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template from KACTL 2024-08-16

template c IntPerm multinomial

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Template (1)	
template.cpp	27 lines
#pragma once	
<pre>#include <bits stdc++.h=""></bits></pre>	

```
#define sz(x) (int)(x).size()
#define all(x) (x).begin(), (x).end()
using namespace std;
using 11 = long long;
using db = long double;
using vi = vector<int>;
using vl = vector<11>;
using vd = vector<db>;
using pii = pair<int, int>;
using pll = pair<11, 11>;
using pdd = pair<db, db>;
const int INF = 0x3ffffffff;
// const int MOD=1000000007;
const int MOD = 998244353;
const 11 LINF = 0x1fffffffffffffffff;
const db DINF = numeric_limits<db>::infinity();
const db EPS = 1e-9;
const db PI = acos(db(-1));
```

1	<pre>int main() { cin.tie(nullptr)->sync_with_stdio(false);</pre>	
1	c.sh	
L	C.SH	2 lines
	g++ -std=gnu++2a -Wall \$1 -o a.out	
	./a.out	

Mathematics (2)

2.1 Goldbatch's Conjecture

- Even number can be written in sum of two primes (Up to 1e12)
- Range of N^{th} prime and $N + 1^{th}$ prime will be less than or equal to 300 (Up to 1e12)

2.2 Divisibility

Number of divisors of N is given by $\prod_{i=1}^{k} (a_i + 1)$ where $N = \prod_{i=1}^{k} p_i^{a_i}$ and p_i are prime factors of N.

Combinatorial (3)

3.1 Permutations

3.1.1 Factorial

n	1 2 3	3 4	5 6	7	8	9	10	
n!	1 2 6	3 24 1	20 720	5040	40320	362880	3628800	-
n	11	12	13	14	15	16	17	
$\overline{n!}$	4.0e7	7 4.8e	$8 6.2 \mathrm{e}{9}$	8.7e1	l0 1.3e	$12 \ 2.1e1$	3 3.6e14	
n	20	25	30	40	50 - 10	00 - 150) 171	
$\overline{n!}$	2e18	2e25	3e32 8	647.3	e64 9e	157 6e26	$52 > \text{DBL}_{-}$	MAX

IntPerm.h

Description: Permutation -> integer conversion. (Not order preserving.) Integer -> permutation can use a lookup table. **Time:** $\mathcal{O}(n)$

int permToInt(vi &v){
 int use = 0, i = 0, r = 0;
 for (int x : v) r = r * ++i + __builtin_popcount(use & -(1 << x)),
 use |= 1 << x; // (note: minus, not ~!)
 return r;
}</pre>

3.1.2 Cycles

Let $g_S(n)$ be the number of *n*-permutations whose cycle lengths all belong to the set S. Then

$$\sum_{n=0}^{\infty} g_S(n) \frac{x^n}{n!} = \exp\left(\sum_{n \in S} \frac{x^n}{n}\right)$$

3.1.3 Derangements

Permutations of a set such that none of the elements appear in their original position.

$$D(n) = (n-1)(D(n-1) + D(n-2)) = nD(n-1) + (-1)^n = \left\lfloor \frac{n!}{e} \right\rfloor$$

3.1.4 Burnside's lemma

Given a group G of symmetries and a set X, the number of elements of X up to symmetry equals

$$\frac{1}{|G|} \sum_{g \in G} |X^g|,$$

where X^g are the elements fixed by g(g.x = x).

If f(n) counts "configurations" (of some sort) of length n, we can ignore rotational symmetry using $G = \mathbb{Z}_n$ to get

$$g(n) = \frac{1}{n} \sum_{k=0}^{n-1} f(\gcd(n,k)) = \frac{1}{n} \sum_{k|n} f(k)\phi(n/k).$$

3.2 Partitions and subsets

3.2.1 Partition function

Number of ways of writing n as a sum of positive integers, disregarding the order of the summands.

$$p(0) = 1, \ p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k+1} p(n - k(3k - 1)/2)$$

3.2.2 Lucas' Theorem

Let n, m be non-negative integers and p a prime. Write $n = n_k p^k + ... + n_1 p + n_0$ and $m = m_k p^k + ... + m_1 p + m_0$. Then $\binom{n}{m} \equiv \prod_{i=0}^k \binom{n_i}{m_i} \pmod{p}$.

3.2.3 Binomials

multinomial.h

return c;

Description: Computes $\binom{k_1+\cdots+k_n}{k_1,k_2,\ldots,k_n} = \frac{(\sum k_i)!}{k_1!k_2!\ldots k_n!}$.

11 multinomial (vi& v) {
 ll c = 1, m = v.empty() ? 1 : v[0];
 rep(i,1,sz(v)) rep(j,0,v[i])
 c = c * ++m / (i+1);

.3 General purpose numbers

3.3.1 Bernoulli numbers

EGF of Bernoulli numbers is $B(t) = \frac{t}{e^t - 1}$ (FFT-able). $B[0, \ldots] = [1, -\frac{1}{2}, \frac{1}{6}, 0, -\frac{1}{30}, 0, \frac{1}{42}, \ldots]$

Sums of powers:

$$\sum_{i=1}^{n} n^{m} = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_{k} \cdot (n+1)^{m+1-k}$$

Euler-Maclaurin formula for infinite sums:

$$\sum_{i=m}^{\infty} f(i) = \int_{m}^{\infty} f(x)dx - \sum_{k=1}^{\infty} \frac{B_{k}}{k!} f^{(k-1)}(m)$$

$$\approx \int_{m}^{\infty} f(x)dx + \frac{f(m)}{2} - \frac{f'(m)}{12} + \frac{f'''(m)}{720} + O(f^{(5)}(m))$$

3.3.2 Stirling numbers of the first kind

Number of permutations on n items with k cycles.

$$c(n,k) = c(n-1,k-1) + (n-1)c(n-1,k), \ c(0,0) = 1$$
$$\sum_{k=0}^{n} c(n,k)x^{k} = x(x+1)\dots(x+n-1)$$

c(8, k) = 8, 0, 5040, 13068, 13132, 6769, 1960, 322, 28, 1 $c(n,2) = 0, 0, 1, 3, 11, 50, 274, 1764, 13068, 109584, \dots$

3.3.3 Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j:s s.t. $\pi(j) > \pi(j+1)$, k+1 j:s s.t. $\pi(j) \geq j$, k j:s s.t. $\pi(j) > j$.

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n,n-1) = 1$$

$$E(n,k) = \sum_{j=0}^{k} (-1)^{j} \binom{n+1}{j} (k+1-j)^{n}$$

3.3.4 Stirling numbers of the second kind

Partitions of n distinct elements into exactly k groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k)$$

$$S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} \binom{k}{j} j^{n}$$

3.3.5 Bell numbers

Total number of partitions of n distinct elements. B(n) = $1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, \dots$ For p prime,

$$B(p^m + n) \equiv mB(n) + B(n+1) \pmod{p}$$

3.3.6 Labeled unrooted trees

on
$$n$$
 vertices: n^{n-2} # on k existing trees of size n_i : $n_1 n_2 \cdots n_k n^{k-2}$ # with degrees d_i : $(n-2)!/((d_1-1)!\cdots(d_n-1)!)$

3.3.7 Catalan numbers

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \binom{2n}{n} - \binom{2n}{n+1} = \frac{(2n)!}{(n+1)!n!}$$

$$C_0 = 1, \ C_{n+1} = \frac{2(2n+1)}{n+2} C_n, \ C_{n+1} = \sum_{n=1}^{\infty} C_i C_{n-n}$$

 $C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, \dots$

- sub-diagonal monotone paths in an $n \times n$ grid.
- strings with n pairs of parenthesis, correctly nested.
- binary trees with with n+1 leaves (0 or 2 children).
- ordered trees with n+1 vertices.
- ways a convex polygon with n+2 sides can be cut into triangles by connecting vertices with straight lines.
- permutations of [n] with no 3-term increasing subseq.

Numerical (4)

4.1 Newton's Method

if
$$F(Q) = 0$$
, then $Q_{2n} \equiv Q_n - \frac{F(Q_n)}{F'(Q_n)} \pmod{x^{2n}}$

$$Q = P^{-1} : Q_{2n} \equiv Q_n \cdot (2 - P \cdot Q_n^2) \pmod{x^{2n}}$$

$$Q = \ln P = \int \frac{P'}{P} dx$$

$$Q = e^p : Q_{2n} \equiv Q_n (1 + P - \ln Q_n) \pmod{x^{2n}}$$

$$Q = \sqrt{P} : Q_{2n} \equiv \frac{1}{2} (Q_n + P \cdot Q_n^{-1}) \pmod{x^{2n}}$$

$$Q = P^k = \alpha^k x^{kt} e^{k \ln T} : P = \alpha \cdot x^t \cdot T \cdot T(0) = 1$$

Group (5)

5.1 Monoid

monoid/MonoidBase.hpp Description: Monoid Base class

e75b74, 6 lines

```
template < class T, T (*combine) (T, T), T (*identity) () >
struct MonoidBase
    using value_type = T;
    static constexpr T op(const T &x,const T &y) {return combine
    static constexpr T unit() {return identity();}
```

Action

action/MonoidActionBase.hpp Description: Monoid Action Base class.

```
template < class MInfo, class MTaq, typename MInfo::value_type
    (*combine) (typename MInfo::value_type,typename MTag::
         value_type)>
struct MonoidActionBase{
    using InfoMonoid = MInfo;
```

```
using TagMonoid = MTag;
    using Info = typename InfoMonoid::value_type;
    using Tag = typename TagMonoid::value type;
    static constexpr Info op(const Info &a,const Tag &b) {
        return combine (a, b);
};
action/DefaultAction.hpp
Description: Default Action class.
                                                      e45000. 10 lines
template < class Monoid>
struct DefaultAction{
    using InfoMonoid = Monoid;
    using TagMonoid = Monoid;
    using Info = typename Monoid::value_type;
    using Tag = typename Monoid::value_type;
    static constexpr Info op(const Info &a,const Tag &b) {
        return Monoid::op(a,b);
};
```

Data Structures (6)

```
OrderedSet.hpp
Description: Ordered Set
```

```
"../template/Header.hpp", <bits/extc++.h>
                                                           1a7f5f, 14 lines
using namespace __gnu_pbds;
template <class T>
using ordered_set = tree<T, null_type, less<T>, rb_tree_tag,
```

```
tree_order_statistics_node_update>;
// can be change to less equal
void usage() {
 ordered set<int> st, st 2;
 st.insert(2);
 st.insert(1);
 cout << st.order_of_key(2);</pre>
 cout << *st.find_by_order(1);</pre>
 st.join(st_2); // merge
```

FenwickTree.hpp

Description: Fenwick / Binary Indexed Tree

43767a, 41 lines

```
template<class T>
struct Fenwick{
    int n,logn;
    vector<T> t:
    Fenwick(){}
    Fenwick(int _n) {init(vector<T>(_n,T{}));}
    template<class U>
    Fenwick(const vector < U > &a) {init(a);}
    template<class U>
    void init(const vector<U> &a) {
        n=(int)a.size();
        logn=31- builtin clz(n);
        t.assign(n+1,T\{\});
        for (int i=1; i<=n; i++) {</pre>
             t[i]=t[i]+a[i-1];
             int j=i+(i\&-i);
             if (j<=n)t[j]=t[j]+t[i];</pre>
    void update(int x,const T &v) {
        for (int i=x+1; i<=n; i+=i&-i) t[i] =t[i]+v;</pre>
```

};

};

```
void update(int 1,int r,const T &v) {
        update (1, v), update (r+1, -v);
   T query(int x) {
       T res{};
        for(int i=x+1;i>0;i-=i&-i)res=res+t[i];
        return res;
    T query(int 1,int r){
        return query (r) -query (1-1);
    int find(const T &k) {
        int x=0;
        T cur{};
        for(int i=1<<logn;i>0;i>>=1)
            if (x+i<=n&&cur+t[x+i]<=k)x+=i,cur=cur+t[x];</pre>
SmallSegmentTree.hpp
Description: Zero-indexed max-tree. Bounds are inclusive to the left and
exclusive to the right. Can be changed by modifying T, f and unit.
Time: \mathcal{O}(\log N)
                                                      0f4bdb, 19 lines
struct Tree {
  typedef int T;
  static constexpr T unit = INT_MIN;
  T f (T a, T b) { return max(a, b); } // (any associative fn)
  vector<T> s; int n;
  Tree(int n = 0, T def = unit) : s(2*n, def), n(n) {}
  void update(int pos, T val) {
    for (s[pos += n] = val; pos /= 2;)
      s[pos] = f(s[pos * 2], s[pos * 2 + 1]);
  T query (int b, int e) { // query [b, e]
   T ra = unit, rb = unit;
    for (b += n, e += n; b < e; b /= 2, e /= 2) {
      if (b % 2) ra = f(ra, s[b++]);
     if (e % 2) rb = f(s[--e], rb);
    return f(ra, rb);
Segment Tree.hpp
Description: Segment Tree
                                                       c51dec. 85 lines
template < class Monoid>
struct SegmentTree{
    using T = typename Monoid::value_type;
    int n;
    vector<T> t;
    SegmentTree(){}
    SegmentTree(int n, function<T(int)> create) {init(n, create);}
    SegmentTree(int n,T v=Monoid::unit()){init(n,[&](int){
         return v; });}
    template < class U>
    SegmentTree(const vector<U> &a){init((int)a.size(),[&](int
         i) {return T(a[i]);});}
    void init(int _n,function<T(int)> create){
        t.assign(4<<(31-__builtin_clz(n)), Monoid::unit());
        function<void(int,int,int)> build=[&](int 1,int r,int i
            if (l==r) return void(t[i]=create(l));
            int m = (1+r)/2;
            build(l, m, i*2);
```

build(m+1, r, i*2+1);

```
pull(i);
    };
    build(0, n-1, 1);
void pull(int i){
    t[i]=Monoid::op(t[i*2],t[i*2+1]);
void modify(int 1,int r,int i,int x,const T &v) {
    if (x<1 | | r<x) return;
    if(l==r)return void(t[i]=v);
    int m = (1+r)/2;
    modify(1, m, i*2, x, v);
    modify (m+1,r,i*2+1,x,v);
    pull(i);
void modify(int x,const T &v) {
    modify (0, n-1, 1, x, v);
template<class U>
void update(int 1, int r, int i, int x, const U &v) {
    if (x<1 | | r<x) return;
    if(l==r)return void(t[i]=Monoid::op(t[i],v));
    int m = (1+r)/2;
    update (1, m, i*2, x, v);
    update (m+1,r,i*2+1,x,v);
    pull(i);
template < class U>
void update(int x,const U &v) {
    update (0, n-1, 1, x, v);
T query(int 1, int r, int i, int x, int y) {
    if (y<1||r<x) return Monoid::unit();</pre>
    if (x<=1&&r<=y) return t[i];
    int m = (1+r)/2;
    return Monoid::op(query(1, m, i*2, x, y), query(m+1, r, i*2+1,
         x,y));
T query(int x,int y) {
    return query (0, n-1, 1, x, y);
template<class F>
int findfirst(int 1, int r, int i, int x, int y, const F &f) {
    if(y<1||r<x||!f(t[i]))return n;
    if (l==r) return 1;
    int m = (1+r)/2;
    int res=findfirst(1, m, i*2, x, y, f);
    if (res==n) res=findfirst (m+1, r, i * 2+1, x, y, f);
    return res;
template < class F>
int findfirst(int x,int y,const F &f) {
    return findfirst(0,n-1,1,x,y,f);
template < class F>
int findlast(int 1,int r,int i,int x,int y,const F &f){
    if (y<1 | | r<x | | ! f (t[i])) return -1;</pre>
    if(l==r)return 1;
    int m = (1+r)/2;
    int res=findlast(m+1,r,i*2+1,x,y,f);
    if(res==-1) res=findlast(1,m,i*2,x,y,f);
    return res;
template < class F>
int findlast(int x,int y,const F &f){
    return findlast(0, n-1, 1, x, y, f);
```

