



206

Significant Structures II

Konstantinos P. Michmizos

Computational Brain Lab

Computer Science | Rutgers University | NJ, USA



Send me an email

(with a positive or negative remark)





Announcements

- Section 4 has *in-person* recitation on Tuesday 5:00 PM 5:55 PM at Livingston TIL-254 (TA: Yuequn Zhang <u>yuequn.zhang@rutgers.edu</u>)
- Section 5 has *in-person* recitation on Thursday 3:00 PM 3:55 PM at Livingston TIL-226 (TA: Vladimir Ivanov vai9@scarletmail.rutgers.edu)
- Section 6 has *in-person* recitation on Thursday 3:00 PM 3:55 PM at Livingston LSH-B115 (TA: James Fu jf980@scarletmail.rutgers.edu)

Office hours for all TAs under home > course information > staff and office hours

What we will cover today

Combinatorics

- Recap
 - Proofs (Direct, Contrapositive)
- Today
 - Proofs
 - Direct
 - Contrapositive
 - Case Analysis
 - Contradiction
 - Induction
 - Counting
 - Partition Method
 - Difference Method

Next Week

- Bijection Rule
- Product Rule



Course Outline

• Part I

- Recap of basics sets, function, proofs, induction
- Basic counting techniques
- Pigeonhole principle
- Generating functions

Part II

- Sample spaces and events
- Basics of probability
- Independence, conditional probability
- Random variables, expectation, variance
- Moment generating functions

• Part III

- Graph Theory
- Machine learning and statistical inference

Proving an Implication via Direct Proof

- To prove: $P \Rightarrow Q$
 - Assume that *P* is true.
 - Show that *Q* logically follows

Direct Proof

• To prove: $P \Rightarrow Q$

The sum of two even numbers is even.

• Assume that *P* is true.

$$x = 2m, y = 2n$$

 $x+y = 2m+2n$

• Show that Q logically follows = 2(m+n)

The product of two odd numbers is odd.

Proof
$$x = 2m+1, y = 2n+1$$

 $xy = (2m+1)(2n+1)$
 $= 4mn + 2m + 2n + 1$
 $= 2(2mn+m+n) + 1$

Example of Proving an Implication

• Theorem: $1 \le x \le 2 \Rightarrow x^2 - 3x + 2 \le 0$

Remember:

$$Assume 1 \leq X \leq 2$$

Step 1: $X^2 - 3X + 2 = (X - 1)(X - 2)$

We start from P

and we construct Q

$$\frac{Step2}{Step3}$$
: $(X-2) \leq 0$

(while adhering to logic)

$$(X-1)$$
 $(X-2)$ $(X-2)$ $(X-2)$ $(X-1)$ $(X-2)$ $(X-2$

Intuition:When x grows, 3x grows faster than x^2 in that range.

Proof by Contrapositive

• To prove: $P \Rightarrow Q$

Remember:

• Prove that $\neg Q \Rightarrow \neg P$.

We reverse the direction

(double flip!)

• Assume $\neg Q$ is true and show that $\neg P$ follows logically.

i.e., we will <u>assume the opposite of our desired conclusion</u> and show that this fancy opposite conclusion could never be true in the first place.

Example of Proof by Contrapositive

• Theorem: *If* r *is irrational, then* \sqrt{r} *is irrational.*

Example of Proof by Contrapositive

• Theorem: *If* r *is irrational, then* \sqrt{r} *is irrational.*

Proof:

We shall prove the contrapositive – "if \sqrt{r} is rational, then r is rational."

Since \sqrt{r} is rational, $\sqrt{r} = a/b$ for some integers a,b.

So $r = a^2/b^2$. Since a,b are integers, a^2,b^2 are integers.

Therefore, r is rational. Q.E.D.

(Q.E.D.)

"which was to be demonstrated", or "quite easily done". © quod erat demonstrandum

What we will cover today

- Recap
 - Sets Venn Functions Proofs (Direct)
- Combinatorics
 - Proofs
 - Direct
 - Contrapositive
 - Case Analysis
 - Contradiction
 - Induction
 - Counting
 - Partition Method
 - Difference Method

Example of Proof by Case Analysis

• Theorem: For all $x \in \mathbb{R}$, $-5 \le |x + 2| - |x - 3| \le 5$

We **hate** absolute values so we want to avoid them as soon as possible

And we have two possible ways:-) (Squaring and...)

One of them is our goal here: To identify all possible cases

Example of Proof by Case Analysis

• Theorem: For all $x \in \mathbb{R}$, $-5 \le |x + 2| - |x - 3| \le 5$

$$|X+2| = X+2$$
, if $X+27/0$ on $X7/-2$
 $|X+2| = -(X+2)$, if $X+27/0$ on $X2/-2$
 $|X+2| = -(X+2)$, if $X=3$
 $|X-3| = X-3$ if $X=3$
 $|X-3| = -(X-3)$ if $X<3$
 $|X-3| = -(X-3)$ if $X<3$

Example of Proof by Case Analysis

• Theorem: For all $x \in \mathbb{R}$, $-5 \le |x + 2| - |x - 3| \le 5$

(46 I:
$$\times 7/3$$
, Want $-5 \le (\times +2) - (\times -3) \le 5$)

(46 I: $\times 7/3$, Want $-5 \le (\times +2) - (\times -3) \le 5$)

(46 II: $-2 \le \times \times 23$, Want $-5 \le (\times +2) - (\times -3) \le 5$)

Want $-5 \le \times \times 24 + \times -3 \le 5$

Want $-5 \le \times \times 24 + \times -3 \le 5$

Want $-5 \le -(\times +2) - -(\times -3) \le 5$

Want $-5 \le -(\times +2) + \times -3 \le 5$

Want $-5 \le -(\times +2) + \times -3 \le 5$

Want $-5 \le -(\times +2) + \times -3 \le 5$

Want $-5 \le -(\times +2) + \times -3 \le 5$

Want $-5 \le -(\times +2) + \times -3 \le 5$

Want $-5 \le -(\times +2) + \times -3 \le 5$

Want $-5 \le -(\times +2) + \times -3 \le 5$

Remember:

Proof by Contradiction

$$\frac{\overline{P} \to \mathbf{F}}{P}$$

The main difference between

Proof by Contradiction and

Proof by Contrapositive is...

that there is no Q here!

