

LINUX OPERATING SYSTEM

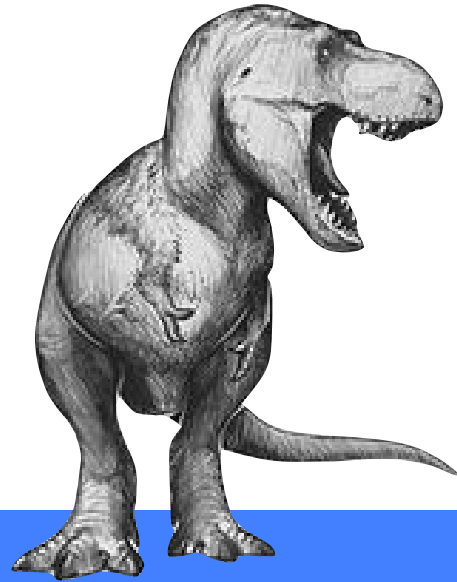
YANG

LINUX操作系统（双语）





双语课→课件内容中英混排



|Lecture 18

File-System Implementation

本讲内容

 文件系统实现

 文件目录

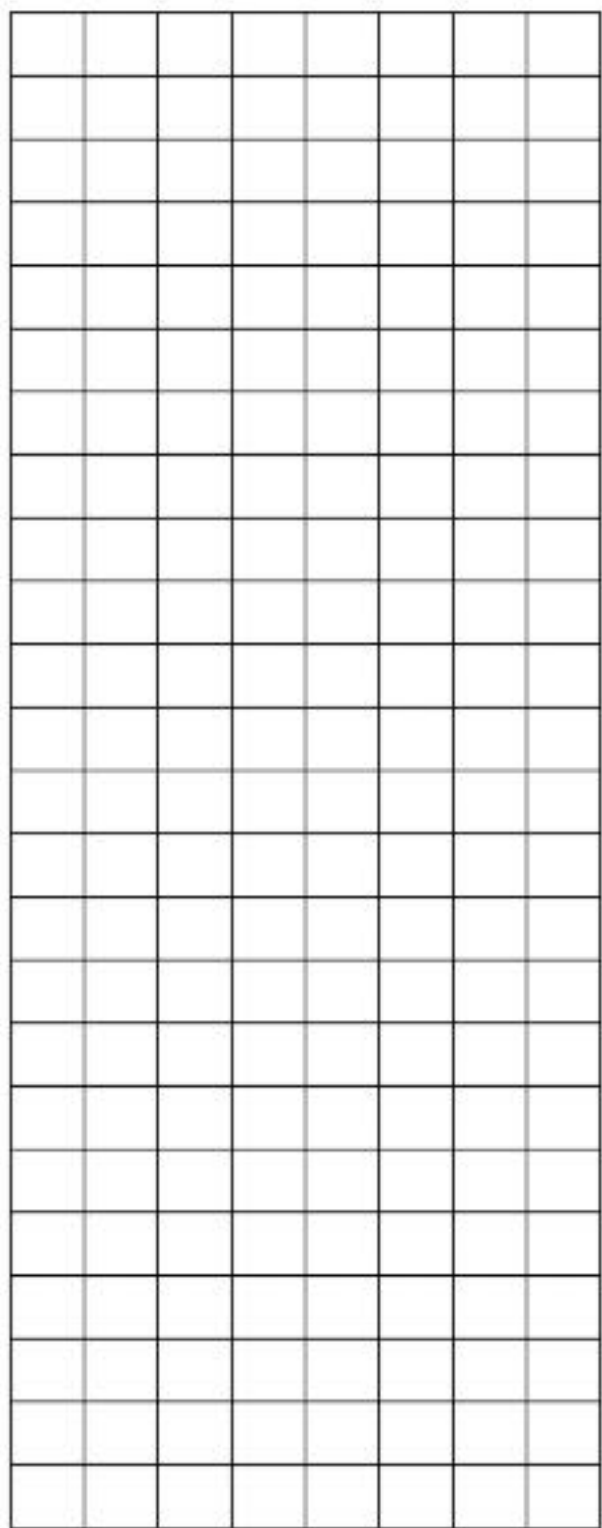
 分配方法

 空闲空间管理

 文件系统结构

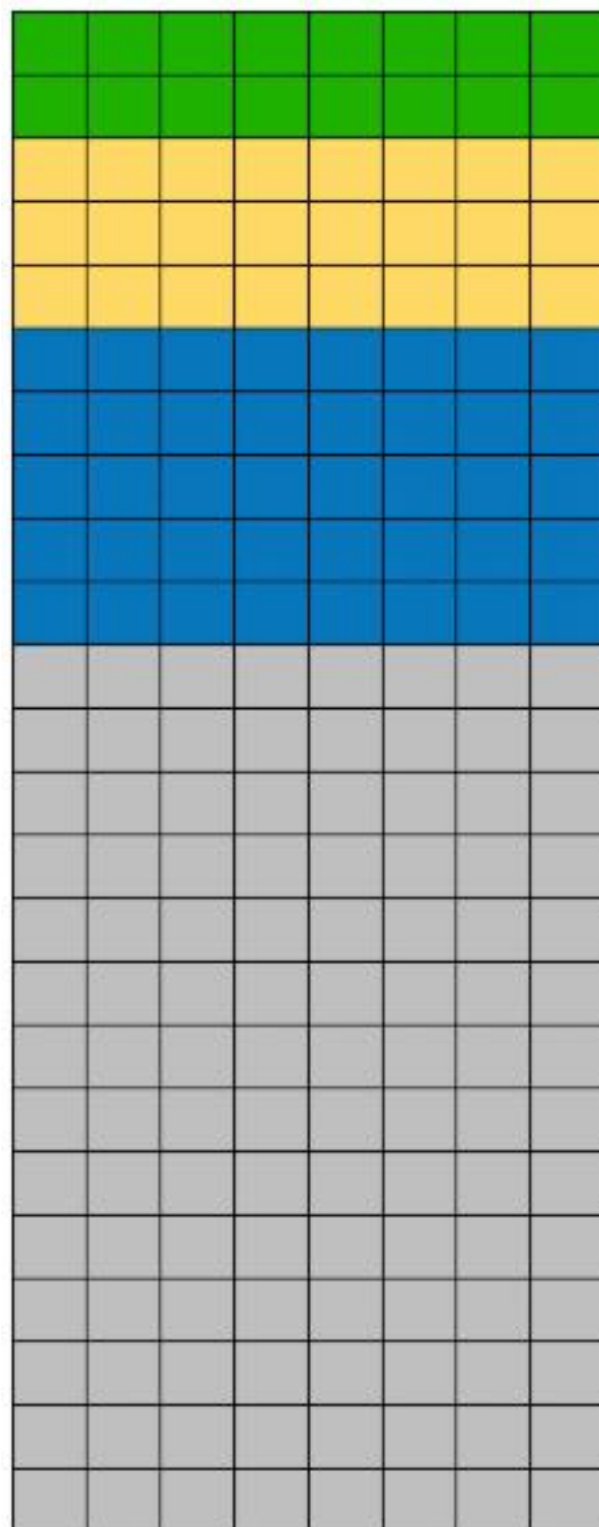
文件系统实现

文件系统要实现什么？



disk

文件系统要实现什么？



disk

partition table

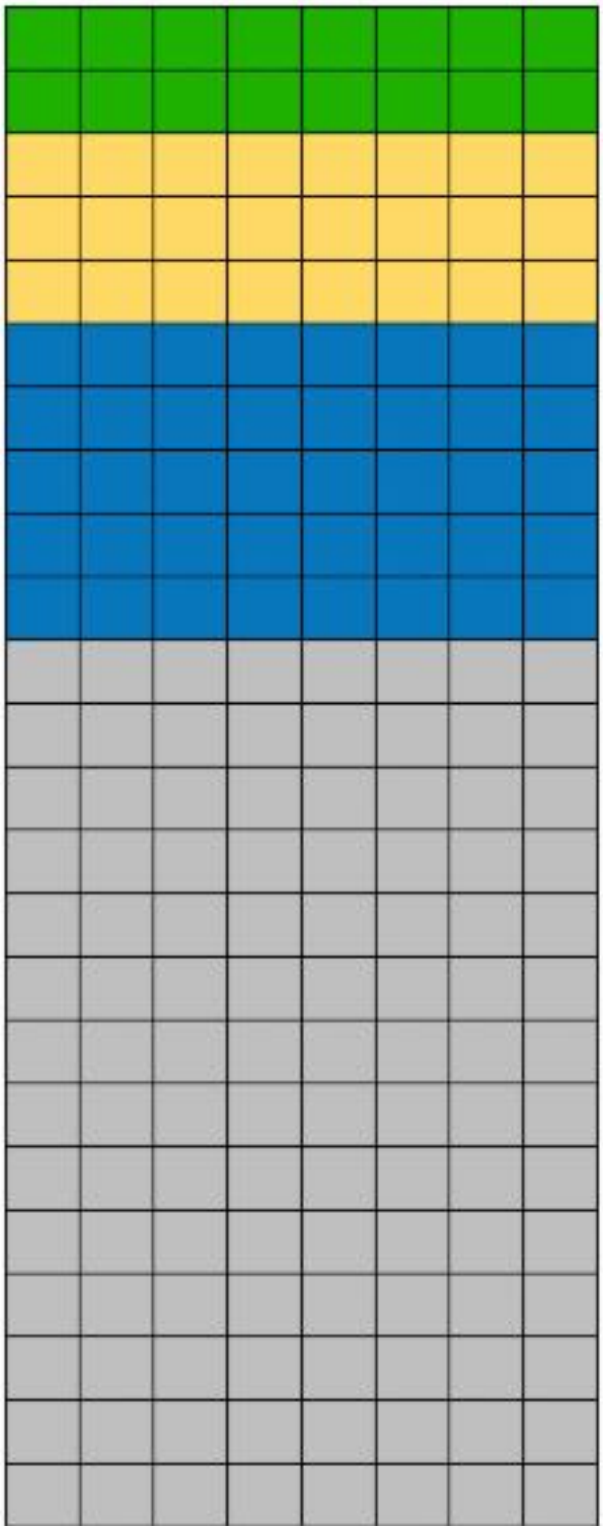
free blocks map

directory entry 1
directory entry 2
directory entry...
directory entry n

File Control Block

- file name
- permission
- date&time
- owner/group/ACL
- file size
- pointer to file data

文件系统要实现什么？



disk

partition table

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directory entry...
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File Control Block
<ul style="list-style-type: none">• file name• permission• date&time• owner/group/ACL• file size• pointer to file data

- 🧠 文件目录
- 🧠 空间分配方法
- 🧠 空闲空间管理

COMMON FILE SYSTEMS

- 👤 Many file systems are in use today, and most operating systems support more than one.
- 👤 For example, most CD-ROMs are written in the **ISO 9660 format**, a standard format agreed on by CD-ROM manufacturers. In addition to removable-media file systems, each operating system has one or more disk-based file systems.
 - 👤 UNIX uses the UNIX file system (**UFS**), which is based on the Berkeley Fast File System (**FFS**).
