    void update(int u) {
        for(;u!=-1;u=stt.par[u])_update(u);
    Path query_all(){
        return path[stt.root];
    Path query_subtree(int u) {
        Path res=path[u];
        while(true) {
            int p=stt.par[u];
            if (p==-1 | | stt.type[p]!=stt.Compress) break;
            if (stt.lch[p] == u) res=TreeDP::compress(path[stt.rch[
                 p]],res);
        return res;
    Path query_reroot(int u) {
        auto rec=[&] (auto &&rec,int u) ->Point{
            int p=stt.par[u];
            Path below=Path::unit(),above=Path::unit();
            while (p!=-1&&stt.type[p]==stt.Compress) {
                int l=stt.lch[p], r=stt.rch[p];
                if(l==u)below=TreeDP::compress(below,path[r]);
                else above=TreeDP::compress(above,rpath[1]);
                u=p;
```

```
p=stt.par[u];
    if(p!=-1){
        u=p;
        p=stt.par[u];
        Point sum=Point::unit();
        while(stt.type[p]==stt.Rake){
            int l=stt.lch[p],r=stt.rch[p];
            sum=TreeDP::rake(sum,u==r?point[1]:point[r
                1);
            p=stt.par[u];
        sum=TreeDP::rake(sum, rec(rec,p));
        above=TreeDP::compress(above,TreeDP::add_vertex
    return TreeDP::rake(TreeDP::add_edge(below), TreeDP
         ::add_edge(above));
};
Point res=rec(rec,u);
if (stt.type[u] == stt.AddVertex) {
    res=TreeDP::rake(res,point[stt.lch[u]]);
return TreeDP::add_vertex(res,u);
```

## Number Theory (7)

### Extended Euclid. hpp

Description: Extended Euclid algorithm for solving diophantine equation (ax + by = gcd(a, b)).

Time:  $\mathcal{O}(\log \max\{a, b\})$ 

```
"../template/Header.hpp"
                                                             229e7c, 13 lines
pair<ll, ll> euclid(ll a, ll b) {
    11 x=1, y=0, x1=0, y1=1;
    while(b!=0){
         11 q=a/b;
         x = q * x1;
         y = q * y1;
         a = q * b;
         swap(x,x1);
         swap(y,y1);
        swap(a,b);
    return {x,y};
```

#### euclid.h

**Description:** Finds two integers x and y, such that  $ax + by = \gcd(a, b)$ . If you just need gcd, use the built in \_\_gcd instead. If a and b are coprime, then x is the inverse of a (mod b).  $x = x_0 + k * (b/g)$   $y = y_0 - k * (a/g)_{\text{ines}}$ 

```
ll euclid(ll a, ll b, ll &x, ll &y) {
 if (!b) return x = 1, y = 0, a;
 ll d = euclid(b, a % b, y, x);
 return y -= a/b * x, d;
```

Description: Chinese Remainder Theorem.

crt (a, m, b, n) computes x such that  $x \equiv a \pmod{m}$ ,  $x \equiv b \pmod{n}$ . If |a| < m and |b| < n, x will obey  $0 \le x < \operatorname{lcm}(m, n)$ . Assumes  $mn < 2^{62}$ . If x0 and y0 is one of the solutions of ax + by = g, then the general solution is x = x0 + k \* (b / g) and y = y0 - k \* (a / g).

```
Time: \log(n)
"euclid.h"
                                                                       04d93a, 7 lines
11 crt(l1 a, l1 m, l1 b, l1 n) {
  if (n > m) swap(a, b), swap(m, n);
  11 x, y, g = euclid(m, n, x, y);
  assert ((a - b) % g == 0); // else no solution
  x = (b - a) % n * x % n / q * m + a;
  return x < 0 ? x + m*n/g : x;
phiFunction.hpp
Description: Euler's \phi function is defined as \phi(n) := \# of positive integers
\leq n that are coprime with n. \phi(1) = 1, p prime \Rightarrow \phi(p^k) = (p-1)p^{k-1}
m, n \text{ coprime} \Rightarrow \phi(mn) = \phi(m)\phi(n). If n = p_1^{k_1} p_2^{k_2} \dots p_r^{k_r} then \phi(n) = p_1^{k_1} p_2^{k_2} \dots p_r^{k_r}
(p_1-1)p_1^{k_1-1}...(p_r-1)p_r^{k_r-1}. \phi(n)=n\cdot\prod_{p\mid n}(1-1/p).
\sum_{d|n} \phi(d) = n, \sum_{1 \le k \le n, \gcd(k,n)=1} k = n\phi(n)/2, n > 1
Euler's thm: a, n coprime \Rightarrow a^{\phi(n)} \equiv 1 \pmod{n}.
Fermat's little thm: p \text{ prime } \Rightarrow a^{p-1} \equiv 1 \pmod{p} \ \forall a.
                                                                       efae90, 10 lines
const int LIM = 5000000;
int phi[LIM];
void calculatePhi() {
  for(int i=0; i<LIM; ++i) phi[i] = i & 1 ? i : i / 2;</pre>
  for (int i = 3; i < LIM; i += 2)</pre>
     if (phi[i] == i)
        for (int j = i; j < LIM; j += i)</pre>
          phi[j] -= phi[j] / i;
```

### FloorSum.hpp

**Description:** Floor sum function.  $f(a,b,c,n) = \sum_{x=0}^{n} \lfloor \frac{ax+b}{c} \rfloor$  becareful when a,b,c are negetive (use custom floor division and mod instead) Time:  $O(\log a)$ 

d088d2, 7 lines 11 floor\_sum(l1 a, l1 b, l1 c, l1 n) { 11 res=n\*(n+1)/2\*(a/c)+(n+1)\*(b/c);a%=c,b%=c; if(a==0) return res; 11 m=(a\*n+b)/c;return res+n\*m-floor\_sum(c,c-b-1,a,m-1);

### 7.1 Prime Numbers

### Miller Rabin.hpp

Description: Deterministic Miller-Rabin primality test. Guaranteed to work for numbers up to 7 · 10<sup>18</sup>; for larger numbers, use Python and ex-

**Time:** 7 times the complexity of  $a^b \mod c$ .

be7e00, 25 lines

```
using ull = uint64 t;
ull modmul(ull a, ull b, ull M) {
  ll ret = a * b - M * ull(1.L / M * a * b);
  return ret + M * (ret < 0) - M * (ret >= (11)M);
ull modpow(ull b, ull e, ull mod) {
 ull ans = 1:
  for (; e; b = modmul(b, b, mod), e /= 2)
   if (e & 1) ans = modmul(ans, b, mod);
  return ans:
bool isPrime(ull n) {
  if (n < 2 || n % 6 % 4 != 1) return (n | 1) == 3;</pre>
 ull A[] = \{2, 325, 9375, 28178, 450775, 9780504, 1795265022\},
      s = builtin ctzll(n-1), d = n >> s;
```

```
for (ull a : A) { // ^ count trailing zeroes
  ull p = modpow(a%n, d, n), i = s;
  while (p != 1 && p != n - 1 && a % n && i--)
   p = modmul(p, p, n);
  if (p != n-1 && i != s) return 0;
return 1;
```

### LinearSieve.hpp

Description: Prime Number Generator in Linear Time Time:  $\mathcal{O}(N)$ 

```
"../template/Header.hpp"
                                                         194fb1, 15 lines
vi linear sieve(int n) {
 vi prime, composite(n + 1);
 for(int i=2; i<=n; ++i) {</pre>
    if(!composite[i]) {
      prime.emplace back(i);
    for(int j=0; j<(int) prime.size() && i*prime[j]<=n; ++j) {</pre>
      composite[i * prime[j]] = true;
      if(i % prime[j] == 0) {
        break;
 return prime;
```

#### Fast Eratosthenes. hpp

return pr;

**Description:** Prime sieve for generating all primes smaller than LIM. Time: LIM=1e9  $\approx 1.5$ s

```
"../template/Header.hpp"
                                                        295b58, 33 lines
const int LIM = 1e6;
bitset<LIM> isPrime:
vi eratosthenes() {
  const int S = (int) round(sqrt(LIM)), R = LIM / 2;
  vi pr = \{2\}, sieve(S + 1);
  pr.reserve(int(LIM/log(LIM) * 1.1));
  vector<pii> cp;
  for(int i=3; i<=S; i+=2) {
    if(!sieve[i]) {
      cp.emplace_back(i, i * i / 2);
      for(int j=i*i; j<=S; j+=2*i) {</pre>
        sieve[i] = 1;
  for(int L=1; L<=R; L+=S) {</pre>
    array<bool, S> block{};
    for(auto &[p, idx]: cp) {
      for(int i=idx; i<S+L; idx=(i+=p)) {</pre>
        block[i - L] = 1;
    for(int i=0; i<min(S, R-L); ++i) {</pre>
      if(!block[i]) {
        pr.emplace_back((L + i) \star 2 + 1);
  for(int i: pr) {
    isPrime[i] = 1;
```

```
GolbatchConjecture.hpp
Description: Find two prime numbers which sum equals s
Time: \mathcal{O}(N \log N)
"FastEratosthenes.hpp"
                                                       88fb23, 18 lines
pair<int, int> goldbatchConjecture(int s, vi pr = {}) {
  if (s <= 2 || s % 2 != 0) {
    return make_pair(-1, -1);
  if (pr.size() == 0) {
    pr = eratosthenes();
  for (auto x : pr) {
    if (x > s / 2) {
     break;
    int d = s - x;
   if (binary_search(pr.begin(), pr.end(), d)) {
      return make_pair(min(x, d), max(x, d));
  return make_pair(-1, -1);
```

## Graph (8)

## 8.1 Matching

HopcroftKarp.hpp

struct HopcroftKarp{

bool dfs(int u) {

Description: Fast bipartite matching algorithm.

```
Time: \mathcal{O}\left(E\sqrt{V}\right)
"../template/Header.hpp
```

```
int n.m:
vi l,r,lv,ptr;
vector<vi> adj;
HopcroftKarp(){}
HopcroftKarp(int _n,int _m) {init(_n,_m);}
void init(int _n,int _m){
    n=_n, m=_m;
    adj.assign(n+m, vi{});
void addEdge(int u,int v) {
    adj[u].emplace_back(v+n);
void bfs() {
    lv=vi(n,-1);
    queue<int> q;
    for (int i=0; i<n; i++) if (1[i]==-1) {</pre>
        lv[i]=0;
        q.emplace(i);
    while(!q.empty()){
        int u=q.front();
        for(int v:adj[u])if(r[v]!=-1&&lv[r[v]]==-1){
            lv[r[v]]=lv[u]+1;
```

q.emplace(r[v]);

int v=adj[u][i];

l[u]=v,r[v]=u;

return true;

for(int &i=ptr[u];i<sz(adj[u]);i++){</pre>

if(r[v] == -1 | | (lv[r[v]] == lv[u] + 1 & & dfs(r[v])))

```
return false:
    int maxMatching() {
        int match=0, cnt=0;
        1=r=vi(n+m,-1);
             ptr=vi(n);
             bfs();
             cnt=0:
             for (int i=0; i < n; i++) if (l[i] ==-1&&dfs(i)) cnt++;</pre>
             match+=cnt;
         }while(cnt);
        return match;
};
```

### Kuhn.hpp

Description: Kuhn Algorithm to find maximum bipartite matching or find augmenting path in bipartite graph.

```
Time: \mathcal{O}(VE)
```

```
"../template/Header.hpp"
                                                                         fc7d17, 15 lines
               vi adi[1010], match(1010, -1);
               bitset<1010> visited;
               bool kuhn (int u) {
                 if(visited[u]) {
                    return false;
                 visited[u] = true;
                 for(auto x: adi[u]) {
                   if(match[x] == -1 \mid \mid kuhn(match[x])) {
                     match[x] = u;
Obd56f, 52 lines
                     return true;
                 return false;
```

### Weighted Matching.hpp

Description: Given a weighted bipartite graph, matches every node on the left with a node on the right such that no nodes are in two matchings and the sum of the edge weights is minimal. Takes cost[N][M], where cost[i][j] = cost for L[i] to be matched with R[j] and returns (min cost, match), where L[i] is matched with R[match[i]]. Negate costs for max cost. Requires  $N \leq M$ .

