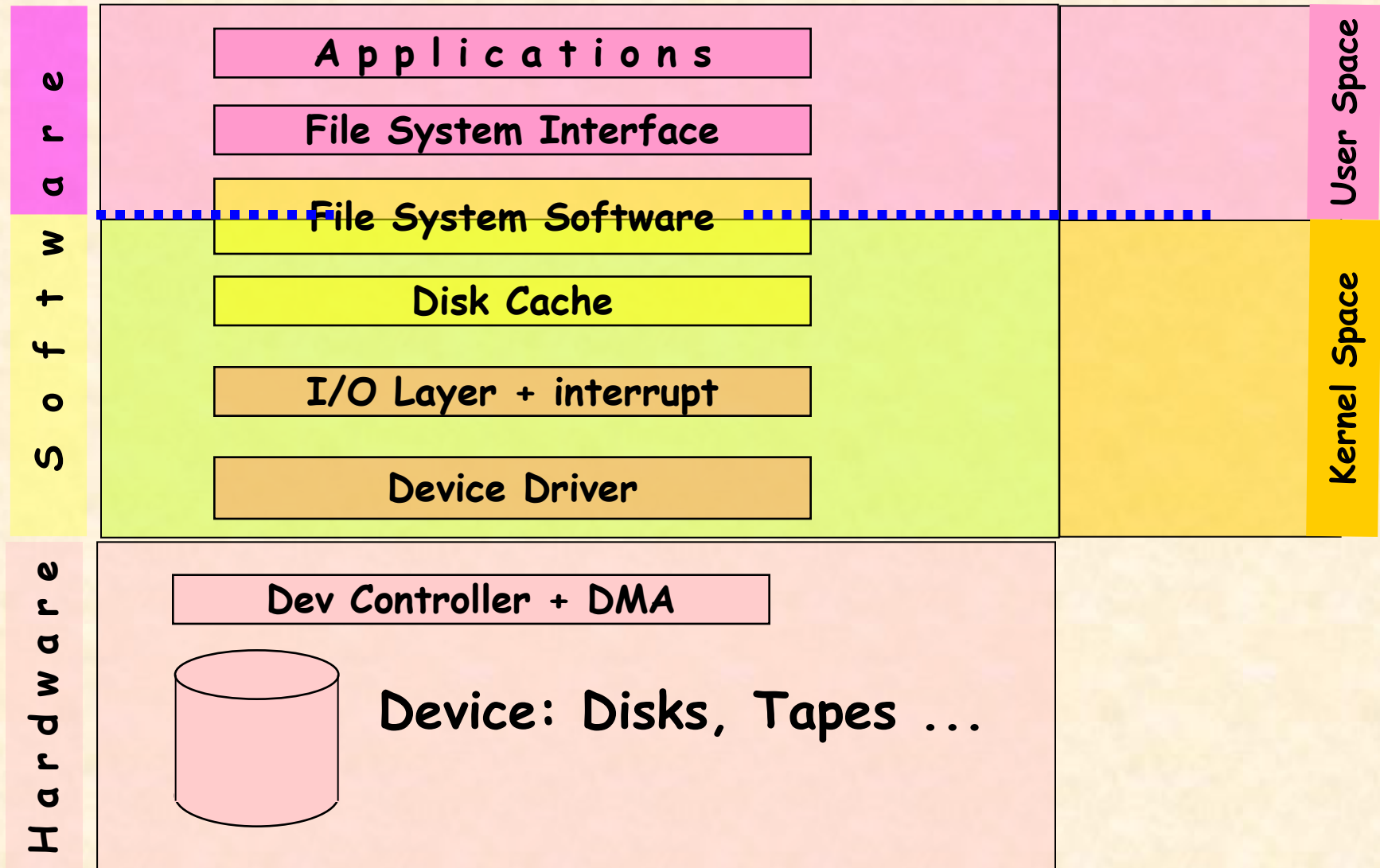


FILE SYSTEMS

File Systems typical layers



File Concept

- ✓ **Contiguous logical address space**
- ✓ **Files are mapped, by the operating system, onto physical devices**
 - **These devices are nonvolatile – contents are persistent through power failures and system reboots**
- ✓ **File** is a named collection of related information that is recorded on secondary storage
- ✓ **Types of information stored in a file :**
 - **Data**
 - ◆ **numeric**
 - ◆ **character**
 - ◆ **binary**
 - **Programs**
 - ◆ **source programs**
 - ◆ **object programs**
 - ◆ **executable programs**

File Structure

- ✓ **Text file :**

- **Is a sequence of characters organized into lines (possibly pages)**

- ✓ **Source file :**

- **Is a sequence of subroutines and functions, which is further organized as declarations followed by executable statements**

- ✓ **Object file :**

- **Is a sequence of bytes organized into code sections that the loader can bring into memory and execute.**

File Attributes

- ✓ **Name** – only information kept in human-readable form
- ✓ **Identifier** – unique tag (number) identifies file within file system
- ✓ **Type** – needed for systems that support different types
- ✓ **Location** – pointer to file location on device
- ✓ **Size** – current file size
- ✓ **Protection** – controls who can do reading, writing, executing
- ✓ **Time, date, and user identification** – Creation, modification, last use, data for protection, security, and usage monitoring
- ✓ **Information about files is kept in the directory structure, which is maintained on the disk**

File Operations

- ✓ File is an abstract data type
- ✓ File Operations
 - **Create** : Allocate space, and create directory entry
 - **Write** : Find file location and write from buffer at write pointer
 - **Read** : Find location and read to buffer from read pointer
 - **Reposition within a file** : Move pointer (Seek)
 - **Delete** : Free space, and free directory entry
 - **Truncating a file** : Erase the contents of a file keeping the attributes
 - **Open (FN)** – search the directory structure on disk for entry FN , and move the content of entry to memory.
 - **Close (FN)** – move the content of entry FN in memory to directory structure on disk

Additional Operations

- ✓ The directory itself is maintained on disk (normally as a file)
- ✓ To avoid the need to search directory for each operation an **open-file table** is maintained
- ✓ Open operation creates an entry in this table and returns pointer to this table
- ✓ All subsequent operations receive this pointer (instead of name) as parameter
- ✓ Close operation releases the entry
- ✓ Open file table is normally split in two levels of internal tables
 - **System-wide table** (**GOFT** - global open file table)
 - **Per-Process table** (**OFT** open file table)

