CLP(B):

Constraint Logic Programming over Boolean Variables

:- use module(library(clpb)).

1 Introduction

This library provides CLP(B), Constraint Logic Programming over Boolean variables. It can be used to model and solve combinatorial problems such as verification, allocation and covering tasks.

CLP(B) is an instance of the general $CLP(\cdot)$ scheme, extending logic programming with reasoning over specialised domains.

The implementation is based on reduced and ordered Binary Decision Diagrams (BDDs).

2 Boolean expressions

A Boolean expression is one of:

0	false
1	true
variable	unknown truth value
atom	universally quantified variable
~ Expr	logical NOT
Expr + Expr	logical OR
Expr * Expr	logical AND
Expr # Expr	exclusive OR
Var ^ Expr	existential quantification
Expr = := Expr	equality
$Expr = \ Expr$	disequality (same as #)
Expr = < Expr	less or equal (implication)
Expr >= Expr	greater or equal
Expr < Expr	less than
Expr > Expr	greater than
card(Is,Exprs)	see below
+(Exprs)	see below
*(Exprs)	see below

where Expr again denotes a Boolean expression.

The Boolean expression card (Is, Exprs) is true iff the number of true expressions in the list Exprs is a member of the list Is of integers and integer ranges of the form From-To.

+ (Exprs) and \star (Exprs) denote, respectively, the disjunction and conjunction of all elements in the list *Exprs* of Boolean expressions.

Atoms denote parametric values that are universally quantified. All universal quantifiers appear implicitly in front of the entire expression. In residual goals, universally quantified variables always appear on the right-hand side of equations. Therefore, they can be used to express functional dependencies on input variables.

3 Interface predicates

 $\mathbf{sat}(+Expr)$ [semidet]

True iff *Expr* is a satisfiable Boolean expression.

taut(+Expr, -T) [semidet]

Succeeds with T = 0 if the Boolean expression Expr cannot be satisfied, and with T = 1 if Expr is always true with respect to the current constraints. Fails otherwise.

labeling(+Vs) [multi]

Assigns truth values to the Boolean variables *Vs* such that all stated constraints are satisfied.

 $\mathbf{sat_count}(+Expr, -N)$ [det]

N is the number of different assignments of truth values to the variables in the Boolean expression *Expr*, such that *Expr* is true and all posted constraints are satisfiable.

weighted_maximum(+Weights, +Vs, -Maximum) [multi]

Maximize a linear objective function over Boolean variables Vs with integer coefficients Weights. This predicate assigns 0 and 1 to the variables in Vs such that all stated constraints are satisfied, and Maximum is the maximum of $\sum w_i \cdot v_i$ over all admissible assignments. On backtracking, all admissible assignments that attain the optimum are generated.

This predicate can also be used to *minimize* a linear Boolean program, since negative integers can appear in *Weights*.