 - 👤 Windows supports disk file-system formats of **FAT**, **FAT32**, and **NTFS**.
 - 👤 Although Linux supports over forty different file systems, the standard Linux file system is known as the extended file system, with the most common versions being **ext3** and **ext4**.
- 👤 Another interesting project is the **FUSE** file system, which allows a user can add a new file system to a variety of operating systems and can use that file system to manage her files.

文件目录

文件控制块

- 🧠 文件系统通过文件控制块（File Control Block）来维护文件结构，FCB包含有关文件的信息，包括所有得、权限、文件内容的位置等。
- 🧠 文件目录用于组织文件，每个目录项对应一个FCB。
- 🧠 文件目录实现的关键
 - 🧠 FCB与文件内容的关联方法
 - 🧠 在目录中“按名”搜索的效率

File Control Block
<ul style="list-style-type: none">• file name• permission• date&time• owner/group/ACL• file size• pointer to file data

INODES

🧠 UFS中的FCB被称作索引结点inode，每个inode都有一个唯一的编号，包含的内容有：

- 🧠 The type of the file
- 🧠 The mode of the file (ACL)
- 🧠 The number of hard links to the file
- 🧠 The user ID of the owner of the file
- 🧠 The group ID to which the file belongs
- 🧠 The number of bytes in the file
- 🧠 **An array of 15 disk-block addresses**
- 🧠 The date and time the file was last accessed
- 🧠 The date and time the file was last modified
- 🧠 The date and time the file was created

inode
<ul style="list-style-type: none">• permission• date&time• owner/group/ACL• file size• ...• an array of 15 disk-block address