Open Files

- **System-wide table** – contains process independent information such as location of the file on disk, file size, access dates
- **Per-Process table** – Tracks all files that a process has open. It has information regarding the use of the file by the process such as current file pointer, access rights
- ✓ Several pieces of data are needed to manage open files:
 - **File pointer:** pointer to last read/write location, per process that has the file open
 - **File-open count:** number of times a file is open – to allow removal of data from open-file table when last process closes it
 - **Disk location of the file:** cache of data access information
 - **Access rights:** per-process access mode information

Open File Locking

- ✓ Facility provided by some operating systems and file systems for locking an open file (or part of a file)
- ✓ Prevents other process from gaining access to a open file
- ✓ File locking may be either **Mandatory** or **Advisory**:
 - **Mandatory** – access is denied depending on locks held and requested. Operating system ensures locking integrity Ex. Windows Systems
 - **Advisory** – processes can find status of locks and decide what to do. It is upto the software developer to ensure locking integrity. Ex. UNIX systems

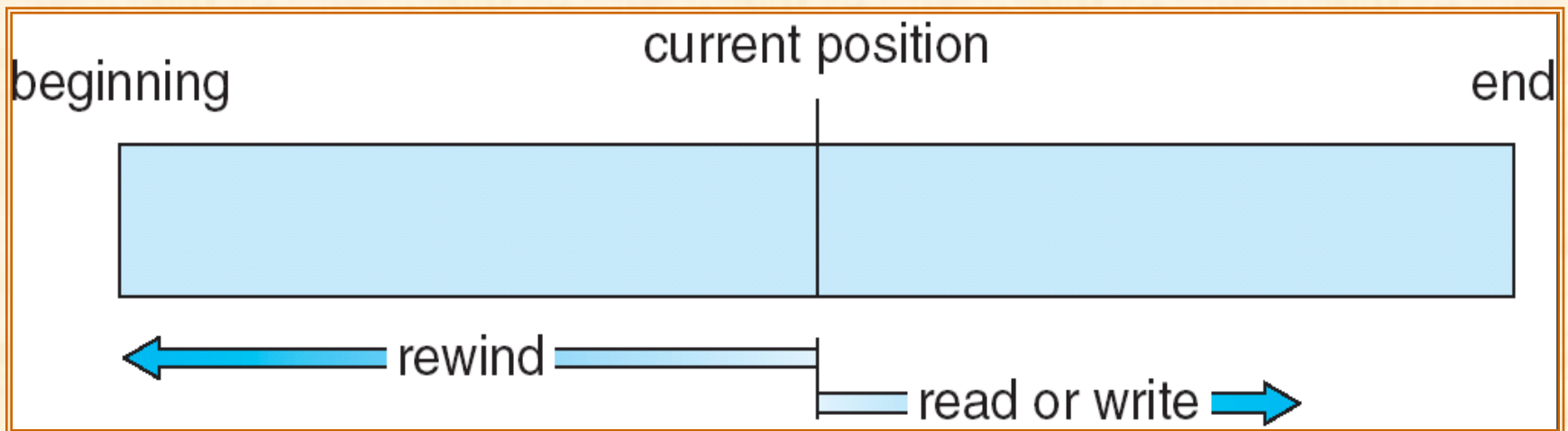
File Types – Name, Extension

| file type | usual extension | function |
|----------------|-----------------------------|--|
| executable | exe, com, bin or none | ready-to-run machine- language program |
| object | obj, o | compiled, machine language, not linked |
| source code | c, cc, java, pas, asm, a | source code in various languages |
| batch | bat, sh | commands to the command interpreter |
| text | txt, doc | textual data, documents |
| word processor | wp, tex, rtf, doc | various word-processor formats |
| library | lib, a, so, dll | libraries of routines for programmers |
| print or view | ps, pdf, jpg | ASCII or binary file in a format for printing or viewing |
| archive | arc, zip, tar | related files grouped into one file, sometimes com- pressed, for archiving or storage |
| multimedia | mpeg, mov, rm, mp3, avi | binary file containing audio or A/V information |

Access Methods

1. Sequential Access :

- simplest access method
 - Information in the file is processed in order, one record after the other
 - Example : editors and compilers access files sequentially
- ✓ Operations on files are : **read next**
 write next
 reset
 no read after last write (rewrite)



Access Methods

2. Direct Access :

- Also called relative access
- Read or write takes in no particular order
- File is viewed as numbered sequence of blocks or records
- Example : databases

✓ Operations on files are :