```
Time: \mathcal{O}\left(N^2M\right)
pair<ll, vector<int>> hungarian(const vector<vector<ll>> &a) {
 if (a.emptv()) return {0, {}};
 int n = a.size() + 1, m = a[0].size() + 1;
 vector<ll> u(n), v(m), p(m);
    vector<int> ans(n - 1);
 for(int i=1;i<n;i++) {</pre>
    p[0] = i;
    int j0 = 0; // add "dummy" worker 0
    vector<ll> dist(m, LLONG_MAX), pre(m, -1);
    vector<bool> done(m + 1);
    do { // dijkstra
      done[j0] = true;
      int i0 = p[j0], j1;
            11 delta = LLONG_MAX;
      for(int j=1; j<m; j++) if (!done[j]) {</pre>
        auto cur = a[i0 - 1][j - 1] - u[i0] - v[j];
        if (cur < dist[j]) dist[j] = cur, pre[j] = j0;</pre>
        if (dist[j] < delta) delta = dist[j], j1 = j;</pre>
      for(int j=0; j<m; j++) {</pre>
        if (done[j]) u[p[j]] += delta, v[j] -= delta;
        else dist[j] -= delta;
```

```
j0 = j1;
  } while (p[j0]);
  while (j0) { // update alternating path
    int j1 = pre[j0];
    p[j0] = p[j1], j0 = j1;
for(int j=1; j<m; j++) if (p[j]) ans[p[j] - 1] = j - 1;</pre>
return {-v[0], ans}; // min cost
```

### Network Flow

#### Dinic.hpp

```
Description: Dinic's Algorithm for finding the maximum flow.
Time: \mathcal{O}(VE \log U) where U is the maximum flow.
template < class T, bool directed = true, bool scaling = true >
struct Dinic{
    static constexpr T INF=numeric_limits<T>::max()/2;
    struct Edge{
        int to;
        T flow, cap;
        Edge(int _to,T _cap):to(_to),flow(0),cap(_cap){}
        T remain() {return cap-flow;}
    int n,s,t;
    T U;
    vector<Edge> e;
    vector<vector<int>> q;
    vector<int> ptr.lv:
    bool calculated;
    T max_flow;
    Dinic(){}
    Dinic(int n, int s, int t) {init(n, s, t);}
    void init(int _n,int _s,int _t){
        n=_n,s=_s,t=_t;
        U=0;
        e.clear();
        g.assign(n, {});
        calculated=false;
    void add edge(int from,int to,T cap){
        assert(0<=from&&from<n&&0<=to&&to<n);
        g[from].emplace_back(e.size());
        e.emplace_back(to,cap);
        g[to].emplace_back(e.size());
        e.emplace_back(from,directed?0:cap);
        U=max(U,cap);
    bool bfs(T scale) {
        lv.assign(n,-1);
        vector<int> q{s};
        lv[s]=0;
        for(int i=0;i<(int)q.size();i++){</pre>
            int u=q[i];
            for(int j:g[u]){
                int v=e[j].to;
                if(lv[v] ==-1&&e[j].remain()>=scale){
                     q.emplace_back(v);
                     lv[v]=lv[u]+1;
        return lv[t]!=-1;
    T dfs(int u, int t, T f) {
        if(u==t||f==0)return f;
        for(int &i=ptr[u];i<(int)q[u].size();i++){</pre>
            int j=g[u][i];
```

```
int v=e[j].to;
             if(lv[v]==lv[u]+1){
                 T res=dfs(v,t,min(f,e[j].remain()));
                 if(res>0){
                     e[j].flow+=res;
                     e[j^1].flow-=res;
                     return res;
        return 0;
    T flow(){
        if(calculated)return max flow;
        calculated=true;
        max flow=0;
        for(T scale=scaling?1LL<<(63-__builtin_clzll(U)):1LL;</pre>
             scale>0;scale>>=1) {
             while(bfs(scale)){
                ptr.assign(n,0);
                 while(true) {
                     T f=dfs(s,t,INF);
                     if (f==0) break;
                     max_flow+=f;
        return max flow;
    pair<T, vector<int>> cut() {
        vector<int> res(n);
        for(int i=0;i<n;i++)res[i]=(lv[i]==-1);</pre>
        return {max_flow,res};
};
MinCostFlow.hpp
Description: minimum-cost flow algorithm.
Time: \mathcal{O}(FE \log V) where F is max flow.
"../template/Header.hpp"
                                                       8ea1d2, 83 lines
template < class F, class C>
struct MinCostFlow{
    struct Edge {
        int to;
        F flow, cap;
        Edge(int _to,F _cap,C _cost):to(_to),flow(0),cap(_cap),
             cost(_cost){}
        F getcap(){
             return cap-flow;
    };
    int n;
    vector<Edge> e;
    vector<vi> adj;
    vector<C> pot, dist;
    vi pre;
    bool neg;
    const F FINF=numeric_limits<F>::max()/2;
    const C CINF=numeric_limits<C>::max()/2;
    MinCostFlow(){}
    MinCostFlow(int _n) {
        init(n);
    void init(int _n){
        n=_n;
        e.clear();
        adj.assign(n,{});
```

```
neg=false;
    void addEdge(int u,int v,F cap,C cost){
         adj[u].emplace_back(sz(e));
         e.emplace_back(v,cap,cost);
         adj[v].emplace_back(sz(e));
         e.emplace_back(u,0,-cost);
         if (cost<0) neg=true;</pre>
    bool dijkstra(int s,int t) {
         using P = pair<C,int>;
         dist.assign(n,CINF);
         pre.assign(n,-1);
         priority_queue<P, vector<P>, greater<P>> pq;
         dist[s]=0;
         pq.emplace(0,s);
         while(!pq.empty()){
              auto [d,u]=pq.top();
              pq.pop();
              if (dist[u] < d) continue;</pre>
              for(int i:adj[u]){
                  int v=e[i].to;
                  C ndist=d+pot[u]-pot[v]+e[i].cost;
                  if(e[i].getcap()>0&&dist[v]>ndist){
                       pre[v]=i;
                       dist[v]=ndist;
                       pg.emplace(ndist, v);
         return dist[t] < CINF;
    pair<F,C> flow(int s,int t) {
         F flow=0;
         C cost=0;
         pot.assign(n,0);
         if (neq) for (int t=0; t< n; t++) for (int i=0; i< sz(e); i++) if (e
               [i].getcap()>0){
              int u=e[i^1].to, v=e[i].to;
              pot[v]=min(pot[v],pot[u]+e[i].cost);
         } // Bellman-Ford
         while (dijkstra(s,t)) {
              for (int i=0; i < n; i++) pot[i] += dist[i];</pre>
              for(int u=t;u!=s;u=e[pre[u]^1].to){
                   aug=min(aug,e[pre[u]].getcap());
              } // find bottleneck
              for(int u=t;u!=s;u=e[pre[u]^1].to){
                  e[pre[u]].flow+=aug;
                  e[pre[u]^1].flow-=aug;
              } // push flow
              flow+=aug;
              cost+=aug*pot[t];
         return {flow,cost};
};
Binary Optimization. hpp
Description: Binary Optimization.
                                          minimize \kappa + \sum_{i} \theta_{i}(x_{i}) +
\textstyle \sum_{i < j} \phi_{ij}(x_i, x_j) + \sum_{i < j < k} \psi_{ijk}(x_i, x_j, x_k) \text{ where } x_i \in \{0, 1\} \text{ and } \phi_{ij}, \psi_{ijk}
are submodular functions. a set function f is submodular if f(S) + f(T) >
f(S \cap T) + f(S \cup T) for all S, T. \phi_{ij}(0,1) + \phi_{ij}(1,0) \ge \phi_{ij}(1,1) + \phi_{ij}(0,0).
template < class T, bool minimize = true >
struct BinaryOptimization{
     static constexpr T INF=numeric_limits<T>::max()/2;
    int n,s,t,node id;
```

```
T base:
map<pair<int,int>,T> edges;
BinaryOptimization(int _n):n(_n),s(n),t(n+1),node_id(n+2),
     base(0){}
void add_edge(int u,int v,T w) {
    assert (w \ge 0);
    if (u==v | |w==0) return;
    auto &e=edges[{u,v}];
    e=min(e+w,INF);
void add0(T w){
    base+=w;
void add1(int i,T a,T b) {
    if(a<=b){
        add0(a);
        add_edge(s,i,b-a);
    }else{
        add0(b);
        add_edge(i,t,a-b);
void add1(int i,T x0,T x1){
    assert(0<=i&&i<n);
    if (!minimize) x0=-x0, x1=-x1;
    \_add1(i, x0, x1);
void _add2(int i,int j,T a,T b,T c,T d) {
    assert (b+c>=a+d);
    add0(a);
    _add1(i,0,c-a);
    _add1(j,0,d-c);
    add_edge(i,j,b+c-a-d);
void add2(int i,int j,T x00,T x01,T x10,T x11){
    assert(i!=j&&0<=i&&i<n&&0<=j&&j<n);
    if(!minimize)x00=-x00,x01=-x01,x10=-x10,x11=-x11;
    _{add2}(i, j, x00, x01, x10, x11);
void _add3(int i,int j,int k,T a,T b,T c,T d,T e,T f,T q,T
    T p=a+d+f+g-b-c-e-h;
    if(p>=0){
        add0(a);
        _add1(i,0,f-b);
        _add1(j,0,g-e);
        add1(k,0,d-c);
        _{add2}(i, j, 0, c+e-a-q, 0, 0);
        add2(i,k,0,0,b+e-a-f,0);
        \_add2(j,k,0,b+c-a-d,0,0);
        int u=node_id++;
        add0(-p);
        add edge(i,u,p);
        add_edge(j,u,p);
        add_edge(k,u,p);
        add_edge(u,t,p);
    }else{
        add0(h);
        _add1(i,c-g,0);
        _add1(j,b-d,0);
        _add1(k,e-f,0);
        _{add2}(i, j, 0, 0, d+f-b-h, 0);
        _{add2}(i,k,0,d+g-c-h,0,0);
        \_add2(j,k,0,0,f+g-e-h,0);
        int u=node_id++;
        add0(p);
        add_edge(s,u,-p);
        add_edge(u,i,-p);
        add_edge(u,j,-p);
```

### KaryOptimization SCC LowLink

```
add_edge(u,k,-p);
    void add3(int i,int j,int k,T x000,T x001,T x010,T x011,T
         x100,T x101,T x110,T x111){
         assert(i!=j&&j!=k&&k!=i&&0<=i&&i<n&&0<=j&&j<n&&0<=k&&k<
        if(!minimize){
             x000=-x000, x001=-x001, x010=-x010, x011=-x011;
             x100=-x100, x101=-x101, x110=-x110, x111=-x111;
         _add3(i, j, k, x000, x001, x010, x011, x100, x101, x110, x111);
    T solve(){
        Dinic<T> dinic(node_id,s,t);
         for(auto &[p,w]:edges){
             auto [u,v]=p;
             dinic.add_edge(u,v,w);
         T ans=dinic.flow()+base;
         return minimize?ans:-ans;
};
KaryOptimization.hpp
Description: k-ary Optimization. minimize \kappa + \sum_{i} \theta_{i}(x_{i}) +
\sum_{i < j} \phi_{ij}(x_i, x_j) where x_i \in \{0, 1, \dots, k-1\} and \phi_{i,j} is monge. A func-
tion f is monge if f(a,c) + f(b,d) \le f(a,d) + f(b,c) for all a < b and
c < d. \phi_{ij}(x-1,y) + \phi_{ij}(x,y+1) \le \phi_{ij}(x-1,y+1) + \phi_{ij}(x,y).
\phi_{ij}(x,y) + \phi_{ij}(x-1,y+1) - \phi_{ij}(x-1,y) - \phi_{ij}(x,y+1) \ge 0.
                                                          422f8a, 88 lines
template < class T, bool minimize = true >
struct K aryOptimization{
    static constexpr T INF=numeric_limits<T>::max()/2;
    int n,s,t,node_id;
    T base;
    vector<int> ks:
    vector<vector<int>> id;
    map<pair<int,int>,T> edges;
    K_aryOptimization(int n, int k) {init(vector<int>(n,k));}
    K_aryOptimization(const vector<int> &_ks) {init(_ks);}
    void init(const vector<int> &_ks){
        ks= ks;
        n=ks.size();
        s=0, t=1, node_id=2;
        base=0:
        id.clear();
         edges.clear();
         for(auto &k:ks) {
             assert (k \ge 1):
             vector<int> a(k+1);
             a[0]=s,a[k]=t;
             for(int i=1;i<k;i++)a[i]=node_id++;</pre>
             id.emplace_back(a);
             for (int i=2; i < k; i++) add_edge(a[i], a[i-1], INF);</pre>
    void add_edge(int u,int v,T w) {
        assert (w \ge 0):
        if (u==v | |w==0) return;
        auto &e=edges[{u,v}];
        e=min(e+w,INF);
    void add0(T w){
        base+=w;
    void _add1(int i, vector<T> cost) {
        add0(cost[0]);
```

