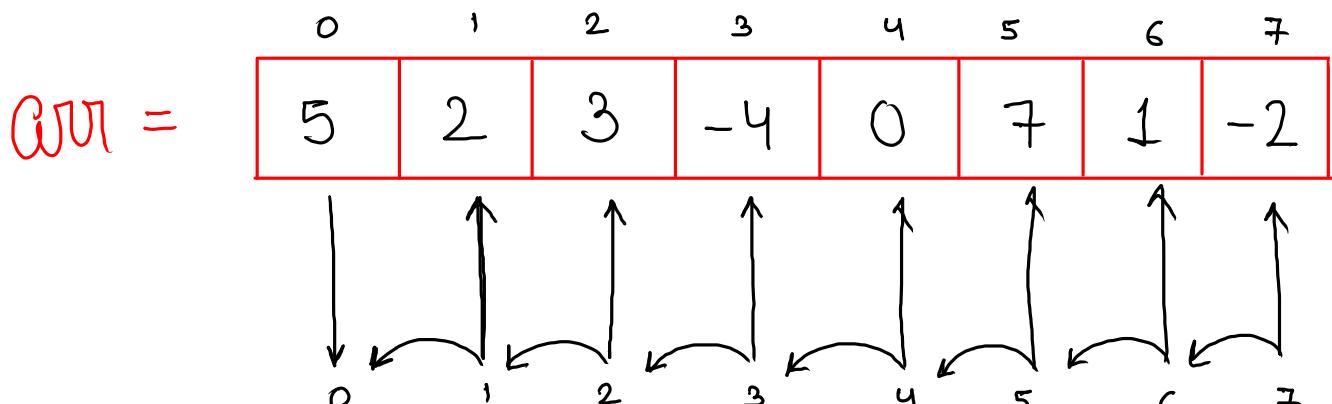


\Rightarrow Prefix Sum (sum of all the elements on the left side including itself)



prefix sum array

eg

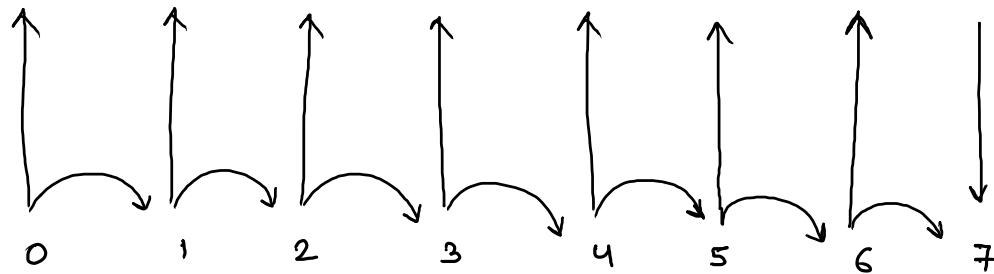
$$\text{prefix}[i] = \text{prefix}[i-1] + \text{arr}[i]$$

\Rightarrow Suffix Sum

(sum of all the elements on
the right side including itself)

$\text{arr} =$

0	1	2	3	4	5	6	7
5	2	3	-4	0	7	1	-2



$\text{suffix} =$
sum
array

0	1	2	3	4	5	6	7
12	7	5	2	6	6	-1	-2

eqn

$$\text{Suffix}[i] = \text{suffix}[i+1] + \text{arr}[i]$$

Code for prefix sum

```
pre[0] = arr[0];
```

```
for( int i=1 ; i<n ; i++ ) {
```

```
    pre[i] = pre[i-1] + arr[i];
```

y

code for suffix sum

$$\text{suffix}[n-1] = \text{arr}[n-1];$$

for (int i = n-2; i >= 0; i--) {

$$\text{suffix}[i] = \text{suffix}[i+1] + \text{arr}[i];$$

y

Dry Run

$\text{suffix}[n-1] = \text{arr}[n-1];$

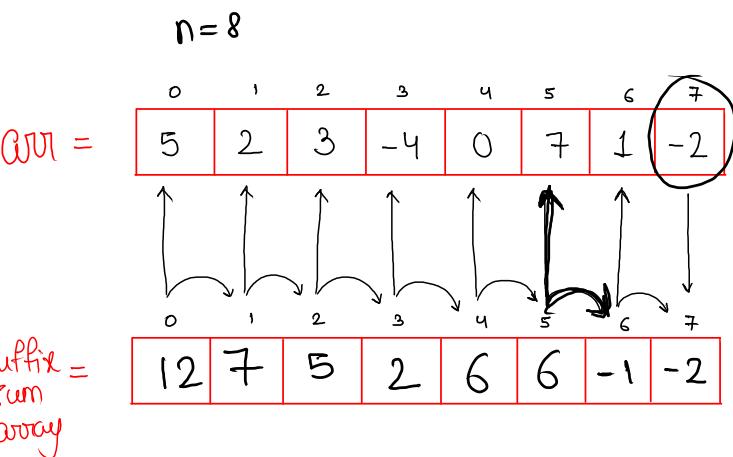
```
for( int i=n-2; i>=0; i-- ) {
    suffix[i] = suffix[i+1] + arr[i];
}