read n

write n

position to n

read next

write next

rewrite n

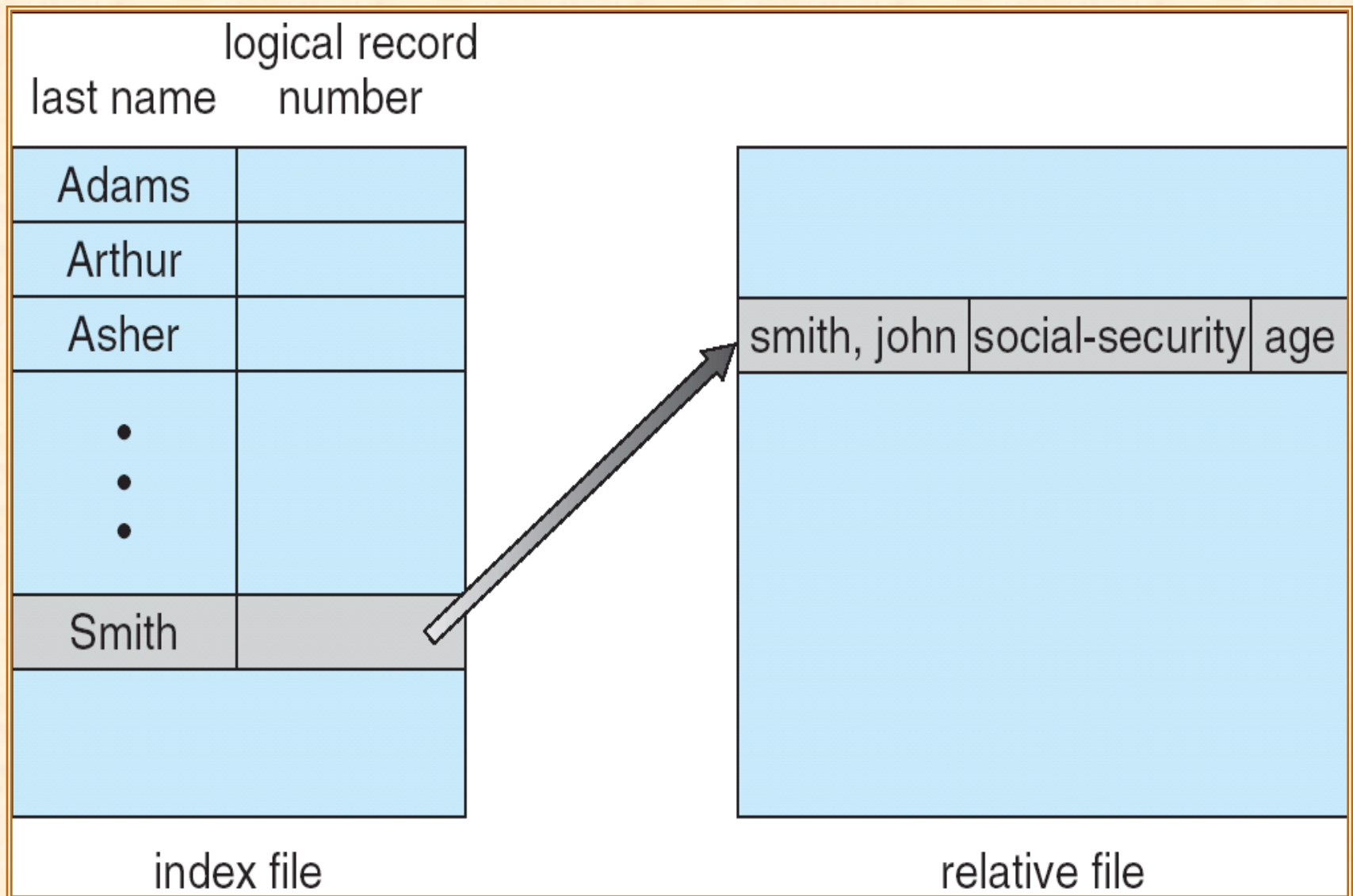
n = relative block number

Simulation of Sequential Access on a Direct-access File

- ✓ **Cp defines our current position**
- ✓ **We can simulate sequential file operations as follows:**

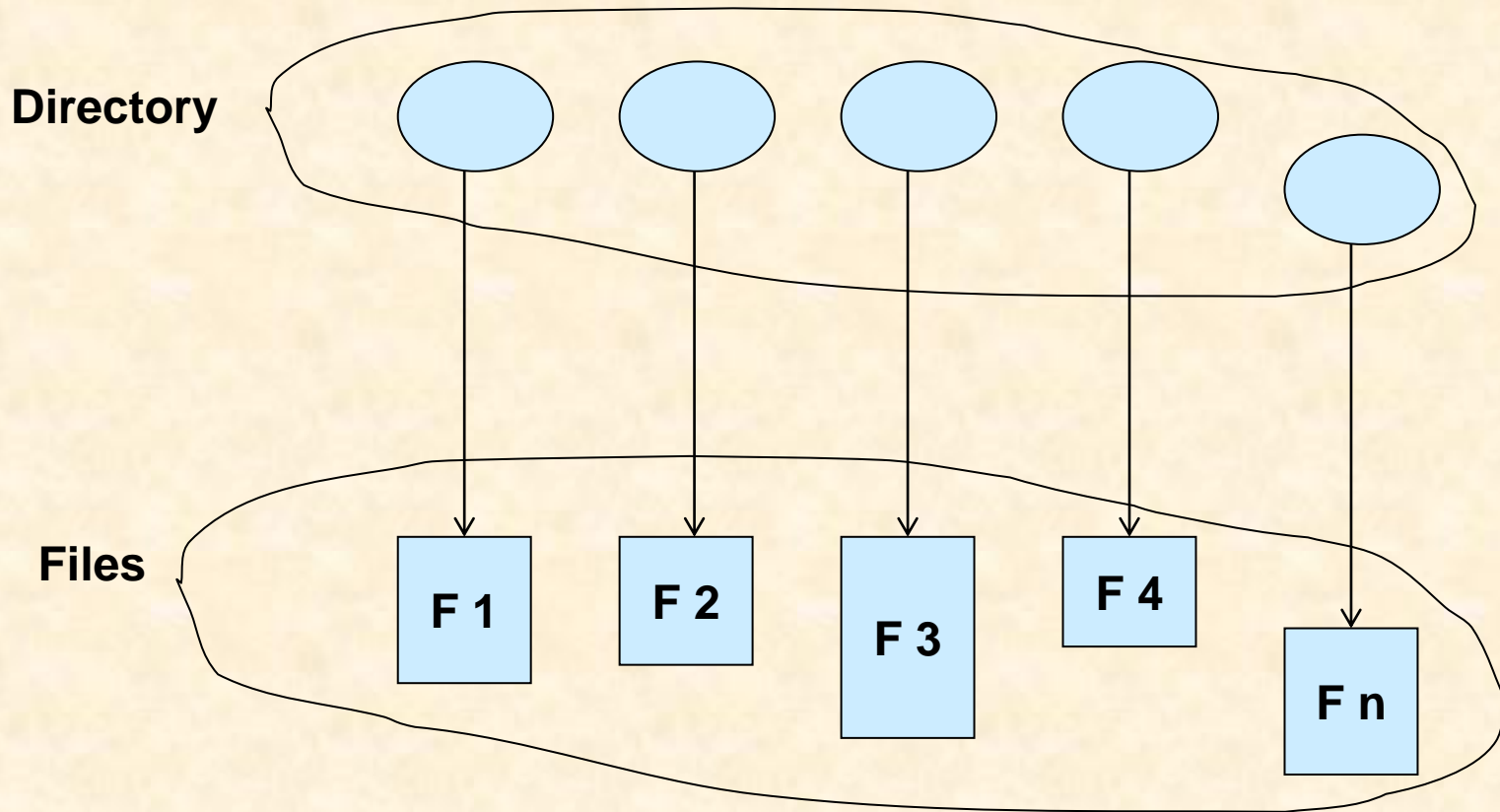
| sequential access | implementation for direct access |
|-------------------|---|
| <i>reset</i> | <i>cp = 0;</i> |
| <i>read next</i> | <i>read cp;</i> <i>cp = cp + 1;</i> |
| <i>write next</i> | <i>write cp;</i> <i>cp = cp + 1;</i> |

