

Chap 10. Signals

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Disclaimer: The slides are borrowed from many sources!

Signals

- Signals are software interrupts from unexpected events
 - o an illegal operation (e.g., divide by 0)
 - o a power failure
 - o an alarm clock
 - the death of a child process
 - a termination request from a user (Ctrl-C)
 - a suspend request from a user (Ctrl-Z)



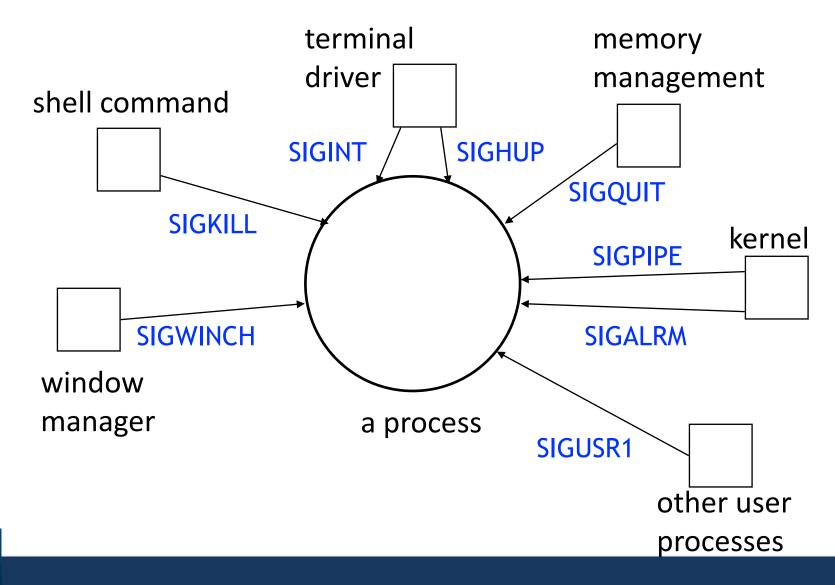


Predefined Signals (1/2)

- 31 signals
- Every signal has a name.
 - o begin with 'SIG'
 - SIGABRT: abort signal from abort()
 - SIGALRM: alarm signal from alarm()
- Actions of the default signal handler
 - terminate the process and generate a core (dump)
 - o ignores and discards the signal (ignore)
 - suspends the process (suspend)
 - resume the process
- Default signal handler can be overridden.



Signal Sources





Predefined Signals (2/2)

```
esjung — esjung@hpclab:/usr/include — ssh hpclab — 84×32
#define SIGTERM
                        15
#define SIGSTKFLT
                        16
#define SIGCHLD
                        17
#define SIGCONT
                        18
#define SIGSTOP
                        19
#define SIGTSTP
                        20
#define SIGTTIN
                        21
#define SIGTTOU
                        22
#define SIGURG
                        23
#define SIGXCPU
                        24
#define SIGXFSZ
                        25
#define SIGVTALRM
                        26
#define SIGPROF
                        27
#define SIGWINCH
                        28
#define SIGIO
                        29
#define SIGPOLL
                        SIGIO
#define SIGLOST
                        29
*/
#define SIGPWR
                        30
#define SIGSYS
                        31
#define SIGUNUSED
                        31
/* These should not be considered constants from userland. */
#define SIGRTMIN
                        32
#define SIGRTMAX
                       NSIG
 * SA FLAGS values:
* SA ONSTACK indicates that a registered stack t will be used.
                                                                   65,1
                                                                                 34%
```



Signal Generation

- Terminal-generated signals
 - CTRL-C → SIGINT
 - CTRL-Z → SIGSTP signal
- Hardware exceptions generate signals
 - o divide by 0 → SIGFPE
 - o invalid memory reference → SIGSEGV
- kill()
 - sends any signal to a process or process group
 - o need to be owner or super-user
- Software conditions
 - SIGALRM: alarm clock expires
 - SIGPIPE: broken pipe
 - SIGURG: out-of-band network data



Handling of Signals

Disposition or action:
 Process has to tell the kernel "if and when this signal occurs, do the following."