};

```
Description: Segment Tree with Lazy Propagation
                                                      91ab0c, 103 lines
template < class MonoidAction >
struct LazySegmentTree{
    using InfoMonoid = typename MonoidAction::InfoMonoid;
    using TagMonoid = typename MonoidAction::TagMonoid;
    using Info = typename MonoidAction::Info;
    using Tag = typename MonoidAction::Tag;
    int n:
    vector<Info> t;
    vector<Tag> lz;
    LazySegmentTree(){}
    LazySegmentTree(int n, function<Info(int)> create) {init(n,
    LazySegmentTree(int n, Info v=InfoMonoid::unit()) {init(n
         , [&] (int) {return v;});}
    template<class T>
    LazySegmentTree(const vector<T> &a) {init((int)a.size(), [&](
         int i) {return Info(a[i]);});}
    void init(int _n,function<Info(int)> create){
        int m=4 << (31- builtin clz(n));
        t.assign(m,InfoMonoid::unit());
        lz.assign(m, TagMonoid::unit());
        function<void(int,int,int)> build=[&](int l,int r,int i
            if (l==r) return void(t[i]=create(l));
            int m = (1+r)/2;
            build(1, m, i * 2);
            build(m+1, r, i*2+1);
            pull(i);
        };
        build(0, n-1, 1);
    void pull(int i) {
        t[i]=InfoMonoid::op(t[i*2],t[i*2+1]);
    void apply(int i,const Tag &v){
        t[i] = MonoidAction::op(t[i], v);
        lz[i]=TagMonoid::op(lz[i],v);
    void push(int i) {
        apply(i*2,lz[i]);
        apply(i*2+1,lz[i]);
        lz[i]=TagMonoid::unit();
    void modify(int l,int r,int i,int x,const Info &v){
        if (x<1 | | r<x) return;
        if(l==r)return void(t[i]=v);
        int m = (1+r)/2;
        push(i);
        modify(1, m, i*2, x, v);
        modify (m+1,r,i*2+1,x,v);
        pull(i);
    void modify(int x,const Info &v){
        modify (0, n-1, 1, x, v);
    void update(int 1,int r,int i,int x,int y,const Tag &v) {
        if (y<1 | | r<x) return;</pre>
        if (x<=l&&r<=y) return apply(i,v);</pre>
        int m = (1+r)/2;
        push(i);
        update(1, m, i*2, x, y, v);
        update (m+1,r,i*2+1,x,y,v);
        pull(i);
    void update(int x,int y,const Tag &v) {
```

LazySegmentTree.hpp

update (0, n-1, 1, x, y, v);

int m = (1+r)/2;

if (x<=1&&r<=y) return t[i];</pre>

Info query(int 1,int r,int i,int x,int y) {

if (y<1||r<x) return InfoMonoid::unit();</pre>

DynamicSegmentTree DSU BinaryTrie

```
push(i);
        return InfoMonoid::op(query(1, m, i*2, x, y), query(m+1, r, i
             *2+1, x, y));
    Info query(int x,int y){
        return query (0, n-1, 1, x, y);
    template < class F>
    int findfirst(int 1,int r,int i,int x,int y,const F &f) {
        if (y<1||r<x||!f(t[i])) return n;</pre>
        if(l==r)return 1;
        int m = (1+r)/2;
        push(i);
        int res=findfirst(1, m, i*2, x, y, f);
        if (res==n) res=findfirst (m+1, r, i*2+1, x, y, f);
        return res;
    template<class F>
    int findfirst(int x,int y,const F &f) {
        return findfirst(0,n-1,1,x,y,f);
    template<class F>
    int findlast(int l,int r,int i,int x,int y,const F &f) {
        if(y<1||r<x||!f(t[i]))return -1;
        if(l==r)return 1;
        int m = (1+r)/2;
        int res=findlast(m+1,r,i*2+1,x,y,f);
        if (res==-1) res=findlast (1, m, i*2, x, y, f);
        return res;
    template<class F>
    int findlast(int x,int y,const F &f) {
        return findlast(0,n-1,1,x,y,f);
};
DynamicSegmentTree.hpp
Description: Dynamic Segment Tree
                                                      e84eeb, 106 lines
template < class MonoidAction >
struct DynamicSegmentTree{
    using InfoMonoid = typename MonoidAction::InfoMonoid;
    using TagMonoid = typename MonoidAction::TagMonoid;
    using Info = typename MonoidAction::Info;
    using Tag = typename MonoidAction::Tag;
    struct Node;
    using Ptr = Node*;
    struct Node {
        Info val:
        Tag lz;
        Ptr l,r;
        Node (Info v): val(v), lz(TagMonoid::unit()), l(nullptr), r(
             nullptr){}
        Node(Info v, Tag t):val(v),lz(t),l(nullptr),r(nullptr){}
    11 lb, ub;
    Ptr rt:
    function<Info(11,11)> create;
    DynamicSegmentTree() {init(0,0);}
    DynamicSegmentTree(ll n) {init(0, n-1);}
    DynamicSegmentTree(ll lb, ll ub) {init(lb, ub);}
```

```
DynamicSegmentTree(11 lb, 11 ub, function<Info(11,11) > create
     ) {init(lb,ub,create);}
void init(ll _lb,ll _ub,function<Info(ll,ll)> _create=[](ll
      1,11 r) {return InfoMonoid::unit();}) {
    lb=_lb, ub=_ub;
    create=_create;
    rt=new Node(create(lb,ub));
Info val(Ptr t){
    return t?t->val:InfoMonoid::unit();
void pull(Ptr &t) {
    t->val=InfoMonoid::op(val(t->1),val(t->r));
void apply(Ptr &t,const Tag &v,ll l,ll r){
    if(!t)t=new Node(create(1,r));
    t->val=MonoidAction::op(t->val, v);
    t->1z=TagMonoid::op(t->1z,v);
void push(Ptr &t, ll 1, ll m, ll r) {
    apply (t->1,t->1z,1,m);
    apply (t->r,t->lz,m+1,r);
    t->lz=TagMonoid::unit();
void modify(ll 1,ll r,Ptr &t,ll x,const Info &v){
    if (x<1 | | r<x) return;
    if (l==r) return void(t->val=v);
    11 m=1+(r-1)/2;
    push(t,l,m,r);
    modify (1, m, t \rightarrow 1, x, v);
    modify (m+1, r, t->r, x, v);
    pull(t);
void modify(ll x,const Info &v) {
    modify(lb,ub,rt,x,v);
void update(ll 1,ll r,Ptr &t,ll x,ll y,const Tag &v) {
    if (y<1 | | r<x) return;
    if (x<=l&&r<=y) return apply (t, v, l, r);</pre>
    11 m=1+(r-1)/2;
    push(t,l,m,r);
    update(1, m, t \rightarrow 1, x, v, v);
    update (m+1,r,t->r,x,y,v);
    pull(t);
void update(ll x,ll y,const Tag &v) {
    update(lb,ub,rt,x,y,v);
Info query(ll 1,11 r,Ptr &t,11 x,11 y){
    if(v<1||r<x)return InfoMonoid::unit();</pre>
    if(x<=1&&r<=y) return t->val;
    11 m=1+(r-1)/2;
    push(t,l,m,r);
    return InfoMonoid::op(query(l,m,t->l,x,y),query(m+1,r,t
         ->r,x,y));
Info query(ll x,ll y){
    return query(lb,ub,rt,x,y);
template < class F >
11 findfirst(ll 1,ll r,Ptr t,ll x,ll y,const F &f) {
    if (y<1 | | r<x | | ! f (t->val) ) return -1;
    if(l==r)return 1;
    11 m=1+(r-1)/2;
    push(t,1,m,r);
    ll res=findfirst(l, m, t->l, x, y, f);
    if (res==-1) res=findfirst (m+1, r, t->r, x, y, f);
    return res;
```

```
template<class F>
    11 findfirst(ll x, ll y, const F &f) {
        return findfirst(lb,ub,rt,x,v,f);
    template < class F>
    11 findlast(ll 1,ll r,Ptr t,ll x,ll y,const F &f) {
        if (y<1||r<x||!t||!f(t->val))return -1;
        if (l==r) return 1;
        11 m=1+(r-1)/2;
        push(t,l,m,r);
        ll res=findlast(m+1,r,t->r,x,y,f);
        if(res==-1)res=findlast(1, m, t->1, x, y, f);
        return res;
    template<class F>
    11 findlast(ll x,ll y,const F &f) {
        return findlast(lb,ub,rt,x,y,f);
};
DSU.hpp
Description: Disjoint Set Union.
struct DSU{
    vector<int> p,sz;
    DSU(int n) {init(n);}
    void init(int n) {
        p.resize(n);
        iota(p.begin(),p.end(),0);
        sz.assign(n,1);
    int find(int u) {
        return p[u] == u?u:p[u] = find(p[u]);
    bool same(int u,int v){
        return find(u) == find(v);
    bool merge(int u,int v) {
        u=find(u), v=find(v);
        if(u==v)return false;
        sz[u]+=sz[v];
        p[v]=u;
        return true;
    int size(int u){
        return sz[find(u)];
};
BinaryTrie.hpp
Description: Binary Trie
                                                      ae5b7a, 66 lines
template<int BIT, class T = uint32_t, class S = int>
struct BinaryTrie{
    struct Node{
        array<int,2> ch;
        S cnt;
        Node():ch\{-1,-1\}, cnt(0){}
    vector<Node> t;
    BinaryTrie():t{Node()}{}
    int new_node(){
        t.emplace back(Node());
        return t.size()-1;
    S size(){
        return t[0].cnt;
```