INODES

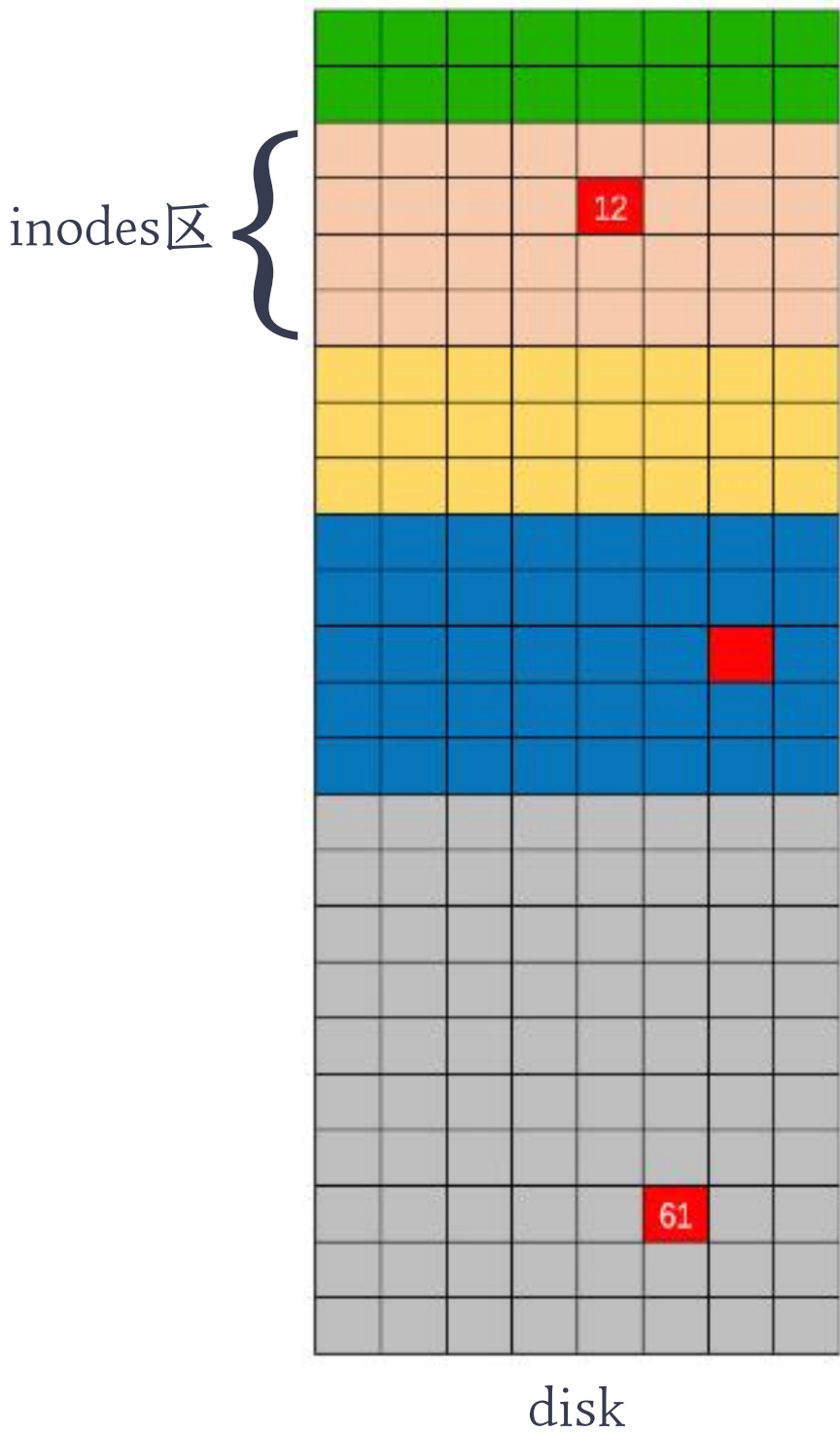
🧠 UFS中的FCB被称作索引结点inode，每个inode都有一个唯一的编号，包含的内容有：

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inode
<ul style="list-style-type: none">• permission• date&time• owner/group/ACL• file size• ...• an array of 15 disk-block address

file name	inode no.
UFS dentry	

INODE区



directory

A.file	13
B.file	11
C.file	12

inode-12
<ul style="list-style-type: none">• permission• date&time• owner/group/ACL• file size• ...• 61(disk-block address)

