



SEARCHING ALGORITHMS

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Searching in Linear Array

- The process of finding a particular element of an array is called “Searching”.
- If the item is not present in the array, then the search is unsuccessful.
- There are two types of search,
 - Linear search
 - Binary Search

Linear Search

- The linear search compares each element of the array with the search key until the search key is found.
- To determine that a value is not in the array, the program must compare the search key to every element in the array.
- It is also called “Sequential Search” because it traverses the data sequentially to locate the element.

Linear Search

- **Algorithm: (Linear Search)**
- **LINEAR (A, SKEY)**
- Here **A** is a Linear Array with **N** elements and **SKEY** is a given item of information to search. This algorithm finds the location of SKEY in A and if successful, it returns its location otherwise it returns -1 for unsuccessful.

- 1. Repeat for $i = 0$ to $N-1$**
- 2. if($A[i] = SKEY$) return i [Successful Search]**
[End of loop]
- 3. return -1 [Un-Successful]**
- 4. Exit.**

Binary Search

- It is useful for the large sorted arrays.
- The binary search algorithm can only be used with ***sorted array and eliminates one half of the elements in the array*** being searched after each comparison.
- The algorithm locates the middle element of the array and compares it to the search key.
- If they are equal, the search key is found and array subscript of that element is returned.
- Otherwise the problem is reduced to searching one half of the array.
- If the search key is less than the middle element of array, the first half of the array is searched.
- If the search key is not the middle element of in the specified sub array, the algorithm is repeated on one quarter of the original array.
- The search continues until the sub array consist of one element that is equal to the search key (search successful).
- But if Search-key not found in the array then the value of END of new selected range will be less than the START of new selected range.

Binary Search Example

Search Key=22	
A[0]	3
A[1]	5
A[2]	9
A[3]	11
A[4]	15
A[5]	17
A[6]	22
A[7]	25
A[8]	32
A[9]	54

Start=0

End = 9

$\text{Mid} = \text{int}(\text{Start} + \text{End}) / 2$

$\text{Mid} = \text{int}(0 + 9) / 2$

Mid=4



Start=4+1 = 5

End = 9

$\text{Mid} = \text{int}(5 + 9) / 2 = 7$



Start = 5

End = 7 - 1 = 6

$\text{Mid} = \text{int}(5 + 6) / 2 = 5$



Start = 5+1 = 6

End = 6

$\text{Mid} = \text{int}(6 + 6) / 2 = 6$

Found at location 6

Successful Search

Binary Search Example

Search Key=8	
A[0]	3
A[1]	5
A[2]	9
A[3]	11
A[4]	15
A[5]	17
A[6]	22
A[7]	25
A[8]	32
A[9]	54

Start=0
End = 9
 $\text{Mid} = \text{int}(\text{Start} + \text{End}) / 2$
 $\text{Mid} = \text{int}(0 + 9) / 2$
Mid=4



Start=0
End = 3
 $\text{Mid} = \text{int}(0 + 3) / 2 = 1$



Start = 1+1 = 2
End = 3
 $\text{Mid} = \text{int}(2 + 3) / 2 = 2.5$



Start = 2
End = 2 - 1 = 1

End is < Start
Un-Successful Search

Binary Search Algorithm

- Here **A** is a sorted Linear Array with **N** elements and **SKEY** is a given item of information to search. This algorithm finds the location of **SKEY** in **A** and if successful, it returns its location otherwise it returns -1 for unsuccessful.
- Binary Search (A, SKEY)
 1. [Initialize segment variables.]
 - Set $START=0$, $END=N-1$ and $MID=INT((START+END)/2)$.
 2. Repeat Steps 3 and 4 while $START \leq END$ and $A[MID] \neq SKEY$.
 3. If $SKEY < A[MID]$. then
 - Set $END=MID-1$.
 - Else
 - Set $START=MID+1$.
 - [End of If Structure.]
 4. Set $MID=INT((START +END)/2)$.
 - [End of Step 2 loop.]
 5. If $A[MID]= SKEY$ then
 - Set $LOC= MID$
 - Else:
 - Set $LOC = -1$
 - [End of IF structure.]
 6. return LOC and Exit

Computational Complexity of Binary Search

- The *Computational Complexity of the Binary Search algorithm is measured* by the maximum (worst case) number of Comparisons it performs for searching operations.
- The searched array is divided by 2 for each comparison/iteration.
- Therefore, the maximum number of comparisons is measured by: *$\log_2(n)$ where n is the size of the array*
- **Example:**
- If a given sorted array 1024 elements, then the maximum number of comparisons required is:
- *$\log_2(1024) = 10$ (only 10 comparisons are enough)*

Computational Complexity of Linear Search

- Note that the *Computational Complexity of the Linear Search is the* maximum number of comparisons you need to search the array.
- As you are visiting all the array elements in the worst case, then, the number of comparisons required is:
- *n (n is the size of the array)*
- **Example:**
- If a given an array of 1024 elements, then the maximum number of comparisons required is:
- *$n-1 = 1023$ (As many as 1023 comparisons may be required)*

Thank You !

Questions ?