Example of Index and Relative Files

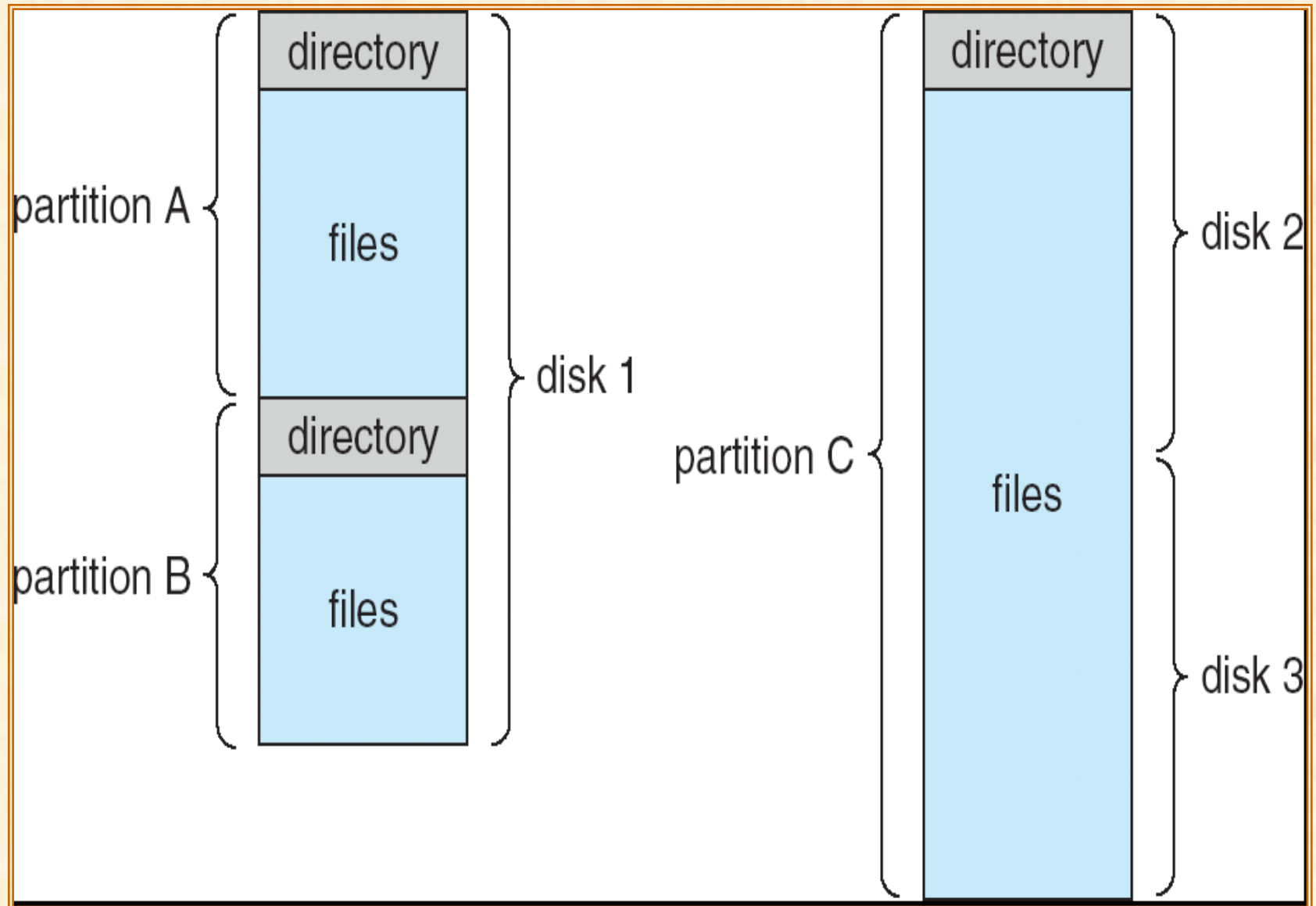


Directory Structure

- ✓ A collection of nodes containing information about all files
- ✓ Both the directory structure and the files reside on disk



A Typical File-system Organization



Operations Performed on Directory

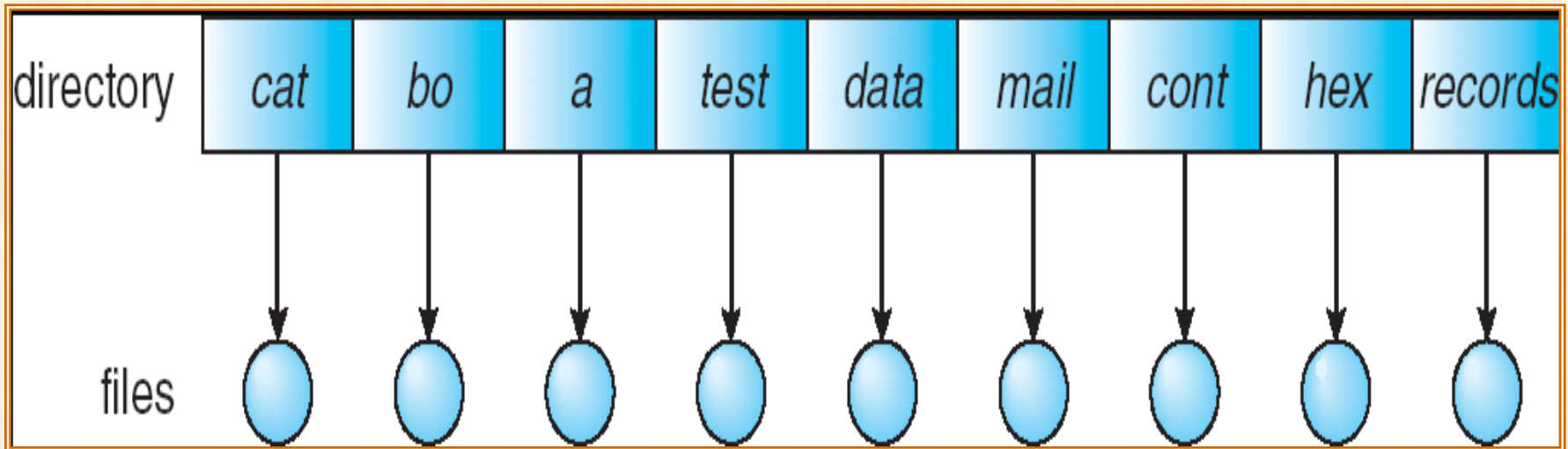
- ✓ **Search for a file**
- ✓ **Create a file**
- ✓ **Delete a file**
- ✓ **List a directory**
- ✓ **Rename a file**
- ✓ **Traverse the file system**