```
for (int j=1; j<ks[i]; j++) {</pre>
             T x=cost[j]-cost[j-1];
             if(x>0)add_edge(id[i][j],t,x);
             if(x<0) add0(x), add_edge(s,id[i][j],-x);
    void add1(int i,vector<T> cost){
         assert (0<=i&&i<n&& (int) cost.size() ==ks[i]);
        if(!minimize)for(auto &x:cost)x=-x;
         add1(i.cost);
    void _add2(int i,int j,vector<vector<T>> cost){
         int h=ks[i], w=ks[j];
         _add1(j,cost[0]);
         for (int x=h-1; x>=0; x--) for (int y=0; y<w; y++) cost [x] [y] -=</pre>
              cost[0][y];
         vector<T> a(h);
         for (int x=0; x<h; x++) a[x]=cost[x][w-1];</pre>
         _add1(i,a);
         for (int x=1; x<h; x++) {</pre>
             for (int y=0; y<w-1; y++) {</pre>
                   \texttt{T} \ \texttt{w} = \texttt{cost}[\texttt{x}][\texttt{y}] + \texttt{cost}[\texttt{x} - 1][\texttt{y} + 1] - \texttt{cost}[\texttt{x} - 1][\texttt{y}] - \texttt{cost} 
                       [x][y+1];
                  assert(w>=0); // monge
                  add_edge(id[i][x],id[j][y+1],w);
    void add2(int i,int j,vector<vector<T>> cost){
         assert(0<=i&&i<n&&0<=j&&j<n&&i!=j);
         assert((int)cost.size() == ks[i]);
         for(auto &v:cost)assert((int)v.size() == ks[j]);
        if(!minimize) for(auto &v:cost) for(auto &x:v) x=-x;
         _add2(i,j,cost);
    pair<T, vector<int>> solve() {
        Dinic<T> dinic(node id,s,t);
         for(auto &[p,w]:edges){
             auto [u,v]=p;
             dinic.add_edge(u,v,w);
         auto [val,cut]=dinic.cut();
        val+=base:
         if(!minimize)val=-val;
         vector<int> ans(n);
         for (int i=0; i<n; i++) {</pre>
             ans[i]=ks[i]-1;
             for(int j=1; j<ks[i]; j++) ans[i] -=cut[id[i][j]];</pre>
         return {val, ans};
};
       Connectivity
8.3
SCC.hpp
Description: Strongly Connected Component.
                                                          82a9d1, 34 lines
"../template/Header.hpp"
template < class G>
pair<int, vector<int>> strongly_connected_component (G &g) {
    static_assert (G::is_directed);
    int n=q.n;
    vector<int> disc(n,-1), low(n), scc(n,-1);
    stack<int> st;
    vector<bool> in_st(n);
    int t=0,scc_cnt=0;
    function<void(int,int)> dfs=[&](int u,int p){
```

```
disc[u]=low[u]=t++;
        st.emplace(u);
        in st[u]=true;
        for(int v:g[u]){
            if(disc[v]==-1){
                 dfs(v,u);
                 low[u]=min(low[u],low[v]);
            }else if(in_st[v]){
                 low[u]=min(low[u],disc[v]);
        if (disc[u] == low[u]) {
             while(true) {
                 int v=st.top();
                 st.pop();
                 in_st[v]=false;
                 scc[v]=scc_cnt;
                 if (v==u) break;
             scc_cnt++;
    };
    for (int i=0; i<n; i++) if (disc[i] ==-1) dfs(i,-1);</pre>
    return {scc_cnt,scc};
LowLink.hpp
Description: Low Link.
                                                        f4ad2f. 33 lines
template<class G>
struct LowLink{
    G &g;
    int n;
    vector<int> disc, low, par, ord;
    vector<pair<int,int>> bridge;
    vector<int> articulation;
    LowLink(G &_g):g(_g), n(g.n), disc(n,-1), low(n), par(n,-1) {
        for (int i=0; i<n; i++) if (disc[i] ==-1) dfs(i);</pre>
    void dfs(int u) {
        disc[u]=low[u]=t++;
        ord.emplace_back(u);
        int child=0;
        bool found_par=false;
        for(int v:q[u]){
             if (disc[v] ==-1) {
                 par[v]=u;
                 dfs(v);
                 low[u] = min(low[u], low[v]);
                 if(low[v]>disc[u])bridge.emplace_back(u,v);
                 if (par[u]!=-1&&low[v]>=disc[u]) articulation.
                      emplace_back(u);
                 child++;
             }else if(v!=par[u]||found_par){
                 low[u]=min(low[u],disc[v]);
                 found_par=true;
        if (par[u] ==-1&&child>1) articulation.emplace_back(u);
};
```

Tree (9)

Description: HLD

HLD.hpp

### HLD CentroidDecom FormalPowerSeries

if(sz[v] \* 2 > t) find cen(v, u, t);

```
cf6882, 45 lines
"../template/Header.hpp"
vector<vi> adj;
vector<int> sz, lvl, hv, hd, p, disc;
void dfs(int u, int parent) {
  sz[u] = 1;
  lvl[u] = lvl[parent] + 1;
  p[u] = parent;
  int c_hv=0, c_max=0;
  for(auto v: adj[u]) {
   if(v == parent) continue;
   dfs(v, u);
   sz[u] += sz[v];
   if(c_max < sz[v])  {
     c_hv = v;
      c_max = sz[v];
 hv[u] = c hv;
void hld(int u, int parent) {
 if(hd[u] == 0) {
   hd[u] = u;
  disc[u] = ++t;
  if(hv[u] != 0) {
   hd[hv[u]] = hd[u];
   hld(hv[u], u);
  for(auto v: adj[u]) {
   if(v == parent || v == hv[u]) {
      continue;
   hld(v, u);
int lca(int u, int v) {
  while(hd[u] != hd[v]) {
   if(lvl[hd[u]] > lvl[hd[v]]) swap(u, v);
    v=p[hd[v]];
 return lvl[u] < lvl[v] ? u: v;</pre>
Centroid Decom. hpp
Description: Centroid
                                                      e46d44, 32 lines
"../template/Header.hpp"
vector<vi> adj;
vi sz;
vector<bool> used;
int find_size(int u, int p) {
  sz[u] = 1;
  for(auto v: adj[u]) {
   if(v == p || used[v]) continue;
    sz[u] += find_size(v, u);
  return sz[u];
int find_cen(int u, int p, int t) {
  for(auto v: adj[u]) {
    if(v == p || used[v]) continue;
```

```
return u;
void decom(int u) {
 u = find_cen(u, 0, find_size(u, 0));
 used[u] = true;
  for(auto v: adj[u]) {
    // dfs do something
  for(auto v: adj[u]) {
   if(used[v]) continue;
    decom(v);
Polynomials (10)
FormalPowerSeries.hpp
Description: basic operations of formal power series
                                                      416433, 136 lines
template < class mint >
struct FormalPowerSeries:vector<mint>{
    using vector<mint>::vector;
    using FPS = FormalPowerSeries;
    FPS & operator += (const FPS & rhs) {
        if (rhs.size()>this->size())this->resize(rhs.size());
        for(int i=0;i<rhs.size();i++)(*this)[i]+=rhs[i];</pre>
        return *this;
    FPS & operator += (const mint &rhs) {
        if (this->empty())this->resize(1);
        (*this) [0] += rhs;
        return *this;
    FPS & operator = (const FPS &rhs) {
        if (rhs.size()>this->size())this->resize(rhs.size());
        for (int i=0; i<rhs.size(); i++) (*this)[i]-=rhs[i];</pre>
        return *this;
    FPS & operator -= (const mint &rhs) {
        if(this->empty())this->resize(1);
        (*this)[0]-=rhs;
        return *this;
    FPS & operator *= (const FPS &rhs) {
        auto res=NTT<mint>()(*this, rhs);
        return *this=FPS(res.begin(),res.end());
    FPS & operator *= (const mint &rhs) {
        for(auto &a:*this)a*=rhs;
        return *this;
    friend FPS operator+(FPS lhs,const FPS &rhs) {return lhs+=
    friend FPS operator+(FPS lhs, const mint &rhs) {return lhs+=
    friend FPS operator+ (const mint &lhs, FPS &rhs) {return rhs+=
    friend FPS operator-(FPS lhs, const FPS &rhs) {return lhs-=
    friend FPS operator-(FPS lhs,const mint &rhs) {return lhs-=
    friend FPS operator-(const mint &lhs, FPS rhs) {return - (rhs-
```

```
friend FPS operator* (FPS lhs, const FPS &rhs) {return lhs*=
friend FPS operator* (FPS lhs, const mint &rhs) {return lhs*=
friend FPS operator* (const mint &lhs, FPS rhs) {return rhs*=
     lhs; }
FPS operator-(){return (*this)*-1;}
FPS rev(){
    FPS res(*this);
    reverse(res.beign(),res.end());
    return res:
FPS pre(int sz) {
    FPS res(this->begin(),this->begin()+min((int)this->size
    if(res.size() < sz) res.resize(sz);</pre>
    return res;
FPS shrink(){
    FPS res(*this);
    while(!res.empty()&&res.back()==mint{})res.pop_back();
    return res;
FPS operator>>(int sz){
    if (this->size() <=sz) return { };</pre>
    FPS res(*this);
    res.erase(res.begin(), res.begin()+sz);
    return res;
FPS operator<<(int sz) {</pre>
    FPS res(*this);
    res.insert(res.begin(),sz,mint{});
    return res:
FPS diff() {
    const int n=this->size();
    FPS res(max(0,n-1));
    for(int i=1; i < n; i++) res[i-1] = (*this)[i] *mint(i);</pre>
    return res;
FPS integral(){
    const int n=this->size();
    FPS res(n+1);
    res[0]=0;
    if(n>0)res[1]=1;
    11 mod=mint::get_mod();
    for(int i=2;i<=n;i++)res[i]=(-res[mod%i])*(mod/i);</pre>
    for (int i=0; i<n; i++) res[i+1] *=(*this)[i];</pre>
    return res;
mint eval(const mint &x) {
    mint res=0,w=1;
    for(auto &a:*this) res+=a*w, w*=x;
    return res:
FPS inv(int deg=-1) {
    assert (!this->empty() && (*this)[0]!=mint(0));
    if (deg==-1) deg=this->size();
    FPS res{mint(1)/(*this)[0]};
    for(int i=2;i>>1<deg;i<<=1){</pre>
        res=(res*(mint(2)-res*pre(i))).pre(i);
    return res.pre(deg);
FPS log(int deg=-1){
    assert(!this->empty()&&(*this)[0] ==mint(1));
```

53856e, 15 lines

### FFT NTT Manacher SuffixArray

```
if (deg==-1) deg=this->size();
         return (pre(deg).diff()*inv(deg)).pre(deg-1).integral()
    FPS exp(int deg=-1) {
         assert (this->empty() | | (*this) [0] == mint(0));
         if (deg==-1) deg=this->size();
        FPS res{mint(1)};
         for(int i=2;i>>1<deg;i<<=1){</pre>
             res=(res*(pre(i)-res.log(i)+mint(1))).pre(i);
         return res.pre(deg);
    FPS pow(ll k, int deg=-1) {
         const int n=this->size();
         if (deg==-1) deg=n;
        if(k==0){
             FPS res(deg);
             if (deg) res[0]=mint(1);
             return res;
         for(int i=0;i<n;i++) {</pre>
             if( int128 t(i)*k>=deg)return FPS(deg,mint(0));
             if ((*this)[i] == mint(0)) continue;
             mint rev=mint(1)/(*this)[i];
             FPS res=(((*this*rev)>>i).log(deg)*k).exp(deg);
             res=((res*binpow((*this)[i],k))<<(i*k)).pre(deg);
             return res;
         return FPS (deg, mint (0));
using FPS=FormalPowerSeries<mint>;
FFT.hpp
Description: Fast Fourier transform
Time: \mathcal{O}(N \log N)
"../template/Header.hpp"
                                                            5d476b, 73 lines
template < class T=11, int mod=0>
struct FFT{
  using vt = vector<T>;
  using cd = complex<db>;
  using vc = vector<cd> ;
  static const bool INT=true;
  static void fft (vc &a) {
    int n=a.size(),L=31-__builtin_clz(n);
    vc rt(n);
    rt[1]=1;
    for (int k=2; k < n; k \neq =2) {
      cd z=polar(db(1),PI/k);
      for(int i=k;i<2*k;i++)rt[i]=i&1?rt[i/2]*z:rt[i/2];</pre>
    vi rev(n);
    for (int i=1; i<n; i++) rev[i] = (rev[i/2] | (i&1) <<L) /2;</pre>
    for (int i=1; i<n; i++) if (i<rev[i]) swap(a[i], a[rev[i]]);</pre>
     \textbf{for}(\textbf{int} \ k=1; k< n; k \star =2) \, \textbf{for}(\textbf{int} \ i=0; i< n; i+=2 \star k) \, \textbf{for}(\textbf{int} \ j=0; j< k) 
         ; j++) {
      cd z=rt[j+k]*a[i+j+k];
      a[i+j+k]=a[i+j]-z;
      a[i+j] += z;
  template<class U>
  static db norm(const U &x) {
    return INT?round(x):x;
  static vt conv(const vt &a,const vt &b) {
```