分配方法

分配方法

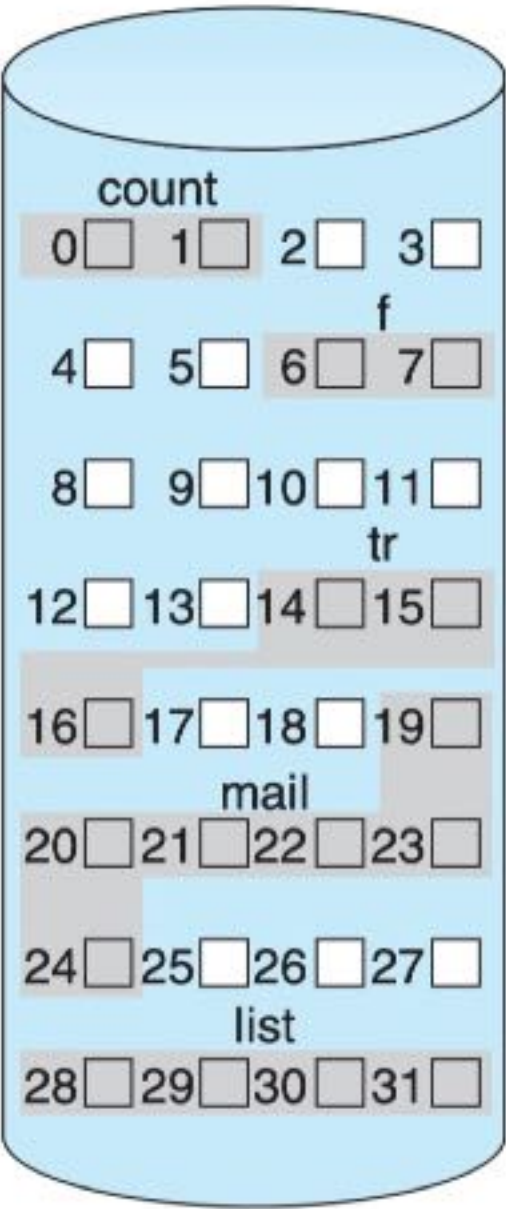
🧠 这里我们讨论如何给文件分配磁盘空间，常用方法有三种：

🧠 连续分配

🧠 链接分配

🧠 索引分配

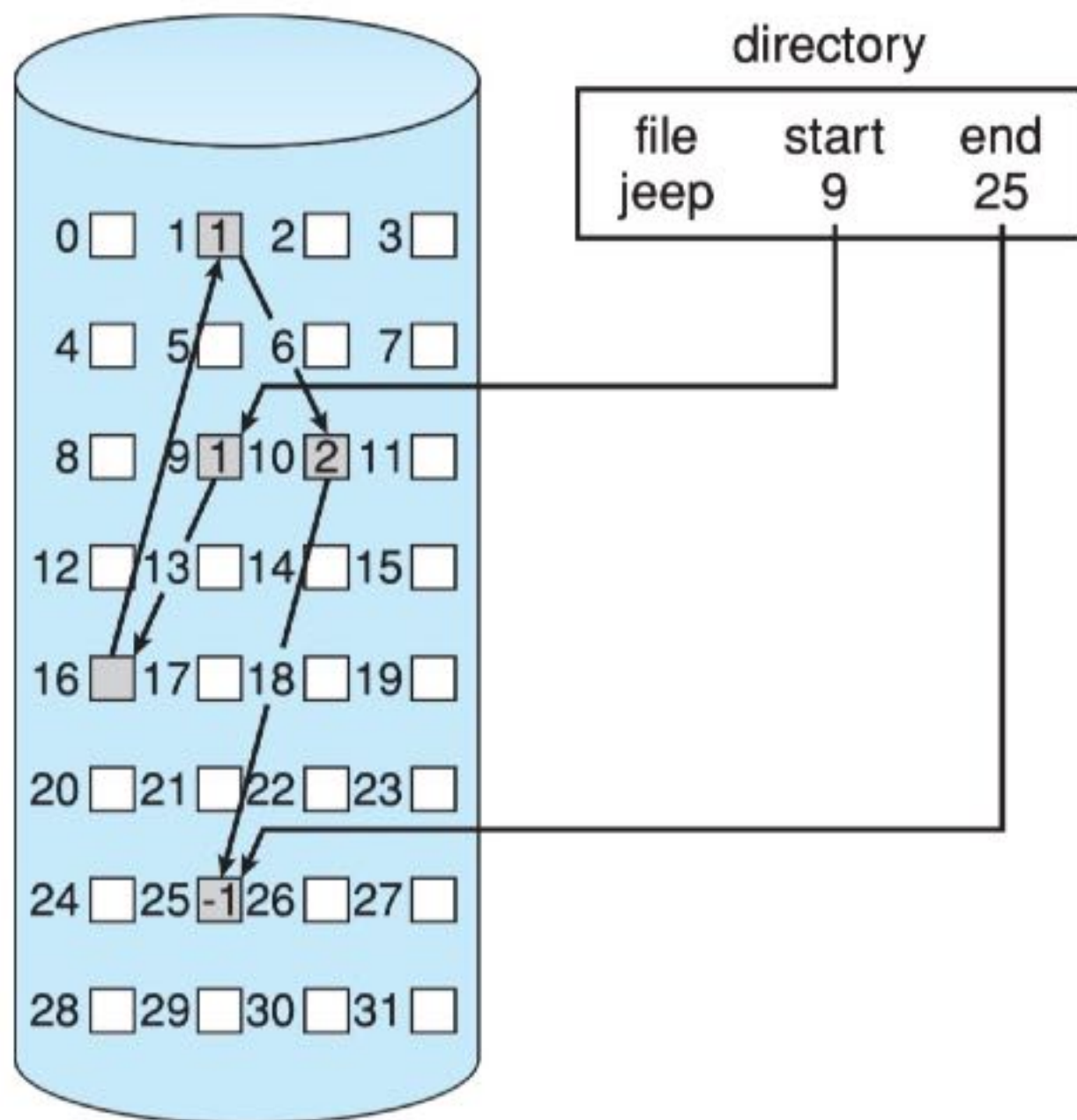
连续分配



directory		
file	start	length
count	0	2
tr	14	3
mail	19	6
list	28	4
f	6	2

- 每个文件在磁盘上占用连续的物理块
- 优点