Ignore the signal:
 all signals can be ignored, except SIGKILL and SIGSTOP

 Let the default action apply: most are to terminate process



Representative UNIX Signals (1/2)

- SIGART: generated by calling the abort function.
- SIGALRM: generated when a timer set with the alarm expires.
- SIGCHLD: whenever a process terminates or stops, the signal is sent to the parent.
- SIGCONT: this signal sent to a stopped process when it is continued.
- SIGFPE: signals an arithmetic exception, such as divide-by-0, floating point overflow, and so on
- SIGILL: indicates that the process has executed an illegal hardware instruction.
- SIGINT: generated by the terminal driver when we type the interrupt key and sent to all processes in the foreground process group.



Representative UNIX Signals (2/2)

- SIGKILL: can't be caught or ignored. a sure way to kill any process.
- **SIGPIPE**: if we write to a pipeline but the reader has terminated, SIGPIPE is generated.
- SIGSEGV: indicates that the process has made an invalid memory reference.
 (→ core dumped)
- SIGTERM: the termination signal sent by the kill(1) command by default.
- SIGSTP: Cntl-Z from the terminal driver which is sent to all processes in the foreground process group.
- SIGUSR1: user defined signal 1
- SIGUSR2: user defined signal 2



signal()

- Signal Handler Registration
- signal(int signo, void(*func)()))
 - \circ specify the action for a signal (signo \rightarrow func)
 - o func
 - SIG_IGN (ignore)
 - SIG_DFL (default)
 - user-defined function
 - Return: the previous func



Example

```
int main()
    signal(SIGINT, foo);
    do usual things until SIGINT */
    return 0;
→void foo( int signo )
             /* deal with SIGINT signal */
   return; /* return to program */
```



Example: alarm2.c (w/ handler) (1/2)

```
#include <stdio.h> // alarm2.c
#include <signal.h>
int alarmFlag=0;
void alarmHandler();
main() {
    signal(SIGALRM, alarmHandler);
    alarm(3);
    printf("Looping ...\n");
    while(!alarmFlag) {
        pause();
    printf("Loop ends due to alarm signal \n");
void alarmHandler() {
    printf("An alarm clock signal was received\n");
    alarmFlag = 1;
```



Example: alarm2.c (w/ handler) (2/2)

Execution

```
[[esjung@hpclab opensource]$ gcc -o alarm2 alarm2.c
[[esjung@hpclab opensource]$ ./alarm2
Looping ...
An alarm clock signal was received
Loop ends due to alarm signal
[esjung@hpclab opensource]$
```



SIGCHLD

 Whenever a process terminates or stops, the signal is sent to the parent.

 When a child process is killed, it sends SGICHILD signal to its parent process



Example: timelimit.c (1/3)

\$ timelimit N command // perform "command" within N seconds

```
#include <stdio.h> // timelimit.c
#include <signal.h>
int delay;
void childHandler();
main(int argc, char *argv[])
{
    int pid;
    sscanf(argv[1], "%d", &delay);
    signal(SIGCHLD, childHandler);
    pid = fork();
    if (pid == 0) { // child
        execvp(argv[2], &argv[2]);
       perror("Limit");
    } else { // parent
        sleep(delay);
        printf("Child %d exceeded limit and is being killed\n", pid);
        kill(pid, SIGINT);
```



Example: timelimit.c (2/3)

```
childHandler() /* Executed if the child dies before the parent */
{
   int childPid, childStatus;
   childPid = wait(&childStatus);
   printf("Child %d terminated within %d seconds\n", childPid, delay);
   exit(0);
}
```



Example: timelimit.c (3/3)

Execution

```
[esjung@hpclab opensource]$ ./timelimit 5 ls
7Process alarm2.c doatexit.c lab3
                                       pipe2.c rdfile timelimit
                                                                      wtfile
alarm2
         a.out
                   lab2
                              midterm
                                       pipe.c
                                                shfile.c timilimit.c
Child 170561 terminated within 5 seconds
[esjung@hpclab opensource]$ ./timelimit 5 sleep 3
Child 170563 terminated within 5 seconds
[esjung@hpclab opensource]$ ./timelimit 5 sleep 100
Child 170573 exceeded limit and is being killed
[esjung@hpclab opensource]$
```



Multiple Signals

- If many signals of the *same* type are waiting to be handled (e.g. two SIGINTs), then most UNIXs will only deliver one of them.
 - the others are thrown away

If many signals of *different* types are waiting to be handled (e.g. a SIGINT, SIGSEGV, SIGUSR1), they are not delivered in any fixed order.