LiChaoTree DynamicLiChaoTree SplayTreeBase

```
bool empty(){
        return size()==0;
    S get_cnt(int i) {
        return i!=-1?t[i].cnt:S(0);
    void insert(T x,S k=1){
        int u=0;
        t[u].cnt+=k;
        for(int i=BIT-1;i>=0;i--){
            int v=x>>i&1;
            if(t[u].ch[v]==-1)t[u].ch[v]=new_node();
            u=t[u].ch[v];
            t[u].cnt+=k;
    void erase(T x,S k=1){
        int u=0;
        assert(t[u].cnt>=k);
        t[u].cnt-=k;
        for(int i=BIT-1; i>=0; i--) {
            int v=x>>i&1;
            u=t[u].ch[v];
            assert (u!=-1&&t[u].cnt>=k);
            t[u].cnt-=k;
    T kth(S k, T x=0) {
        assert(k<size());
        int u=0;
        T res=0;
        for(int i=BIT-1; i>=0; i--) {
            int v=x>>i&1;
            if(k<get_cnt(t[u].ch[v])){
                u=t[u].ch[v];
            }else{
                res|=T(1)<<i;
                k-=get_cnt(t[u].ch[v]);
                u=t[u].ch[v^1];
        return res;
    T min(T x) {
        return kth(0,x);
    T max(T x){
        return kth(size()-1,x);
};
LiChaoTree.hpp
Description: Li-Chao Tree (minimize).
                                                      4ab713, 52 lines
template<class T>
struct LiChaoTree{
    static const T INF=numeric_limits<T>::max()/2;
    struct Line{
        T m,c;
        Line(T _m, T _c):m(_m),c(_c){}
        inline T eval(T x)const{return m*x+c;}
    };
    vector<T> xs;
    vector<Line> t;
    LiChaoTree(){}
    LiChaoTree (const vector<T> &x):xs(x) {init(x);}
    LiChaoTree(int n):xs(n){
        vector<T> x(n);
        iota(x.begin(), x.end(), 0);
```

```
init(x);
    void init(const vector<T> &x) {
        sort(xs.begin(),xs.end());
        xs.erase(unique(xs.begin(),xs.end()),xs.end());
        t.assign(4<<(31-__builtin_clz(xs.size())),Line(0,INF));
    void insert(int 1, int r, int i, Line v) {
        int m = (1+r)/2;
        if (v.eval(xs[m]) <t[i].eval(xs[m])) swap(t[i],v);</pre>
        if (v.eval(xs[1]) <t[i].eval(xs[1]))insert(1, m, i*2, v);</pre>
        if (v.eval(xs[r]) <t[i].eval(xs[r]))insert(m+1,r,i*2+1,v)</pre>
    inline void insert(T m,T c){
        insert (0, (int) xs.size()-1,1, Line(m,c));
    void insert_range(int 1, int r, int i, T x, T y, Line v) {
        if (y<xs[1] | |xs[r]<x) return;
        if (x<=xs[1]&&xs[r]<=y) return insert(1,r,i,v);</pre>
        int m = (1+r)/2;
        insert_range(1, m, i * 2, x, y, v);
        insert_range(m+1, r, i*2+1, x, y, v);
    inline void insert_range(T m, T c, T x, T y) {
        insert_range(0, (int)xs.size()-1,1,x,y,Line(m,c));
    T query(int 1,int r,int i,T x){
        if(l==r)return t[i].eval(x);
        int m = (1+r)/2;
        if (x \le xs[m]) return min (t[i].eval(x), query(1, m, i \times 2, x));
        return min(t[i].eval(x), query(m+1, r, i*2+1, x));
    inline T query(T x){
        return query(0,(int)xs.size()-1,1,x);
};
DynamicLiChaoTree.hpp
Description: Dynamic Li-Chao Tree (minimize).
                                                        b8af36, 50 lines
template<class T>
struct DvnamicLiChaoTree{
    static const T INF=numeric_limits<T>::max()/2;
    struct Line{
        T m,c;
        Line(T _m, T _c):m(_m),c(_c){}
        inline T eval(T x)const{return m*x+c;}
    struct Node;
    using Ptr = Node*;
    struct Node {
        Line v:
        Ptr l,r;
        Node():v(0,INF),l(nullptr),r(nullptr){}
        Node(Line _v):v(_v),l(nullptr),r(nullptr){}
    };
    11 1b, ub;
    Ptr root:
    DynamicLiChaoTree(ll _lb,ll _ub):lb(_lb),ub(_ub),root(
         nullptr) {}
    void insert(T 1,T r,Ptr &t,Line v){
        if(!t)return void(t=new Node(v));
        T m=1+(r-1)/2;
        if(v.eval(m) < t->v.eval(m)) swap(t->v,v);
        if (v.eval(1) <t->v.eval(1)) insert(1, m, t->1, v);
        if (v.eval(r) <t->v.eval(r)) insert(m+1, r, t->r, v);
```

```
inline void insert(T m, T c) {
        insert(lb, ub, root, Line(m, c));
    void insert_range(T 1,T r,Ptr &t,T x,T y,Line v){
        if (y<1 | | r<x) return;</pre>
        if(!t)t=new Node();
        if (x<=l&&r<=v) return insert(l,r,t,v);</pre>
        T m=1+(r-1)/2;
        insert_range(l, m, t->l, x, y, v);
        insert_range(m+1, r, t->r, x, y, v);
    inline void insert_range(T m, T c, T x, T y) {
        insert_range(lb, ub, root, x, y, Line(m, c));
    T query(T 1,T r,Ptr t,T x) {
        if(!t)return INF;
        T m=1+(r-1)/2;
        if (x<=m) return min(t->v.eval(x), query(1, m, t->1, x));
        return min(t->v.eval(x), query(m+1, r, t->r, x));
    inline T query(T x){
        return query(lb,ub,root,x);
};
SplayTreeBase.hpp
Description: Splay Tree. splay(u) will make node u be the root of the tree
in amortized O(log n) time.
                                                       cc90a9, 113 lines
template < class Node >
struct SplayTreeBase{
    using Ptr = Node*;
    bool is root (Ptr t) {
         return ! (t->p) | | (t->p->1!=t&&t->p->r!=t);
    } // The parent of the root stores the path parant in link
    int size(Ptr t){
        return t?t->size:0;
    virtual void push(Ptr t){};
    virtual void pull(Ptr t){};
    int pos(Ptr t){
        if(t->p){
             if(t->p->1==t)return -1;
             if(t->p->r==t)return 1;
        return 0;
    void rotate(Ptr t) {
        Ptr x=t->p, y=x->p;
        if (pos(t) ==-1) {
             if ((x->l=t->r))t->r->p=x;
             t->r=x, x->p=t;
        }else{
             if((x->r=t->1))t->1->p=x;
             t->l=x, x->p=t;
        pull(x),pull(t);
        if((t->p=y)){
             if (y->1==x) y->1=t;
             if (y->r==x) y->r=t;
    void splay(Ptr t) {
        if(!t)return;
        push(t);
        while(!is_root(t)){
             Ptr x=t->p;
             if(is root(x)){
```

```
push(x), push(t);
            rotate(t);
            Ptr y=x->p;
            push(y),push(x),push(t);
            if(pos(x) == pos(t)) rotate(x), rotate(t);
            else rotate(t), rotate(t);
Ptr get_first(Ptr t){
    while (t->1) push (t), t=t->1;
    splay(t);
    return t;
Ptr get_last(Ptr t){
    while (t->r) push (t), t=t->r;
    splay(t);
    return t:
Ptr merge(Ptr 1,Ptr r) {
    splay(1), splay(r);
    if(!1)return r;
    if(!r)return 1;
    l=get_last(1);
    1->r=r;
    r->p=1;
    pull(1);
    return 1;
pair<Ptr,Ptr> split(Ptr t,int k) {
    if(!t)return {nullptr,nullptr};
    if (k==0) return {nullptr,t};
    if(k==size(t))return {t,nullptr};
    push(t);
    if(k<=size(t->1)){
        auto x=split(t->1,k);
        t->l=x.second;
        t->p=nullptr;
        if (x.second) x.second->p=t;
        pull(t);
        return {x.first,t};
        auto x=split(t->r,k-size(t->l)-1);
        t->r=x.first;
        t->p=nullptr;
        if(x.first)x.first->p=t;
        pull(t);
        return {t,x.second};
void insert(Ptr &t,int k,Ptr v) {
    splav(t);
    auto x=split(t,k);
    t=merge(merge(x.first,v),x.second);
void erase(Ptr &t,int k) {
    splav(t);
    auto x=split(t,k);
    auto v=split(x.second,1);
    // delete y.first;
    t=merge(x.first,y.second);
template<class T>
Ptr build(const vector<T> &v) {
    if(v.empty())return nullptr;
    function<Ptr(int,int)> build=[&](int 1,int r){
        if(l==r)return new Node(v[1]);
        int m = (1+r)/2;
```

```
return merge(build(1,m),build(m+1,r));
        return build(0, v.size()-1);
};
LazyReversibleBBST.hpp
Description: Lazy Reversible BBST Base.
                                                     904708, 81 lines
template < class Tree, class Node, class MonoidAction>
struct LazyReversibleBBST:Tree{
    using Tree::merge;
    using Tree::split;
    using typename Tree::Ptr;
    using InfoMonoid = typename MonoidAction::InfoMonoid;
    using TagMonoid = typename MonoidAction::TagMonoid;
    using Info = typename MonoidAction::Info;
    using Tag = typename MonoidAction::Tag;
    LazyReversibleBBST() = default;
    Info sum(Ptr t) {
        return t?t->sum:InfoMonoid::unit();
   void pull(Ptr t) {
       if(!t)return;
       push(t);
       t->size=1;
       t->sum=t->val;
       t->revsum=t->val;
       if(t->1){
            t->size+=t->l->size;
            t->sum=InfoMonoid::op(t->l->sum,t->sum);
            t->revsum=InfoMonoid::op(t->revsum,t->l->revsum);
       if(t->r){
            t->size+=t->r->size;
            t->sum=InfoMonoid::op(t->sum,t->r->sum);
            t->revsum=InfoMonoid::op(t->r->revsum,t->revsum);
    void push (Ptr t) {
       if(!t)return;
       if(t->rev){
            toggle (t->1);
            toggle(t->r);
            t->rev=false;
       if(t->lz!=TagMonoid::unit()){
            propagate(t->1,t->1z);
            propagate(t->r,t->lz);
            t->1z=TagMonoid::unit();
    void toggle(Ptr t) {
        if(!t)return;
        swap(t->1,t->r);
        swap(t->sum,t->revsum);
       t->rev^=true;
   void propagate(Ptr t,const Tag &v){
       if(!t)return;
       t->val=MonoidAction::op(t->val,v);
       t->sum=MonoidAction::op(t->sum, v);
        t->revsum=MonoidAction::op(t->revsum,v);
        t->1z=TagMonoid::op(t->1z,v);
    void apply(Ptr &t,int 1,int r,const Tag &v) {
        if (1>r) return;
```

```
auto x=split(t,1);
        auto y=split(x.second,r-l+1);
        propagate(y.first,v);
        t=merge(x.first, merge(y.first, y.second));
    Info query(Ptr &t,int 1,int r){
        if(l>r)return InfoMonoid::unit();
        auto x=split(t,1);
        auto y=split(x.second, r-1+1);
        Info res=sum(y.first);
        t=merge(x.first, merge(y.first, y.second));
        return res;
    void reverse(Ptr &t,int 1,int r) {
        if(1>r)return;
        auto x=split(t,1);
        auto y=split(x.second, r-1+1);
        toggle(y.first);
        t=merge(x.first, merge(y.first, y.second));
};
LazyReversibleSplayTree.hpp
Description: Lazy Reversible Splay Tree.
"SplayTreeBase.hpp", "LazyReversibleBBST.hpp"
                                                     b8455b, 23 lines
template < class MonoidAction >
struct LazyReversibleSplayTreeNode{
    using Ptr = LazyReversibleSplayTreeNode*;
    using InfoMonoid = typename MonoidAction::InfoMonoid;
    using TagMonoid = typename MonoidAction::TagMonoid;
    using Info = typename MonoidAction::Info;
    using Tag = typename MonoidAction::Tag;
    using value_type = Info;
    Ptr l,r,p;
    Info val, sum, revsum;
    Tag lz:
    int size;
    LazyReversibleSplayTreeNode(const Info &_val=InfoMonoid::
         unit(),const Tag & lz=TagMonoid::unit())
        :1(),r(),p(),val(_val),sum(_val),revsum(_val),lz(_lz),
             size(1), rev(false) {}
};
template < class MonoidAction >
struct LazyReversibleSplayTree
    : LazyReversibleBBST<SplayTreeBase<
         LazyReversibleSplayTreeNode<MonoidAction>>,
      LazyReversibleSplayTreeNode<MonoidAction>, MonoidAction>{
    using Node = LazyReversibleSplayTreeNode<MonoidAction>;
LinkCutTreeBase.hpp
Description: Link Cut Tree Base.
Usage: evert(u): make u be the root of the tree.
link(u,v): attach u to v.
cut(u,v): remove edge between u and v.
get_root(u): get the root of the tree containing u.
lca(u,v): get the lowest common ancestor of u and v.
fold(u,v): get the value of the path from u to v. _{\rm b432c3,\ 59\ lines}
template<class Splay>
struct LinkCutTreeBase:Splay{
    using Node = typename Splay::Node;
    using Ptr = Node*;
    using T = typename Node::value_type;
    Ptr expose(Ptr t) {
        Ptr pc=nullptr; // preferred child
        for(Ptr cur=t;cur;cur=cur->p) {
```

```
this->splay(cur);
        cur->r=pc;
        this->pull(cur);
        pc=cur;
    this->splay(t);
    return pc;
void evert(Ptr t) { // make t be the root of the tree
    expose(t);
    this->toggle(t);
    this->push(t);
void link(Ptr u,Ptr v){ // attach u to v
    evert(u);
    expose (v);
    u->p=v;
void cut(Ptr u,Ptr v){ // cut edge between u and v
    evert(u);
    expose(v);
    assert (u->p==v);
    v->1=u->p=nullptr;
    this->pull(v);
Ptr get_root(Ptr t){
    expose(t);
    while (t->1) this->push (t), t=t->1;
    this->splay(t);
    return t;
Ptr lca(Ptr u,Ptr v) {
    if(get_root(u)!=get_root(v))return nullptr;
    expose(u);
    return expose(v);
void set_val(Ptr t,const T &val){
    this->evert(t);
    t->val=val;
    this->pull(t);
T get val(Ptr t){
    this->evert(t);
    return t->val;
T fold (Ptr u, Ptr v) {
    evert(u);
    expose (v);
    return v->sum;
```

```
Usage: using Lct = LazyLinkCutTree<Action>;
using Ptr = Lct::Ptr;
using Node = Lct:: Node;
vector<Ptr> ptr(n);
for(int i=0;i<n;i++)ptr[i]=new Node(val[i]);</pre>
auto link=[&] (int u,int v) {
Lct::link(ptr[u],ptr[v]);
auto cut=[&](int u,int v){
Lct::cut(ptr[u],ptr[v]);
auto update=[&](int u,int v,const Action:: Tag &val){
Lct: apply(ptr[u],ptr[v],val);
};
auto query=[&](int u,int v){
return Lct::fold(ptr[u],ptr[v]);
                                                     ead3da, 12 lines
"LazyReversibleSplayTree.hpp", "LinkCutTreeBase.hpp"
template < class MonoidAction >
struct LazyLinkCutTree:LinkCutTreeBase<LazyReversibleSplayTree</pre>
    MonoidAction>>{
    using base = LinkCutTreeBase<LazyReversibleSplayTree<
        MonoidAction>>;
    using Ptr = typename base::Ptr;
    using Tag = typename MonoidAction::Tag;
    void apply (Ptr u, Ptr v, const Tag &val) {
        this->evert(u);
        this->expose(v);
        this->propagate(v, val);
};
Number Theory (7)
```

```
ExtendedEuclid.hpp
```

Description: Extended Euclid algorithm for solving diophantine equation (ax + by = gcd(a, b)).

Time: $\mathcal{O}(\log \max\{a, b\})$

```
"../template/Header.hpp"
                                                            229 e7c, 13 lines
pair<11,11> euclid(11 a,11 b) {
    11 x=1, v=0, x1=0, v1=1;
    while(b!=0){
         11 q=a/b;
         x = q * x1;
         y = q * y1;
         a-=q*b;
         swap(x,x1);
         swap(y,y1);
         swap(a,b);
    return {x,y};
```

euclid.h

Description: Finds two integers x and y, such that $ax + by = \gcd(a, b)$. If you just need gcd, use the built in gcd instead. If a and b are coprime, then x is the inverse of a (mod b). $x = x_0 + k * (b/g)$ $y = y_0 - k * (a/g)_{\text{ines}}$

```
ll euclid(ll a, ll b, ll &x, ll &y) {
 if (!b) return x = 1, y = 0, a;
 11 d = euclid(b, a % b, y, x);
 return y -= a/b * x, d;
```

```
Description: Chinese Remainder Theorem.
crt (a, m, b, n) computes x such that x \equiv a \pmod{m}, x \equiv b \pmod{n}. If
|a| < m and |b| < n, x will obey 0 \le x < \operatorname{lcm}(m, n). Assumes mn < 2^{62}
Time: \log(n)
"euclid.h"
                                                                     04d93a, 7 lines
ll crt(ll a, ll m, ll b, ll n) {
  if (n > m) swap(a, b), swap(m, n);
  ll x, y, q = euclid(m, n, x, y);
  assert((a - b) % q == 0); // else no solution
  x = (b - a) % n * x % n / g * m + a;
  return x < 0 ? x + m*n/q : x;
phiFunction.hpp
Description: Euler's \phi function is defined as \phi(n) := \# of positive integers
< n that are coprime with n. \phi(1) = 1, p prime \Rightarrow \phi(p^k) = (p-1)p^{k-1},
m, n \text{ coprime } \Rightarrow \phi(mn) = \phi(m)\phi(n). If n = p_1^{k_1} p_2^{k_2} ... p_r^{k_r} then \phi(n) = p_1^{k_1} p_2^{k_2} ... p_r^{k_r}
(p_1 - 1)p_1^{k_1 - 1}...(p_r - 1)p_r^{k_r - 1}...\phi(n) = n \cdot \prod_{p|n} (1 - 1/p).
\sum_{d|n} \phi(d) = n, \ \sum_{1 \le k \le n, \gcd(k,n)=1} k = n\phi(n)/2, n > 1
Euler's thm: a, n coprime \Rightarrow a^{\phi(n)} \equiv 1 \pmod{n}.
Fermat's little thm: p \text{ prime } \Rightarrow a^{p-1} \equiv 1 \pmod{p} \ \forall a.
                                                                     efae90, 10 lines
const int LIM = 5000000;
int phi[LIM];
void calculatePhi() {
  for(int i=0; i<LIM; ++i) phi[i] = i & 1 ? i : i / 2;</pre>
  for (int i = 3; i < LIM; i += 2)</pre>
     if (phi[i] == i)
        for (int j = i; j < LIM; j += i)</pre>
          phi[j] -= phi[j] / i;
7.1 Prime Numbers
```

LinearSieve.hpp

Description: Prime Number Generator in Linear Time

```
Time: \mathcal{O}(N)
```

```
"../template/Header.hpp"
                                                         194fb1, 15 lines
vi linear sieve(int n) {
 vi prime, composite(n + 1);
  for(int i=2; i<=n; ++i) {</pre>
    if(!composite[i]) {
      prime.emplace back(i);
    for(int j=0; j<(int) prime.size() && i*prime[j]<=n; ++j) {</pre>
      composite[i * prime[j]] = true;
      if(i % prime[j] == 0) {
        break:
 return prime;
```

FastEratosthenes.hpp

Description: Prime sieve for generating all primes smaller than LIM.

Time: LIM=1e9 ≈ 1.5 s

vector<pii> cp;

```
"../template/Header.hpp"
                                                        295b58, 33 lines
const int LIM = 1e6;
bitset<LIM> isPrime;
vi eratosthenes() {
  const int S = (int) round(sqrt(LIM)), R = LIM / 2;
  vi pr = \{2\}, sieve(S + 1);
  pr.reserve(int(LIM/log(LIM) * 1.1));
```

};

```
for(int i=3; i<=S; i+=2) {</pre>
    if(!sieve[i]) {
      cp.emplace back(i, i * i / 2);
      for(int j=i*i; j<=S; j+=2*i) {</pre>
        sieve[j] = 1;
  for(int L=1; L<=R; L+=S) {</pre>
    arrav<bool, S> block{};
    for(auto &[p, idx]: cp) {
      for(int i=idx; i<S+L; idx=(i+=p)) {</pre>
        block[i - L] = 1;
    for(int i=0; i<min(S, R-L); ++i) {</pre>
      if(!block[i]) {
        pr.emplace_back((L + i) \star 2 + 1);
  for(int i: pr) {
    isPrime[i] = 1;
  return pr;
GolbatchConjecture.hpp
Description: Find two prime numbers which sum equals s
Time: \mathcal{O}(N \log N)
"FastEratosthenes.hpp"
                                                          88fb23, 18 lines
pair<int, int> goldbatchConjecture(int s, vi pr = {}) {
  if (s <= 2 || s % 2 != 0) {
    return make_pair(-1, -1);
  if (pr.size() == 0) {
    pr = eratosthenes();
  for (auto x : pr) {
    if (x > s / 2) {
      break;
    int d = s - x;
    if (binary_search(pr.begin(), pr.end(), d)) {
      return make_pair(min(x, d), max(x, d));
  return make_pair(-1, -1);
HopcroftKarp.hpp
Description: Fast bipartite matching algorithm.
Time: \mathcal{O}\left(E\sqrt{V}\right)
"../template/Header.hpp'
                                                          0bd56f, 52 lines
    int n,m;
    vi l,r,lv,ptr;
```

Graph (8)

8.1 Matching

```
struct HopcroftKarp{
    vector<vi> adj;
    HopcroftKarp(){}
   HopcroftKarp(int _n,int _m) {init(_n,_m);}
   void init(int _n,int _m) {
       n=_n, m=_m;
       adj.assign(n+m, vi{});
```

```
void addEdge(int u,int v) {
        adj[u].emplace back(v+n);
    void bfs() {
        lv=vi(n,-1);
        queue<int> q;
         for (int i=0; i<n; i++) if (1[i]==-1) {
             lv[i]=0;
             q.emplace(i);
         while(!q.empty()){
             int u=q.front();
             q.pop();
             for(int v:adj[u])if(r[v]!=-1&&lv[r[v]]==-1){
                 lv[r[v]]=lv[u]+1;
                 q.emplace(r[v]);
    bool dfs(int u) {
         for(int &i=ptr[u];i<sz(adj[u]);i++) {</pre>
             int v=adj[u][i];
             if (r[v] ==-1||(lv[r[v]] == lv[u] +1&&dfs(r[v]))){
                 l[u]=v,r[v]=u;
                  return true;
         return false;
    int maxMatching() {
         int match=0, cnt=0;
         l=r=vi(n+m,-1);
             ptr=vi(n);
             bfs();
             cnt=0:
             for (int i=0; i<n; i++) if (l[i] ==-1&&dfs(i)) cnt++;</pre>
             match+=cnt;
         }while(cnt);
         return match;
};
Description: Kuhn Algorithm to find maximum bipartite matching or find
augmenting path in bipartite graph.
Time: \mathcal{O}(VE)
"../template/Header.hpp"
                                                         fc7d17, 15 lines
vi adj[1010], match(1010, -1);
bitset<1010> visited;
bool kuhn(int u) {
  if(visited[u]) {
    return false;
  visited[u] = true;
  for(auto x: adj[u]) {
    if (match[x] == -1 \mid \mid kuhn(match[x]))  {
      match[x] = u;
      return true;
 return false;
```

8.2 Network Flow

Dinic.hpp

Description: Dinic's Algorithm for finding the maximum flow.