- 缺点

链接分配

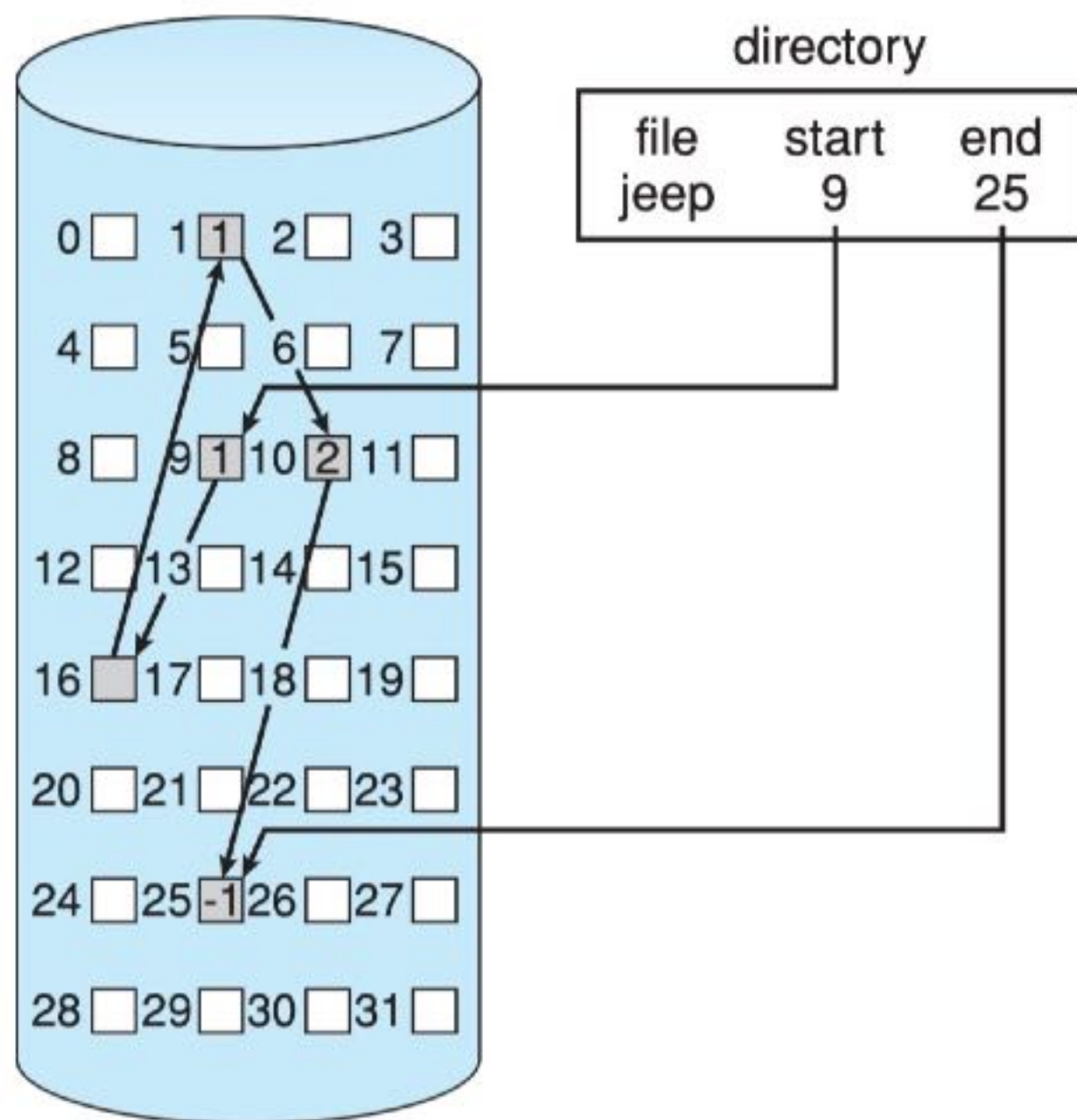


🧠 文件所占用的物理块分散在磁盘的不同位置，通过指针将它们链接起来。

🧠 优点

🧠 缺点

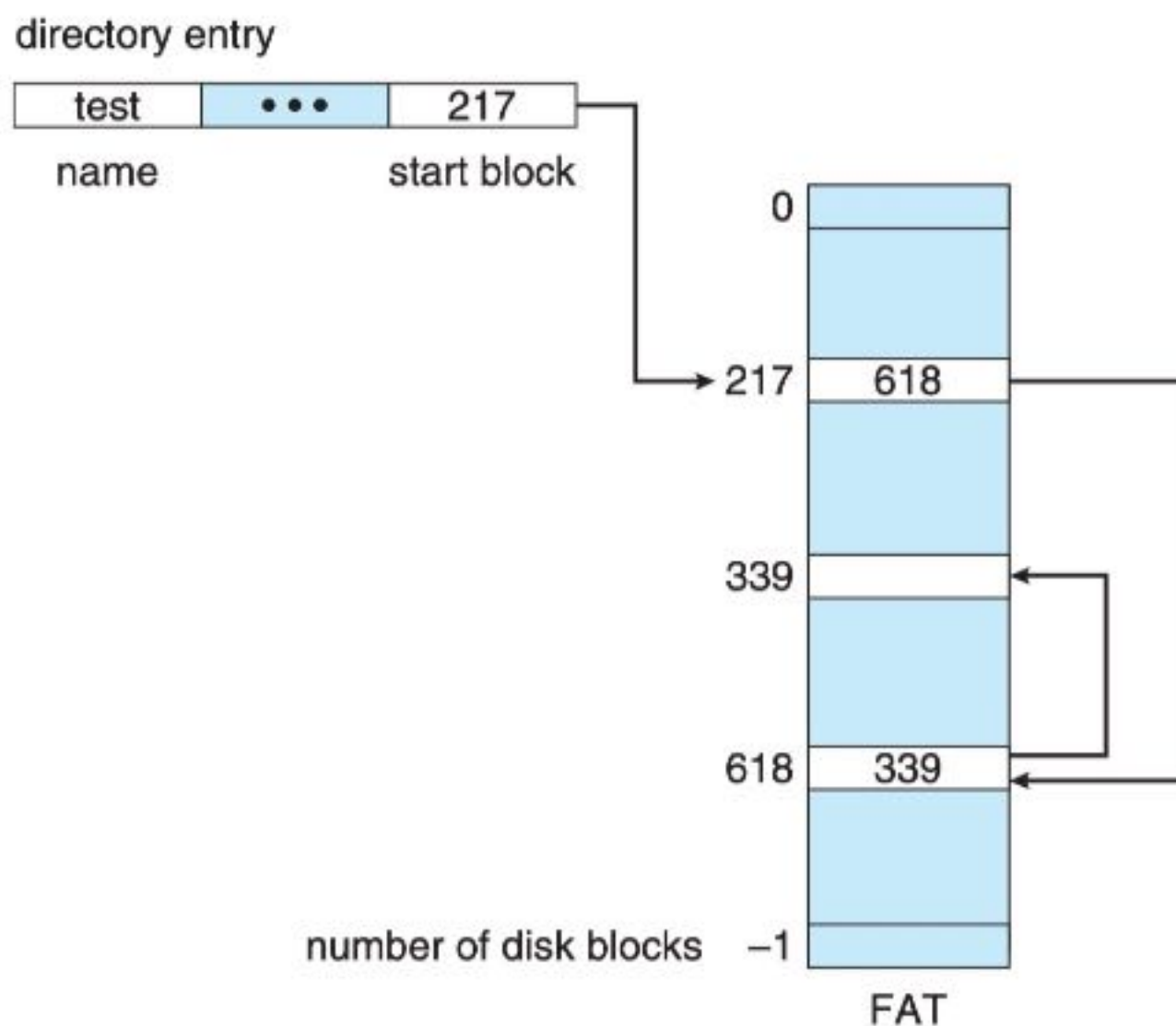
簇 CLUSTER



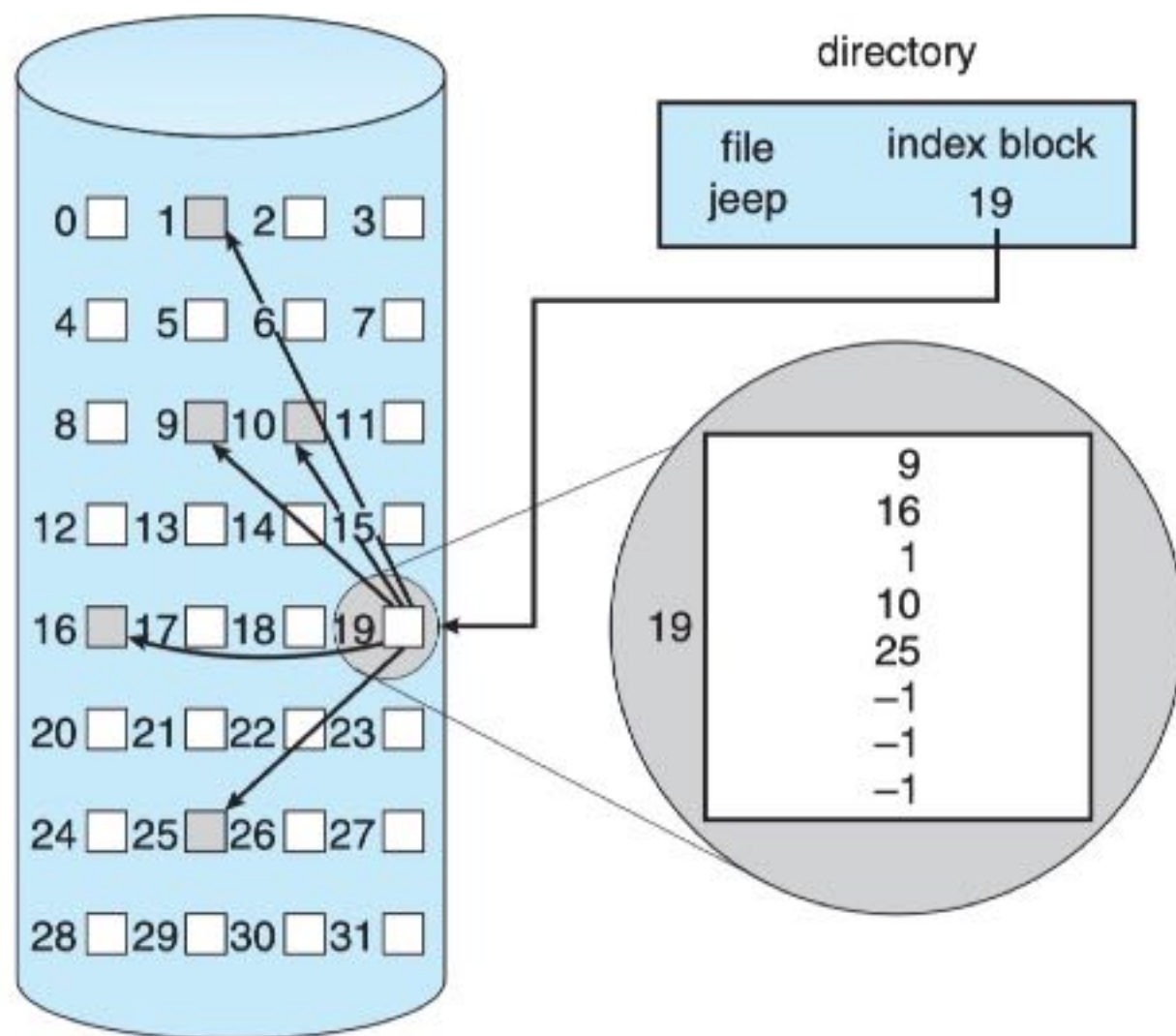
💡 **簇 (Cluster)** 指一组物理块的集合。如果以簇作为分配单位，可以节省指针占用的空间比例。