The Reset Problem in early System V UNIX

- In Linux (and many other UNIXs), the signal disposition in a process is reset to its default action immediately after the signal has been delivered.
- Must call signal() again to reinstall the signal handler function.



Reset Problem Example

```
int main()
   signal(SIGINT, foo);
   /* do usual things until SIGINT */
void foo(int signo)
   signal(SIGINT, foo); /* reinstall */
   return;
```



Reset Problem

```
void ouch( int sig )
       printf( "OUCH! - I got signal %d\n", sig );
       (void) signal(SIGINT, ouch);
                                   Problem: from the time that the interrupt
                                    Function starts to just before the
int main()
                                    signal handler is re-established
                                   the signal will not be handled.
       (void) signal( SIGINT, ouch );
       while (1)
              printf("Hello World!\n");
              sleep(1);
```

To keep catching the signal with this function, must call the signal system call again.

If another SIGINT signal is received during this time, default behavior will be done, i.e., program will terminate.

Re-installation may be too slow!

- There is a (very) small time period in foo() when a new SIGINT signal will cause the default action to be carried out -- process termination.
- POSIX, BSD signal functions solve it (and some other later UNIXs)



Modification in BSD 4.x signal environment

- Persistent Handlers
 - Signal handlers remain installed even after the signal occurs and do not need to be explicitly reinstalled.
- Masking
 - Signals are blocked for the duration of a signal handler (i.e. recursive signals are not normally allowed).
 - A "signal mask" can be set to block most signals during critical regions.
- Signal handlers normally remain installed during and after signal delivery.



kill(), raise()

• kill - sends a signal to a process or a group of process

• raise - function allows a process to send a signal to itself



kill()

pid means

- o pid > 0 : signal to the process whose process ID is pid
- pid == 0: signal to the processes whose process group ID equals that of sender
- o pid < 0 : signal to the processes whose process group ID equals abs. of pid
- pid == -1 : unspecified (used as a broadcast signal in SVR4, 4.3 + BSD)

Permission to send signals

- The super-user can send a signal to any process.
- The real or effective user ID of the sender has to equal the real or effective user ID of the receiver.



alarm()

```
#include <unistd.h>
unsigned int alarm (unsigned int seconds);
    Returns: 0 or number of seconds until previously set alarm
```

- alarm() sets a timer to expire at a specified time in future.
 - when timer expires, SIGALRM signal is generated,
 - default action of the signal is to terminate the process.
- Only one alarm clock per process
 - previously registered alarm clock is replaced by the new value.
- if alarm(0), a previous unexpired alarm is cancelled.



pause()

- suspends the calling process until a signal is caught.
- returns only if a signal handler is executed and that handler returns.
 - If signal handler is not registered, just quit
 - If signal handler is registered, return after the handler is processed.



abort()

```
#include <stdlib.h>
void abort(void);
    This function never returns
```

Causes abnormal program termination.

This function sends the SIGABRT signal to the process.

 SIGABRT signal handler to perform any cleanup that it wants to do, before the process terminated.



sleep()

- This function causes the calling process to be suspended until either
 - The amount of wall clock time specified by second has elapsed (returns 0)
 - A signal is caught by the process and the signal handler returns (returns the number of unslept seconds)



Process group

 Group of related processes that should receive a common signal for certain events

Identified by Process group ID

- Example
 - grp = setpgrp();
 - Initialize process group id, the same as its pid
 - After fork(), child process will have parent's process group id



setpgrp()

```
#include <signal.h>
#include <stdio.h>
main(){
    register int i;
    setpgrp();
    for(i=0;i<10;i++){
        if(fork()==0){
            /* child process */
            if(i&1)
                setpgrp();
            printf("pid = %d pgrp= %d\n", getpid(),getpgrp());
            pause();
    kill(0,SIGINT);
```