Time: $\mathcal{O}(VE \log U)$ where U is the maximum flow. 2b9ab1, 88 lines

```
template<class T, bool directed=true, bool scaling=true>
struct Dinic{
    static constexpr T INF=numeric limits<T>::max()/2;
    struct Edge{
        int to:
        T flow, cap;
        Edge(int _to,T _cap):to(_to),flow(0),cap(_cap){}
        T remain() { return cap-flow; }
   int n,s,t;
    T U;
    vector<Edge> e;
    vector<vector<int>> q;
    vector<int> ptr.lv;
    bool calculated;
    T max_flow;
    Dinic(){}
    Dinic(int n, int s, int t) {init(n, s, t);}
    void init(int _n,int _s,int _t){
        n=_n,s=_s,t=_t;
        U=0:
        e.clear();
        g.assign(n,{});
        calculated=false;
    void add_edge(int from,int to,T cap){
        assert(0<=from&&from<n&&0<=to&&to<n);
        g[from].emplace_back(e.size());
        e.emplace_back(to,cap);
        g[to].emplace_back(e.size());
        e.emplace_back(from,directed?0:cap);
        U=max(U,cap);
    bool bfs(T scale) {
        lv.assign(n,-1);
        vector<int> q{s};
        lv[s]=0;
        for (int i=0; i < (int) q.size(); i++) {</pre>
            int u=q[i];
            for(int j:q[u]){
                int v=e[j].to;
                if(lv[v] ==-1&&e[j].remain()>=scale){
                     q.emplace_back(v);
                     lv[v]=lv[u]+1;
        return lv[t]!=-1;
    T dfs(int u,int t,T f) {
        if (u==t | |f==0) return f;
        for(int &i=ptr[u];i<(int)q[u].size();i++){</pre>
            int j=q[u][i];
            int v=e[j].to;
            if(lv[v]==lv[u]+1){
                T res=dfs(v,t,min(f,e[j].remain()));
                if(res>0){
                     e[i].flow+=res;
                     e[j^1].flow-=res;
                     return res;
        return 0;
```

MinCostFlow BinaryOptimization

```
T flow(){
        if(calculated)return max flow;
        calculated=true:
        max flow=0;
        for(T scale=scaling?1LL<<(63-__builtin_clzll(U)):1LL;</pre>
             scale>0;scale>>=1){
             while (bfs(scale)) {
                ptr.assign(n,0);
                 while(true) {
                     T f=dfs(s,t,INF);
                     if (f==0) break;
                     max_flow+=f;
        return max_flow;
    pair<T, vector<int>> cut() {
        flow();
        vector<int> res(n);
        for(int i=0;i<n;i++)res[i]=(lv[i]==-1);</pre>
        return {max flow,res};
};
MinCostFlow.hpp
Description: minimum-cost flow algorithm.
Time: \mathcal{O}(FE \log V) where F is max flow.
"../template/Header.hpp"
                                                       8ea1d2, 83 lines
template < class F, class C>
struct MinCostFlow{
    struct Edge {
        int to;
        F flow, cap;
        C cost;
        Edge(int _to,F _cap,C _cost):to(_to),flow(0),cap(_cap),
             cost(cost){}
        F getcap(){
             return cap-flow;
    };
    int n;
    vector<Edge> e;
    vector<vi> adj;
    vector<C> pot, dist;
    vi pre;
    bool neg;
    const F FINF=numeric limits<F>::max()/2;
    const C CINF=numeric_limits<C>::max()/2;
    MinCostFlow() {}
    MinCostFlow(int _n) {
        init(_n);
    void init(int _n){
        e.clear();
        adj.assign(n,{});
        neg=false;
    void addEdge(int u,int v,F cap,C cost){
        adj[u].emplace_back(sz(e));
        e.emplace_back(v,cap,cost);
        adj[v].emplace back(sz(e));
        e.emplace_back(u,0,-cost);
        if(cost<0)neg=true;</pre>
    bool dijkstra(int s,int t) {
        using P = pair<C,int>;
```

```
dist.assign(n,CINF);
        pre.assign(n,-1);
        priority_queue<P, vector<P>, greater<P>> pq;
        dist[s]=0;
        pq.emplace(0,s);
         while(!pq.empty()){
             auto [d,u]=pq.top();
             pq.pop();
             if (dist[u] < d) continue;</pre>
             for(int i:adj[u]){
                  int v=e[i].to;
                  C ndist=d+pot[u]-pot[v]+e[i].cost;
                  if(e[i].getcap()>0&&dist[v]>ndist){
                       pre[v]=i;
                       dist[v]=ndist;
                       pg.emplace(ndist, v);
         return dist[t] < CINF;
    pair<F,C> flow(int s,int t) {
        F flow=0;
        C cost=0;
        pot.assign(n,0);
         if (neq) for (int t=0; t< n; t++) for (int i=0; i< sz(e); i++) if (e
              [i].getcap()>0){
              int u=e[i^1].to, v=e[i].to;
             pot[v]=min(pot[v],pot[u]+e[i].cost);
         } // Bellman-Ford
         while(dijkstra(s,t)){
             for (int i=0; i<n; i++) pot[i] += dist[i];</pre>
             F aug=FINF;
              for (int u=t; u!=s; u=e[pre[u]^1].to) {
                  aug=min(aug,e[pre[u]].getcap());
              } // find bottleneck
             for(int u=t;u!=s;u=e[pre[u]^1].to){
                  e[pre[u]].flow+=aug;
                  e[pre[u]^1].flow-=aug;
             } // push flow
              flow+=aug;
             cost+=aug*pot[t];
         return {flow,cost};
};
BinaryOptimization.hpp
Description: Binary Optimization.
                                           minimize \kappa + \sum_{i} \theta_{i}(x_{i}) +
\sum_{i < j} \phi_{ij}(x_i, x_j) + \sum_{i < j < k} \psi_{ijk}(x_i, x_j, x_k) \text{ where } x_i \in \{0, 1\} \text{ and } \phi_{ij}, \psi_{ijk}
are submodular functions. a set function f is submodular if f(S) + f(T) >
f(S \cap T) + f(S \cup T) for all S, T. \phi_{ij}(0,1) + \phi_{ij}(1,0) \ge \phi_{ij}(1,1) + \phi_{ij}(0,0).
template < class T, bool minimize = true >
struct BinaryOptimization{
    static constexpr T INF=numeric_limits<T>::max()/2;
    int n,s,t,node_id;
    T base;
    map<pair<int,int>,T> edges;
    BinaryOptimization(int _n):n(_n),s(n),t(n+1),node_id(n+2),
         base(0){}
    void add_edge(int u,int v,T w) {
        assert (w \ge 0);
        if (u==v | |w==0) return;
         auto &e=edges[{u,v}];
        e=min(e+w,INF);
    void add0(T w){
```

```
base+=w:
void add1(int i,T a,T b){
    if(a<=b){
        add0(a);
        add_edge(s,i,b-a);
    }else{
        add0(b);
        add_edge(i,t,a-b);
void add1(int i,T x0,T x1){
    assert(0<=i&&i<n);
    if (!minimize) x0=-x0, x1=-x1;
    _{add1(i,x0,x1);}
void _add2(int i,int j,T a,T b,T c,T d) {
    assert (b+c>=a+d);
    add0(a);
    _add1(i,0,c-a);
    _add1(j,0,d-c);
    add_edge(i,j,b+c-a-d);
void add2(int i,int j,T x00,T x01,T x10,T x11){
    assert(i!=j&&0<=i&&i<n&&0<=j&&j<n);
    if (!minimize) x00=-x00, x01=-x01, x10=-x10, x11=-x11;
    _{add2}(i, j, x00, x01, x10, x11);
void _add3(int i,int j,int k,T a,T b,T c,T d,T e,T f,T q,T
    T p=a+d+f+g-b-c-e-h;
    if(p>=0){
        add0(a);
        _add1(i,0,f-b);
        _add1(j,0,g-e);
        _add1(k,0,d-c);
        _{add2}(i, j, 0, c+e-a-q, 0, 0);
        \_add2(i,k,0,0,b+e-a-f,0);
        add2(j,k,0,b+c-a-d,0,0);
        int u=node id++;
        add0(-p);
        add edge(i,u,p);
        add_edge(j,u,p);
        add edge(k,u,p);
        add_edge(u,t,p);
    }else{
        add0(h);
        _add1(i,c-q,0);
        add1(i,b-d,0);
        add1(k,e-f,0);
        _{add2}(i, j, 0, 0, d+f-b-h, 0);
        _{add2}(i,k,0,d+g-c-h,0,0);
        add2(j,k,0,0,f+g-e-h,0);
        int u=node_id++;
        add0(p);
        add_edge(s,u,-p);
        add_edge(u,i,-p);
        add_edge(u,j,-p);
        add_edge(u,k,-p);
void add3(int i,int j,int k,T x000,T x001,T x010,T x011,T
     x100, T x101, T x110, T x111) {
    assert (i!=j\&\&j!=k\&\&k!=i\&\&0<=i\&\&i<n\&\&0<=j\&\&j<n\&\&0<=k\&\&k<
         n);
    if(!minimize){
        x000=-x000, x001=-x001, x010=-x010, x011=-x011;
        x100=-x100, x101=-x101, x110=-x110, x111=-x111;
```

KaryOptimization SCC LowLink HLD

```
_add3(i,j,k,x000,x001,x010,x011,x100,x101,x110,x111);
    T solve(){
        Dinic<T> dinic(node_id,s,t);
         for(auto &[p,w]:edges){
             auto [u, v]=p;
             dinic.add_edge(u,v,w);
        T ans=dinic.flow()+base;
         return minimize?ans:-ans;
};
KaryOptimization.hpp
Description: k-ary Optimization.
                                          minimize \kappa + \sum_{i} \theta_{i}(x_{i}) +
\sum_{i < j} \phi_{ij}(x_i, x_j) where x_i \in \{0, 1, \dots, k-1\} and \phi_{i,j} is monge. A func-
tion f is monge if f(a,c) + f(b,d) \leq f(a,d) + f(b,c) for all a < b and
c < d. \phi_{ij}(x-1,y) + \phi_{ij}(x,y+1) \le \phi_{ij}(x-1,y+1) + \phi_{ij}(x,y).
\phi_{ij}(x,y) + \phi_{ij}(x-1,y+1) - \phi_{ij}(x-1,y) - \phi_{ij}(x,y+1) \ge 0.
                                                           422f8a, 88 lines
template < class T, bool minimize = true >
struct K aryOptimization{
    static constexpr T INF=numeric_limits<T>::max()/2;
    int n,s,t,node_id;
    T base;
    vector<int> ks;
    vector<vector<int>> id;
    map<pair<int,int>,T> edges;
    K aryOptimization(int n.int k) {init(vector<int>(n.k));}
    K_aryOptimization(const vector<int> &_ks) {init(_ks);}
    void init(const vector<int> &_ks){
        ks= ks;
        n=ks.size();
        s=0, t=1, node id=2;
        base=0;
        id.clear();
        edges.clear();
         for(auto &k:ks) {
             assert (k>=1);
             vector<int> a(k+1);
             a[0]=s, a[k]=t;
             for (int i=1; i < k; i++) a[i] = node_id++;</pre>
             id.emplace back(a);
             for(int i=2;i<k;i++)add_edge(a[i],a[i-1],INF);</pre>
    void add_edge(int u,int v,T w) {
        assert (w \ge 0):
        if (u==v | |w==0) return;
        auto &e=edges[{u,v}];
        e=min(e+w,INF);
    void add0(T w){
        base+=w:
    void _add1(int i, vector<T> cost) {
        add0(cost[0]);
         for(int j=1; j<ks[i]; j++) {</pre>
             T x=cost[j]-cost[j-1];
             if(x>0)add_edge(id[i][j],t,x);
             if (x<0) add0 (x), add_edge(s,id[i][j],-x);</pre>
    void add1(int i, vector<T> cost) {
        assert (0<=i&&i<n&& (int) cost.size() ==ks[i]);
        if(!minimize) for(auto &x:cost) x=-x;
         _add1(i,cost);
```

```
void _add2(int i,int j,vector<vector<T>> cost){
         int h=ks[i], w=ks[j];
         _add1(j,cost[0]);
         for (int x=h-1; x>=0; x--) for (int y=0; y<w; y++) cost[x][y]-=</pre>
              cost[0][y];
         vector<T> a(h);
         for (int x=0; x<h; x++) a[x] = cost[x][w-1];</pre>
         _add1(i,a);
         for (int x=0; x<h; x++) for (int y=0; y<w; y++) cost[x][y]-=a[x</pre>
              1;
         for (int x=1; x<h; x++) {</pre>
             for (int y=0; y<w-1; y++) {</pre>
                  T = cost[x][y] + cost[x-1][y+1] - cost[x-1][y] - cost
                       [x][y+1];
                 assert (w>=0); // monge
                 add_edge(id[i][x],id[j][y+1],w);
        }
    void add2(int i,int j,vector<vector<T>> cost){
        assert (0 \le i \& i \le n \& \& 0 \le j \& \& j \le n \& \& i! = j);
         assert((int)cost.size() == ks[i]);
         for(auto &v:cost)assert((int)v.size()==ks[i]);
        if(!minimize) for(auto &v:cost) for(auto &x:v) x=-x;
         _add2(i,j,cost);
    pair<T, vector<int>> solve() {
        Dinic<T> dinic(node id,s,t);
         for(auto &[p,w]:edges){
             auto [u, v] =p;
             dinic.add_edge(u,v,w);
         auto [val,cut]=dinic.cut();
        val+=base;
        if(!minimize) val=-val;
        vector<int> ans(n);
         for (int i=0; i < n; i++) {</pre>
             ans[i]=ks[i]-1;
             for(int j=1; j<ks[i]; j++) ans[i]-=cut[id[i][j]];</pre>
         return {val, ans};
};
8.3
        Connectivity
SCC.hpp
Description: Strongly Connected Component.
"../template/Header.hpp"
                                                          82a9d1, 34 lines
template<class G>
pair<int, vector<int>>> strongly_connected_component(G &g) {
    static_assert(G::is_directed);
    int n=q.n;
    vector<int> disc(n,-1), low(n), scc(n,-1);
    stack<int> st;
    vector<bool> in_st(n);
    int t=0,scc_cnt=0;
    function<void(int,int)> dfs=[&](int u,int p){
        disc[u]=low[u]=t++;
        st.emplace(u);
         in_st[u]=true;
        for(int v:q[u]){
             if(disc[v] ==-1) {
                 dfs(v,u);
                  low[u] = min(low[u], low[v]);
             }else if(in_st[v]){
                  low[u]=min(low[u],disc[v]);
```

```
if(disc[u] == low[u]) {
             while(true){
                 int v=st.top();
                 st.pop();
                 in_st[v]=false;
                 scc[v]=scc_cnt;
                 if (v==u) break;
            scc_cnt++;
    };
    for(int i=0;i<n;i++)if(disc[i]==-1)dfs(i,-1);</pre>
    return {scc_cnt,scc};
LowLink.hpp
Description: Low Link.
                                                       f4ad2f, 33 lines
template<class G>
struct LowLink{
    G &q;
    int n;
    vector<int> disc, low, par, ord;
    vector<pair<int,int>> bridge;
    vector<int> articulation;
    int t=0:
    LowLink(G &_g):g(_g), n(g.n), disc(n,-1), low(n), par(n,-1) {
        for (int i=0; i<n; i++) if (disc[i]==-1) dfs(i);</pre>
    void dfs(int u) {
        disc[u]=low[u]=t++;
        ord.emplace_back(u);
        int child=0;
        bool found_par=false;
        for(int v:q[u]){
             if (disc[v] ==-1) {
                 par[v]=u;
                 dfs(v);
                 low[u]=min(low[u],low[v]);
                 if(low[v]>disc[u])bridge.emplace_back(u,v);
                 if(par[u]!=-1&&low[v]>=disc[u])articulation.
                      emplace_back(u);
                 child++;
             }else if(v!=par[u]||found par){
                 low[u]=min(low[u],disc[v]);
             }else{
                 found par=true;
        if (par[u] ==-1&&child>1) articulation.emplace_back(u);
};
Tree (9)
HLD.hpp
Description: HLD
"../template/Header.hpp"
                                                       cf6882, 45 lines
vector<vi> adj;
vector<int> sz, lvl, hv, hd, p, disc;
int t;
void dfs(int u, int parent) {
  sz[u] = 1;
  lvl[u] = lvl[parent] + 1;
  p[u] = parent;
  int c_hv=0, c_max=0;
  for(auto v: adj[u]) {
```