Organize the Directory (Logically) to obtain

- ✓ **Efficiency** – locating a file quickly
- ✓ **Naming** – convenient to users
 - Two users can have same name for different files
 - The same file can have several different names
- ✓ **Grouping** – logical grouping of files by properties, (e.g., all Java programs, all games, ...)

Single-Level Directory

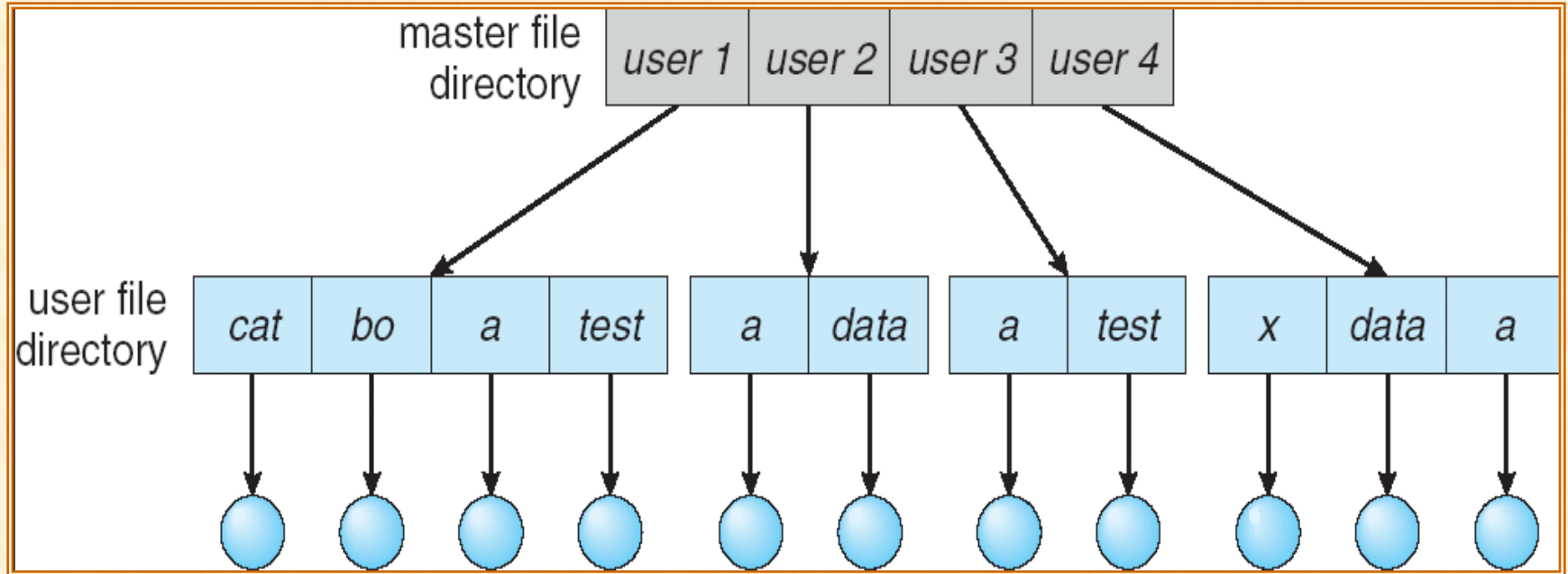
- ✓ Simplest directory structure
- ✓ A single directory for all users
- ✓ All files are contained in the same directory



- ✓ Naming problem
- ✓ Grouping problem

Two-Level Directory

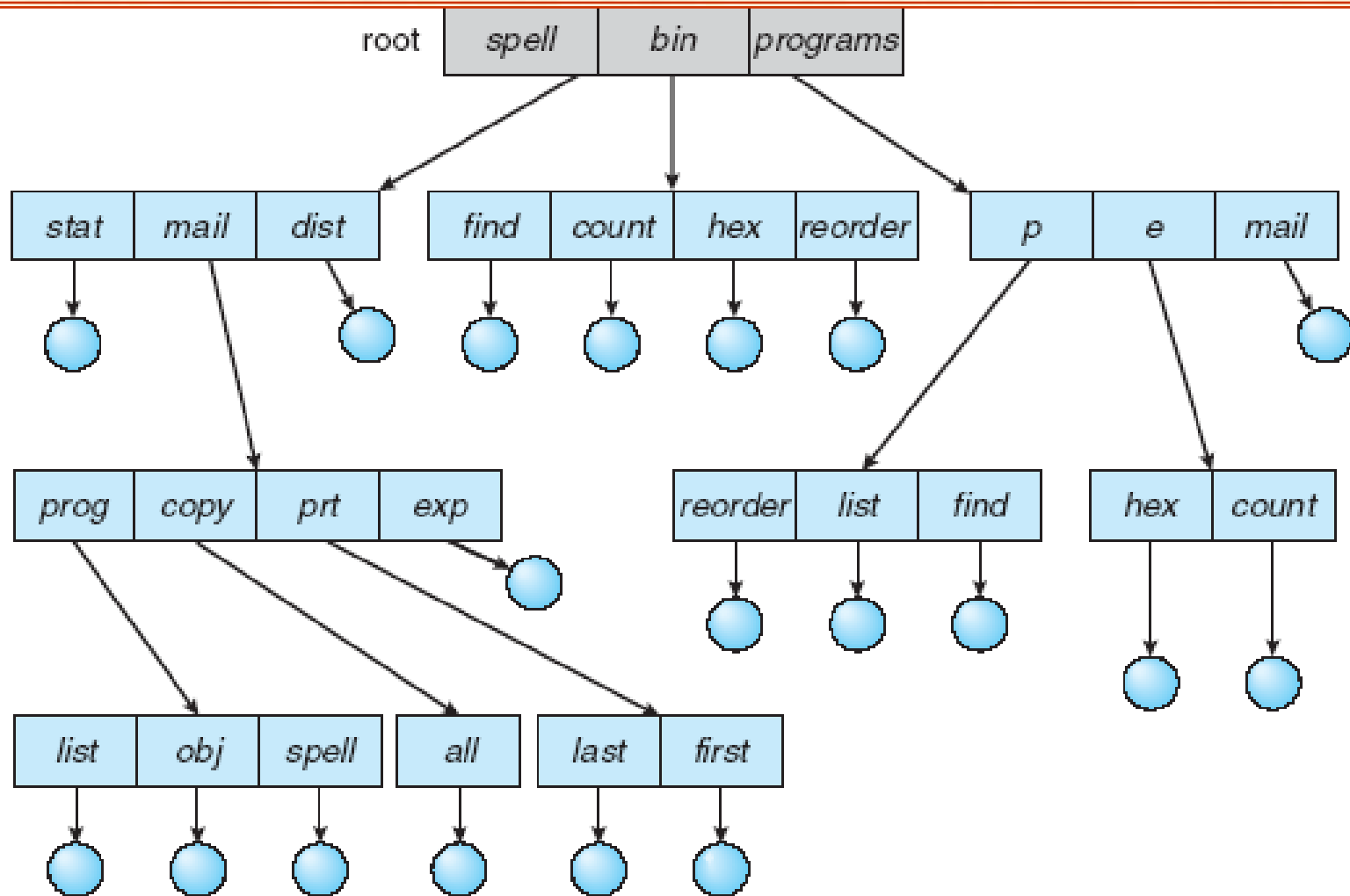
- ✓ **Separate directory for each user**