```
if(a.empty()||b.empty())return {};
    vt res(a.size()+b.size()-1);
    int L=32-__builtin_clz(res.size()), n=1<<L;</pre>
    vc in(n), out(n);
    copy(a.begin(),a.end(),in.begin());
    for (int i=0; i < b.size(); i++) in[i].imag(b[i]);</pre>
    fft(in);
    for(auto &x:in)x*=x;
    for (int i=0; i < n; i++) out [i] = in [-i&(n-1)] - conj(in[i]);</pre>
    fft(out);
    for(int i=0;i<res.size();i++)res[i]=norm(imag(out[i])/(4*n)</pre>
        );
    return res;
  static vl convMod(const vl &a,const vl &b) {
    assert (mod>0);
    if(a.empty()||b.empty())return {};
    vl res(a.size()+b.size()-1);
    int L=32-__builtin_clz(res.size()), n=1<<L;</pre>
    11 cut=int(sqrt(mod));
    vc in1(n),in2(n),out1(n),out2(n);
    for(int i=0;i<a.size();i++)in1[i]=cd(l1(a[i])/cut,l1(a[i])%</pre>
         cut); // a1 + i * a2
    for(int i=0;i<b.size();i++)in2[i]=cd(ll(b[i])/cut,ll(b[i])%</pre>
         cut); // b1 + i * b2
    fft(in1),fft(in2);
    for (int i=0; i<n; i++) {</pre>
      int j=-i&(n-1);
      out1[j]=(in1[i]+conj(in1[j]))*in2[i]/(2.1*n); // f1 * (g1)
            + i * q2) = f1 * q1 + i f1 * q2
      out2[j]=(in1[i]-conj(in1[j]))*in2[i]/cd(0.1,2.1*n); //f2
             * (q1 + i * q2) = f2 * q1 + i f2 * q2
    fft(out1),fft(out2);
    for(int i=0;i<res.size();i++){</pre>
      11 x=round(real(out1[i])), y=round(imag(out1[i]))+round(
           real(out2[i])), z=round(imag(out2[i]));
      res[i]=((xmod*cut+y)mod*cut+z)mod; // a1*b1*cut^2
           + (a1 * b2 + a2 * b1) * cut + a2 * b2
    return res;
 vt operator()(const vt &a,const vt &b){
    return mod>0?conv(a,b):convMod(a,b);
};
template<>
struct FFT<db>{
 static const bool INT=false;
NTT.hpp
Description: Number theoretic transform
Time: \mathcal{O}(N \log N)
"../template/Header.hpp", "../modular-arithmetic/BinPow.hpp",
"../modular-arithmetic/MontgomeryModInt.hpp"
                                                        2b2392, 39 lines
template < class mint = mint >
struct NTT{
 using vm = vector<mint>;
  static constexpr mint root=mint::get_root();
    static_assert(root!=0);
  static void ntt(vm &a) {
    int n=a.size(),L=31-__builtin_clz(n);
    vm rt(n);
    rt[1]=1;
    for (int k=2, s=2; k < n; k *=2, s++) {
      mint z[]={1,binpow(root,MOD>>s)};
```

```
for (int i=k; i<2*k; i++) rt[i]=rt[i/2]*z[i&1];</pre>
    vi rev(n);
    for (int i=1; i<n; i++) rev[i] = (rev[i/2] | (i&1) <<L) /2;</pre>
    for (int i=1; i < n; i++) if (i < rev[i]) swap (a[i], a[rev[i]]);</pre>
    for (int k=1; k < n; k *=2) for (int i=0; i < n; i+=2*k) for (int j=0; j < k
      mint z=rt[j+k]*a[i+j+k];
      a[i+j+k]=a[i+j]-z;
      a[i+j]+=z;
  static vm conv(const vm &a,const vm &b) {
    if(a.empty()||b.empty())return {};
    int s=a.size()+b.size()-1, n=1<<(32-__builtin_clz(s));</pre>
    mint inv=mint(n).inv();
    vm in1(a),in2(b),out(n);
    in1.resize(n),in2.resize(n);
    ntt(in1),ntt(in2);
    for (int i=0; i<n; i++) out [-i&(n-1)]=in1[i]*in2[i]*inv;</pre>
    ntt(out);
    return vm(out.begin(),out.begin()+s);
  vm operator()(const vm &a,const vm &b){
    return conv(a,b);
};
```

## Strings (11)

"../template/Header.hpp"

#### | Manacher.hpp

**Description:** Manacher's Algorithm. pal[i] := the length of the longest palindrome centered at i/2.

```
template < class STR>
vector < int > manacher (const STR &s) {
    int n = (int) s. size();
    if (n==0) return {};
    vector < int > pal (2*n-1);
    for (int p=0, l=-1, r=-1; p<2*n-1; p++) {
        int i = (p+1) >> 1, j=p>>1;
        int k = (i>=r?0:min (r-i, pal [2*(l+r)-p]));
        while (j+k+l < n&&i-k-1>=0&&s[j+k+1] ==s[i-k-1])k++;
        pal [p] = k;
        if (j+k>r) l=i-k, r=j+k;
    }
    for (int i=0; i<2*n-1; i++) pal[i] = pal[i] <<1| (i&1^1);
    return pal;
}</pre>
```

#### SuffixArray.hpp

Description: Suffix Automaton.

```
"../data-structure/SparseTable.hpp", "../group/monoid/Min.hpp"
                                                        b9cfb1, 39 lines
template<class STR>
struct SuffixArray{
    int n:
    vector<int> sa,isa,lcp;
    SparseTable<MinMonoid<int>> st;
    SuffixArray(){}
    SuffixArray(const STR &s){init(s);}
    void init(const STR &s) {
        n=(int)s.size();
        sa=isa=lcp=vector<int>(n+1);
        sa[0]=n;
        iota(sa.begin()+1, sa.end(), 0);
        sort(sa.begin()+1, sa.end(), [&] (int i, int j) {return s[i
              ]<s[j];});
```

```
for (int i=1; i<=n; i++) {</pre>
             int x=sa[i-1], y=sa[i];
             isa[y]=i>1&&s[x]==s[y]?isa[x]:i;
        for(int len=1;len<=n;len<<=1) {</pre>
             vector<int> ps(sa),pi(isa),pos(n+1);
             iota(pos.begin(),pos.end(),0);
             for (auto i:ps) if ((i-=len)>=0) sa[pos[isa[i]]++]=i;
             for (int i=1; i<=n; i++) {</pre>
                 int x=sa[i-1], y=sa[i];
                 isa[y]=pi[x]==pi[y]\&pi[x+len]==pi[y+len]?isa[x
         for (int i=0, k=0; i<n; i++) {</pre>
             for(int j=sa[isa[i]-1]; j+k<n&&s[j+k]==s[i+k]; k++);</pre>
             lcp[isa[i]]=k;
             if(k)k--;
        st.init(lcp);
    int get_lcp(int i,int j){
        if (i== j) return n-i;
        auto [l,r]=minmax(isa[i],isa[j]);
        return st.query(1+1,r);
};
```

### ZAlgo.hp

**Description:** Z Algorithm. z[i] := the length of the longest common prefix between s and <math>s[i:].

"../template/Header.hpp" b93726, 12

```
template < class STR>
vector < int > z_algorithm (const STR &s) {
    int n = (int) s.size();
    vector < int > z(n);
    z[0] = n;
    for (int i = 1, 1 = 0, r = 1; i < n; i + +) {
        if (i < r) z[i] = min (r - i, z[i - 1]);
        while (i + z[i] < n & & s[z[i]] = s[i + z[i]]) z[i] + +;
        if (i + z[i] > r) l = i, r = i + z[i];
    }
    return z;
}
```

### PrefixFunction.hpp

**Description:** Prefix function. pi[i] := the length of the longest proper prefix of <math>s[0:i] which is also a suffix of s[0:i].

```
template < class STR>
vector < int > prefix_function(const STR &s) {
    int n= (int) s. size();
    vector < int > pi(n);
    for(int i=1, j=0; icn; i++) {
        while (j>0&&s[i]!=s[j]) j=pi[j-1];
        if(s[i]==s[j]) j++;
        pi[i]=j;
    }
    return pi;
}
```

### Suffix Automaton.hpp

Description: Suffix Automaton.

Find whether a string t is a substring of a string s by traversing the automaton.

Find whether a string t is a suffix of a string s by checking whether the last node is a terminal node.

Find the number of distinct substrings of a string s by calculating the number of distinc path using DP.

Count the number of occurrences of string t in string s. Let p be the node we end up at after traversing t in the automaton. The answer is the number of paths from p to terminal nodes.

Find first occurrence of string t in string s by calculating the longest path in the automaton after reaching node p.

```
a50940, 49 lines
template<class STR>
struct SuffixAutomaton{
    using T = typename STR::value type;
    struct Node {
        map<T,int> nxt;
        int link,len;
        Node (int link, int len): link(link), len(len) {}
   };
   vector<Node> nodes;
    int last:
    SuffixAutomaton():nodes{Node(-1,0)},last(0){}
    SuffixAutomaton(const STR &s):SuffixAutomaton() {
        for(auto c:s)extend(c);
    int new_node(int link,int len){
        nodes.emplace_back(Node(link,len));
        return (int) nodes.size()-1;
    void extend(T c){
        int cur=new_node(0, nodes[last].len+1);
        int p=last;
        while (p!=-1&&!nodes[p].nxt.count(c)) {
            nodes[p].nxt[c]=cur;
            p=nodes[p].link;
        if(p!=-1){
            int q=nodes[p].nxt[c];
            if (nodes[p].len+1==nodes[q].len) {
                nodes[cur].link=q;
            }else{
                int r=new_node(nodes[q].link,nodes[p].len+1);
                nodes[r].nxt=nodes[q].nxt;
                while (p! = -1 \& & nodes [p] . nxt [c] = = q) {
                     nodes[p].nxt[c]=r;
                     p=nodes[p].link;
                nodes[q].link=nodes[cur].link=r;
        last=cur;
    11 distinct_substrings() {
        11 res=0;
        for(int i=1;i<(int)nodes.size();i++){</pre>
            res+=nodes[i].len-nodes[nodes[i].link].len;
        return res;
};
```

# Geometry (12)

### 12.1 Geometric primitives

#### Point.h

**Description:** Class to handle points in the plane. T can be e.g. double or long long. (Avoid int.)

```
template <class T> int sgn(T x) \{ return (x > 0) - (x < 0); \}
template<class T>
struct Point {
  typedef Point P;
  T x, y;
  explicit Point(T x=0, T y=0) : x(x), y(y) {}
  bool operator<(P p) const { return tie(x,y) < tie(p.x,p.y); }</pre>
  bool operator==(P p) const { return tie(x,y)==tie(p.x,p.y); }
  P operator+(P p) const { return P(x+p.x, y+p.y); }
  P operator-(P p) const { return P(x-p.x, y-p.y); }
  P operator*(T d) const { return P(x*d, y*d); }
  P operator/(T d) const { return P(x/d, y/d); }
  T dot(P p) const { return x*p.x + y*p.y; }
  T cross(P p) const { return x*p.y - y*p.x; }
  T cross(P a, P b) const { return (a-*this).cross(b-*this); }
  T dist2() const { return x*x + y*y; }
  double dist() const { return sqrt((double)dist2()); }
  // angle to x-axis in interval [-pi, pi]
  double angle() const { return atan2(y, x); }
  P unit() const { return *this/dist(); } // makes dist()=1
  P perp() const { return P(-y, x); } // rotates +90 degrees
  P normal() const { return perp().unit(); }
  // returns point rotated 'a' radians ccw around the origin
  P rotate (double a) const {
    return P(x*cos(a)-y*sin(a),x*sin(a)+y*cos(a)); }
  friend ostream& operator<<(ostream& os, P p) {</pre>
    return os << "(" << p.x << "," << p.v << ")"; }
```

#### lineDistance.h

#### Description:

Returns the signed distance between point p and the line containing points a and b. Positive value on left side and negative on right as seen from a towards b. a==b gives nan. P is supposed to be Point<T> or Point3D<T> where T is e.g. double or long long. It uses products in intermediate steps so watch out for overflow if using int or long long. Using Point3D will always give a non-negative distance. For Point3D, call .dist on the result of the cross product.

"Point.h"



f6bf6b, 4 lines

```
template<class P>
double lineDist(const P& a, const P& b, const P& p) {
  return (double) (b-a).cross(p-a)/(b-a).dist();
}
```

### Segment Distance.h

#### Description:

Returns the shortest distance between point p and the line segment from point s to e.