if(v == parent) continue;

if(used[v]) continue;

decom(v);

```
dfs(v, u);
    sz[u] += sz[v];
    if(c_max < sz[v]) {</pre>
      c hv = v;
      c_max = sz[v];
  hv[u] = c_hv;
void hld(int u, int parent) {
  if(hd[u] == 0) {
   hd[u] = u;
  disc[u] = ++t;
  if(hv[u] != 0) {
   hd[hv[u]] = hd[u];
   hld(hv[u], u);
  for(auto v: adj[u]) {
   if(v == parent || v == hv[u]) {
      continue;
   hld(v, u);
int lca(int u, int v) {
  while(hd[u] != hd[v]) {
   if(lvl[hd[u]] > lvl[hd[v]]) swap(u, v);
    v=p[hd[v]];
  return lvl[u] < lvl[v] ? u: v;</pre>
Centroid Decom. hpp
Description: Centroid
"../template/Header.hpp"
                                                      e46d44. 32 lines
vector<vi> adj;
vi sz;
vector<bool> used;
int find_size(int u, int p) {
  sz[u] = 1;
  for(auto v: adj[u]) {
    if(v == p || used[v]) continue;
    sz[u] += find_size(v, u);
  return sz[u];
int find_cen(int u, int p, int t) {
  for(auto v: adj[u]) {
   if(v == p || used[v]) continue;
    if(sz[v] * 2 > t) find_cen(v, u, t);
  return u;
void decom(int u) {
  u = find_cen(u, 0, find_size(u, 0));
  used[u] = true;
  for(auto v: adi[u]) {
    // dfs do something
  for(auto v: adj[u]) {
```

```
Polynomials (10)
FormalPowerSeries.hpp
Description: basic operations of formal power series
                                                      416433, 136 lines
template < class mint >
struct FormalPowerSeries:vector<mint>{
    using vector<mint>::vector;
    using FPS = FormalPowerSeries;
    FPS & operator += (const FPS &rhs) {
        if (rhs.size()>this->size())this->resize(rhs.size());
        for(int i=0;i<rhs.size();i++)(*this)[i]+=rhs[i];</pre>
        return *this;
    FPS & operator += (const mint &rhs) {
        if (this->empty())this->resize(1);
        (*this) [0] += rhs;
        return *this;
    FPS & operator -= (const FPS & rhs) {
        if (rhs.size()>this->size())this->resize(rhs.size());
        for(int i=0;i<rhs.size();i++)(*this)[i]-=rhs[i];</pre>
        return *this;
    FPS & operator -= (const mint &rhs) {
        if (this->empty())this->resize(1);
        (*this) [0]-=rhs;
        return *this;
    FPS & operator *= (const FPS & rhs) {
        auto res=NTT<mint>()(*this,rhs);
        return *this=FPS(res.begin(),res.end());
    FPS & operator *= (const mint &rhs) {
        for(auto &a:*this)a*=rhs;
        return *this;
    friend FPS operator+ (FPS lhs, const FPS &rhs) {return lhs+=
    friend FPS operator+(FPS lhs, const mint &rhs) {return lhs+=
    friend FPS operator+(const mint &lhs, FPS &rhs) {return rhs+=
    friend FPS operator-(FPS lhs,const FPS &rhs) {return lhs-=
    friend FPS operator-(FPS lhs, const mint &rhs) {return lhs-=
    friend FPS operator-(const mint &lhs, FPS rhs) {return -(rhs-
         lhs);}
    friend FPS operator* (FPS lhs, const FPS &rhs) {return lhs*=
    friend FPS operator* (FPS lhs, const mint &rhs) {return lhs*=
    friend FPS operator* (const mint &lhs, FPS rhs) {return rhs*=
         lhs; }
    FPS operator-() {return (*this) *-1;}
    FPS rev(){
```

FPS res(*this);

return res;

reverse(res.beign(), res.end());

```
FPS pre(int sz) {
    FPS res(this->begin(),this->begin()+min((int)this->size
    if(res.size() < sz) res.resize(sz);</pre>
    return res;
FPS shrink(){
    FPS res(*this);
    while(!res.empty()&&res.back() ==mint{})res.pop_back();
    return res:
FPS operator>>(int sz){
    if (this->size()<=sz) return { };</pre>
    FPS res(*this);
    res.erase(res.begin(), res.begin()+sz);
    return res;
FPS operator<<(int sz) {</pre>
    FPS res(*this);
    res.insert(res.begin(),sz,mint{});
    return res;
FPS diff(){
    const int n=this->size();
    FPS res(max(0,n-1));
    for(int i=1; i < n; i++) res[i-1] = (*this)[i] *mint(i);</pre>
    return res;
FPS integral(){
    const int n=this->size();
    FPS res(n+1);
    res[0]=0;
    if(n>0)res[1]=1;
    11 mod=mint::get_mod();
    for(int i=2;i<=n;i++)res[i]=(-res[mod%i])*(mod/i);</pre>
    for (int i=0; i<n; i++) res[i+1] *=(*this)[i];</pre>
    return res;
mint eval(const mint &x) {
    mint res=0,w=1;
    for(auto &a:*this) res+=a*w, w*=x;
    return res;
FPS inv(int deg=-1){
    assert(!this->empty()&&(*this)[0]!=mint(0));
    if (deg==-1) deg=this->size();
    FPS res{mint(1)/(*this)[0]};
    for(int i=2;i>>1<deq;i<<=1) {</pre>
        res=(res*(mint(2)-res*pre(i))).pre(i);
    return res.pre(deg);
FPS log(int deg=-1){
    assert (!this->empty() && (*this)[0] ==mint(1));
    if (deg==-1) deg=this->size();
    return (pre(deg).diff()*inv(deg)).pre(deg-1).integral()
FPS exp(int deg=-1) {
    assert (this->empty() | | (*this) [0] == mint(0));
    if (deg==-1) deg=this->size();
    FPS res{mint(1)};
    for(int i=2;i>>1<deq;i<<=1){</pre>
        res=(res*(pre(i)-res.log(i)+mint(1))).pre(i);
    return res.pre(deg);
FPS pow(ll k, int deg=-1) {
```

FFT NTT Manacher SuffixArray

```
const int n=this->size();
        if (deg==-1) deg=n;
        if(k==0){
             FPS res(deg);
             if (deg) res[0]=mint(1);
             return res;
        for(int i=0;i<n;i++){</pre>
             if (__int128_t(i) *k>=deg) return FPS(deg, mint(0));
             if ((*this)[i] ==mint(0)) continue;
             mint rev=mint(1)/(*this)[i];
             FPS res=(((*this*rev)>>i).log(deg)*k).exp(deg);
             res=((res*binpow((*this)[i],k))<<(i*k)).pre(deg);
             return res;
        return FPS(deg,mint(0));
};
using FPS=FormalPowerSeries<mint>;
FFT.hpp
Description: Fast Fourier transform
Time: \mathcal{O}(N \log N)
"../template/Header.hpp"
                                                         5d476b, 73 lines
template<class T=11,int mod=0>
struct FFT{
  using vt = vector<T>;
  using cd = complex<db>;
  using vc = vector<cd>;
  static const bool INT=true;
  static void fft(vc &a){
    int n=a.size(),L=31- builtin clz(n);
    vc rt(n);
    rt[1]=1;
    for(int k=2; k<n; k*=2) {</pre>
      cd z=polar(db(1),PI/k);
      for(int i=k;i<2*k;i++)rt[i]=i&1?rt[i/2]*z:rt[i/2];</pre>
    vi rev(n);
    for(int i=1;i<n;i++)rev[i]=(rev[i/2]|(i&1)<<L)/2;</pre>
    for (int i=1; i<n; i++) if (i<rev[i]) swap(a[i], a[rev[i]]);</pre>
    for (int k=1; k < n; k \neq 2) for (int i=0; i < n; i+2 \neq k) for (int j=0; j < k
         ; j++) {
      cd z=rt[j+k]*a[i+j+k];
      a[i+j+k]=a[i+j]-z;
      a[i+j]+=z;
  template<class U>
  static db norm (const U &x) {
    return INT?round(x):x;
  static vt conv(const vt &a,const vt &b) {
    if(a.empty()||b.empty())return {};
    vt res(a.size()+b.size()-1);
    int L=32-__builtin_clz(res.size()),n=1<<L;</pre>
    vc in(n), out(n);
    copy(a.begin(),a.end(),in.begin());
    for(int i=0;i<b.size();i++)in[i].imag(b[i]);</pre>
    fft (in):
    for(auto &x:in)x*=x;
    for (int i=0; i<n; i++) out[i]=in[-i&(n-1)]-conj(in[i]);</pre>
    fft (out);
    for(int i=0;i<res.size();i++)res[i]=norm(imag(out[i])/(4*n)</pre>
    return res;
```

```
static vl convMod(const vl &a,const vl &b) {
    assert (mod>0);
    if(a.emptv()||b.emptv())return {};
    vl res(a.size()+b.size()-1);
    int L=32-__builtin_clz(res.size()), n=1<<L;</pre>
    11 cut=int(sqrt(mod));
    vc in1(n),in2(n),out1(n),out2(n);
    for(int i=0;i<a.size();i++)in1[i]=cd(ll(a[i])/cut,ll(a[i])%</pre>
         cut); // a1 + i * a2
    for (int i=0; i < b. size(); i++) in2[i] = cd(ll(b[i]) / cut, ll(b[i]) %</pre>
         cut); // b1 + i * b2
    fft(in1), fft(in2);
    for (int i=0; i<n; i++) {</pre>
      int j=-i&(n-1);
      out1[j]=(in1[i]+conj(in1[j]))*in2[i]/(2.1*n); // f1 * (g1)
             + i * q2) = f1 * q1 + i f1 * q2
      out2[j]=(in1[i]-conj(in1[j]))*in2[i]/cd(0.1,2.1*n); // f2
             * (q1 + i * q2) = f2 * q1 + i f2 * q2
    fft(out1),fft(out2);
    for(int i=0;i<res.size();i++){</pre>
      11 x=round(real(out1[i])), y=round(imag(out1[i]))+round(
            real(out2[i])), z=round(imag(out2[i]));
      res[i]=((xmod*cut+y)mod*cut+z)mod; // a1*b1*cut^2
            + (a1 * b2 + a2 * b1) * cut + a2 * b2
    return res;
  vt operator()(const vt &a,const vt &b){
    return mod>0?conv(a,b):convMod(a,b);
};
template<>
struct FFT<db>{
  static const bool INT=false;
NTT.hpp
Description: Number theoretic transform
Time: \mathcal{O}(N \log N)
"../template/Header.hpp", "../modular-arithmetic/BinPow.hpp",
"../modular-arithmetic/MontgomeryModInt.hpp"
                                                         2b2392, 39 lines
template < class mint = mint >
struct NTT{
  using vm = vector<mint>;
  static constexpr mint root=mint::get_root();
    static assert(root!=0);
  static void ntt(vm &a) {
    int n=a.size(),L=31-__builtin_clz(n);
    vm rt(n);
    rt[1]=1;
    for (int k=2, s=2; k < n; k *=2, s++) {
      mint z[]={1,binpow(root,MOD>>s)};
      for(int i=k;i<2*k;i++)rt[i]=rt[i/2]*z[i&1];</pre>
    vi rev(n);
    for (int i=1; i<n; i++) rev[i] = (rev[i/2] | (i&1) <<L) /2;</pre>
    for (int i=1; i < n; i++) if (i < rev[i]) swap (a[i], a[rev[i]]);</pre>
    for (int k=1; k < n; k *=2) for (int i=0; i < n; i+=2*k) for (int j=0; j < k
         ; j++) {
      mint z=rt[j+k]*a[i+j+k];
      a[i+j+k]=a[i+j]-z;
      a[i+j] += z;
  static vm conv(const vm &a,const vm &b) {
    if(a.empty()||b.empty())return {};
```

```
int s=a.size()+b.size()-1,n=1<<(32-_builtin_clz(s));
   mint inv=mint(n).inv();
   vm inl(a),in2(b),out(n);
   inl.resize(n),in2.resize(n);
   ntt(inl),ntt(in2);
   for(int i=0;i<n;i++)out[-i&(n-1)]=in1[i]*in2[i]*inv;
   ntt(out);
   return vm(out.begin(),out.begin()+s);
}
vm operator()(const vm &a,const vm &b){
   return conv(a,b);
}
</pre>
```

Strings (11)

Manacher.hpp

Description: Manacher's Algorithm. pal[i] := the length of the longest palindrome centered at i/2.
"../template/Header.hpp"

53856e. 15 lines

```
template < class STR>
vector < int > manacher (const STR &s) {
    int n= (int) s. size();
    if (n==0) return {};
    vector < int > pal (2*n-1);
    for (int p=0, l=-1, r=-1; p<2*n-1; p++) {
        int i= (p+1) >> 1, j=p>> 1;
        int k= (i>=r?0:min (r-i, pal [2*(l+r)-p]));
        while (j+k+1<n&&i-k-1>=0&&s[j+k+1]==s[i-k-1])k++;
        pal[p]=k;
        if (j+k>r) l=i-k, r=j+k;
    }
    for (int i=0; i<2*n-1; i++) pal[i]=pal[i] <<1| (i&1^1);
    return pal;
}</pre>
```

SuffixArray.hpp

Description: Suffix Automaton.

```
"../data-structure/SparseTable.hpp", "../group/monoid/Min.hpp" b9cfb1, 39 lines
```

```
template<class STR>
struct SuffixArray{
    int n;
    vector<int> sa,isa,lcp;
    SparseTable<MinMonoid<int>> st;
    SuffixArray(){}
    SuffixArray(const STR &s) {init(s);}
    void init(const STR &s){
        n=(int)s.size();
        sa=isa=lcp=vector<int>(n+1);
        sa[0]=n;
        iota(sa.begin()+1,sa.end(),0);
        sort(sa.begin()+1, sa.end(), [&](int i, int j) {return s[i
             ]<s[j];});
        for (int i=1; i<=n; i++) {</pre>
             int x=sa[i-1],y=sa[i];
             isa[y]=i>1&&s[x]==s[y]?isa[x]:i;
        for (int len=1; len<=n; len<<=1) {</pre>
             vector<int> ps(sa),pi(isa),pos(n+1);
             iota(pos.begin(),pos.end(),0);
             for(auto i:ps) if((i-=len)>=0) sa[pos[isa[i]]++]=i;
             for(int i=1;i<=n;i++) {</pre>
                 int x=sa[i-1],y=sa[i];
                 isa[y]=pi[x]==pi[y]\&\&pi[x+len]==pi[y+len]?isa[x
                     ]:i;
```

ZAlgo.hpp

Description: Z Algorithm. z[i] := the length of the longest common prefix between s and <math>s[i:].

