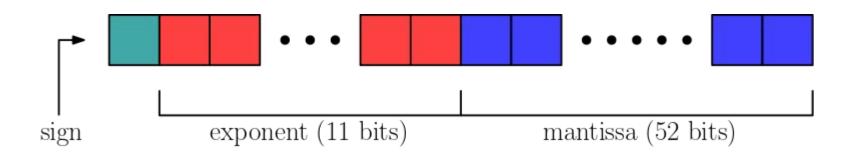
Floating-point expansions

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double-precision floating-point

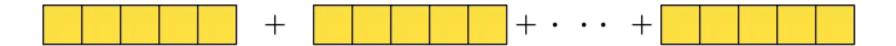




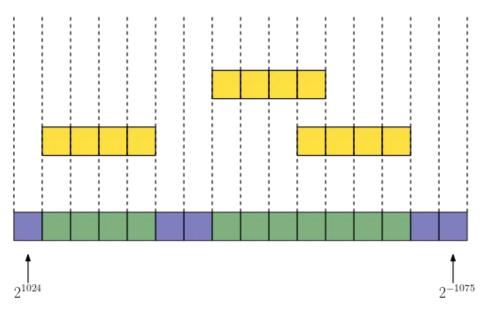
52 bits of precision

how do we increase precision?

represent numbers as unevaluated sums of double-precision floating-point numbers



we use whole range of doubles as precision

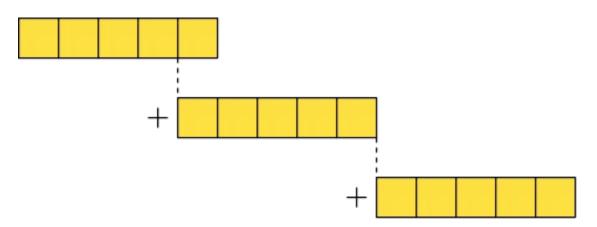




over 2000 bits of precision!

Non-overlapping expansions

Definition 2.6. An expansion $u_0, u_1, \ldots, u_{n-1}$ is ulp-nonoverlapping if for all 0 < i < n, we have $|u_i| \le \text{ulp}(u_{i-1})$.



Cost analysis

- Cost measure: floating point operations (flops)
 - Idexp & frexp counted as 1 flop
- Timing: RDTSC, chrono, QueryPerformanceCounter

Experimental setup

- Compiler: MSVC cl.exe on Windows
 - Also compared to g++ on Linux

- Optimizations: everything except compiler vectorization; fast floating point operations turned off
 - MSVC flags: /O2 /GL
 - **g++** flags: -O3 -fno-tree-vectorize -ffp-contract=off -march=native

Addition - baseline

```
void addition(double *a, double *b, double *s, int length_a, int length_b, int length_result)
   double *tmp = (double *)alloca((length_a + length_b) * sizeof(double));
   merge(a, b, tmp, length_a, length_b);
   renormalizationalgorithm(tmp, length_a + length_b, s, length_result);
   return;
```

- Adds two floating point expansions
- Components: Merge + renormalization

Addition - baseline

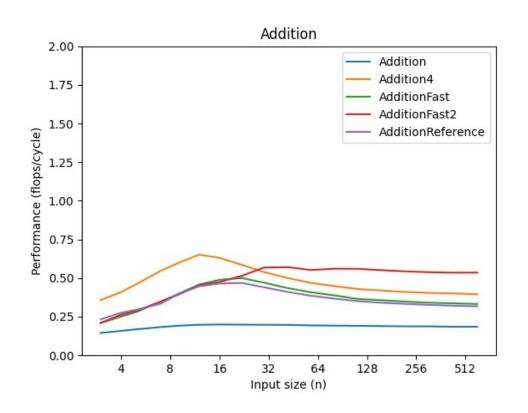
```
void renormalizationalgorithm(double* x, int size of x, double* f, int m)
       double *err = (double *)alloca((size of x) * sizeof(double));
       double *f_tmp = (double *)alloca((m + 1) * sizeof(double));
       for (int i = 0; i <= m; i++)
            f_tmp[i] = 0;
       vecSum(x, err, size_of_x);
       vecSumErrBranch(err, size_of_x, m + 1, f_tmp);
       for (int i = 0; i <= (m - 2); i++)
           vecSumErr(&(f_tmp[i]), m - i + 1, &(f_tmp[i]));
            f[i] = f tmp[i];
       f[m - 1] = f tmp[m - 1];
```

Renormalization:normalizes input vector toULP-nonoverlapping

Components:

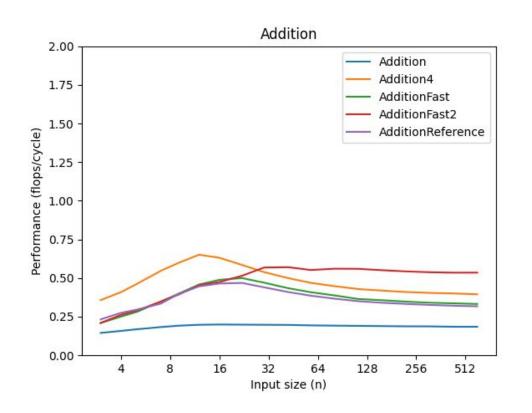
- VecSum
- VecSumErrBranch
- VecSumErr

Addition - Performance



- Addition: baseline
- Addition4: includes optimized (and vectorized) components
- AdditionFast: uses fast two sum (performance is lower, but runtime also)
- AdditionFast2: uses vectorized loop of VecSumFrr calls