- ✓ **Path name**
- ✓ **Can have the same file name for different user**
- ✓ **Efficient searching**
- ✓ **No grouping capability**

Tree-Structured Directories

Most common directory structure

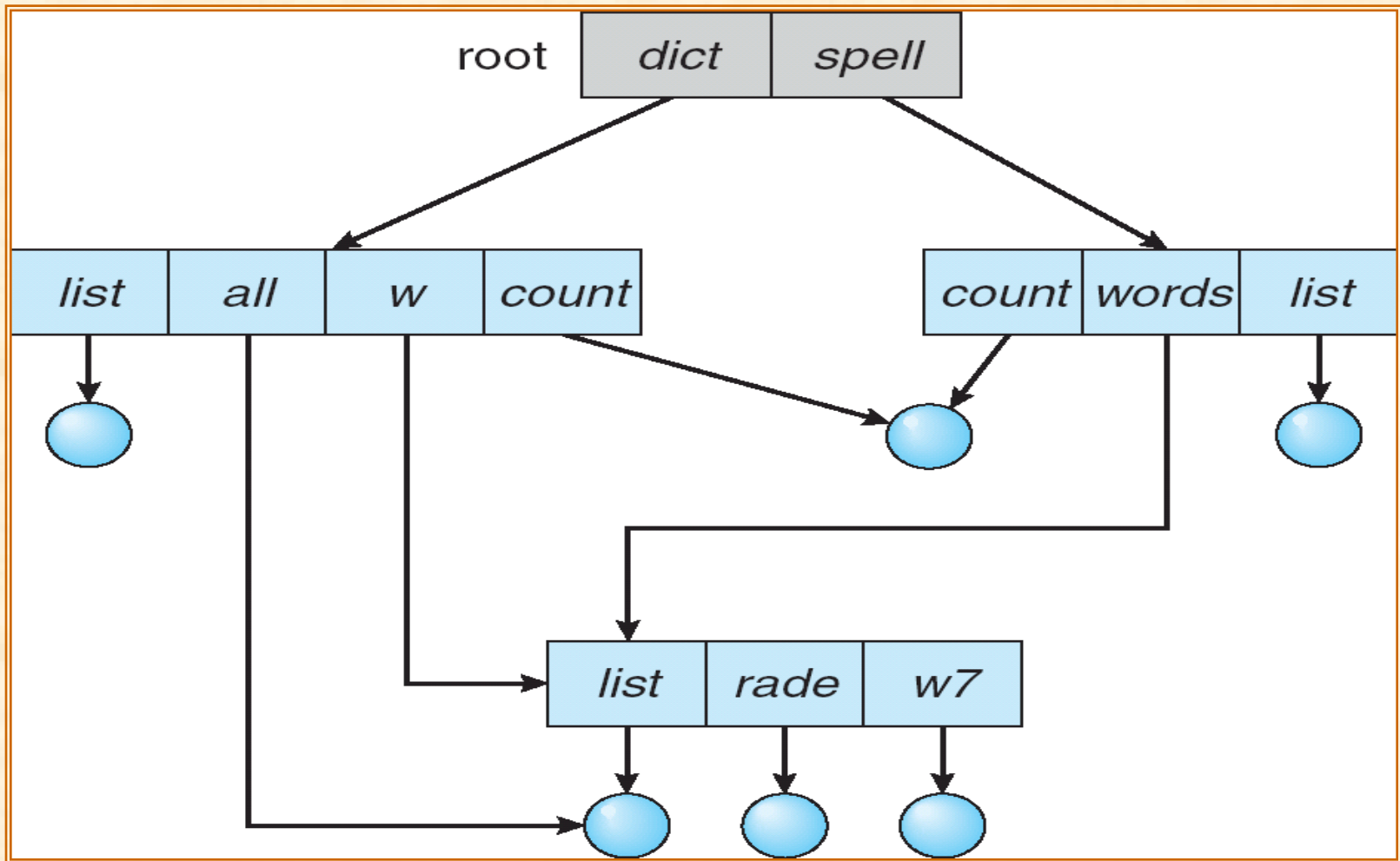


Path names

- ✓ **Path names can be of 2 types:**
 - **Absolute**
 - ◆ **Begins at the root and follows a path down to the specified file**
 - **Relative**
 - ◆ **Defines path from current directory**
- ✓ **Advantages of Tree-Structured Directories**
 - **Efficient searching**
 - **Grouping Capabilities**

Acyclic-Graph Directories

- ✓ **Allows directories to share subdirectories and files**



Acyclic-Graph Directories (Cont.)

