

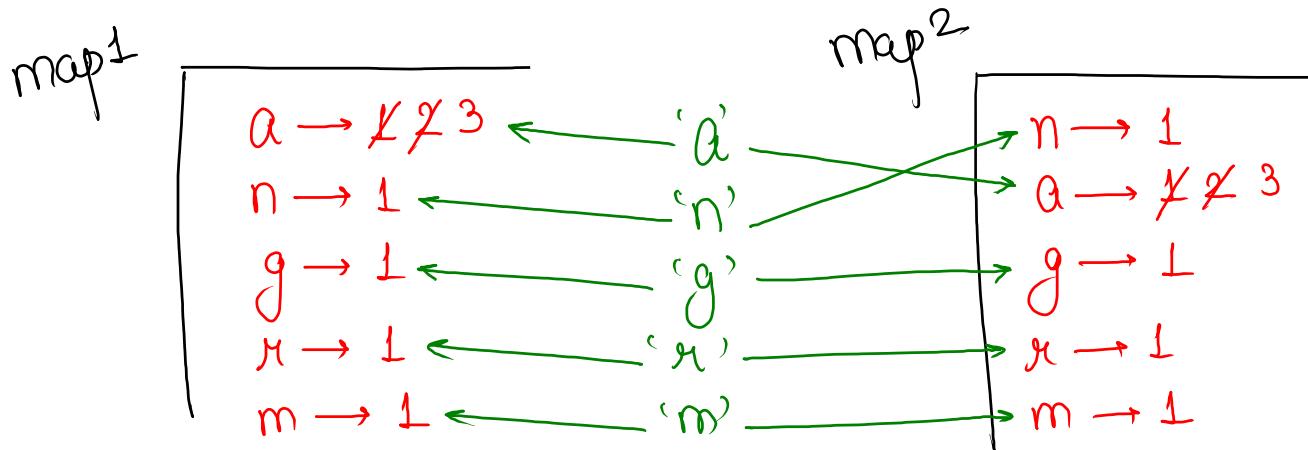
## Valid Anagram 5

→ Anagram :- when 2 strings are have same character , appearing same no. of time

Idea

$s = \text{"anagram"}$

$t = \text{"nagaram"}$



## pseudo code

- 1) declare map1
- 2) traverse in first string
  - 2.1) check if curr char is present in map  
then find oldfreq  
and put curr element back with oldfreq+1
  - 2.2) else  
put curr element with freq 1
- 3) declare map2
- 4) traverse in second string
  - 4.1) check if curr char is present in map  
then find oldfreq  
and put curr element back with oldfreq+1
  - 4.2) else  
put curr element with freq 1
- 5) traverse in map1  
if check if curr key is present in map2  
and freq of both should be same  
else return false
- 6) return true.

Code

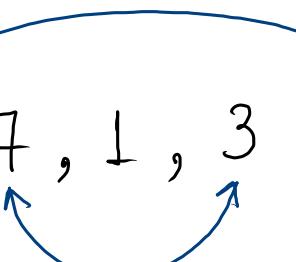
```
public static boolean checkAnagram(String s, String t) {  
    HashMap<Character, Integer> map1 = new HashMap<>();  
    for (char c : s.toCharArray()) {  
        if (map1.containsKey(c)) {  
            int oldFreq = map1.get(c);  
            map1.put(c, oldFreq + 1);  
        } else {  
            map1.put(c, 1);  
        }  
    }  
  
    HashMap<Character, Integer> map2 = new HashMap<>();  
    for (char c : t.toCharArray()) {  
        if (map2.containsKey(c)) {  
            int oldFreq = map2.get(c);  
            map2.put(c, oldFreq + 1);  
        } else {  
            map2.put(c, 1);  
        }  
    }  
  
    for (Map.Entry<Character, Integer> e : map1.entrySet()) {  
        char ch = e.getKey();  
        int freq = e.getValue();  
        if (!map2.containsKey(ch)) {  
            return false;  
        }  
        if (freq != map2.get(ch)) {  
            return false;  
        }  
    }  
    return true;  
}
```

$$T.C = O(n+m)$$

$$S.C = O(n+m)$$

## Max Number of K-Sum Pairs

$\text{arr} = [4, 7, 1, 3, 6, 8, 5]$ ,  $K = 10$



Count = 2

(Note:- no repeated elements)

$$\boxed{\text{arr}[i] + \text{arr}[j] == K}$$

$$\text{arr}[j] == K - \text{arr}[i]$$

Solution :- Consider every element  $\text{arr}$  as  $\text{arr}[i]$ , and try to find  $\text{arr}[j]$  using hashmap.