Addition - Performance



- Many dependencies
 between instructions, so
 difficult to parallelize
- Many memory operations



- low performance (< 0.65 flops/cycle)
- ~6.1x overall speedup

Multiplication 1 - baseline

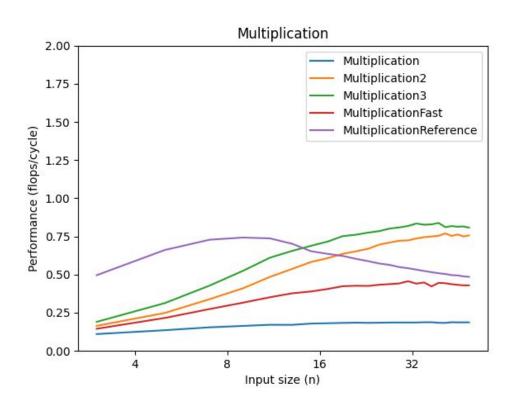
```
1 void multiplication(double *a, double *b, double *r, const int sizea, const int sizeb, const int sizer)
       twoMultFMA(a[0], b[0], &(r_ext[0]), &(err[0]));
        for (int n = 1; n <= (k - 1); n++)
           for (int i = 0; i <= n; i++)
               twoMultFMA(a[i], b[n - i], &(p[i]), &(e tmp[i]));
           vecSum(tmp, tmp1, (n * n + n));
       renormalizationalgorithm(r_ext, k + 1, r, sizea);
```

Multiplies two floating point expansions

Components:

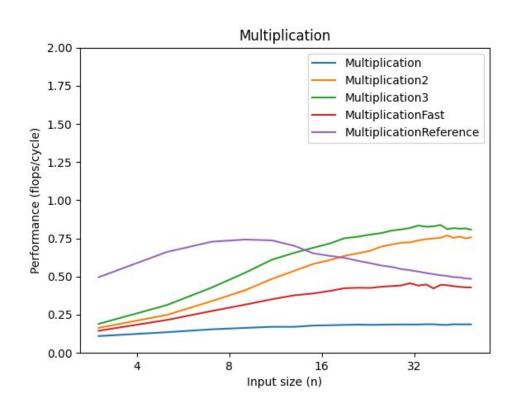
- TwoMultFMA
- VecSum
- Renormalization

Multiplication - performance



- Multiplication: baseline
- Multiplication3: Some vectorizations
- MultiplicationFast: uses fast two sum and vectorized loop of VecSumFrr calls

Multiplication - performance

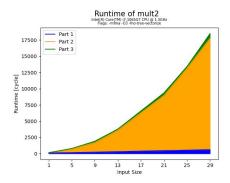


- Even more memory operations than Addition, so very low performance (< 0.75 f/c)</p>
- ~4.8x speedup

 Created a specialized implementation (input size of 4), but have not analyzed it yet

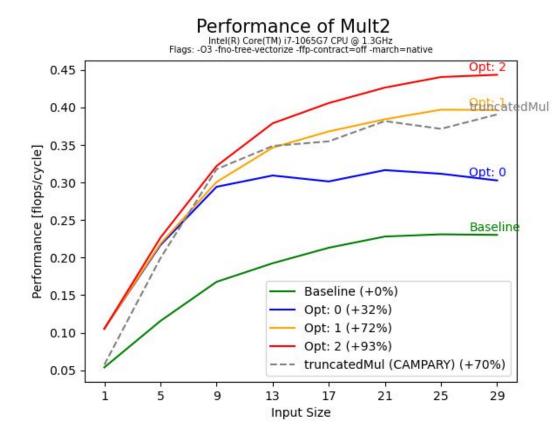
Multiplication Algo 2

```
Input: ulp-nonoverlapping FP expansions x = x_0 + ... +
     x_{n-1}; y = y_0 + \ldots + y_{m-1}.
Output: ulp-nonoverlapping FP expansion \pi = \pi_0 + \ldots +
     \pi_{r-1}.
    e \leftarrow e_{x_0} + e_{y_0}
  2: for i \leftarrow 0 to |r \cdot p/b| + 1 do
        B_i \leftarrow 1.5 \cdot 2^{e - (i+1)b + p - 1}
     end for
     for i \leftarrow 0 to min(n-1,r) do
        for j \leftarrow 0 to min(m-1, r-1-i) do
           (P, E) \leftarrow 2 \text{MultFMA}(x_i, y_i)
           \ell \leftarrow e - e_{x_i} - e_{y_i}
            sh \leftarrow |\ell/b|
           \ell \leftarrow \ell - sh \cdot b
           B \leftarrow \text{Accumulate}(P, E, B, sh, \ell)
        end for
        if j < m - 1 then
           P \leftarrow x_i \cdot y_i
           \ell \leftarrow e - e_{x_i} - e_{y_i}
           sh \leftarrow |\ell/b|
           \ell \leftarrow \ell - sh \cdot b
           B \leftarrow Accumulate(P, 0., B, sh, \ell)
        end if
     end for
     for i \leftarrow 0 to |r \cdot p/b| + 1 do
         B_i \leftarrow B_i - 1.5 \cdot 2^{e - (i+1)b + p - 1}
     end for
     \pi[0:r-1] \leftarrow \text{Renormalize}(B[0:|r\cdot p/b|+1])
     return FP expansion \pi = \pi_0 + \ldots + \pi_{r-1}.
```