```
Usage: Point \langle double\rangle a, b(2,2), p(1,1); bool onSegment = segDist(a,b,p) \langle 1e-10;
```

84d6d3, 11 lines

### SegmentIntersection.h

#### Description:

If a unique intersection point between the line segments going from s1 to e1 and from s2 to e2 exists then it is returned. If no intersection point exists an empty vector is returned. If infinitely many exist a vector with 2 elements is returned, containing the endpoints of the common line segment. The wrong position will be returned if P is Point<ll> and the intersection point does not have integer coordinates. Products of three coordinates are used in intermediate steps so watch out for overflow if using int or long long.



```
Usage: vector<P> inter = segInter(s1,e1,s2,e2);
if (sz(inter) == 1)
cout << "segments intersect at " << inter[0] << endl;</pre>
"Point.h", "OnSegment.h"
                                                         9d57f2, 13 lines
```

```
template < class P > vector < P > segInter (P a, P b, P c, P d) {
 auto oa = c.cross(d, a), ob = c.cross(d, b),
       oc = a.cross(b, c), od = a.cross(b, d);
  // Checks if intersection is single non-endpoint point.
  if (sgn(oa) * sgn(ob) < 0 && sgn(oc) * sgn(od) < 0)
   return { (a * ob - b * oa) / (ob - oa) };
  set<P> s:
  if (onSegment(c, d, a)) s.insert(a);
  if (onSegment(c, d, b)) s.insert(b);
  if (onSegment(a, b, c)) s.insert(c);
 if (onSegment(a, b, d)) s.insert(d);
  return {all(s)};
```

#### lineIntersection.h

### Description:

If a unique intersection point of the lines going through s1,e1 and s2,e2 exists {1, point} is returned. If no intersection point exists  $\{0, (0,0)\}$  is returned and if infinitely many exists  $\{-1, e^2\}$ (0,0)} is returned. The wrong position will be returned if P is Point<|ll> and the intersection point does not have integer coordinates. Products of three coordinates are used in inter- s1 mediate steps so watch out for overflow if using int or ll. Usage: auto res = lineInter(s1,e1,s2,e2);



```
if (res.first == 1)
cout << "intersection point at " << res.second << endl;</pre>
template<class P>
```

```
pair<int, P> lineInter(P s1, P e1, P s2, P e2) {
  auto d = (e1 - s1).cross(e2 - s2);
 if (d == 0) // if parallel
   return {-(s1.cross(e1, s2) == 0), P(0, 0)};
  auto p = s2.cross(e1, e2), q = s2.cross(e2, s1);
  return {1, (s1 * p + e1 * q) / d};
```

#### sideOf.h

**Description:** Returns where p is as seen from s towards e.  $1/0/-1 \Leftrightarrow \text{left/on}$ line/right. If the optional argument eps is given 0 is returned if p is within distance eps from the line. P is supposed to be Point<T> where T is e.g. double or long long. It uses products in intermediate steps so watch out for overflow if using int or long long.

```
Usage: bool left = sideOf(p1,p2,q) ==1;
```

```
"Point.h"
                                                       3af81c, 9 lines
template<class P>
int sideOf(P s, P e, P p) { return sgn(s.cross(e, p)); }
template<class P>
int sideOf(const P& s, const P& e, const P& p, double eps) {
  auto a = (e-s).cross(p-s);
  double 1 = (e-s).dist()*eps;
  return (a > 1) - (a < -1);
```

### OnSegment.h

**Description:** Returns true iff p lies on the line segment from s to e. Use (segDist(s,e,p) <=epsilon) instead when using Point < double >.

```
"Point.h"
template < class P > bool on Segment (P s, P e, P p) {
 return p.cross(s, e) == 0 \&\& (s - p).dot(e - p) <= 0;
```

#### linearTransformation.h Description:

Apply the linear transformation (translation, rotation and scaling) which takes line p0-p1 to line q0-q1 to point r.



```
typedef Point<double> P;
P linearTransformation(const P& p0, const P& p1,
    const P& q0, const P& q1, const P& r) {
 P dp = p1-p0, dq = q1-q0, num(dp.cross(dq), dp.dot(dq));
 return q0 + P((r-p0).cross(num), (r-p0).dot(num))/dp.dist2();
```

#### Angle.h

"Point.h"

Description: A class for ordering angles (as represented by int points and a number of rotations around the origin). Useful for rotational sweeping. Sometimes also represents points or vectors.

Usage: vector $\langle Angle \rangle$  v = {w[0], w[0].t360() ...}; // sorted int j = 0; rep(i,0,n) { while (v[j] < v[i].t180()) ++j; } // sweeps j such that (j-i) represents the number of positively oriented triangles with vertices at 0 and i 0f0602, 35 lines

```
struct Angle {
 int x, y;
  Angle(int x, int y, int t=0) : x(x), y(y), t(t) {}
  Angle operator-(Angle b) const { return {x-b.x, y-b.y, t}; }
  int half() const {
    assert(x || y);
    return v < 0 || (v == 0 && x < 0);
  Angle t90() const { return \{-y, x, t + (half() \&\& x >= 0)\}; \}
  Angle t180() const { return {-x, -y, t + half()}; }
 Angle t360() const { return {x, y, t + 1}; }
bool operator<(Angle a, Angle b) {</pre>
  // add a.dist2() and b.dist2() to also compare distances
  return make_tuple(a.t, a.half(), a.y * (ll)b.x) <</pre>
         make_tuple(b.t, b.half(), a.x * (11)b.y);
  Given two points, this calculates the smallest angle between
  them, i.e., the angle that covers the defined line segment.
pair<Angle, Angle> segmentAngles(Angle a, Angle b) {
  if (b < a) swap(a, b);
  return (b < a.t180() ?
          make_pair(a, b) : make_pair(b, a.t360()));
```

Angle operator+(Angle a, Angle b) { // point a + vector b

Angle angleDiff(Angle a, Angle b) {  $// angle b - angle a}$ 

return  $\{a.x*b.x + a.y*b.y, a.x*b.y - a.y*b.x, tu - (b < a)\};$ 

Angle r(a.x + b.x, a.y + b.y, a.t);

return r.t180() < a ? r.t360() : r;</pre>

int tu = b.t - a.t; a.t = b.t;

**if** (a.t180() < r) r.t--;

### 12.2 Circles

### CircleIntersection.h.

**Description:** Computes the pair of points at which two circles intersect. Returns false in case of no intersection. "Point.h"

```
typedef Point<double> P;
bool circleInter(P a, P b, double r1, double r2, pair < P, P >* out) {
  if (a == b) { assert(r1 != r2); return false; }
  P \text{ vec} = b - a;
  double d2 = vec.dist2(), sum = r1+r2, dif = r1-r2,
          p = (d2 + r1*r1 - r2*r2)/(d2*2), h2 = r1*r1 - p*p*d2;
  if (sum*sum < d2 || dif*dif > d2) return false;
  P \text{ mid} = a + \text{vec*p, per} = \text{vec.perp()} * \text{sqrt(fmax(0, h2) / d2);}
  *out = {mid + per, mid - per};
  return true;
```

### CircleTangents.h

Description: Finds the external tangents of two circles, or internal if r2 is negated. Can return 0, 1, or 2 tangents - 0 if one circle contains the other (or overlaps it, in the internal case, or if the circles are the same); 1 if the circles are tangent to each other (in which case .first = .second and the tangent line is perpendicular to the line between the centers). .first and .second give the tangency points at circle 1 and 2 respectively. To find the tangents of a circle with a point set r2 to 0. "Point.h"

```
template<class P>
vector<pair<P, P>> tangents(P c1, double r1, P c2, double r2) {
 P d = c2 - c1;
  double dr = r1 - r2, d2 = d.dist2(), h2 = d2 - dr * dr;
  if (d2 == 0 || h2 < 0) return {};</pre>
  vector<pair<P, P>> out;
  for (double sign : {-1, 1}) {
    P v = (d * dr + d.perp() * sqrt(h2) * sign) / d2;
    out.push_back(\{c1 + v * r1, c2 + v * r2\});
 if (h2 == 0) out.pop back();
 return out;
```

### CirclePolygonIntersection.h

Description: Returns the area of the intersection of a circle with a ccw polygon.

Time:  $\mathcal{O}(n)$ 

alee63, 19 lines "../../content/geometry/Point.h" typedef Point<double> P; #define arg(p, q) atan2(p.cross(q), p.dot(q)) double circlePoly(P c, double r, vector<P> ps) { **auto** tri = [&] (P p, P q) { **auto** r2 = r \* r / 2; P d = q - p;auto a = d.dot(p)/d.dist2(), b = (p.dist2()-r\*r)/d.dist2(); **auto** det = a \* a - b; if (det <= 0) return arg(p, q) \* r2;</pre> **auto** s = max(0., -a-sqrt(det)), t = min(1., -a+sqrt(det));if (t < 0 || 1 <= s) return arg(p, q) \* r2;</pre> P u = p + d \* s, v = p + d \* t;**return** arg(p,u) \* r2 + u.cross(v)/2 + arg(v,q) \* r2; **auto** sum = 0.0; rep(i, 0, sz(ps))sum += tri(ps[i] - c, ps[(i + 1) % sz(ps)] - c);return sum;

#### circumcircle.h

### Description:

The circumcirle of a triangle is the circle intersecting all three vertices. ccRadius returns the radius of the circle going through points A, B and C and ccCenter returns the center of the same circle.



"Point.h" 1caa3a, 9 lines

```
typedef Point<double> P;
double ccRadius(const P& A, const P& B, const P& C) {
  return (B-A).dist()*(C-B).dist()*(A-C).dist()/
    abs((B-A).cross(C-A))/2;
}
P ccCenter(const P& A, const P& B, const P& C) {
  P b = C-A, c = B-A;
  return A + (b*c.dist2()-c*b.dist2()).perp()/b.cross(c)/2;
}
```

### MinimumEnclosingCircle.h

**Description:** Computes the minimum circle that encloses a set of points. **Time:** expected  $\mathcal{O}(n)$ 

```
"circumcircle.h"
                                                     09dd0a, 17 lines
pair<P, double> mec(vector<P> ps) {
  shuffle(all(ps), mt19937(time(0)));
 P \circ = ps[0];
  double r = 0, EPS = 1 + 1e-8;
  rep(i,0,sz(ps)) if ((o - ps[i]).dist() > r * EPS) {
   o = ps[i], r = 0;
   rep(j, 0, i) if ((o - ps[j]).dist() > r * EPS) {
     o = (ps[i] + ps[j]) / 2;
     r = (o - ps[i]).dist();
     rep(k,0,j) if ((o - ps[k]).dist() > r * EPS) {
       o = ccCenter(ps[i], ps[j], ps[k]);
       r = (o - ps[i]).dist();
   }
  return {o, r};
```

## 12.3 Polygons

### InsidePolygon.h

**Description:** Returns true if p lies within the polygon. If strict is true, it returns false for points on the boundary. The algorithm uses products in intermediate steps so watch out for overflow.

```
Usage: vector\langle P \rangle v = {P{4,4}, P{1,2}, P{2,1}};
bool in = inPolygon(v, P{3, 3}, false);
Time: \mathcal{O}(n)
```

"Point.h", "OnSegment.h", "SegmentDistance.h" 2bf504, 11 lines

```
template < class P >
bool inPolygon(vector < P > & p, P a, bool strict = true) {
   int cnt = 0, n = sz(p);
   rep(i,0,n) {
      P q = p[(i + 1) % n];
      if (onSegment(p[i], q, a)) return !strict;
      //or: if (segDist(p[i], q, a) <= eps) return !strict;
   cnt ^= ((a.y < p[i].y) - (a.y < q.y)) * a.cross(p[i], q) > 0;
   }
   return cnt;
}
```

#### PolygonArea.h

**Description:** Returns twice the signed area of a polygon. Clockwise enumeration gives negative area. Watch out for overflow if using int as T!