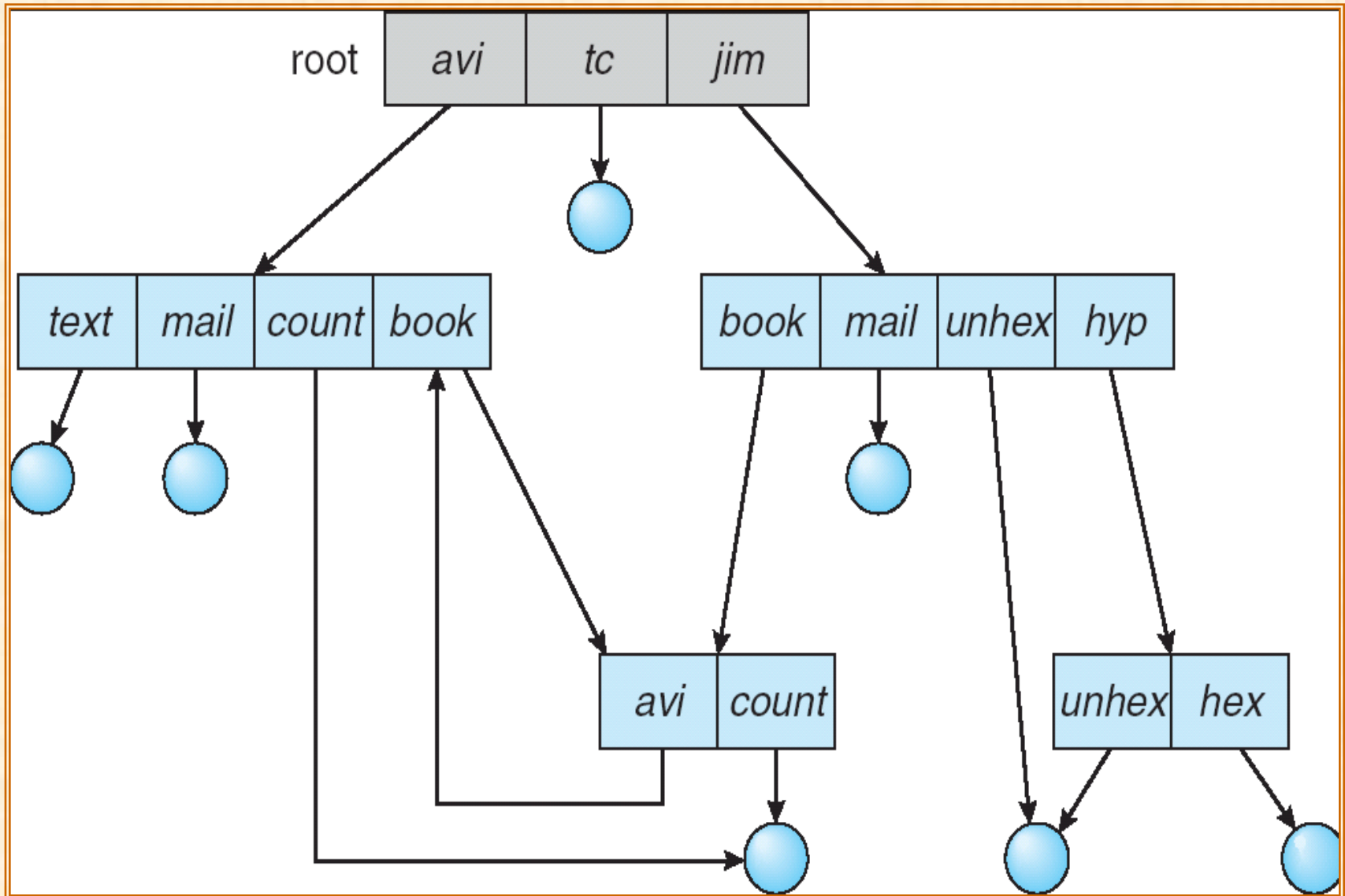
- ✓ **Tree structure prohibits the sharing of files or directories**
 - **Acyclic graph allows directories to have shared subdirectories and files**
- ✓ **Shared files and subdirectories can be implemented in several ways: a common way is :**
 - **Create a new directory called a *link***
 - ◆ **Link is effectively a pointer to another file or sub directory**
 - ◆ **Link may be implemented as absolute or relative path name (a symbolic link)**
 - **Duplicate all the information about them in both sharing directories**
 - ◆ **A link is clearly different from the original directory entry; thus 2 are not equal**
 - ◆ **It makes original and the copy indistinguishable**
 - ◆ **Problem:**
 - ★ **Maintaining consistency if the files is modified**

Acyclic-Graph Directories (Cont.)

✓ Disadvantages:

- File may now have multiple absolute path names, distinct file names may refer to the same file
 - ◈ Trying to traverse the entire file system becomes significant, since we do not want to traverse shared structures more than once
- Deletion
 - ◈ When can the space allocated to a shared file be deallocated and reused ?
 - ◈ Remove the file whenever anyone deletes it, this may leave **dangling pointers** to the now-nonexistent file
 - ◈ Different approaches:
 - ★ search for the links and remove them
 - ★ preserve the file until all references to it are deleted

General Graph Directory



General Graph Directory (Cont.)

- ✓ **Problem with using acyclic graph structure is ensuring that there are no cycles**
- ✓ **How do we guarantee no cycles?**
 - **Allow only links to file not subdirectories**
 - **Garbage collection**
 - **Every time a new link is added use a cycle detection algorithm to determine whether it is OK**

Protection

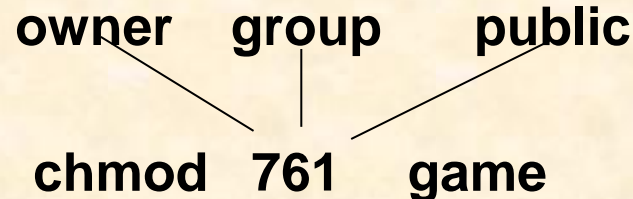
- ✓ When information is kept in computer system , a major concern is its:
 - Protection from both physical damage (reliability) and
 - Improper access (protection)
 - Reliability is provided by duplicate copies of files
 - Protection can be provided in many ways
 - ◆ For a small single user system we might provide protection by physically removing the floppy disks and locking them in a desk drawer or file cabinet
- ✓ Types of access
 - Read, Write , Execute , Append, Delete, List