- (l:1-4) initialize bins (B)
- (I:7-11) multiply each x_i with **each** y_j while i, j smaller than r
 - -> add product and overflow to correct bin
- (l:13-18) multiply each x_i with last y_j while i smaller than r
 -> add only product to correct bin
- (1:21-24) make bins non-overlapping

Multiplication Algo 2



Optimizations:

- Opt 0: Precompute
- Opt 1: Opt 0 + Replace complex with simpler functions
- Opt 2: Opt 1 + Remove function calls

Why has CAMPARY lower Performance?

- Missing optimizations (e.g. divisions by const.)
- has less FLOPS
- probably optimized for GPU

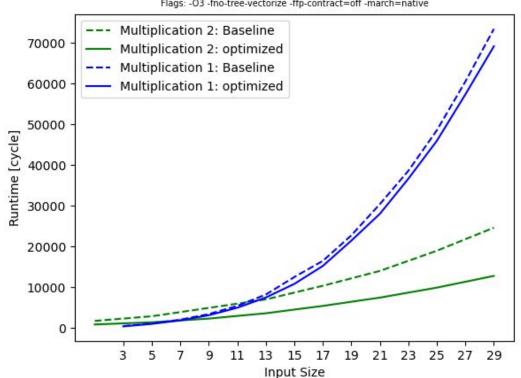
Remarks:

■ Loop unrolling didn't help

Comparison Mult1 & Mult2

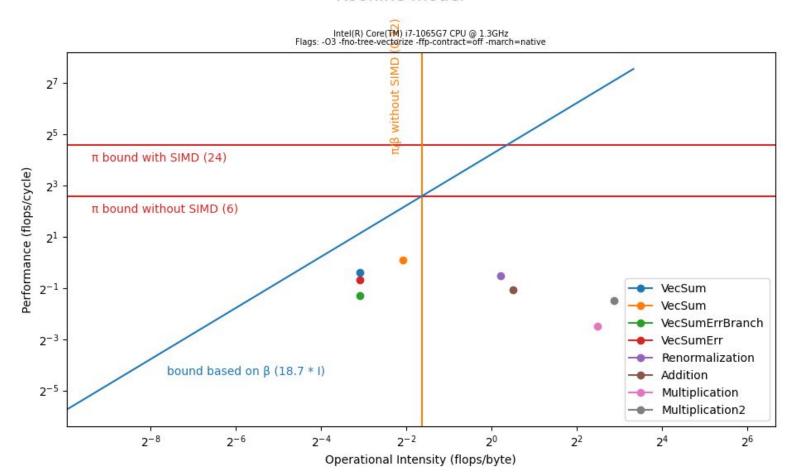


Intel(R) Core(TM) i7-1065G7 CPU @ 1.3GHz Flags: -O3 -fno-tree-vectorize -ffp-contract=off -march=native



- Multiplication 2 is faster for n > 7
- Reasons:
 - Different input requirements
 - Number of flops

Roofline model



References

- Multiplication & Addition: https://ieeexplore.ieee.org/document/7118139
- Multipliation 2: https://ieeexplore.ieee.org/document/7563270

General Remarks

- Pay attention to the length: e.g., 10 minutes typically means 7–8 slides
- Use proper visuals as much as possible, avoid text-only bullet slides
- Don't put an overview or organization slide the talk is too short
- For the very motivated, check out this small guide http://people.inf.ethz.ch/markusp/teaching/guides/guide-presentations.pdf

Typical Organization I

- Algorithm that you consider (maybe 2 slides)
 - State problem that it solves (input:..., output: ...)
 - If possible visualize how it works or show high-level pseudocode
 - State asymptotic runtime
- Cost analysis (cost measure, exact count)
- Baseline implementation (briefly explain), maybe show already performance plot and extract percentage of peak
- Optimizations you performed
 - Briefly discuss major optimizations/code versions
 - Maybe explain the most interesting in a bit greater detail

Typical Organization II

- Experimental results
 - Very brief: Experimental setup (platform, compiler)
 - Performance plot over a range of sizes with different code versions
 - Make sure you also push input size to the limit in the experiments
 - Extract overall speedup

- Every project is different so adapt as needed
- Focus on the most interesting things, don't explain everything that will be in the final report.

Try to Make Nice Plots

DFT 2ⁿ (single precision) on Pentium 4, 2.53 GHz