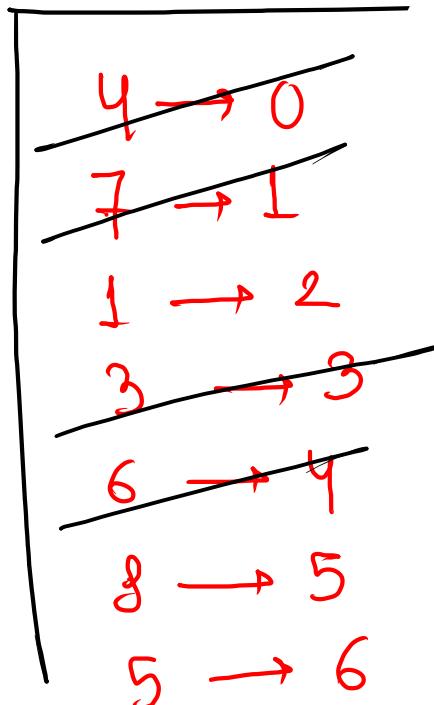
$\text{arr} = [ \begin{smallmatrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 \\ 4, 7, 1, 3, 6, 8, 5, 2, 9, 0 \end{smallmatrix} ]$ ,  $\underline{\underline{k}} = 10$

$\uparrow \quad \uparrow \quad \uparrow \quad \uparrow \quad \uparrow \quad \uparrow \quad \uparrow$

$i$

~~map~~

$\text{Count} = \cancel{0} \cancel{+} 2$



$\text{arr}[i]$	$\frac{\text{arr}[j]}{\text{arr}[i]} (k - \text{arr}[i])$
$i = 0, 4$	6
$i = 1, 7$	3
$i = 2, 1$	9
$i = 3, 3$	7
$i = 4, 6$	4
$i = 5, 8$	2
$i = 6, 5$	5

Code

```
public static void main(String[] args) {
    Scanner scn = new Scanner(System.in);
    int n = scn.nextInt();
    int[] arr = new int[n];
    for (int i = 0; i < n; i++) {
        arr[i] = scn.nextInt();
    }
    int k = scn.nextInt();
    int ans = maxKpairs(arr, n, k);
    System.out.println(ans);
}

public static int maxKpairs(int[] arr, int n, int k) {
    HashMap<Integer, Integer> map = new HashMap<>();
    for (int i = 0; i < n; i++) {
        map.put(arr[i], i);
    }

    int count = 0;
    for (int i = 0; i < n; i++) {
        int num1 = arr[i];
        int num2 = k - num1;
        if (map.containsKey(num2)) {
            if (i != map.get(num2)) {
                count++;
                map.remove(num1);
                map.remove(num2);
            }
        }
    }
    return count;
}
```

$$T.C = O(n)$$

$$S.C = O(n)$$

arr = [ 7, 2, 3, -2, 5, 12 ], k = 10

map	Count = 0 ✓ 2
7 → 0	i = 0, (7, 3) ✓
2 → 1	i = 1, (2, 8) ✗
3 → 2	i = 2, (3, 7) ✗
-2 → 3	i = 3, (-2, 12) ✓
5 → 4	i = 4, (5, 5) ✗
12 → 5	i = 5, (12, -2) ✗

~~→ Queue~~

[FIFO:- first in first out]

(dynamic)

doesn't have any indexing

structure

(open from both ends)



Note:- we can add elements from rear end &  
remove elements from front end only

## Syntax

Queue<DataType> que = new LinkedList<>()

Queue<DataType> que = new ArrayList<>()

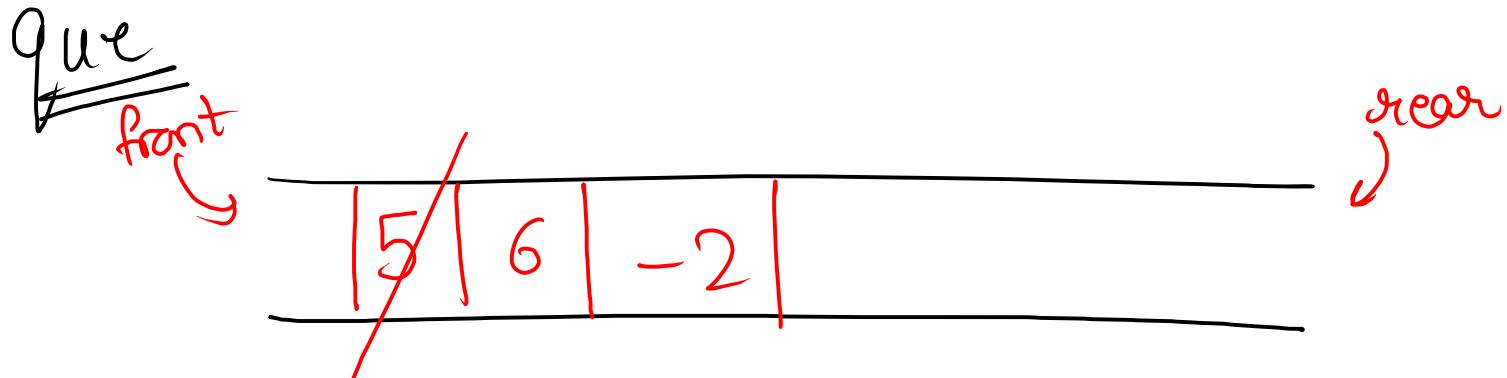
inbuilt

que.add(value); // to insert value from rear  
end

{ que.remove(); // to delete and return element  
que.poll(); // in front of queue

que.peek(); // return front element from queue  
without removing it

que.size() / que.isEmpty() } evergreen



que.add(5);

que.add(6);

que.add(-2);

que.peek() → 5

que.poll() → 5

que.poll() → 6

que.isEmpty() → false.