

Find the index of Rotation

$n=5$

$$\text{arr} = [5, 1, 2, 3, 4]$$

↳ If it is sorted then it would be like
 this $\rightarrow [1, 2, 3, 4, 5]$

Rotated index $\rightarrow \underline{0}$

Explanation

Ex1: $[5, 1, 2, 3, 4]$

$^0 \quad 1 \quad 2 \quad 3 \quad 4$

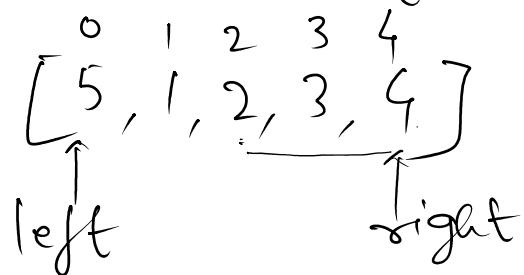
$$\text{arr}[0] > \text{arr}[1]$$

Ex2: $[4, 6, 8, 1, 2, 3]$

Output: -2.

Solution:-

(i.) We have determine where the value on its right is smaller.



$$\text{mid} = 2$$

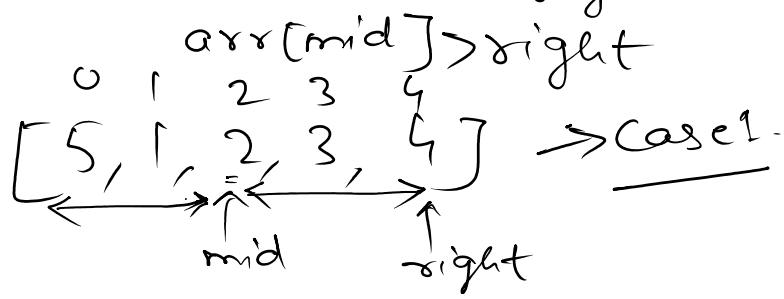
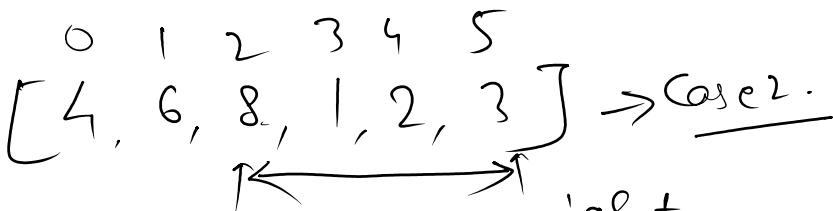
$$\text{arr}[2] = 2.$$

$$\text{arr}[\text{right}] = 4$$

$\text{arr}[right] = 4$

Case1. $\text{arr}[right] > \text{arr}[mid]$

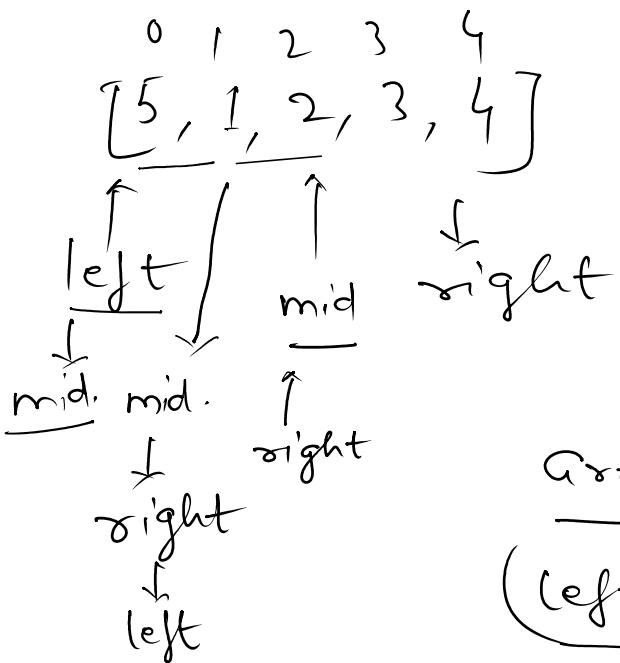
We have to check in left, because right side, it will sorted.



$right = mid;$

Case2:- $\text{arr}[mid] > \text{arr}[right]$

$(left = mid + 1);$



$\text{arr}[mid] > \text{arr}[right]$
 $(left < right)$

0 1 2 3 4 5

0 1 2 3 4 5
[4, 5, 6, 1, 2, 3]

Dry Run

1. left = 0
right = 5

mid = 2

arr[mid] > arr[right]

6 > 3

left = mid

2. left = 2
right = 5

0 1 2 3 4 5
[4, 5, 6, 1, 2, 3]

mid = 3

arr[mid] < arr[right]

1 < 3.

right = 3

3. left = 2

right = 3

mid = 2

arr[mid] > arr[right]

6 > 1.