```
"Point.h" f12300, 6 lines
```

```
template<class T>
T polygonArea2(vector<Point<T>>& v) {
```

```
T = v.back().cross(v[0]);
  rep(i,0,sz(v)-1) a += v[i].cross(v[i+1]);
 return a;
PolygonCenter.h
Description: Returns the center of mass for a polygon.
Time: \mathcal{O}(n)
"Point.h"
                                                        9706dc, 9 lines
typedef Point<double> P;
P polygonCenter(const vector<P>& v) {
 P res(0, 0); double A = 0;
 for (int i = 0, j = sz(v) - 1; i < sz(v); j = i++) {
    res = res + (v[i] + v[j]) * v[j].cross(v[i]);
    A += v[j].cross(v[i]);
 return res / A / 3;
PolygonCut.h
Description:
Returns a vector with the vertices of a polygon with every-
thing to the left of the line going from s to e cut away.
Usage: vector<P> p = ...;
p = polygonCut(p, P(0,0), P(1,0));
"Point.h", "lineIntersection.h"
                                                        f2b7d4, 13 lines
typedef Point < double > P;
vector<P> polygonCut(const vector<P>& poly, P s, P e) {
 vector<P> res;
 rep(i, 0, sz(poly)) {
    P cur = poly[i], prev = i ? poly[i-1] : poly.back();
    bool side = s.cross(e, cur) < 0;</pre>
    if (side != (s.cross(e, prev) < 0))</pre>
     res.push_back(lineInter(s, e, cur, prev).second);
    if (side)
      res.push_back(cur);
```

### ConvexHull.h

return res:

#### Description:

"Point.h"

Returns a vector of the points of the convex hull in counterclockwise order. Points on the edge of the hull between two other points are not considered part of the hull.  $\mathbf{Time} : \mathcal{O}(n \log n)$ 

countereen two

typedef Point<11> P;
vector<P> convexHull(vector<P> pts) {
 if (sz(pts) <= 1) return pts;
 sort(all(pts));
 vector<P> h(sz(pts)+1);
 int s = 0, t = 0;
 for (int it = 2; it--; s = --t, reverse(all(pts)))
 for (P p: pts) {
 while (t >= s + 2 && h[t-2].cross(h[t-1], p) <= 0) t--;
 h[t++] = p;
 }
 return {h.begin(), h.begin() + t - (t == 2 && h[0] == h[1])};</pre>

#### HullDiameter.h

**Description:** Returns the two points with max distance on a convex hull (ccw, no duplicate/collinear points).

#### Time: $\mathcal{O}\left(n\right)$

```
"Point.h" c571b8, 12 lines
```

```
typedef Point<1l> P;
```

```
array<P, 2> hullDiameter(vector<P> S) {
  int n = sz(S), j = n < 2 ? 0 : 1;
  pair<11, array<P, 2>> res({0, {S[0], S[0]}});
  rep(i,0,j)
  for (;; j = (j + 1) % n) {
    res = max(res, {(S[i] - S[j]).dist2(), {S[i], S[j]}});
    if ((S[(j + 1) % n] - S[j]).cross(S[i + 1] - S[i]) >= 0)
      break;
  }
  return res.second;
}
```

#### PointInsideHull.h

**Description:** Determine whether a point t lies inside a convex hull (CCW order, with no collinear points). Returns true if point lies within the hull. If strict is true, points on the boundary aren't included.

### Time: $\mathcal{O}\left(\log N\right)$

#### LineHullIntersection.h

return sqn(l[a].cross(l[b], p)) < r;</pre>

**Description:** Line-convex polygon intersection. The polygon must be ccw and have no collinear points. lineHull(line, poly) returns a pair describing the intersection of a line with the polygon:  $\bullet$  (-1,-1) if no collision,  $\bullet$  (i,-1) if touching the corner i,  $\bullet$  (i,i) if along side (i,i+1),  $\bullet$  (i,j) if crossing sides (i,i+1) and (j,j+1). In the last case, if a corner i is crossed, this is treated as happening on side (i,i+1). The points are returned in the same order as the line hits the polygon. extrVertex returns the point of a hull with the max projection onto a line.

#### Time: $\mathcal{O}(\log n)$

```
"Point.h"
              #define cmp(i, j) sqn(dir.perp().cross(poly[(i)%n]-poly[(j)%n]))
              #define extr(i) cmp(i + 1, i) >= 0 && cmp(i, i - 1 + n) < 0
              template <class P> int extrVertex(vector<P>& poly, P dir) {
310954, 13 lines
               int n = sz(poly), lo = 0, hi = n;
               if (extr(0)) return 0;
                while (lo + 1 < hi) {
                  int m = (lo + hi) / 2;
                  if (extr(m)) return m;
                  int ls = cmp(lo + 1, lo), ms = cmp(m + 1, m);
                  (1s < ms \mid | (1s == ms \&\& 1s == cmp(1o, m)) ? hi : 1o) = m;
               return lo;
              #define cmpL(i) sqn(a.cross(poly[i], b))
              template <class P>
              array<int, 2> lineHull(P a, P b, vector<P>& poly) {
               int endA = extrVertex(poly, (a - b).perp());
                int endB = extrVertex(poly, (b - a).perp());
               if (cmpL(endA) < 0 \mid | cmpL(endB) > 0)
                  return {-1, -1};
                array<int, 2> res;
                rep(i, 0, 2) {
```

int lo = endB, hi = endA, n = sz(poly);

### ClosestPair kdTree FastDelaunay PolyhedronVolume

```
while ((lo + 1) % n != hi) {
   int m = ((lo + hi + (lo < hi ? 0 : n)) / 2) % n;
   (cmpL(m) == cmpL(endB) ? lo : hi) = m;
 res[i] = (lo + !cmpL(hi)) % n;
 swap (endA, endB);
if (res[0] == res[1]) return {res[0], -1};
if (!cmpL(res[0]) && !cmpL(res[1]))
 switch ((res[0] - res[1] + sz(poly) + 1) % sz(poly)) {
   case 0: return {res[0], res[0]};
   case 2: return {res[1], res[1]};
return res;
```

### 12.4 Misc. Point Set Problems

#### Closest Pair h

Description: Finds the closest pair of points.

Time:  $\mathcal{O}(n \log n)$ 

```
"Point.h"
                                                      ac41a6, 17 lines
typedef Point<11> P;
pair<P, P> closest(vector<P> v) {
  assert (sz(v) > 1);
  set<P> S;
  sort(all(v), [](P a, P b) { return a.y < b.y; });</pre>
  pair<11, pair<P, P>> ret{LLONG_MAX, {P(), P()}};
  int j = 0;
  for (P p : v) {
   P d{1 + (ll)sqrt(ret.first), 0};
    while (v[j].v \le p.v - d.x) S.erase(v[j++]);
    auto lo = S.lower_bound(p - d), hi = S.upper_bound(p + d);
    for (; lo != hi; ++lo)
      ret = min(ret, \{(*lo - p).dist2(), \{*lo, p\}\});
    S.insert(p);
  return ret.second;
kdTree.h
```

**Description:** KD-tree (2d, can be extended to 3d)

bac5b0, 63 lines typedef long long T; typedef Point<T> P; const T INF = numeric\_limits<T>::max(); bool on\_x(const P& a, const P& b) { return a.x < b.x; }</pre> bool on\_y(const P& a, const P& b) { return a.y < b.y; }</pre> struct Node { P pt; // if this is a leaf, the single point in it T x0 = INF, x1 = -INF, y0 = INF, y1 = -INF; // bounds Node \*first = 0, \*second = 0; T distance (const P& p) { // min squared distance to a point T x = (p.x < x0 ? x0 : p.x > x1 ? x1 : p.x);T y = (p.y < y0 ? y0 : p.y > y1 ? y1 : p.y);return (P(x,y) - p).dist2(); Node (vector<P>&& vp) : pt(vp[0]) { for (P p : vp) { x0 = min(x0, p.x); x1 = max(x1, p.x);y0 = min(y0, p.y); y1 = max(y1, p.y);**if** (vp.size() > 1) { // split on x if width  $>= h \, eight \, (not \, ideal...)$  $sort(all(vp), x1 - x0 >= y1 - y0 ? on_x : on_y);$ 

```
// divide by taking half the array for each child (not
      // best performance with many duplicates in the middle)
      int half = sz(vp)/2;
      first = new Node({vp.begin(), vp.begin() + half});
      second = new Node({vp.begin() + half, vp.end()});
 }
};
struct KDTree {
  Node* root:
  KDTree(const vector<P>& vp) : root(new Node({all(vp)})) {}
 pair<T, P> search(Node *node, const P& p) {
    if (!node->first) {
      // uncomment if we should not find the point itself:
      // if (p = node \rightarrow pt) return \{INF, P()\};
      return make_pair((p - node->pt).dist2(), node->pt);
    Node *f = node->first, *s = node->second;
    T bfirst = f->distance(p), bsec = s->distance(p);
    if (bfirst > bsec) swap(bsec, bfirst), swap(f, s);
    // search closest side first, other side if needed
    auto best = search(f, p);
    if (bsec < best.first)</pre>
     best = min(best, search(s, p));
    return best;
  // find nearest point to a point, and its squared distance
  // (requires an arbitrary operator for Point)
  pair<T, P> nearest (const P& p) {
    return search(root, p);
};
FastDelaunav.h
Description: Fast Delaunay triangulation. Each circumcircle contains none
of the input points. There must be no duplicate points. If all points are on a
line, no triangles will be returned. Should work for doubles as well, though
there may be precision issues in 'circ'. Returns triangles in order {t[0][0],
t[0][1], t[0][2], t[1][0], \ldots\}, all counter-clockwise.
Time: \mathcal{O}(n \log n)
"Point.h"
                                                        eefdf5, 88 lines
typedef Point<11> P;
typedef struct Ouad* O;
typedef __int128_t 111; // (can be ll if coords are < 2e4)
P arb(LLONG_MAX, LLONG_MAX); // not equal to any other point
struct Quad {
 Q rot, o; P p = arb; bool mark;
 P& F() { return r()->p; }
 O& r() { return rot->rot; }
```

Q prev() { return rot->o->rot; }

Q next() { return r()->prev(); }

Q makeEdge(P orig, P dest) {

H = r -> 0; r -> r() -> r() = r;

r->p = orig; r->F() = dest;

111 p2 = p.dist2(), A = a.dist2()-p2,

B = b.dist2()-p2, C = c.dist2()-p2;

bool circ(P p, P a, P b, P c) { // is p in the circumcircle?

Q r = H ? H : new Quad{new Quad{new Quad{new Quad{0}}}};

```
void splice(0 a, 0 b) {
                                                                    swap(a->o->rot->o, b->o->rot->o); swap(a->o, b->o);
                                                                  Q connect(Q a, Q b) {
                                                                    Q = makeEdge(a->F(), b->p);
                                                                    splice(q, a->next());
                                                                    splice(q->r(), b);
                                                                    return q;
                                                                  pair<Q,Q> rec(const vector<P>& s) {
                                                                    if (sz(s) <= 3) {
                                                                       Q a = makeEdge(s[0], s[1]), b = makeEdge(s[1], s.back());
                                                                       if (sz(s) == 2) return { a, a->r() };
                                                                       splice(a->r(), b);
                                                                      auto side = s[0].cross(s[1], s[2]);
                                                                      Q c = side ? connect(b, a) : 0;
                                                                       return {side < 0 ? c->r() : a, side < 0 ? c : b->r() };
                                                                  #define H(e) e->F(), e->p
                                                                  #define valid(e) (e->F().cross(H(base)) > 0)
                                                                    Q A, B, ra, rb;
                                                                    int half = sz(s) / 2;
                                                                    tie(ra, A) = rec({all(s) - half});
                                                                    tie(B, rb) = rec({sz(s) - half + all(s)});
                                                                    while ((B->p.cross(H(A)) < 0 && (A = A->next())) | |
                                                                            (A->p.cross(H(B)) > 0 && (B = B->r()->o)));
                                                                    O base = connect(B->r(), A);
                                                                    if (A->p == ra->p) ra = base->r();
                                                                    if (B->p == rb->p) rb = base;
                                                                  #define DEL(e, init, dir) Q e = init->dir; if (valid(e)) \
                                                                       while (circ(e->dir->F(), H(base), e->F())) { \
                                                                        Q t = e->dir; \setminus
                                                                        splice(e, e->prev()); \
                                                                        splice(e->r(), e->r()->prev()); \
                                                                         e->o = H; H = e; e = t; \setminus
                                                                    for (;;) {
                                                                      DEL(LC, base->r(), o); DEL(RC, base, prev());
                                                                       if (!valid(LC) && !valid(RC)) break;
                                                                      if (!valid(LC) || (valid(RC) && circ(H(RC), H(LC))))
                                                                        base = connect(RC, base->r());
                                                                        base = connect(base->r(), LC->r());
                                                                    return { ra, rb };
                                                                  vector<P> triangulate(vector<P> pts) {
                                                                    sort(all(pts)); assert(unique(all(pts)) == pts.end());
                                                                    if (sz(pts) < 2) return {};
                                                                    Q e = rec(pts).first;
                                                                    vector<Q> q = \{e\};
                                                                    int qi = 0;
                                                                    while (e->o->F().cross(e->F(), e->p) < 0) e = e->o;
                                                                   #define ADD { Q c = e; do { c->mark = 1; pts.push_back(c->p); \
                                                                    q.push_back(c->r()); c = c->next(); } while (c != e); }
                                                                    ADD; pts.clear();
                                                                    while (qi < sz(q)) if (!(e = q[qi++]) \rightarrow mark) ADD;
return p.cross(a,b)*C + p.cross(b,c)*A + p.cross(c,a)*B > 0;
                                                                    return pts;
rep(i,0,4) r = r \rightarrow rot, r \rightarrow p = arb, r \rightarrow o = i & 1 ? <math>r : r \rightarrow r();
```

return r;

### $12.5 \ 3D$

#### PolyhedronVolume.h

Description: Magic formula for the volume of a polyhedron. Faces should point outwards. 305<u>8c3, 6 lines</u>