"../template/Header.hpp" b93726, 12 line

```
template < class STR>
vector < int > z_algorithm(const STR &s) {
    int n = (int) s.size();
    vector < int > z(n);
    z[0] = n;
    for (int i = 1, 1 = 0, r = 1; i < n; i + +) {
        if (i < r) z[i] = min(r - i, z[i - 1]);
        while (i + z[i] < n&&s[z[i]] = = s[i + z[i]]) z[i] + +;
        if (i + z[i] > r) l = i, r = i + z[i];
    }
    return z;
}
```

PrefixFunction.hpp

Description: Prefix function. pi[i] := the length of the longest proper prefix of <math>s[0:i] which is also a suffix of s[0:i].

```
template < class STR>
vector < int > prefix_function(const STR &s) {
    int n=(int)s.size();
    vector < int > pi(n);
    for(int i=1,j=0;i<n;i++) {
        while(j>0&&s[i]!=s[j])j=pi[j-1];
        if(s[i]==s[j])j++;
        pi[i]=j;
    }
    return pi;
}
```

SuffixAutomaton.hpp

Description: Suffix Automaton. Find whether a string t is a substring of a string s by traversing the automaton. Find whether a string t is a suffix of a string s by checking whether the last node is a terminal node. Find the number of distinct substrings of a string s by calculating the number of distinc path using DP. Count the number of occurrences of string t in string s. Let p be the node we end up at after traversing t in the automaton. The answer is the number of paths from p to terminal nodes. Find first occurrence of string t in string s by calculating the longest path in the automaton after reaching node p.

```
template < class STR>
struct SuffixAutomaton{
    using T = typename STR::value_type;
    struct Node{
        map<T,int> nxt;
        int link,len;
        Node(int link,int len):link(link),len(len){}
};
    vector<Node> nodes;
int last;
SuffixAutomaton():nodes{Node(-1,0)},last(0){}
```

```
SuffixAutomaton(const STR &s):SuffixAutomaton() {
        for(auto c:s)extend(c);
    int new node(int link,int len) {
        nodes.emplace_back(Node(link,len));
        return (int) nodes.size()-1;
    void extend(T c){
        int cur=new_node(0, nodes[last].len+1);
        int p=last;
        while (p!=-1&&!nodes[p].nxt.count(c)) {
            nodes[p].nxt[c]=cur;
            p=nodes[p].link;
        if(p!=-1){
            int q=nodes[p].nxt[c];
            if (nodes[p].len+1==nodes[q].len) {
                nodes[cur].link=q;
                int r=new_node(nodes[q].link,nodes[p].len+1);
                nodes[r].nxt=nodes[q].nxt;
                while (p! =-1&&nodes [p].nxt [c] ==q) {
                    nodes[p].nxt[c]=r;
                    p=nodes[p].link;
                nodes[q].link=nodes[cur].link=r;
        last=cur;
    11 distinct_substrings(){
        11 res=0;
        for(int i=1;i<(int)nodes.size();i++){</pre>
            res+=nodes[i].len-nodes[nodes[i].link].len;
        return res;
};
```

Geometry (12)

12.1 Geometric primitives

Point.h

Description: Class to handle points in the plane. T can be e.g. double or long long. (Avoid int.)

```
template \langle class T \rangle int sqn(T x) \{ return (x > 0) - (x < 0); \}
template<class T>
struct Point {
 typedef Point P;
 explicit Point (T x=0, T y=0) : x(x), y(y) {}
 bool operator<(P p) const { return tie(x,y) < tie(p.x,p.y); }</pre>
 bool operator==(P p) const { return tie(x,y)==tie(p.x,p.y); }
 P operator+(P p) const { return P(x+p.x, y+p.y); }
 P operator-(P p) const { return P(x-p.x, y-p.y); }
 P operator*(T d) const { return P(x*d, y*d); }
 P operator/(T d) const { return P(x/d, y/d); }
 T dot(P p) const { return x*p.x + y*p.y; }
 T cross(P p) const { return x*p.y - y*p.x; }
 T cross(P a, P b) const { return (a-*this).cross(b-*this); }
 T dist2() const { return x*x + y*y; }
 double dist() const { return sqrt((double)dist2()); }
  // angle to x-axis in interval [-pi, pi]
 double angle() const { return atan2(y, x); }
 P unit() const { return *this/dist(); } // makes dist()=1
 P perp() const { return P(-y, x); } // rotates +90 degrees
```

```
P normal() const { return perp().unit(); }
// returns point rotated 'a' radians ccw around the origin
P rotate(double a) const {
  return P(x*cos(a)-y*sin(a),x*sin(a)+y*cos(a)); }
friend ostream& operator<<(ostream& os, P p) {
  return os << "(" << p.x << "," << p.y << ")"; }
;</pre>
```

lineDistance.h

Description:

Returns the signed distance between point p and the line containing points a and b. Positive value on left side and negative on right as seen from a towards b. a==b gives nan. P is supposed to be Point<T> or Point3D<T> where T is e.g. double or long long. It uses products in intermediate steps so watch out for overflow if using int or long long. Using Point3D will always give a non-negative distance. For Point3D, call .dist on the result of the cross product.



f6bf6b, 4 lines

```
template < class P>
double lineDist(const P& a, const P& b, const P& p) {
  return (double) (b-a).cross(p-a)/(b-a).dist();
}
```

Segment Distance.h

Description:

Returns the shortest distance between point p and the line segment from point s to e.

Usage: Point < double > a, b(2,2), p(1,1);

bool onSegment = segDist(a,b,p) < 1e-10;
"Point.h"

5c88f4, 6 lines

```
typedef Point<double> P;
double segDist(P& s, P& e, P& p) {
   if (s==e) return (p-s).dist();
   auto d = (e-s).dist2(), t = min(d,max(.0,(p-s).dot(e-s)));
   return ((p-s)*d-(e-s)*t).dist()/d;
}
```

SegmentIntersection.h

Description:

If a unique intersection point between the line segments going from s1 to e1 and from s2 to e2 exists then it is returned. If no intersection point exists an empty vector is returned. If infinitely many exist a vector with 2 elements is returned, containing the endpoints of the common line segment. The wrong position will be returned if P is Point<11> and the intersection point does not have integer coordinates. Products of three coordinates are used in intermediate steps so watch out for overflow if using int or long long.



```
Usage: vector<P> inter = segInter(s1,e1,s2,e2);
if (sz(inter)==1)
cout << "segments intersect at " << inter[0] << endl;</pre>
"Point.h", "OnSegment.h"
                                                     9d57f2, 13 lines
template<class P> vector<P> segInter(P a, P b, P c, P d) {
 auto oa = c.cross(d, a), ob = c.cross(d, b),
       oc = a.cross(b, c), od = a.cross(b, d);
  // Checks if intersection is single non-endpoint point.
  if (sgn(oa) * sgn(ob) < 0 && sgn(oc) * sgn(od) < 0)
    return { (a * ob - b * oa) / (ob - oa) };
  set<P> s:
  if (onSegment(c, d, a)) s.insert(a);
  if (onSegment(c, d, b)) s.insert(b);
 if (onSegment(a, b, c)) s.insert(c);
 if (onSegment(a, b, d)) s.insert(d);
 return {all(s)};
```

lineIntersection.h

Description:

If a unique intersection point of the lines going through s1,e1 and s2,e2 exists {1, point} is returned. If no intersection point exists {0, (0,0)} is returned and if infinitely many exists {-1, (0,0)} is returned. The wrong position will be returned if P is Point<|| > and the intersection point does not have integer coordinates. Products of three coordinates are used in intermediate steps so watch out for overflow if using int or ll.



```
template < class P >
pair < int, P > lineInter(P s1, P e1, P s2, P e2) {
   auto d = (e1 - s1).cross(e2 - s2);
   if (d == 0) // if parallel
      return {-(s1.cross(e1, s2) == 0), P(0, 0)};
   auto p = s2.cross(e1, e2), q = s2.cross(e2, s1);
   return {1, (s1 * p + e1 * q) / d};
}
```

sideOf.h

Description: Returns where p is as seen from s towards e. $1/0/-1 \Leftrightarrow left/on line/right$. If the optional argument eps is given 0 is returned if p is within distance eps from the line. P is supposed to be Point<T> where T is e.g. double or long long. It uses products in intermediate steps so watch out for overflow if using int or long long.

```
template < class P >
int sideOf(P s, P e, P p) { return sgn(s.cross(e, p)); }

template < class P >
int sideOf(const P& s, const P& e, const P& p, double eps) {
    auto a = (e-s).cross(p-s);
    double 1 = (e-s).dist()*eps;
    return (a > 1) - (a < -1);</pre>
```

OnSegment.h

Description: Returns true iff p lies on the line segment from s to e. Use (segDist(s,e,p)<=epsilon) instead when using Point<double>.

"Point.h" c597e8, 3 lines

```
template<class P> bool onSegment(P s, P e, P p) {
   return p.cross(s, e) == 0 && (s - p).dot(e - p) <= 0;
}</pre>
```

linearTransformation.h Description:

Apply the linear transformation (translation, rotation and scaling) which takes line p0-p1 to line q0-q1 to point r.

```
d po res ql
```

```
typedef Point<double> P;
P linearTransformation(const P& p0, const P& p1,
    const P& q0, const P& q1, const P& r) {
    P dp = p1-p0, dq = q1-q0, num(dp.cross(dq), dp.dot(dq));
    return q0 + P((r-p0).cross(num), (r-p0).dot(num))/dp.dist2();
}
```

Angle.h

"Point.h"

Description: A class for ordering angles (as represented by int points and a number of rotations around the origin). Useful for rotational sweeping. Sometimes also represents points or vectors.

```
Usage: vector < Angle > v = \{w[0], w[0].t360() ...\}; // sorted
int j = 0; rep(i,0,n) { while (v[j] < v[i].t180()) ++j; }
// sweeps j such that (j-i) represents the number of positively
oriented triangles with vertices at 0 and i
struct Angle {
  int x, v;
  int t;
  Angle(int x, int y, int t=0) : x(x), y(y), t(t) {}
  Angle operator-(Angle b) const { return {x-b.x, y-b.y, t}; }
  int half() const {
    assert(x || v);
    return y < 0 || (y == 0 && x < 0);
  Angle t90() const { return \{-y, x, t + (half() \&\& x >= 0)\}; \}
  Angle t180() const { return {-x, -y, t + half()}; }
  Angle t360() const { return {x, y, t + 1}; }
bool operator < (Angle a, Angle b) {
  // add a.dist2() and b.dist2() to also compare distances
  return make_tuple(a.t, a.half(), a.y * (ll)b.x) <</pre>
         make_tuple(b.t, b.half(), a.x * (ll)b.y);
// Given two points, this calculates the smallest angle between
// them, i.e., the angle that covers the defined line segment.
pair<Angle, Angle> segmentAngles(Angle a, Angle b) {
  if (b < a) swap(a, b);
  return (b < a.t180() ?
          make_pair(a, b) : make_pair(b, a.t360()));
Angle operator+(Angle a, Angle b) { // point a + vector b
  Angle r(a.x + b.x, a.v + b.v, a.t);
  if (a.t180() < r) r.t--;</pre>
  return r.t180() < a ? r.t360() : r;
Angle angleDiff(Angle a, Angle b) { // angle b - angle a}
 int tu = b.t - a.t; a.t = b.t;
  return {a.x*b.x + a.y*b.y, a.x*b.y - a.y*b.x, tu - (b < a)};
```

12.2 Circles

CircleIntersection.h

Description: Computes the pair of points at which two circles intersect. Returns false in case of no intersection.

CircleTangents.h

Description: Finds the external tangents of two circles, or internal if r2 is negated. Can return 0, 1, or 2 tangents - 0 if one circle contains the other (or overlaps it, in the internal case, or if the circles are the same); 1 if the circles are tangent to each other (in which case .first = .second and the tangent line is perpendicular to the line between the centers). .first and .second give the tangency points at circle 1 and 2 respectively. To find the tangents of a circle with a point set r2 to 0.

```
"Point.h" b0153d, 13 lines
```

```
template<class P>
```

Description: Returns the area of the intersection of a circle with a ccw polygon.

```
Time: \mathcal{O}\left(n\right)
```

```
"../../content/geometry/Point.h" alee63, 19 lines
```

```
typedef Point<double> P;
#define arg(p, q) atan2(p.cross(q), p.dot(q))
double circlePoly(P c, double r, vector<P> ps) {
 auto tri = [&] (P p, P q) {
    auto r2 = r * r / 2;
    P d = q - p;
    auto a = d.dot(p)/d.dist2(), b = (p.dist2()-r*r)/d.dist2();
    auto det = a * a - b;
    if (det <= 0) return arg(p, q) * r2;</pre>
    auto s = max(0., -a-sqrt(det)), t = min(1., -a+sqrt(det));
    if (t < 0 || 1 <= s) return arg(p, q) * r2;</pre>
    P u = p + d * s, v = p + d * t;
    return arg(p,u) * r2 + u.cross(v)/2 + arg(v,q) * r2;
 auto sum = 0.0;
  rep(i, 0, sz(ps))
   sum += tri(ps[i] - c, ps[(i + 1) % sz(ps)] - c);
 return sum:
```

circumcircle.h

Description:

"Point.h"

The circumcirle of a triangle is the circle intersecting all three vertices. ccRadius returns the radius of the circle going through points A, B and C and ccCenter returns the center of the same circle.



14

```
typedef Point<double> P;
double ccRadius(const P& A, const P& B, const P& C) {
  return (B-A).dist()*(C-B).dist()*(A-C).dist()/
      abs((B-A).cross(C-A))/2;
}

P ccCenter(const P& A, const P& B, const P& C) {
  P b = C-A, c = B-A;
```

return A + (b*c.dist2()-c*b.dist2()).perp()/b.cross(c)/2;

Minimum Enclosing Circle.h

Description: Computes the minimum circle that encloses a set of points. **Time:** expected $\mathcal{O}(n)$

ac41a6, 17 lines

```
r = (o - ps[i]).dist();
    rep(k, 0, j) if ((o - ps[k]).dist() > r * EPS) {
     o = ccCenter(ps[i], ps[j], ps[k]);
      r = (o - ps[i]).dist();
return {o, r};
```

12.3 Polygons

InsidePolygon.h

Description: Returns true if p lies within the polygon. If strict is true, it returns false for points on the boundary. The algorithm uses products in intermediate steps so watch out for overflow. Usage: vector $\langle P \rangle$ v = {P{4,4}, P{1,2}, P{2,1}};

```
bool in = inPolygon(v, P{3, 3}, false);
Time: \mathcal{O}\left(n\right)
"Point.h", "OnSegment.h", "SegmentDistance.h"
                                                                     2bf504, 11 lines
template<class P>
```

```
bool inPolygon(vector<P> &p, P a, bool strict = true) {
 int cnt = 0, n = sz(p);
 rep(i,0,n) {
   P q = p[(i + 1) % n];
   if (onSegment(p[i], q, a)) return !strict;
   //or: if (segDist(p[i], q, a) \le eps) return !strict;
   cnt ^= ((a.y<p[i].y) - (a.y<q.y)) * a.cross(p[i], q) > 0;
 return cnt;
```

PolygonArea.h

Description: Returns twice the signed area of a polygon. Clockwise enumeration gives negative area. Watch out for overflow if using int as T!

```
f12<u>300, 6 lines</u>
template<class T>
T polygonArea2(vector<Point<T>>& v) {
 T = v.back().cross(v[0]);
 rep(i,0,sz(v)-1) a += v[i].cross(v[i+1]);
 return a:
```

PolygonCenter.h

Description: Returns the center of mass for a polygon. Time: $\mathcal{O}(n)$

```
9706dc, 9 lines
typedef Point <double > P;
P polygonCenter(const vector<P>& v) {
 P res(0, 0); double A = 0;
  for (int i = 0, j = sz(v) - 1; i < sz(v); j = i++) {
   res = res + (v[i] + v[j]) * v[j].cross(v[i]);
   A += v[j].cross(v[i]);
  return res / A / 3;
```

PolygonCut.h Description:

Returns a vector with the vertices of a polygon with everything to the left of the line going from s to e cut away.

```
Usage: vector<P> p = ...;
p = polygonCut(p, P(0,0), P(1,0));
```

f2b7d4, 13 lines

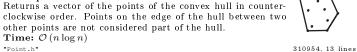
```
"Point.h", "lineIntersection.h"
typedef Point < double > P;
vector<P> polygonCut(const vector<P>& poly, P s, P e) {
  vector<P> res;
```

```
rep(i,0,sz(polv)) {
  P cur = poly[i], prev = i ? poly[i-1] : poly.back();
  bool side = s.cross(e, cur) < 0;</pre>
  if (side != (s.cross(e, prev) < 0))</pre>
    res.push_back(lineInter(s, e, cur, prev).second);
  if (side)
    res.push_back(cur);
return res:
```

Convex Hull, h.

Description:

clockwise order. Points on the edge of the hull between two other points are not considered part of the hull.



```
"Point.h"
typedef Point<11> P;
vector<P> convexHull(vector<P> pts) {
 if (sz(pts) <= 1) return pts;</pre>
 sort(all(pts));
 vector<P> h(sz(pts)+1);
 int s = 0, t = 0;
 for (int it = 2; it--; s = --t, reverse(all(pts)))
   for (P p : pts) {
     while (t \ge s + 2 \&\& h[t-2].cross(h[t-1], p) \le 0) t--;
 return {h.begin(), h.begin() + t - (t == 2 && h[0] == h[1])};
```

HullDiameter.h

Description: Returns the two points with max distance on a convex hull (ccw, no duplicate/collinear points).