Access Lists and Groups

- ✓ Mode of access: read, write, execute
- ✓ Three classes of users

| | | | RWX |
|------------------|---|---|-------|
| a) owner access | 7 | ⇒ | 1 1 1 |
| b) group access | 6 | ⇒ | 1 1 0 |
| c) public access | 1 | ⇒ | 0 0 1 |

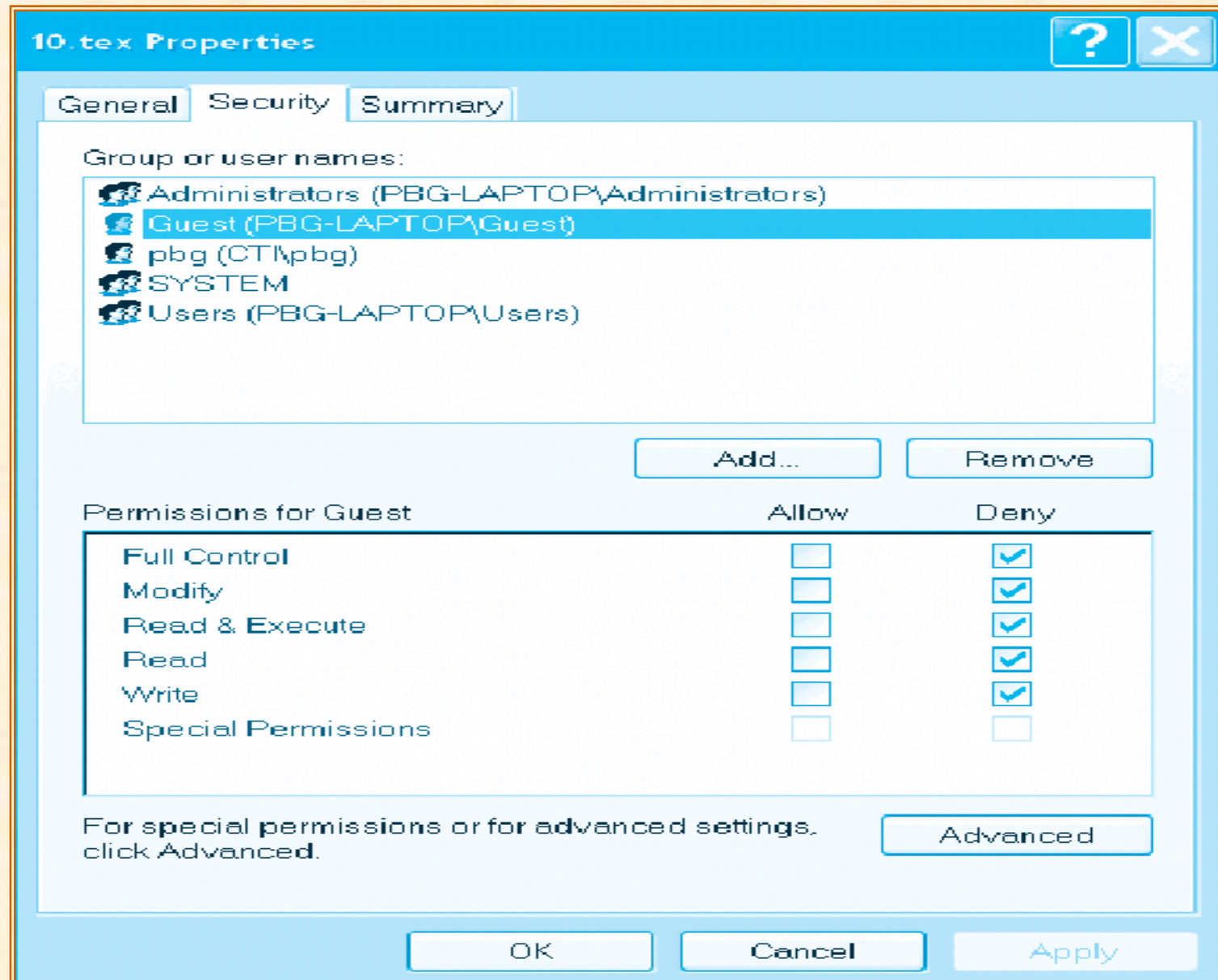
- ✓ Ask manager to create a group (unique name), say G, and add some users to the group.
- ✓ For a particular file (say *game*) or subdirectory, define an appropriate access.



Attach a group to a file

chgrp G game

Windows XP Access-control List Management



A Sample UNIX Directory Listing

| | | | | | |
|------------|-------|---------|-------|--------------|---------------|
| -rw-rw-r-- | 1 pbg | staff | 31200 | Sep 3 08:30 | intro.ps |
| drwx----- | 5 pbg | staff | 512 | Jul 8 09:33 | private/ |
| drwxrwxr-x | 2 pbg | staff | 512 | Jul 8 09:35 | doc/ |
| drwxrwx--- | 2 pbg | student | 512 | Aug 3 14:13 | student-proj/ |
| -rw-r--r-- | 1 pbg | staff | 9423 | Feb 24 2003 | program.c |
| -rwxr-xr-x | 1 pbg | staff | 20471 | Feb 24 2003 | program |
| drwx--x--x | 4 pbg | faculty | 512 | Jul 31 10:31 | lib/ |
| drwx----- | 3 pbg | staff | 1024 | Aug 29 06:52 | mail/ |
| drwxrwxrwx | 3 pbg | staff | 512 | Jul 8 09:35 | test/ |

Consistency Semantics

- ✓ Defines what happens when number of users try to access file simultaneously
- ✓ **Session** - set of reads and writes within one open and close file
- ✓ **UNIX**
 - A write to a file is immediately visible to other users
 - A number of users may share same read/write pointer
 - Implemented by a single image for each file
- ✓ **AFS** - Andrew File System (distributed)
 - A write is not seen immediately by other users
 - When file is closed, updates visible for sessions which start after file was closed
 - There are number of file images

Consistency Semantics

- ✓ **Immutable shared files semantics:**
 - **Once a file is declared as shared by its creator , it cannot be modified**

- ✓ **2 important properties:**
 - **Its name may not be reused and**
 - **Its contents may not be altered**

End of the Chapter