Results

```
[esjung@hpclab opensource]$ ./a.out

pid = 100672 pgrp= 100671

pid = 100673 pgrp= 100673

pid = 100674 pgrp= 100671

pid = 100676 pgrp= 100675

pid = 100675 pgrp= 100675

pid = 100677 pgrp= 100677

pid = 100678 pgrp= 100671

pid = 100680 pgrp= 100671
```



Treatment of Signals

- Sending a signal to a process
 - issig() : check for receipt of signals

- Handling a signal
 - psig(): handle signals after recognizing their existence



How to Send a Signal to Process

Send a Signal

- sets a bit in the signal field in process table entry.
- process can remember different type of signal.
- process can not remember how many signals it receives of particular type.
- u area contains an array of signal-handler fields.
 kernel stores the address of the user-function in the field.

Check for Signal when

- about to return from kernel mode to user mode
- enters or leaves the sleep state at a suitably low scheduling priority.

Handle Signals

only when returns from kernel mode to user mode

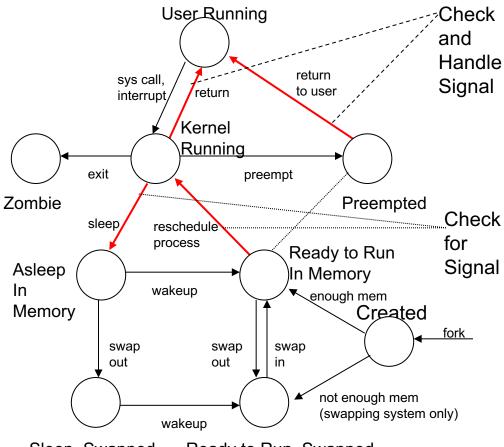


How to Send a Signal to Process(Cont.)

no effect on a process running in the kernel mode

a process never executes in user mode before handling outstanding

signals





Ready to Run, Swapped



Algorithm for issig

```
algorithm issig
                        /* test for receipt of signals */
input : none
output: true, if process receives signals that it does not ignore
          false otherwise
{
  while (received signal field in process table entry not 0)
   {
        find a signal number sent to the process;
        if ( signal is death of child )
                if (ignoring death of child signals)
                        free process table entries of zombie children;
                else if (catching death of child signals)
                        return(true);
        else if (not ignoring signal)
                return(true);
        turn off signal bit in received signal field in process table;
   return(false);
```

Algorithm for Handling Signal

- 1. determines signal type
- 2. turns off signal bit in the process table entry
- 3. if receives a signal to ignore
 - continues as if the signal has never occurred.
- 4. If signal handling function is set to its default value,
 - kernel dumps core image for signals that imply something is wrong with process and exit
 - kernel does not dump core for signals that do not imply error.
- 5. If receives a signal to catch,
 - accesses the user saved register context, and find the program counter and stack pointer
 - clears the signal handler field in u area(undesirable side-effects)
 - kernel creates a new stack frame and writes a program counter and stack pointer from user saved register context
 - kernel changes the user register context: program counter to the address of signal catcher function, and stack pointer to account for the growth of the user stack.



Algorithm for psig

```
algorithm psig
                        /* handle signals after recognizing their
  existence */
input : none
output : none
  get signal number set in process table entry;
  clear signal number in process table entry;
   if (user had called signal sys call to ignore this signal)
        return:
                       /*done*/
   if (user specified function to handle this signal)
       get user virtual address of signal catcher stored in u area;
       /*the next statement has undesirable side-effects*/
        clear u area entry that stored address of signal catcher;
       modify user level context;
                artificially create user stack frame to mimic
                call to signal catcher function;
       modify system level context;
                write address of signal catcher into program
                counter filed of user saved register context;
        return:
```



Algorithm for psig (cont.)

```
if (signal is type that system should dump core image of
  process)
{
    create file named "core" in current directory;
    write contents of user level context to file "core";
}
invoke exit algorithm immediately;
}
```