```
template<class V, class L>
double signedPolyVolume(const V& p, const L& trilist) {
  double v = 0:
  for (auto i : trilist) v += p[i.a].cross(p[i.b]).dot(p[i.c]);
 return v / 6:
```

#### Point3D.h

**Description:** Class to handle points in 3D space. T can be e.g. double or

```
template<class T> struct Point3D {
  typedef Point3D P;
  typedef const P& R;
  T x, y, z;
  explicit Point3D(T x=0, T y=0, T z=0) : x(x), y(y), z(z) {}
  bool operator<(R p) const {</pre>
    return tie(x, y, z) < tie(p.x, p.y, p.z); }
  bool operator==(R p) const {
    return tie(x, y, z) == tie(p.x, p.y, p.z); }
  P operator+(R p) const { return P(x+p.x, y+p.y, z+p.z); }
  P operator-(R p) const { return P(x-p.x, y-p.y, z-p.z); }
  P operator*(T d) const { return P(x*d, y*d, z*d); }
  P operator/(T d) const { return P(x/d, y/d, z/d); }
  T dot(R p) const { return x*p.x + y*p.y + z*p.z; }
  P cross(R p) const {
    return P(y*p.z - z*p.y, z*p.x - x*p.z, x*p.y - y*p.x);
  T dist2() const { return x*x + y*y + z*z; }
  double dist() const { return sgrt((double) dist2()); }
  //Azimuthal angle (longitude) to x-axis in interval [-pi, pi]
  double phi() const { return atan2(y, x); }
  //Zenith angle (latitude) to the z-axis in interval [0, pi]
  double theta() const { return atan2(sgrt(x*x+y*y),z); }
  P unit() const { return *this/(T)dist(); } //makes dist()=1
  //returns unit vector normal to *this and p
  P normal(P p) const { return cross(p).unit(); }
  //returns point rotated 'angle' radians ccw around axis
  P rotate (double angle, P axis) const {
    double s = sin(angle), c = cos(angle); P u = axis.unit();
    return u*dot(u)*(1-c) + (*this)*c - cross(u)*s;
};
```

**Description:** Computes all faces of the 3-dimension hull of a point set. \*No four points must be coplanar\*, or else random results will be returned. All faces will point outwards.

```
Time: \mathcal{O}\left(n^2\right)
```

"Point3D.h" 5b45fc, 49 lines

```
typedef Point3D<double> P3;
struct PR {
  void ins(int x) { (a == -1 ? a : b) = x; }
  void rem(int x) { (a == x ? a : b) = -1; }
  int cnt() { return (a != -1) + (b != -1); }
 int a, b;
struct F { P3 q; int a, b, c; };
vector<F> hull3d(const vector<P3>& A) {
  assert (sz(A) >= 4);
```

```
vector<vector<PR>> E(sz(A), vector<PR>(sz(A), {-1, -1}));
#define E(x,y) E[f.x][f.y]
 vector<F> FS;
 auto mf = [&](int i, int j, int k, int l) {
   P3 q = (A[j] - A[i]).cross((A[k] - A[i]));
   if (q.dot(A[1]) > q.dot(A[i]))
     q = q * -1;
   F f{q, i, j, k};
   E(a,b).ins(k); E(a,c).ins(j); E(b,c).ins(i);
   FS.push back(f);
 rep(i,0,4) rep(j,i+1,4) rep(k,j+1,4)
   mf(i, j, k, 6 - i - j - k);
 rep(i,4,sz(A)) {
   rep(j,0,sz(FS)) {
     F f = FS[j];
     if(f.q.dot(A[i]) > f.q.dot(A[f.a])) {
       E(a,b).rem(f.c);
       E(a,c).rem(f.b);
       E(b,c).rem(f.a);
       swap(FS[j--], FS.back());
       FS.pop back();
   int nw = sz(FS);
   rep(j,0,nw) {
     F f = FS[i];
#define C(a, b, c) if (E(a,b).cnt() != 2) mf(f.a, f.b, i, f.c);
     C(a, b, c); C(a, c, b); C(b, c, a);
 for (F& it : FS) if ((A[it.b] - A[it.a]).cross(
   A[it.c] - A[it.a]).dot(it.q) <= 0) swap(it.c, it.b);
 return FS:
```

### sphericalDistance.h

Description: Returns the shortest distance on the sphere with radius radius between the points with azimuthal angles (longitude) f1 ( $\phi_1$ ) and f2 ( $\phi_2$ ) from x axis and zenith angles (latitude) t1  $(\theta_1)$  and t2  $(\theta_2)$  from z axis (0 =north pole). All angles measured in radians. The algorithm starts by converting the spherical coordinates to cartesian coordinates so if that is what you have you can use only the two last rows. dx\*radius is then the difference between the two points in the x direction and d\*radius is the total distance between the points. 611f07, 8 lines

```
double sphericalDistance (double f1, double t1,
    double f2, double t2, double radius) {
  double dx = \sin(t2) \cdot \cos(f2) - \sin(t1) \cdot \cos(f1);
 double dy = sin(t2) * sin(f2) - sin(t1) * sin(f1);
 double dz = cos(t2) - cos(t1);
 double d = sqrt(dx*dx + dy*dy + dz*dz);
 return radius*2*asin(d/2);
```

# Dynamic Programming (13)

### DVC.hpp

**Description:** Optimize  $O(N^2K)$  to  $O(NK \log N)$ 

```
"../template/Header.hpp"
                                                           aa5ddf, 19 lines
vector<vl> cst, dp;
11 cost(int 1, int r) {
 return cst[l][r];
```

```
void divide(int 1, int r, int opt_1, int opt_r, int c) {
  if(1 > r) return ;
  int mid = (1 + r) / 2;
  pair<11, int> best = make_pair(INF, -1);
  for(int k=opt_1; k<=min(mid, opt_r); ++k) {</pre>
    best = min(best, make_pair(dp[c - 1][k] + cost(k + 1, mid),
  dp[c][mid] = best.first;
  divide(1, mid - 1, opt_1, best.second, c);
  divide(mid + 1, r, best.second, opt_r, c);
// for(int c=1; c \leq K; ++c) divide(1, N, 1, N, c);
SlopeTrick.hpp
Description: Absolute Smth
"../template/Header.hpp"
                                                      f62f9a, 36 lines
11 extending value;
struct slope trick {
  multiset<11> ms_1, ms_r;
  11 min v = 011, 1z 1 = 011, 1z r = 011;
  bool extending = false;
  void add_line(ll v) {
    if(extending) {
      lz 1 -= extending value;
      lz r -= extending value:
    extending = true;
    if(ms_l.empty() && ms_r.empty()) {
      ms l.emplace(v);
      ms_r.emplace(v);
    else if(v <= *ms l.rbegin() + lz l) {
      min_y += (*ms_l.rbegin() + lz_l) - v;
      ms_r.emplace(*ms_l.rbegin() + lz_l - lz_r);
      ms_l.erase(--ms_l.end());
      ms_l.emplace(v - lz_l);
      ms_l.emplace(v - lz_l);
    else if(v >= *ms_r.begin() + lz_r) {
      \min v += v - (*ms r.begin() + lz r);
      ms_1.emplace(*ms_r.begin() + lz_r - lz_l);
      ms_r.erase(ms_r.begin());
      ms_r.emplace(v - lz_r);
      ms_r.emplace(v - lz_r);
    else {
      ms_l.emplace(v - lz_l);
      ms_r.emplace(v - lz_r);
```

## Convolutions (14)

### And Convolution.hpp

**Description:** Bitwise AND Convolution. Superset Zeta Transform: A'[S] = $\sum_{T \supset S} A[T]$ . Superset Mobius Transform:  $A[T] = \sum_{S \supset T} (-1)^{|S-T|} A'[S]$ . Time:  $\mathcal{O}(N \log N)$ .

```
"../template/Header.hpp"
template<class T>
void superset_zeta(vector<T> &a){
    int n=(int)a.size();
    assert (n==(n\&-n));
     for (int i=1; i < n; i < <=1) {</pre>
```

```
for(int j=0; j<n; j++) {</pre>
             if (j&i) {
                 a[j^i]+=a[j];
template < class T>
void superset_mobius(vector<T> &a){
    int n=(int)a.size();
    assert (n==(n\&-n));
    for(int i=n;i>>=1;) {
         for(int j=0; j<n; j++) {</pre>
             if(j&i){
                 a[j^i] -= a[j];
template<class T>
vector<T> and convolution(vector<T> a, vector<T> b) {
    superset_zeta(a);
    superset_zeta(b);
    for (int i=0; i < (int) a.size(); i++) a[i] *=b[i];</pre>
    superset_mobius(a);
    return a;
GCDConvolution.hpp
Description: GCD Convolution. Multiple Zeta Transform: A'[n] =
\sum_{n|m} A[m]. Multiple Mobius Transform: A[n] = \sum_{n|m} \mu(m/n) A'[m].
Time: \mathcal{O}(N \log \log N).
"../template/Header.hpp"
template<class T>
void multiple_zeta(vector<T> &a) {
    int n=(int)a.size();
    vector<bool> is_prime(n,true);
    for(int p=2;p<n;p++){
         if(!is prime[p])continue;
         for (int i= (n-1)/p; i>=1; i--) {
             is_prime[i*p]=false;
             a[i] += a[i * p];
template < class T>
void multiple_mobius(vector<T> &a){
    int n=(int)a.size();
    vector<bool> is_prime(n,true);
    for (int p=2; p<n; p++) {</pre>
        if(!is_prime[p])continue;
         for (int i=1; i*p<n; i++) {</pre>
             is_prime[i*p]=false;
             a[i] -= a[i*p];
template<class T>
vector<T> gcd_convolution(vector<T> a, vector<T> b) {
    multiple_zeta(a);
    multiple_zeta(b);
    for(int i=0;i<(int)a.size();i++)a[i]*=b[i];</pre>
    multiple_mobius(a);
    return a;
```

```
LCMConvolution.hpp
Description: LCM Convolution. Divisor Zeta Transform: A'[n] =
\sum_{d|n} \hat{A}[d]. Divisor Mobius Transform: A[n] = \sum_{d|n} \mu(n/d) A'[d].
Time: \mathcal{O}(N \log \log N).
"../template/Header.hpp"
                                                            41fe9d, 34 lines
template<class T>
void divisor_zeta(vector<T> &a) {
    int n=(int)a.size();
    vector<bool> is prime(n,true);
    for(int p=2;p<n;p++) {</pre>
        if(!is_prime[p])continue;
         for (int i=1;i*p<n;i++) {</pre>
             is_prime[i*p]=false;
             a[i*p]+=a[i];
template < class T>
void divisor_mobius(vector<T> &a) {
    int n=(int)a.size();
    vector<bool> is_prime(n,true);
    for (int p=2; p<n; p++) {</pre>
        if(!is_prime[p])continue;
         for (int i=(n-1)/p;i>=1;i--) {
             is_prime[i*p]=false;
             a[i*p]-=a[i];
template < class T>
vector<T> lcm_convolution(vector<T> a, vector<T> b) {
    divisor zeta(a);
    divisor_zeta(b);
    for(int i=0;i<(int)a.size();i++)a[i]*=b[i];</pre>
    divisor_mobius(a);
    return a;
ORConvolution.hpp
Description: Bitwise OR Convolution. Subset Zeta Transform: A'[S] =
\sum_{T \subset S} A[T]. Subset Mobius Transform: A[T] = \sum_{S \subset T} (-1)^{|T-S|} A'[S].
Time: \mathcal{O}(N \log N).
"../template/Header.hpp"
                                                            c58b77, 34 lines
template<class T>
void subset_zeta(vector<T> &a) {
    int n=(int)a.size();
    assert (n==(n\&-n));
    for (int i=1; i<n; i<<=1) {</pre>
         for (int j=0; j<n; j++) {</pre>
             if(j&i){
                  a[j] += a[j^i];
template < class T>
void subset mobius(vector<T> &a){
    int n=(int)a.size();
    assert (n==(n\&-n));
    for(int i=n;i>>=1;) {
         for (int j=0; j<n; j++) {</pre>
             if(j&i){
```