Time: $\mathcal{O}\left(n\right)$

```
"Point.h"
                                                      c571b8, 12 lines
typedef Point<ll> P;
array<P, 2> hullDiameter(vector<P> S) {
 int n = sz(S), j = n < 2 ? 0 : 1;
 pair<11, array<P, 2>> res({0, {S[0], S[0]}});
 rep(i,0,j)
   for (;; j = (j + 1) % n) {
      res = \max(\text{res}, \{(S[i] - S[j]).dist2(), \{S[i], S[j]\}\});
      if ((S[(j+1) % n] - S[j]).cross(S[i+1] - S[i]) >= 0)
        break;
 return res.second;
```

PointInsideHull.h

Description: Determine whether a point t lies inside a convex hull (CCW order, with no collinear points). Returns true if point lies within the hull. If strict is true, points on the boundary aren't included.

Time: $\mathcal{O}(\log N)$

"Point.h", "sideOf.h", "OnSegment.h"

```
typedef Point<11> P;
bool inHull(const vector<P>& 1, P p, bool strict = true) {
 int a = 1, b = sz(1) - 1, r = !strict;
 if (sz(1) < 3) return r && onSegment(1[0], 1.back(), p);</pre>
 if (sideOf(1[0], 1[a], 1[b]) > 0) swap(a, b);
 if (sideOf(1[0], 1[a], p) >= r || sideOf(1[0], 1[b], p) <= -r)</pre>
    return false;
  while (abs(a - b) > 1) {
    int c = (a + b) / 2;
    (sideOf(1[0], 1[c], p) > 0 ? b : a) = c;
```

```
return sqn(l[a].cross(l[b], p)) < r;</pre>
```

LineHullIntersection.h

Description: Line-convex polygon intersection. The polygon must be ccw and have no collinear points. lineHull(line, poly) returns a pair describing the intersection of a line with the polygon: \bullet (-1,-1) if no collision, \bullet (i,-1)if touching the corner i, \bullet (i,i) if along side (i,i+1), \bullet (i,j) if crossing sides (i, i+1) and (j, j+1). In the last case, if a corner i is crossed, this is treated as happening on side (i, i+1). The points are returned in the same order as the line hits the polygon. extrVertex returns the point of a hull with the max projection onto a line.

```
Time: \mathcal{O}(\log n)
```

```
"Point.h"
                                                     7cf45b, 39 lines
#define cmp(i, j) sqn(dir.perp().cross(poly[(i)%n]-poly[(j)%n]))
#define extr(i) cmp(i + 1, i) >= 0 && cmp(i, i - 1 + n) < 0
template <class P> int extrVertex(vector<P>& poly, P dir) {
  int n = sz(poly), lo = 0, hi = n;
  if (extr(0)) return 0;
  while (10 + 1 < hi) {
   int m = (10 + hi) / 2;
   if (extr(m)) return m;
    int 1s = cmp(1o + 1, 1o), ms = cmp(m + 1, m);
    (1s < ms \mid | (1s == ms \&\& 1s == cmp(1o, m)) ? hi : 1o) = m;
 return lo:
#define cmpL(i) sqn(a.cross(poly[i], b))
template <class P>
array<int, 2> lineHull(P a, P b, vector<P>& poly) {
  int endA = extrVertex(poly, (a - b).perp());
  int endB = extrVertex(poly, (b - a).perp());
 if (cmpL(endA) < 0 || cmpL(endB) > 0)
    return {-1, -1};
  array<int, 2> res;
  rep(i, 0, 2) {
    int lo = endB, hi = endA, n = sz(poly);
    while ((lo + 1) % n != hi) {
      int m = ((lo + hi + (lo < hi ? 0 : n)) / 2) % n;
      (cmpL(m) == cmpL(endB) ? lo : hi) = m;
    res[i] = (lo + !cmpL(hi)) % n;
    swap (endA, endB);
 if (res[0] == res[1]) return {res[0], -1};
  if (!cmpL(res[0]) && !cmpL(res[1]))
    switch ((res[0] - res[1] + sz(poly) + 1) % sz(poly)) {
      case 0: return {res[0], res[0]};
      case 2: return {res[1], res[1]};
  return res;
```

12.4 Misc. Point Set Problems

ClosestPair.h

"Point.h"

71446b, 14 lines

Description: Finds the closest pair of points.

Time: $\mathcal{O}(n \log n)$

```
typedef Point<ll> P;
pair<P, P> closest (vector<P> v) {
  assert (sz(v) > 1);
  set<P> S:
  sort(all(v), [](P a, P b) { return a.y < b.y; });
  pair<11, pair<P, P>> ret{LLONG_MAX, {P(), P()}};
  int \dot{j} = 0;
  for (P p : v) {
    P d{1 + (ll) sqrt (ret.first), 0};
```

kdTree FastDelaunay PolyhedronVolume Point3D

```
ret = min(ret, \{(*lo - p).dist2(), \{*lo, p\}\});
    S.insert(p);
  return ret.second;
kdTree h
Description: KD-tree (2d, can be extended to 3d)
                                                     bac5b0, 63 lines
typedef long long T;
typedef Point<T> P:
const T INF = numeric_limits<T>::max();
bool on_x(const P& a, const P& b) { return a.x < b.x; }</pre>
bool on v(const P& a, const P& b) { return a.v < b.v; }
struct Node {
  P pt; // if this is a leaf, the single point in it
  T x0 = INF, x1 = -INF, y0 = INF, y1 = -INF; // bounds
  Node *first = 0, *second = 0;
  T distance (const P& p) { // min squared distance to a point
   T x = (p.x < x0 ? x0 : p.x > x1 ? x1 : p.x);
    T y = (p.y < y0 ? y0 : p.y > y1 ? y1 : p.y);
    return (P(x,y) - p).dist2();
  Node(vector<P>&& vp) : pt(vp[0]) {
    for (P p : vp) {
      x0 = min(x0, p.x); x1 = max(x1, p.x);
      y0 = min(y0, p.y); y1 = max(y1, p.y);
    if (vp.size() > 1) {
      // split on x if width >= height (not ideal...)
      sort(all(vp), x1 - x0 >= y1 - y0 ? on_x : on_y);
      // divide by taking half the array for each child (not
      // best performance with many duplicates in the middle)
      int half = sz(vp)/2;
      first = new Node({vp.begin(), vp.begin() + half});
      second = new Node({vp.begin() + half, vp.end()});
};
struct KDTree {
  Node* root:
  KDTree(const vector<P>& vp) : root(new Node({all(vp)})) {}
  pair<T, P> search(Node *node, const P& p) {
    if (!node->first) {
      // uncomment if we should not find the point itself:
      // if (p = node \rightarrow pt) return \{INF, P()\};
      return make_pair((p - node->pt).dist2(), node->pt);
    Node *f = node -> first, *s = node -> second;
    T bfirst = f->distance(p), bsec = s->distance(p);
    if (bfirst > bsec) swap(bsec, bfirst), swap(f, s);
    // search closest side first, other side if needed
    auto best = search(f, p);
    if (bsec < best.first)</pre>
     best = min(best, search(s, p));
    return best;
```

while (v[j].y <= p.y - d.x) S.erase(v[j++]);</pre>

for (; lo != hi; ++lo)

auto lo = S.lower_bound(p - d), hi = S.upper_bound(p + d);

```
// find nearest point to a point, and its squared distance
  // (requires an arbitrary operator< for Point)
  pair<T, P> nearest(const P& p) {
    return search (root, p);
};
FastDelaunay.h
Description: Fast Delaunay triangulation. Each circumcircle contains none
of the input points. There must be no duplicate points. If all points are on a
line, no triangles will be returned. Should work for doubles as well, though
there may be precision issues in 'circ'. Returns triangles in order {t[0][0],
t[0][1], t[0][2], t[1][0], \dots\}, all counter-clockwise.
Time: \mathcal{O}(n \log n)
"Point.h"
typedef Point<ll> P;
typedef struct Quad* Q;
typedef int128 t 111; // (can be ll if coords are < 2e4)
P arb(LLONG_MAX, LLONG_MAX); // not equal to any other point
struct Quad {
  Q rot, o; P p = arb; bool mark;
 P& F() { return r()->p; }
  Q& r() { return rot->rot; }
  Q prev() { return rot->o->rot; }
  0 next() { return r()->prev(); }
bool circ(P p, P a, P b, P c) { // is p in the circumcircle?
  111 p2 = p.dist2(), A = a.dist2()-p2,
      B = b.dist2()-p2, C = c.dist2()-p2;
  return p.cross(a,b) \starC + p.cross(b,c) \starA + p.cross(c,a) \starB > 0;
O makeEdge(P orig, P dest) {
  O r = H ? H : new Ouad{new Ouad{new Ouad{0}}};
  H = r -> 0; r -> r() -> r() = r;
  rep(i,0,4) r = r - rot, r - p = arb, r - o = i & 1 ? <math>r : r - r();
  r->p = orig; r->F() = dest;
  return r;
void splice(Q a, Q b) {
  swap(a->o->rot->o, b->o->rot->o); swap(a->o, b->o);
Q connect(Q a, Q b) {
  Q = makeEdge(a->F(), b->p);
  splice(q, a->next());
  splice(q->r(), b);
  return q;
pair<0.0> rec(const vector<P>& s) {
  if (sz(s) \le 3) {
    Q = makeEdge(s[0], s[1]), b = makeEdge(s[1], s.back());
    if (sz(s) == 2) return { a, a->r() };
    splice(a->r(), b);
    auto side = s[0].cross(s[1], s[2]);
    Q c = side ? connect(b, a) : 0;
    return {side < 0 ? c->r() : a, side < 0 ? c : b->r() };
#define H(e) e->F(), e->p
#define valid(e) (e->F().cross(H(base)) > 0)
 Q A, B, ra, rb;
 int half = sz(s) / 2;
  tie(ra, A) = rec({all(s) - half});
  tie(B, rb) = rec({sz(s) - half + all(s)});
  while ((B->p.cross(H(A)) < 0 \&& (A = A->next())) | |
          (A->p.cross(H(B)) > 0 && (B = B->r()->o)));
  Q base = connect(B \rightarrow r(), A);
```

```
if (A->p == ra->p) ra = base->r();
  if (B->p == rb->p) rb = base;
#define DEL(e, init, dir) Q e = init->dir; if (valid(e)) \
    while (circ(e->dir->F(), H(base), e->F())) {
      0 t = e \rightarrow dir; \
      splice(e, e->prev()); \
      splice(e->r(), e->r()->prev()); \
      e->o = H; H = e; e = t; \setminus
  for (;;) {
    DEL(LC, base->r(), o); DEL(RC, base, prev());
    if (!valid(LC) && !valid(RC)) break;
    if (!valid(LC) || (valid(RC) && circ(H(RC), H(LC))))
      base = connect(RC, base->r());
    else
      base = connect(base->r(), LC->r());
  return { ra, rb };
vector<P> triangulate(vector<P> pts) {
 sort(all(pts)); assert(unique(all(pts)) == pts.end());
 if (sz(pts) < 2) return {};
 Q e = rec(pts).first;
 vector<Q> q = \{e\};
 int qi = 0;
  while (e->o->F().cross(e->F(), e->p) < 0) e = e->o;
#define ADD { Q c = e; do { c->mark = 1; pts.push_back(c->p); \
 q.push\_back(c->r()); c = c->next(); } while (c != e); }
 ADD; pts.clear();
  while (qi < sz(q)) if (!(e = q[qi++]) -> mark) ADD;
  return pts;
```

$12.5 \quad 3D$

PolyhedronVolume.h

Description: Magic formula for the volume of a polyhedron. Faces should point outwards.

```
template < class V, class L>
double signedPolyVolume(const V& p, const L& trilist) {
  double v = 0;
  for (auto i : trilist) v += p[i.a].cross(p[i.b]).dot(p[i.c]);
  return v / 6;
}
```

Point 3D.h

Description: Class to handle points in 3D space. T can be e.g. double or long long. 8058ae, 32 lines

```
template < class T > struct Point3D {
  typedef Point3D P;
  typedef const P& R;
 T x, y, z;
  explicit Point3D(T x=0, T y=0, T z=0) : x(x), y(y), z(z) {}
 bool operator<(R p) const {</pre>
    return tie(x, y, z) < tie(p.x, p.y, p.z); }
  bool operator==(R p) const {
    return tie(x, y, z) == tie(p.x, p.y, p.z); }
  P operator+(R p) const { return P(x+p.x, y+p.y, z+p.z); }
  P operator-(R p) const { return P(x-p.x, y-p.y, z-p.z); }
  P operator*(T d) const { return P(x*d, y*d, z*d); }
  P operator/(T d) const { return P(x/d, y/d, z/d); }
 T dot(R p) const { return x*p.x + y*p.y + z*p.z; }
 P cross(R p) const {
    return P(y*p.z - z*p.y, z*p.x - x*p.z, x*p.y - y*p.x);
 T dist2() const { return x*x + y*y + z*z; }
```

```
double dist() const { return sqrt((double)dist2()); }
  //Azimuthal angle (longitude) to x-axis in interval [-pi, pi]
  double phi() const { return atan2(y, x); }
  //Zenith angle (latitude) to the z-axis in interval [0, pi]
  double theta() const { return atan2(sqrt(x*x+y*y),z); }
  P unit() const { return *this/(T)dist(); } //makes dist()=1
  //returns unit vector normal to *this and p
  P normal(P p) const { return cross(p).unit(); }
  //returns point rotated 'angle' radians ccw around axis
  P rotate (double angle, P axis) const {
   double s = sin(angle), c = cos(angle); P u = axis.unit();
    return u*dot(u)*(1-c) + (*this)*c - cross(u)*s;
};
```

3dHull.h

Description: Computes all faces of the 3-dimension hull of a point set. *No four points must be coplanar*, or else random results will be returned. All faces will point outwards.

Time: $\mathcal{O}\left(n^2\right)$

"Point3D.h" 5b45fc, 49 lines typedef Point3D<double> P3; struct PR { **void** ins(**int** x) { $(a == -1 ? a : b) = x; }$ **void** rem(**int** x) { $(a == x ? a : b) = -1; }$ int cnt() { return (a != -1) + (b != -1); } **struct** F { P3 q; **int** a, b, c; }; vector<F> hull3d(const vector<P3>& A) { assert (sz(A) >= 4); vector<vector<PR>> $E(sz(A), vector<PR>(sz(A), \{-1, -1\}));$ #define E(x,y) E[f.x][f.y] vector<F> FS; auto mf = [&](int i, int j, int k, int l) { P3 q = (A[j] - A[i]).cross((A[k] - A[i]));if (q.dot(A[1]) > q.dot(A[i])) q = q * -1;F f{q, i, j, k}; E(a,b).ins(k); E(a,c).ins(j); E(b,c).ins(i);FS.push_back(f); rep(i,0,4) rep(j,i+1,4) rep(k,j+1,4)mf(i, j, k, 6 - i - j - k);rep(i,4,sz(A)) { rep(j, 0, sz(FS)) { F f = FS[j];**if**(f.q.dot(A[i]) > f.q.dot(A[f.a])) { E(a,b).rem(f.c);E(a,c).rem(f.b);E(b,c).rem(f.a);swap(FS[j--], FS.back()); FS.pop_back(); int nw = sz(FS); rep(j,0,nw) { F f = FS[j];**#define** C(a, b, c) **if** (E(a,b).cnt() != 2) mf(f.a, f.b, i, f.c); C(a, b, c); C(a, c, b); C(b, c, a);for (F& it : FS) if ((A[it.b] - A[it.a]).cross(A[it.c] - A[it.a]).dot(it.q) <= 0) swap(it.c, it.b);return FS;

sphericalDistance.h

Description: Returns the shortest distance on the sphere with radius radius between the points with azimuthal angles (longitude) f1 (ϕ_1) and f2 (ϕ_2) from x axis and zenith angles (latitude) t1 (θ_1) and t2 (θ_2) from z axis (0 =north pole). All angles measured in radians. The algorithm starts by converting the spherical coordinates to cartesian coordinates so if that is what you have you can use only the two last rows. dx*radius is then the difference between the two points in the x direction and d*radius is the total distance between the points. 611f07, 8 lines

```
double sphericalDistance(double f1, double t1,
    double f2, double t2, double radius) {
 double dx = \sin(t2) \cdot \cos(f2) - \sin(t1) \cdot \cos(f1);
 double dy = sin(t2) * sin(f2) - sin(t1) * sin(f1);
 double dz = cos(t2) - cos(t1);
 double d = sqrt(dx*dx + dy*dy + dz*dz);
 return radius*2*asin(d/2);
```