Code.

```
int left=0, right=n-1;  
while (left < right) {  
    int mid=(left+right)/2;  
    if (arr[mid] < arr[right]) {  
        right=mid;  
    } else {  
        left=mid+1;  
    }  
}  
S.O.P(n[left-1]);
```

arr = [6, 7, 1, 2, 3, 4]
 0 1 2 3 4 5

Output → 1

Dry Run -

Dry Run

1. $\text{left} = 0,$
 $\text{right} = 5$

$\text{mid} = 2$

$\text{arr}[\text{mid}] = 1 < \text{arr}[\text{right}]$

$\text{right} = \text{mid} = 2$

0 1 2 3 4 5
[6, 7, 1, 2, 3, 4]

2. $\text{left} = 0$
 $\text{right} = 2$

$\text{mid} = 1$

$\text{arr}[\text{mid}] = 7 > \text{arr}[\text{right}]$

$\text{left} = 2.$

$\text{right} = 2$

$(\text{left} < \text{right}) \times$

S.O.Pn(left - 1);

The banana challenge

$$n=4. \quad 0 \quad 1 \quad 2 \quad 3.$$

$$arr = [3, 6, 7, 11]$$

$$h=8.$$

Output $\rightarrow 4.$

Let's assume $k=3$

1 \rightarrow 1st hour \rightarrow 1 pile $\rightarrow 3 \rightarrow$ all will be eaten

2nd hour \rightarrow 2nd pile $\rightarrow 6 \rightarrow$ 3 banana can be eaten

3rd hour \rightarrow 3rd pile \rightarrow 3 banana from 2nd pile

4th hour \rightarrow 4th pile $\rightarrow 3$ from 3rd

$$3rd \rightarrow 4$$

$$4th \rightarrow 11$$

$$5th \rightarrow 3, \text{ remaining} \rightarrow 12$$

$$6th \rightarrow 3, \text{ rem} \rightarrow 9$$

$$7th \rightarrow 3, \text{ rem} \rightarrow 6$$

$$8th \rightarrow 3, \text{ rem} \rightarrow 3$$

$$9th \rightarrow 3, \text{ rem} \rightarrow 0$$

$$\underline{\text{hours: 9.}}$$

That's why, speed can not be $3/\text{hour}$

$$[3, 6, 7, 11]$$

$$\underline{k=4}$$

Hour	Pile	Remaining
------	------	-----------

1	1st	0	[3, 6, 7, 11]
2	2	2	
3	2	0 (2 eaten)	
4	3	3	
5	3	0	
6	4	7	
7	4	3	
8	4	0	
[3, 6, 7, 11]			

$$K=4$$

$$\begin{aligned} 3/4 &= 0.75 = \underline{\underline{1}} \\ 6/4 &= 1.5 = \underline{\underline{2}} \\ 7/4 &= 1.75 = \underline{\underline{2}} \\ 11/4 &= 2.75 = \underline{\underline{3}} \end{aligned} \quad \left. \begin{array}{l} \\ \\ \end{array} \right\} \textcircled{8} =$$

To determine, Speed we need to find
max and min speed

maxspeed:- No. of bananas available in
largest file.

3, 6, 7, 11

11 bananas / hour

minspeed:- 1 banana / hour

Speed will lie between 1 and 11.

Two Approach:-

Two Approach:-

1. Using for loop, checking will every speed between 1 and 11.

$$i=1, i \leq 11$$

for any value of i , we have to calculate min hour to finish eating all banana

if min hour to finish all banana is less than equal to given hour, then that value of i , which is speed will be my answer.

2nd Approach.

$$\text{left} = 1, \text{right} = 11 \quad (6)$$

$$\text{mid} = (\text{left} + \text{right}) / 2$$

↓

If with speed = mid, no. of hours taken to finish banana $\leq h$, then
 $\frac{8}{8}$

with $K = 6$,

1. total hours taken ≤ 8

2. total hours taken > 8 .

1st possibility

total hours ≤ 8

$$\underline{\underline{K=6}}$$

ans = 6.

if any speed $< k$, total hours ≤ 8 .
ans = that speed

2nd possibility

total hours > 8

if this case, we should increase speed
so, we should take values in sight of mid

```
public static int hoursTaken(int piles[], int k) {
    int totalHour = 0
    for (int i=0; i<piles.length; i++) {
        int hour = Math.ceil ((double) piles[i]/k);
        totalHour += hour;
    }
    return totalHour;
}
```

```
p.s. v. m (String [] args) {
    int h = sc.nextInt();
    int left = 1;
    int max = Integer.MIN_VALUE;
    for (int i=0; i<piles.length; i++) {
        max = Math.max (max, piles[i]);
    }
```

```
for (int i = 1; i < piles.length; i++) {
    max = Math.max(max, piles[i]);
}

int right = max; int ans = -1;
while (left <= right) {
    int mid = (left + right) / 2;
    if (hoursTaken(piles, mid) <= h) {
        ans = mid;
        right = mid - 1;
    } else {
        left = mid + 1;
    }
}
System.out.println(ans);
}
```