FILE ALLOCATION TABLE

💡 文件分配表FAT是一个典型的链接分配方案，不过它没有^①在物理块或簇的尾部加入指针，而是用一张表来记录文件占用物理块号的顺序。



索引分配

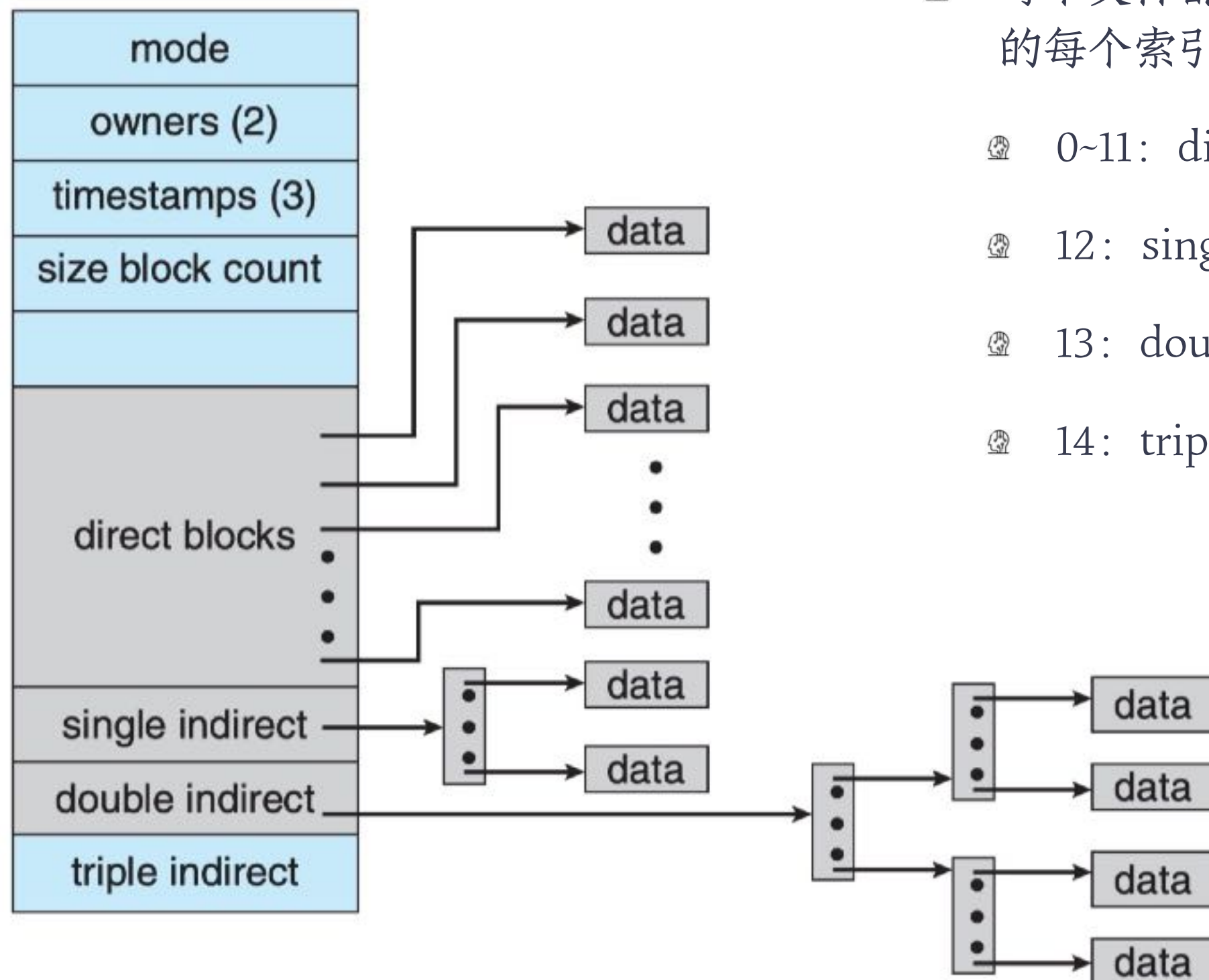


将文件占用的所有物理块号按逻辑顺序保存在一张索引表中，存有索引表的物理块称索引块（index block）

优点

缺点

UFS的索引块



💡 每个文件都要有一个索引块，UFS的每个索引块有15个物理块地址

💡 0~11: direct block

💡 12: single indirect block

💡 13: double indirect block

💡 14: triple indirect block

空闲空间管理

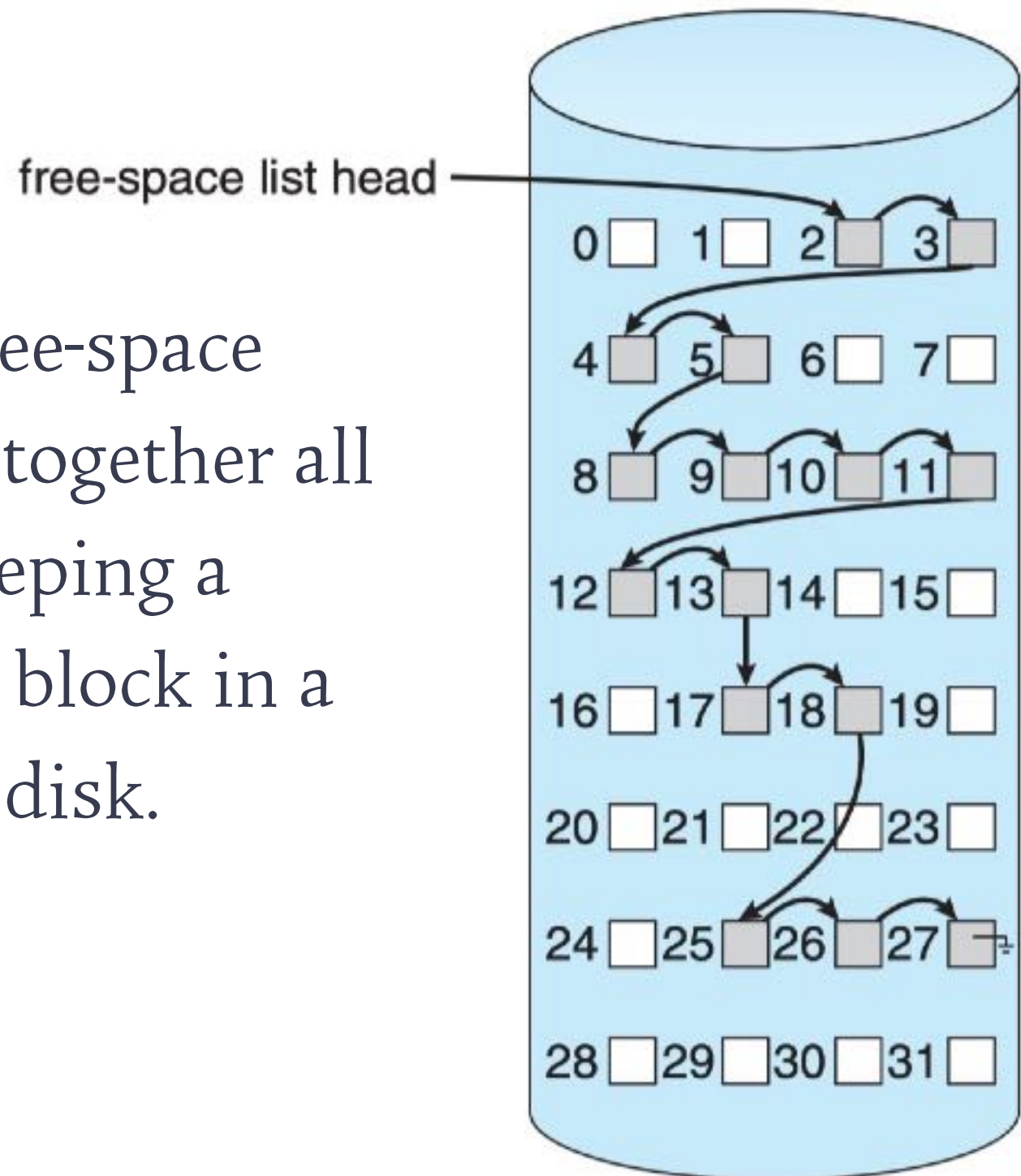
BIT MAP