Dynamic Programming (13)

DVC.hpp

Description: Optimize $O(N^2K)$ to $O(NK \log N)$

"../template/Header.hpp" aa5ddf, 19 lines

```
vector<vl> cst, dp;
11 cost(int 1, int r) {
 return cst[l][r];
void divide(int 1, int r, int opt 1, int opt r, int c) {
 if(1 > r) return ;
 int mid = (1 + r) / 2;
 pair<11, int> best = make_pair(INF, -1);
 for(int k=opt_1; k<=min(mid, opt_r); ++k) {</pre>
   best = min(best, make_pair(dp[c - 1][k] + cost(k + 1, mid),
 dp[c][mid] = best.first;
  divide(1, mid - 1, opt_1, best.second, c);
 divide (mid + 1, r, best.second, opt_r, c);
// for(int c=1; c \leq K; ++c) divide(1, N, 1, N, c);
```

SlopeTrick.hpp

Description: Absolute Smth

"../template/Header.hpp" f62f9a, 36 lines

```
11 extending_value;
struct slope_trick {
 multiset<11> ms_1, ms_r;
 11 \text{ min_y} = 011, 1z_1 = 011, 1z_r = 011;
 bool extending = false;
 void add line(ll v) {
    if(extending) {
      lz_l -= extending_value;
      lz_r -= extending_value;
    extending = true;
    if(ms_l.empty() && ms_r.empty()) {
      ms_l.emplace(v);
      ms_r.emplace(v);
    else if(v <= *ms_l.rbegin() + lz_l) {</pre>
```

```
min_y += (*ms_l.rbegin() + lz_l) - v;
     ms_r.emplace(*ms_l.rbegin() + lz_l - lz_r);
     ms l.erase(--ms l.end());
     ms_l.emplace(v - lz_l);
     ms_1.emplace(v - lz_1);
    else if(v >= *ms_r.begin() + lz_r) {
     min_y += v - (*ms_r.begin() + lz_r);
     ms_l.emplace(*ms_r.begin() + lz_r - lz_l);
     ms_r.erase(ms_r.begin());
     ms_r.emplace(v - lz_r);
     ms_r.emplace(v - lz_r);
    else {
     ms_l.emplace(v - lz_l);
     ms_r.emplace(v - lz_r);
};
```

Convolutions (14)

And Convolution.hpp

Description: Bitwise AND Convolution. Superset Zeta Transform: A'[S] = $\sum_{T \supset S} A[T]$. Superset Mobius Transform: $A[T] = \sum_{S \supset T} (-1)^{|S-T|} A'[S]$. Time: $\mathcal{O}(N \log N)$.

```
"../template/Header.hpp"
```

```
template<class T>
void superset_zeta(vector<T> &a){
    int n=(int)a.size();
    assert (n==(n\&-n));
    for (int i=1; i<n; i<<=1) {</pre>
         for (int j=0; j<n; j++) {</pre>
             if(j&i){
                 a[j^i] += a[j];
template<class T>
void superset mobius(vector<T> &a){
    int n=(int)a.size();
    assert (n==(n\&-n));
    for(int i=n;i>>=1;){
         for (int j=0; j<n; j++) {</pre>
             if(i&i){
                 a[j^i]-=a[j];
template<class T>
vector<T> and_convolution(vector<T> a, vector<T> b) {
    superset zeta(a);
    superset zeta(b);
    for(int i=0;i<(int)a.size();i++)a[i]*=b[i];</pre>
    superset_mobius(a);
    return a;
```

GCDConvolution.hpp

Description: GCD Convolution. Multiple Zeta Transform: A'[n] = $\sum_{n|m} A[m]$. Multiple Mobius Transform: $A[n] = \sum_{n|m} \mu(m/n)A'[m]$. Time: $\mathcal{O}(N \log \log N)$. "../template/Header.hpp 7f6c2d, 34 lines

```
template<class T>
void multiple zeta(vector<T> &a) {
    int n=(int)a.size();
    vector<bool> is_prime(n,true);
    for(int p=2;p<n;p++){
        if(!is prime[p])continue;
        for (int i=(n-1)/p;i>=1;i--) {
            is prime[i*p]=false;
            a[i]+=a[i*p];
template<class T>
void multiple mobius(vector<T> &a){
    int n=(int)a.size();
    vector<bool> is_prime(n, true);
    for (int p=2; p<n; p++) {</pre>
        if(!is_prime[p])continue;
        for (int i=1; i*p<n; i++) {</pre>
            is_prime[i*p]=false;
            a[i] -= a[i * p];
template < class T>
vector<T> gcd_convolution(vector<T> a, vector<T> b) {
    multiple_zeta(a);
    multiple_zeta(b);
    for(int i=0;i<(int)a.size();i++)a[i]*=b[i];</pre>
    multiple_mobius(a);
    return a;
LCMConvolution.hpp
Description: LCM Convolution. Divisor Zeta Transform: A'[n] =
\sum_{d|n} A[d]. Divisor Mobius Transform: A[n] = \sum_{d|n} \mu(n/d) A'[d].
Time: \mathcal{O}(N \log \log N).
"../template/Header.hpp"
                                                        41fe9d, 34 lines
template<class T>
void divisor_zeta(vector<T> &a){
    int n=(int)a.size();
    vector<bool> is_prime(n,true);
    for(int p=2;p<n;p++){
        if(!is_prime[p])continue;
        for (int i=1; i*p<n; i++) {</pre>
            is_prime[i*p]=false;
            a[i*p]+=a[i];
template < class T>
void divisor_mobius(vector<T> &a) {
    int n=(int)a.size();
    vector<bool> is_prime(n,true);
    for(int p=2;p<n;p++){
        if(!is_prime[p])continue;
        for(int i=(n-1)/p;i>=1;i--){
            is_prime[i*p]=false;
            a[i*p]-=a[i];
template<class T>
vector<T> lcm_convolution(vector<T> a, vector<T> b) {
```

```
divisor zeta(a);
     divisor_zeta(b);
    for (int i=0; i < (int) a.size(); i++) a[i] *=b[i];</pre>
    divisor mobius(a);
     return a;
ORConvolution.hpp
Description: Bitwise OR Convolution. Subset Zeta Transform: A'[S] =
\sum_{T \subseteq S} A[T]. Subset Mobius Transform: A[T] = \sum_{S \subseteq T} (-1)^{|T-S|} A'[S].
\mathbf{Time} : \mathcal{O}(N \log N).
"../template/Header.hpp"
                                                              c58b77, 34 lines
template<class T>
void subset_zeta(vector<T> &a) {
    int n=(int)a.size();
    assert (n==(n\&-n));
    for (int i=1; i < n; i < <=1) {</pre>
         for (int j=0; j<n; j++) {</pre>
              if(j&i){
                  a[j]+=a[j^i];
template < class T>
void subset_mobius(vector<T> &a) {
    int n=(int)a.size();
    assert (n==(n\&-n));
    for(int i=n;i>>=1;) {
         for (int j=0; j<n; j++) {</pre>
              if(j&i){
                  a[j] -= a[j^i];
    }
template<class T>
vector<T> or_convolution(vector<T> a, vector<T> b) {
    subset_zeta(a);
    subset zeta(b);
    for(int i=0;i<(int)a.size();i++)a[i]*=b[i];</pre>
    subset_mobius(a);
    return a;
XORConvolution.hpp
Description: Bitwise XOR Convolution. Fast Walsh-Hadamard Transform:
A'[S] = \sum_{T} (-1)^{|S\&T|} A[T].
Time: \mathcal{O}(N \log N).
"../template/Header.hpp"
                                                              05848d, 29 lines
template < class T>
void fwht(vector<T> &a){
    int n=(int)a.size();
    assert (n==(n\&-n));
    for (int i=1; i < n; i < <=1) {</pre>
         for (int j=0; j<n; j++) {</pre>
              if(j&i){
                   T &u=a[j^i],&v=a[j];
                   tie (u, v) = make_pair(u+v, u-v);
template<class T>
```

```
vector<T> xor convolution(vector<T> a, vector<T> b) {
    int n=(int)a.size();
    fwht(a);
    fwht(b);
    for (int i=0; i<n; i++) a[i] *=b[i];</pre>
    fwht(a);
   T div=T(1)/T(n);
    if (div==T(0)) {
        for(auto &x:a) x/=n;
    }else{
        for(auto &x:a)x*=div;
    return a;
```

Various (15)

Gaussian Elimination.hpp Description: Gaussian Elimination

e89ecb, 34 lines

```
"../template/Header.hpp"
struct Gauss {
 int n, sz;
 vector<ll> basis;
  Gauss(int n = 0) {
   init(n);
 void init(int n) {
    n = _n, sz = 0;
   basis.assign(n, 0);
 void insert(ll x) {
    for (int i = n - 1; i >= 0; i--)
      if (x >> i & 1) {
        if (!basis[i]) {
          basis[i] = x;
          sz++;
          return;
        x ^= basis[i];
 11 \text{ getmax}(11 \text{ k} = 0)  {
    11 tot = 111 << sz, res = 0;
    for (int i = n - 1; i >= 0; i--)
      if (basis[i]) {
        tot >>= 1;
        if ((k >= tot && res >> i & 1) || (k < tot && res >> i
             & 1 ^ 1))
          res ^= basis[i];
        if (k >= tot)
          k -= tot;
    return res;
};
```

Binary Trie.hpp

Description: Binary Trie "../template/Header.hpp"

525bf4, 59 lines using node_t = array<int, 2>;

```
template<size_t S>
struct binary_trie {
 vector<node_t> t = {node_t()};
 vector<int> cnt = {0};
 int cnt_nodes = 0;
 void insert(int v) {
   int cur = 0;
```

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```
cnt[0]++;
    for(int i=S-1; i>=0; --i) {
      int b = (v \& (1 << i)) ? 1: 0;
      if(!t[cur][b]) {
       t[cur][b] = ++cnt_nodes;
       t.emplace_back(node_t());
       cnt.emplace back(0);
      cnt[t[cur][b]]++;
     cur = t[curl[b];
  void remove(int v) {
   int cur = 0;
    cnt[0]--;
    for(int i=S-1; i>=0; --i) {
     int b = (v & (1 << i)) ? 1: 0;</pre>
     cnt[t[cur][b]]--;
      cur = t[cur][b];
  int get_min(int v) {
    int cur = 0, res = 0;
    for(int i=(int) S-1; i>=0; --i) {
     int b = (v \& (1 << i)) ? 1 : 0;
     if(t[cur][b] && cnt[t[cur][b]]) {
       cur = t[cur][b];
      else {
        res |= (1 << i);
        cur = t[cur][!b];
    return res;
  int get_max(int v) {
   int cur = 0, res = 0;
    for(int i=(int) S-1; i>=0; --i) {
     int b = (v \& (1 << i)) ? 1 : 0;
     if(t[cur][!b] && cnt[t[cur][!b]]) {
       res |= (1 << i);
        cur = t[cur][!b];
     else {
        cur = t[cur][b];
    return res;
};
```

InfixPropostfix.hpp Description: Infix to Pro-Postfix

"../template/Header.hpp"
stack<char> opr;

517f57, 47 lines

```
stack<char> opr;
stack<int> val;

bool isOpr(char x) {
   return x == '+' || x == '*';
}

int prio(char x) {
   if(x == '(') return -1;
   if(x == '+') return 1;
   if(x == '*') return 2;
   return 0;
}
```

```
int do opr(int 1, int r, char o) {
 if (0 == '+') {
   return 1 + r;
 return 1 * r;
void pop_stack() {
 int rhs = val.top(); val.pop();
 int lhs = val.top(); val.pop();
 int new_val = do_opr(lhs, rhs, opr.top());
 val.emplace(new_val);
 opr.pop();
int cal(string s) {
 for(auto x: s) {
   if(isdigit(x)) val.emplace(x - '0');
   else if(x == '(') opr.emplace('(');
    else if(x == ')') {
     while(!opr.empty() && opr.top() != '(')
       pop_stack();
     opr.pop();
    else {
     while(!opr.empty() && prio(opr.top()) >= prio(x))
       pop_stack();
     opr.emplace(x);
 while(!opr.empty()) pop_stack();
 return val.top();
```

15.1 Optimization tricks

__builtin_ia32_ldmxcsr(40896); disables denormals (which make floats 20x slower near their minimum value).

15.1.1 Bit hacks

- x & -x is the least bit in x.
- for (int x = m; x;) { --x &= m; ... } loops over all subset masks of m (except m itself).
- c = x&-x, r = x+c; (((r^x) >> 2)/c) | r is the next number after x with the same number of bits set.
- rep(b,0,K) rep(i,0,(1 << K))
 if (i & 1 << b) D[i] += D[i^(1 << b)];
 computes all sums of subsets.</pre>

15.1.2 Pragmas

- #pragma GCC optimize ("Ofast") will make GCC auto-vectorize loops and optimizes floating points better.
- #pragma GCC target ("avx2") can double performance of vectorized code, but causes crashes on old machines.
- #pragma GCC optimize ("trapv") kills the program on integer overflows (but is really slow).

Competitive Programming Topics



topics.txt

Bitonic cycle

159 lines

Recursion Divide and conquer Finding interesting points in N log N Algorithm analysis Master theorem Amortized time complexity Greedy algorithm Scheduling Max contiguous subvector sum Invariants Huffman encoding Graph theory Dynamic graphs (extra book-keeping) Breadth first search Depth first search * Normal trees / DFS trees Dijkstra's algorithm MST: Prim's algorithm Bellman-Ford Konig's theorem and vertex cover Min-cost max flow Lovasz toggle Matrix tree theorem Maximal matching, general graphs Hopcroft-Karp Hall's marriage theorem Graphical sequences Flovd-Warshall Euler cycles Flow networks * Augmenting paths * Edmonds-Karp Bipartite matching Min. path cover Topological sorting Strongly connected components Cut vertices, cut-edges and biconnected components Edge coloring * Trees Vertex coloring * Bipartite graphs (=> trees) * 3^n (special case of set cover) Diameter and centroid K'th shortest path Shortest cycle Dynamic programming Knapsack Coin change Longest common subsequence Longest increasing subsequence Number of paths in a dag Shortest path in a dag Dynprog over intervals Dynprog over subsets Dynprog over probabilities Dynprog over trees 3^n set cover Divide and conquer Knuth optimization Convex hull optimizations RMQ (sparse table a.k.a 2^k-jumps)

Log partitioning (loop over most restricted) Combinatorics Computation of binomial coefficients Pigeon-hole principle Inclusion/exclusion Catalan number Pick's theorem Number theory Integer parts Divisibility Euclidean algorithm Modular arithmetic * Modular multiplication * Modular inverses * Modular exponentiation by squaring Chinese remainder theorem Fermat's little theorem Euler's theorem Phi function Frobenius number Quadratic reciprocity Pollard-Rho Miller-Rabin Hensel lifting Vieta root jumping Game theory Combinatorial games Game trees Mini-max Nim Games on graphs Games on graphs with loops Grundy numbers Bipartite games without repetition General games without repetition Alpha-beta pruning Probability theory Optimization Binary search Ternary search Unimodality and convex functions Binary search on derivative Numerical methods Numeric integration Newton's method Root-finding with binary/ternary search Golden section search Matrices Gaussian elimination Exponentiation by squaring Sorting Radix sort Geometry Coordinates and vectors * Cross product * Scalar product Convex hull Polygon cut Closest pair Coordinate-compression Quadtrees KD-trees All segment-segment intersection Discretization (convert to events and sweep) Angle sweeping Line sweeping Discrete second derivatives Strings

Longest common substring Palindrome subsequences Knuth-Morris-Pratt Tries Rolling polynomial hashes Suffix array Suffix tree Aho-Corasick Manacher's algorithm Letter position lists Combinatorial search Meet in the middle Brute-force with pruning Best-first (A*) Bidirectional search Iterative deepening DFS / A* Data structures LCA (2^k-jumps in trees in general) Pull/push-technique on trees Heavy-light decomposition Centroid decomposition Lazy propagation Self-balancing trees Convex hull trick (wcipeg.com/wiki/Convex_hull_trick) Monotone queues / monotone stacks / sliding queues Sliding queue using 2 stacks Persistent segment tree

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