```

Note :-

$T.C = O(n)$

$S.C = O(n)$



$$i=6, \quad \text{suf}[6] = \text{suf}[7] + \text{arr}[6]$$

$$i=5, \quad \text{suf}[5] = \text{suf}[6] + \text{arr}[5]$$

$$i=4, \quad \text{suf}[4] = \text{suf}[5] + \text{arr}[4]$$

$$i=3, \quad \text{suf}[3] = \text{suf}[4] + \text{arr}[3]$$

$$i=2, \quad \text{suf}[2] = \text{suf}[3] + \text{arr}[2]$$

$$i=1, \quad \text{suf}[1] = \text{suf}[2] + \text{arr}[1]$$

$$i=0, \quad \text{suf}[0] = \text{suf}[1] + \text{arr}[0]$$

Code

```
public static void main(String[] args) {  
    int[] arr = {5, -4, 0, -1, 3};  
    int n = arr.length;  
    int[] prefix = new int[n];  
    prefix[0] = arr[0];  
    for (int i = 1; i < n; i++) {  
        prefix[i] = prefix[i - 1] + arr[i];  
    }  
    // print  
    for (int i = 0; i < n; i++) {  
        System.out.print(prefix[i] + " ");  
    }  
}
```

Print Prefix Sum

```
public static void main(String[] args) {  
    Scanner scn = new Scanner(System.in);  
    int n = scn.nextInt();  
    int[] arr = new int[n];  
    for (int i = 0; i < n; i++) {  
        arr[i] = scn.nextInt();  
    }  
    int[] ans = prefixSum(arr, n);  
    for (int i = 0; i < n; i++) {  
        System.out.println(ans[i]);  
    }  
}  
public static int[] prefixSum(int[] arr, int n) {  
    int[] pre = new int[n];  
    pre[0] = arr[0];  
    for (int i = 1; i < n; i++) {  
        pre[i] = pre[i - 1] + arr[i];  
    }  
    return pre;  
}
```

$$T.C = O(n)$$

$$S.C = O(n)$$

where n is size of array

Greatest Till Me (maximum till the current index)

$$\text{arr} = [5, 2, 3, 7, -4, 9, 0]$$

0 1 2 3 4 5 6

prefix max array = $[5, 5, 5, 7, 7, 9, 9]$

0 1 2 3 4 5 6

~~egⁿ)~~

$$\text{pre}[i] = \text{Math.max}(\text{pre}[i-1], \text{arr}[i]);$$

Code

```
public static void main(String[] args) {
    Scanner scn = new Scanner(System.in);
    int n = scn.nextInt();
    int[] arr = new int[n];
    for (int i = 0; i < n; i++) {
        arr[i] = scn.nextInt();
    }
    int[] ans = greatestTillMe(arr, n);
    for (int i = 0; i < n; i++) {
        System.out.println(ans[i]);
    }
}
public static int[] greatestTillMe(int[] arr, int n) {
    int[] pre = new int[n];
    pre[0] = arr[0];
    for (int i = 1; i < n; i++) {
        pre[i] = Math.max( pre[i - 1], arr[i] );
    }
    return pre;
}
```

dry run

$$\text{arr} = [5, 2, 3, 7, -4, 9, 0]$$
$$\text{pre} = \underline{\underline{[5, 5, 5, 7, 7, 9, 9]}}$$

- $i=1, \text{pre}[i] = \max(5, 2);$
- $i=2, \text{pre}[i] = \max(5, 3);$
- $i=3, \text{pre}[i] = \max(5, 7);$
- $i=4, \text{pre}[i] = \max(7, -4);$
- $i=5, \text{pre}[i] = \max(7, 9);$
- $i=6, \text{pre}[i] = \max(9, 0);$

Find Pivot Index 1

(9mp)

ans =

0	1	2	3	4	5	6	7
5	-3	2	7	1	1	1	1

$$\begin{array}{cccc} \downarrow & \downarrow & \downarrow & \downarrow \\ l=0 & l=5 & l=2 & l=4 \\ r=10 & r=13 & r=11 & r=4 \end{array}$$

pivot
index

ans = 3

$\text{arr} =$

0	1	2	3	4	5	6	7
5	-3	2	7	1	1	1	1

$\text{prefix sum array (pre)}$ =

0	1	2	3	4	5	6	7
5	2	4	11	12	13	14	15

$\text{suffix sum array (suf)}$ =

0	1	2	3	4	5	6	7
15	10	13	11	4	3	2	1

pseudo
code

1) create prefix array

$$\text{pre}[0] = \text{arr}[0];$$

loop from 1 to n

$$\text{pre}[i] = \text{pre}[i-1] + \text{arr}[i]$$

2) create suffix array

$$\text{suf}[n-1] = \text{arr}[n-1]$$

loop from (n-2) to 0

$$\text{suf}[i] = \text{suf}[i+1] + \text{arr}[i]$$

3) loop from 0 to n

check if $\text{pre}[i] == \text{suf}[i]$

return i;

return -1;

Code

```
public static int pivotIndex(int[] arr, int n) {  
    int[] pre = new int[n];  
    pre[0] = arr[0];  
    for (int i = 1; i < n; i++) {  
        pre[i] = pre[i - 1] + arr[i];  
    }  
  
    int[] suf = new int[n];  
    suf[n - 1] = arr[n - 1];  
    for (int i = n - 2; i >= 0; i--) {  
        suf[i] = suf[i + 1] + arr[i];  
    }  
  
    for (int i = 0; i < n; i++) {  
        if (pre[i] == suf[i]) {  
            return i;  
        }  
    }  
    return -1;  
}
```

Operations = $3 \times n$
 $T.C = O(n)$

$m/m = 2n$
 $S.C = O(n)$

where n is
size of array