- ☞ Frequently, the free-space list is implemented as a bit map or bit vector. Each block is represented by 1 bit. If the block is free, the bit is 1; if the block is allocated, the bit is 0.
- ☞ A bit map sample is following:

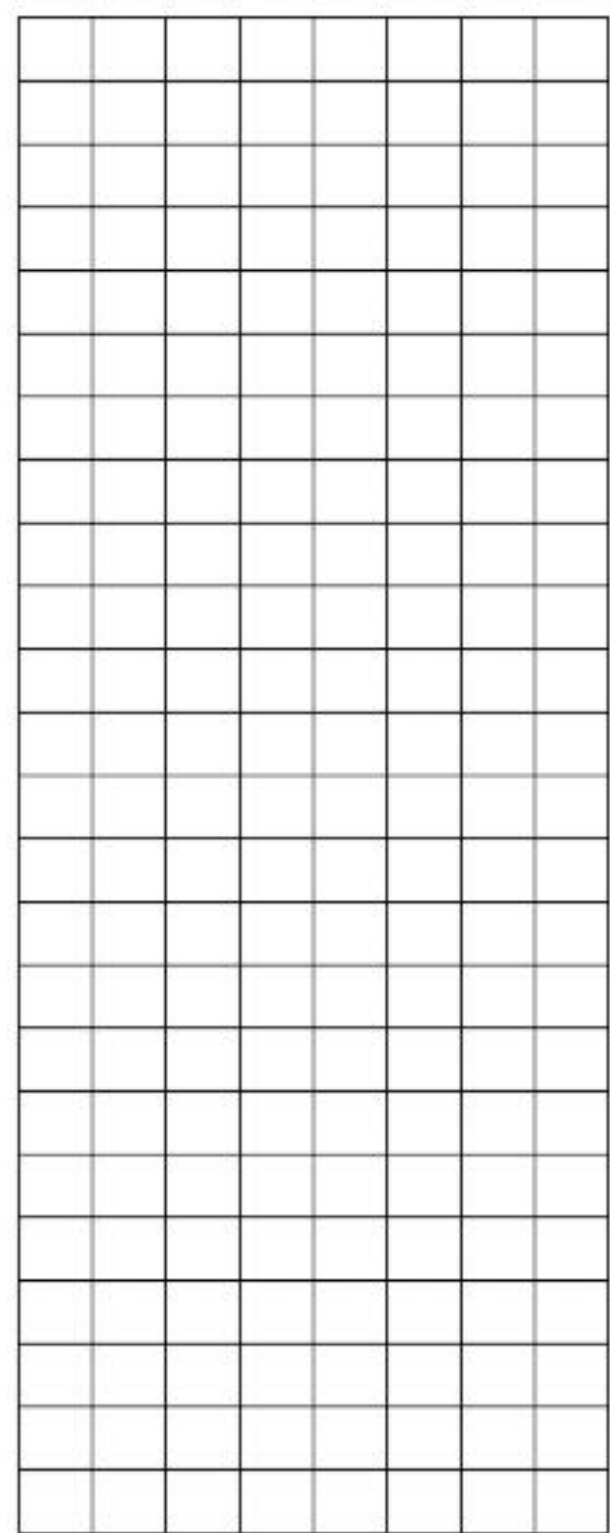
001111001111110001100000011100000 ...

LINKED LIST

- Another approach to free-space management is to link together all the free disk blocks, keeping a pointer to the first free block in a special location on the disk.