```
template < class T>
vector<T> or_convolution(vector<T> a, vector<T> b) {
    subset_zeta(a);
     subset zeta(b);
    for (int i=0; i < (int) a.size(); i++) a[i] *=b[i];</pre>
    subset mobius(a);
    return a:
XORConvolution.hpp
Description: Bitwise XOR Convolution. Fast Walsh-Hadamard Transform:
A'[S] = \sum_{T} (-1)^{|S\&T|} A[T].
Time: \mathcal{O}(N \log N).
"../template/Header.hpp'
                                                           058<u>48d, 29 lines</u>
template<class T>
void fwht(vector<T> &a) {
    int n=(int)a.size();
    assert (n==(n\&-n));
    for (int i=1; i<n; i<<=1) {</pre>
         for(int j=0; j<n; j++) {
             if(j&i){
                  T &u=a[j^i],&v=a[j];
                  tie(u, v) = make_pair(u+v, u-v);
template < class T>
vector<T> xor_convolution(vector<T> a, vector<T> b) {
    int n=(int)a.size();
     fwht(a);
     fwht(b);
     for(int i=0;i<n;i++)a[i] *=b[i];</pre>
    fwht(a);
    T \text{ div}=T(1)/T(n);
    if(div==T(0)){
         for(auto &x:a) x/=n;
     }else{
         for(auto &x:a)x*=div;
     return a;
MaxPlusConvolution.hpp
Description: Max Plus Convolution. Find C[k] = \max_{i+j=k} \{A[i] + B[j]\}.
Time: \mathcal{O}(N).
                                                           7176a2, 94 lines
// SMAWCK algorithm for finding row-wise maxima.
f/f(i,j,k) checks if M[i]/[j] \le M[i]/[k].
// f(i,j,k) checks if M[i][k] is at least as good as M[i][j].
// higher is better.
template<class F>
vector<int> smawck(const F &f,const vector<int> &rows,const
     vector<int> &cols) {
    int n=(int)rows.size(), m=(int)cols.size();
    if (max(n,m) \le 2) {
         vector<int> ans(n,-1);
         for (int i=0; i<n; i++) {</pre>
              for(int j:cols){
                  if (ans[i] ==-1||f(rows[i], ans[i], j)){
                       ans[i]=j;
```

 $a[j] -= a[j^i];$ 

```
Chula[) ? - ? )]
```

```
Gaussian Elimination Binary Trie Infix Propostfix
```

```
21
```

```
return ans;
    if(n<m){
        // reduce
        vector<int> st:
        for(int j:cols){
            while(true) {
                if(st.empty()){
                     st.emplace_back(j);
                 }else if(f(rows[(int)st.size()-1],st.back(),j))
                     st.pop_back();
                 }else if(st.size()<n){</pre>
                     st.emplace_back(j);
                     break:
                 }else{
                     break
        return smawck(f,rows,st);
    vector<int> ans(n,-1);
    vector<int> new rows;
    for (int i=1; i < n; i += 2) {</pre>
        new_rows.emplace_back(rows[i]);
    auto res=smawck(f,new_rows,cols);
    for(int i=0;i<new_rows.size();i++) {</pre>
        ans[2*i+1]=res[i];
    for (int i=0, l=0, r=0; i<n; i+=2) {</pre>
        if (i+1==n) r=m;
        while(r<m&&cols[r]<=ans[i+1])r++;
        ans[i]=cols[l++];
        for(;1<r;1++){
            if(f(rows[i], ans[i], cols[l])){
                 ans[i]=cols[1];
        1--:
    return ans;
template < class F>
vector<int> smawck(const F &f,int n,int m) {
    vector<int> rows(n),cols(m);
    iota(rows.begin(),rows.end(),0);
    iota(cols.begin(),cols.end(),0);
    return smawck(f,rows,cols);
// Max Plus Convolution.
// b must be convex, i.e. b[i]-b[i-1]>=b[i+1]-b[i].
template < class T>
vector<T> max_plus_convolution_arbitary_convex(vector<T> a,
    const vector<T> &b) {
    if(a.empty()||b.empty())return {};
    if((int)b.size()==1){
        for (auto &x:a) x+=b[0];
        return a;
    int n=(int)a.size(),m=(int)b.size();
    auto f=[&](int i,int j){
```

```
return a[j]+b[i-j];
    };
    auto cmp=[&](int i,int j,int k) {
        if(i<k)return false;</pre>
        if (i-j>=m) return true;
        return f(i, j) <= f(i, k);
    };
    auto best=smawck(cmp,n+m-1,n);
    vector<T> ans(n+m-1);
    for(int i=0;i<n+m-1;i++) {</pre>
        ans[i]=f(i,best[i]);
    return ans;
Various (15)
Gaussian Elimination.hpp
Description: Gaussian Elimination
"../template/Header.hpp"
                                                       e89ecb, 34 lines
struct Gauss {
 int n, sz;
 vector<ll> basis;
 Gauss(int n = 0) {
   init(n);
 void init(int n) {
   n = _n, sz = 0;
   basis.assign(n, 0);
 void insert(ll x) {
    for (int i = n - 1; i >= 0; i--)
     if (x >> i & 1) {
        if (!basis[i]) {
          basis[i] = x;
          sz++;
          return:
        x ^= basis[i];
 11 \text{ getmax}(11 \text{ k} = 0)  {
   11 tot = 111 << sz, res = 0;
    for (int i = n - 1; i >= 0; i--)
     if (basis[i]) {
       tot >>= 1;
        if ((k >= tot && res >> i & 1) || (k < tot && res >> i
             & 1 ^ 1))
          res ^= basis[i];
        if (k >= tot)
          k -= tot;
    return res;
BinaryTrie.hpp
Description: Binary Trie
"../template/Header.hpp"
                                                       525bf4, 59 lines
using node_t = array<int, 2>;
template<size_t S>
```

struct binary\_trie {

int cur = 0;

vector<int> cnt = {0}; int cnt\_nodes = 0;

void insert(int v) {

vector<node\_t> t = {node\_t()};

```
cnt[0]++;
    for(int i=S-1; i>=0; --i) {
      int b = (v \& (1 << i)) ? 1: 0;
      if(!t[cur][b]) {
        t[cur][b] = ++cnt_nodes;
        t.emplace_back(node_t());
        cnt.emplace_back(0);
      cnt[t[cur][b]]++;
      cur = t[curl[b];
  void remove(int v) {
    int cur = 0;
    cnt[0]--;
    for(int i=S-1; i>=0; --i) {
      int b = (v & (1 << i)) ? 1: 0;
      cnt[t[cur][b]]--;
      cur = t[curl[b];
  int get_min(int v) {
    int cur = 0, res = 0;
    for(int i=(int) S-1; i>=0; --i) {
      int b = (v \& (1 << i)) ? 1 : 0;
      if(t[cur][b] && cnt[t[cur][b]]) {
        cur = t[cur][b];
      else {
        res |= (1 << i);
        cur = t[cur][!b];
    return res;
  int get_max(int v) {
    int cur = 0, res = 0;
    for(int i=(int) S-1; i>=0; --i) {
      int b = (v \& (1 << i)) ? 1 : 0;
      if(t[cur][!b] && cnt[t[cur][!b]]) {
        res |= (1 << i);
        cur = t[cur][!b];
        cur = t[cur][b];
    return res;
};
InfixPropostfix.hpp
Description: Infix to Pro-Postfix
"../template/Header.hpp"
                                                      517f57, 47 lines
stack<char> opr;
stack<int> val;
bool isOpr(char x) {
  return x == '+' || x == '*';
int prio(char x) {
  if(x == '(') return -1;
  if(x == '+') return 1;
  if(x == '*') return 2;
  return 0;
```

```
int do_opr(int 1, int r, char o) {
  if(o == '+') {
   return 1 + r;
 return 1 * r;
void pop_stack() {
 int rhs = val.top(); val.pop();
  int lhs = val.top(); val.pop();
  int new_val = do_opr(lhs, rhs, opr.top());
  val.emplace(new_val);
  opr.pop();
int cal(string s) {
  for(auto x: s) {
   if(isdigit(x)) val.emplace(x - '0');
   else if(x == '(') opr.emplace('(');
    else if(x == ')') {
     while(!opr.empty() && opr.top() != '(')
       pop_stack();
      opr.pop();
    else {
      while(!opr.empty() && prio(opr.top()) >= prio(x))
       pop_stack();
      opr.emplace(x);
  while(!opr.empty()) pop_stack();
 return val.top();
```

### 15.1 Optimization tricks

\_\_builtin\_ia32\_ldmxcsr(40896); disables denormals (which make floats 20x slower near their minimum value).

### 15.1.1 Bit hacks

Chula[) ? - ? )]

- x & -x is the least bit in x.
- for (int x = m; x; ) { --x &= m; ... } loops over all subset masks of m (except m itself).
- c = x&-x, r = x+c; (((r^x) >> 2)/c) | r is the next number after x with the same number of bits set.
- rep(b,0,K) rep(i,0,(1 << K))</li>
   if (i & 1 << b) D[i] += D[i^(1 << b)];</li>
   computes all sums of subsets.

### 15.1.2 Pragmas

- #pragma GCC optimize ("Ofast") will make GCC auto-vectorize loops and optimizes floating points better.
- #pragma GCC target ("avx2") can double performance of vectorized code, but causes crashes on old machines.
- #pragma GCC optimize ("trapv") kills the program on integer overflows (but is really slow).

# Competitive Programming Topics



topics.txt

Bitonic cycle

159 lines

Recursion Divide and conquer Finding interesting points in N log N Algorithm analysis Master theorem Amortized time complexity Greedy algorithm Scheduling Max contiguous subvector sum Invariants Huffman encoding Graph theory Dynamic graphs (extra book-keeping) Breadth first search Depth first search \* Normal trees / DFS trees Dijkstra's algorithm MST: Prim's algorithm Bellman-Ford Konig's theorem and vertex cover Min-cost max flow Lovasz toggle Matrix tree theorem Maximal matching, general graphs Hopcroft-Karp Hall's marriage theorem Graphical sequences Floyd-Warshall Euler cycles Flow networks \* Augmenting paths \* Edmonds-Karp Bipartite matching Min. path cover Topological sorting Strongly connected components Cut vertices, cut-edges and biconnected components Edge coloring \* Trees Vertex coloring \* Bipartite graphs (=> trees) \* 3^n (special case of set cover) Diameter and centroid K'th shortest path Shortest cycle Dynamic programming Knapsack Coin change Longest common subsequence Longest increasing subsequence Number of paths in a dag Shortest path in a dag Dynprog over intervals Dynprog over subsets Dynprog over probabilities Dynprog over trees 3^n set cover Divide and conquer Knuth optimization Convex hull optimizations RMQ (sparse table a.k.a 2^k-jumps)

Log partitioning (loop over most restricted) Combinatorics Computation of binomial coefficients Pigeon-hole principle Inclusion/exclusion Catalan number Pick's theorem Number theory Integer parts Divisibility Euclidean algorithm Modular arithmetic \* Modular multiplication \* Modular inverses \* Modular exponentiation by squaring Chinese remainder theorem Fermat's little theorem Euler's theorem Phi function Frobenius number Quadratic reciprocity Pollard-Rho Miller-Rabin Hensel lifting Vieta root jumping Game theory Combinatorial games Game trees Mini-max Nim Games on graphs Games on graphs with loops Grundy numbers Bipartite games without repetition General games without repetition Alpha-beta pruning Probability theory Optimization Binary search Ternary search Unimodality and convex functions Binary search on derivative Numerical methods Numeric integration Newton's method Root-finding with binary/ternary search Golden section search Matrices Gaussian elimination Exponentiation by squaring Sorting Radix sort Geometry Coordinates and vectors \* Cross product \* Scalar product Convex hull Polygon cut Closest pair Coordinate-compression Quadtrees KD-trees All segment-segment intersection Discretization (convert to events and sweep) Angle sweeping Line sweeping Discrete second derivatives Strings

Longest common substring Palindrome subsequences Knuth-Morris-Pratt Tries Rolling polynomial hashes Suffix array Suffix tree Aho-Corasick Manacher's algorithm Letter position lists Combinatorial search Meet in the middle Brute-force with pruning Best-first (A\*) Bidirectional search Iterative deepening DFS / A\* Data structures LCA (2^k-jumps in trees in general) Pull/push-technique on trees Heavy-light decomposition Centroid decomposition Lazy propagation Self-balancing trees Convex hull trick (wcipeg.com/wiki/Convex\_hull\_trick) Monotone queues / monotone stacks / sliding queues Sliding queue using 2 stacks Persistent segment tree