To prove P, you prove that **not P** would lead to a ridiculous result,

and so P must be true.

I am working 20 hours per day.

If I had won the lottery, then I would not be working 20 hours per day.

• I have not won the lottery.

Proof by Contradiction – Work Chart

- To prove *P*
 - Assume *P* is false.
 - Logically deduce something that is known to be false.

Proof by Contradiction – Work Chart

To prove a proposition P by contradiction:

- 1. Write, "We use proof by contradiction."
- 2. Write, "Suppose P is false."
- 3. Deduce something known to be false (a logical contradiction).
- 4. Write, "This is a contradiction. Therefore, P must be true."

Example of Proof by Contradiction

• Theorem: *There are infinitely many primes*

A **prime number** (or a prime) is a natural number greater than 1 that cannot be formed by multiplying two smaller natural numbers.

This is one of the most famous, most often quoted, and most beautiful proofs in all of mathematics. Its origins date back more than 2000 years to Euclid

Example of Proof by Contradiction

• Theorem: *There are infinitely many primes*

Assume: There are finitely many primes – And let $p_1, p_2, ..., p_N$ be all the primes.

Now we construct a new number, $p = p_1 p_2 \dots p_N + 1$

Clearly, *p* is larger than any of the primes, so it does not equal one of them.

Therefore it cannot be prime and must be **composite**, i.e., <u>divisible by at least one of the primes</u>.

But our assumption was that *p* is not prime and therefore divisible by any prime number.

On the other hand, we know that any number must be divisible by *some* prime (*fundamental theorem* of arithmetic or the unique factorization theorem or the unique-prime-factorization theorem)

This leads to a **contradiction**, and therefore the assumption must be false.

So there must be infinitely many primes.

Take a Break



Induction

• Let P(m) be a predicate of non-negative integers

• You want to prove that P(m) is true for all non-negative integers.

• Step 1: Prove that P(0) is true

• Step 2: Prove that $P(n) \Rightarrow P(n + 1)$ for all non-negative integers.

Example of Induction

• Theorem: For all $n \in \mathbb{N}$,

•
$$1 + 2 + \dots + n = \frac{n(n+1)}{2}$$

Example of Induction

• Theorem: $For \ all \ n \in \mathbb{N}$,

•
$$1 + 2 + \dots + n = \frac{n(n+1)}{2}$$
 $P(n)$

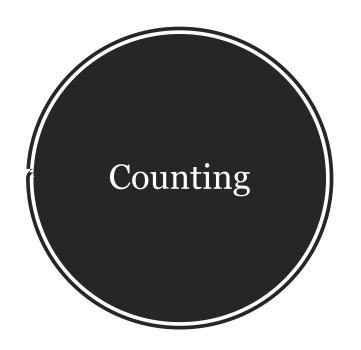
Base (ase: $P(o)$ is three, for $P(o)$ $P(o)$

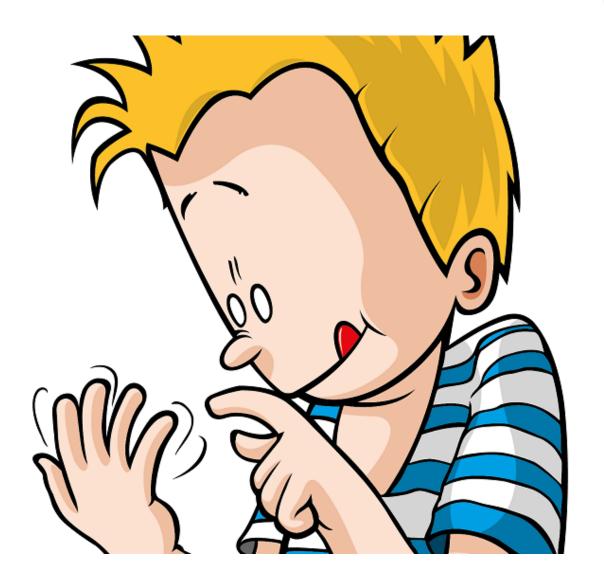
Inductive $P(o)$ $P(o)$ $P(o)$

Assume $P(o)$ is three. $P(o)$

Want: $P(o)$ $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$
 $P(o)$

Intuition: During induction, my goal is to construct what I have assumed as true





Counting

- Basic Question: What is the size of a given set?
- Easy when the set is explicitly defined.
 - $X = \{1,2,3,4\}$, what is |X|?
- Tricky when set is implicit or a defined via set operations.
 - How many ways to get flush in the game of poker?
 - How many ways to assign time slots to courses at Rutgers?
 - How many operations before my algorithm terminates?

Counting

- In the next few lectures
 - Fundamental tools and techniques for counting
 - Sum Rule
 - Product Rule
 - Difference Method
 - Bijection Method
 - Permutations/Combinations
 - Inclusion Exclusion
 - Binomial/Multinomial coefficients

-> Intermediate

-> Advanced