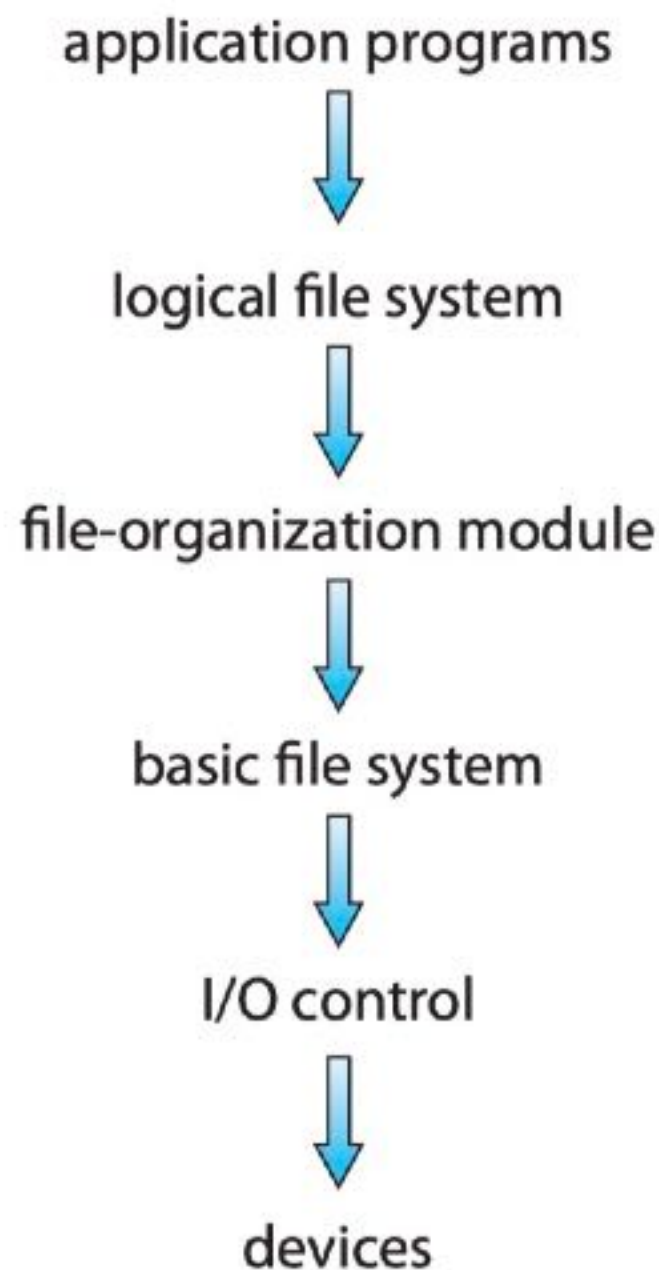
文件系统的实现工作



disk

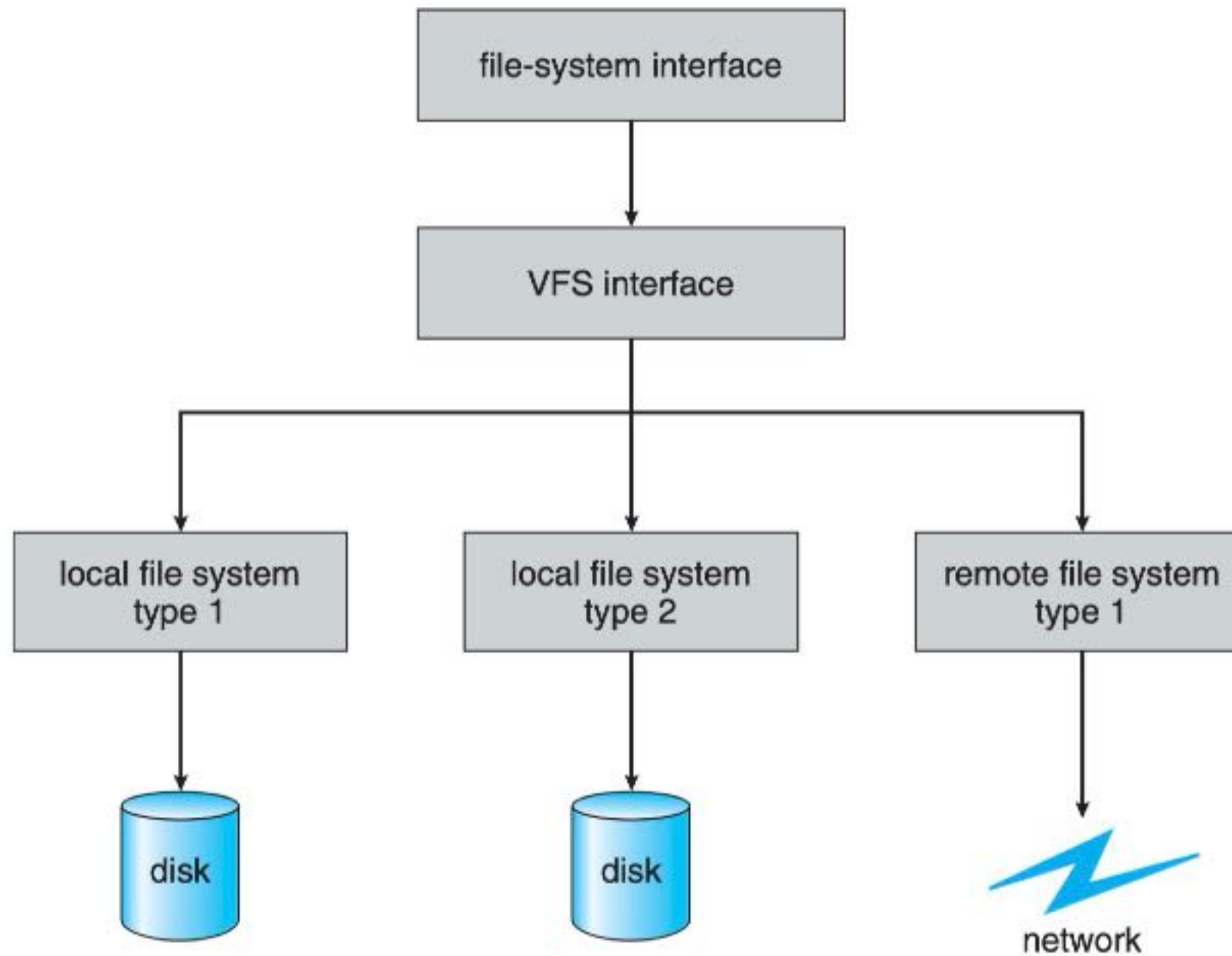
文件系统结构

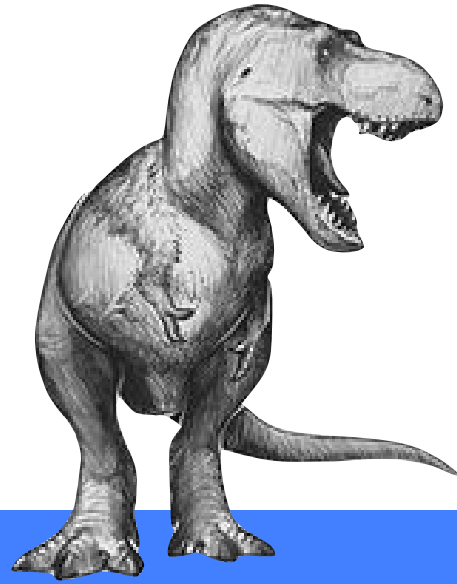
LAYERED FILE SYSTEM



你发出一个读写A文件X位置数据的请求，系统从目录中找到这个文件并读出相应的FCB，按照既有的分配方案计算X位置所在的物理块号，编制一个对该物理块号的读写请求，然后发给磁盘控制器。

VIRTUAL FILE SYSTEM





|Lecture 18

The End

实践5 文件链接

硬链接和软链接

硬链接

-  输入命令：ln file1 file1Hardlink，为file1文件创建一个硬链接，使用ls -il命令观察file1和file1Hardlink有什么联系

软链接

-  输入命令：ln -s file2 file2Softlink，为file1文件创建一个软链接（符号链接），使用ls -il命令观察这两个文件有什么联系

