



IEEE 802.11 Distribute Coordinator Function Discrete–Event Simulator against State–of–the–Art Analytical Model

Simulation and Performance Evaluation project

Luca Mosetti, Shandy Darma

November, 2023



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Analytical Models

1 Related works

IEEE 802.11—Saturation Throughput Analysis

George Bianchi

Abstract. To study the throughput of the IEEE 802.11 medium access control (MAC) protocol, we consider a simplified model of the MAC protocol. In this model, the stations are assumed to be saturated, i.e., they always have a packet to transmit. We derive an analytical expression for the throughput of the MAC protocol. The results are compared with those obtained by simulation.



1. INTRODUCTION

THE IEEE 802.11 MAC protocol is a distributed MAC protocol. It is designed to be used in a wireless LAN environment. The protocol is based on the CSMA/CD principle. The stations in the network are assumed to be saturated, i.e., they always have a packet to transmit. We derive an analytical expression for the throughput of the MAC protocol. The results are compared with those obtained by simulation.

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For more information on this topic, see the paper by Bianchi et al. [1].

Performance Analysis of the IEEE 802.11 Distributed Coordination Function

George Bianchi

Abstract. Bianchi et al. have developed the IEEE 802.11 MAC protocol. In this paper, we consider a simplified model of the MAC protocol. We derive an analytical expression for the throughput of the MAC protocol. The results are compared with those obtained by simulation.

1. INTRODUCTION

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Remarks on IEEE 802.11 DCF Performance Analysis

Giuseppe Bianchi and Rocco Tassi

Abstract. This paper presents a simplified model of the IEEE 802.11 MAC protocol. In this model, the stations are assumed to be saturated, i.e., they always have a packet to transmit. We derive an analytical expression for the throughput of the MAC protocol. The results are compared with those obtained by simulation.

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Let us denote with P the event that a station is being transmitted at a given time t . We denote with $P(t)$ the probability that a station is being transmitted at a given time t . We denote with $P(t, \tau)$ the probability that a station is being transmitted at a given time t and that the next transmission occurs at a given time τ .

2. ANALYSIS

In this section, we analyze the performance of the IEEE 802.11 MAC protocol. We derive an analytical expression for the throughput of the MAC protocol. The results are compared with those obtained by simulation.

For more information on this topic, see the paper by Bianchi et al. [1].

Refinements on IEEE 802.11 Distributed Coordination Function Model Approaches

Rocco Tassi, Giuseppe Bianchi, and Rocco Tassi

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IEEE TRANSACTIONS ON COMMUNICATIONS, VOL. 48, NO. 1, JANUARY 2000

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Introduction

2 Distribute Coordinator Function

- DCF follows Carrier Sense Multiple Access / Collision Avoidance principles
- Stations (STAs) use a “listen-before-talk” approach
- Key mechanisms: carrier sensing, random backoff, ACK

ACK

- Positive acknowledgment after successful reception
- Transmitted at the end of the packet after a SIFS
- If no ACK, reschedule transmission

Carrier Sensing Mechanism

- STAs sense the medium before transmission
- If a STA senses the channel idle for DIFS, it generates a random backoff period



Backoff strategy

2 Distribute Coordinator Function

- Backoff period expressed in backoff slots
 - Decreases when the medium is sensed idle
 - “Frozen” during a detected transmission
-
- At each transmission, backoff period uniformly chosen in $[0, CW)$
 - After unsuccessful transmission, $CW = \min \left(2 \cdot CW, 2^m \cdot W \right)$
 - Resets to W after success or retry limit R reached



BAS & RTS/CTS

2 Distribute Coordinator Function

BAS mode



- Default mechanism
- Transmit directly the entire packet

RTS/CTS mode

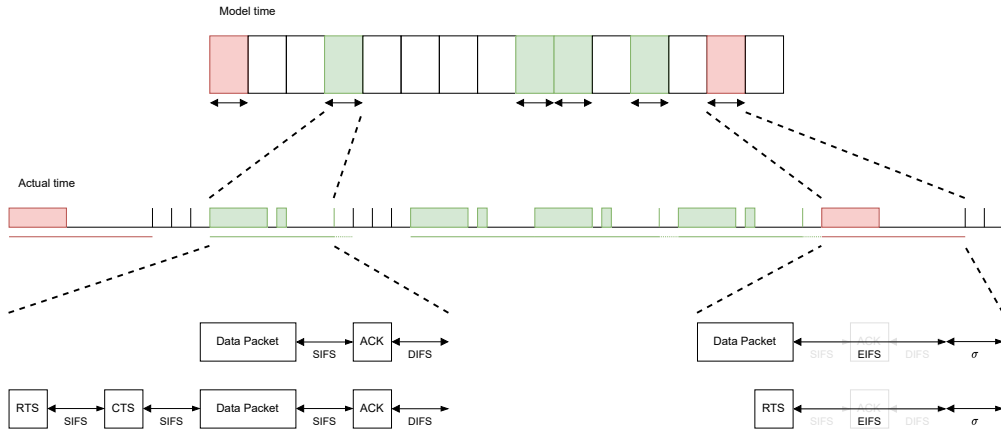


- Four-way handshaking mechanism
- Request-to-Send (RTS)
- Clear-to-Send (CTS)



Times

2 Distribute Coordinator Function





Times

2 Distribute Coordinator Function

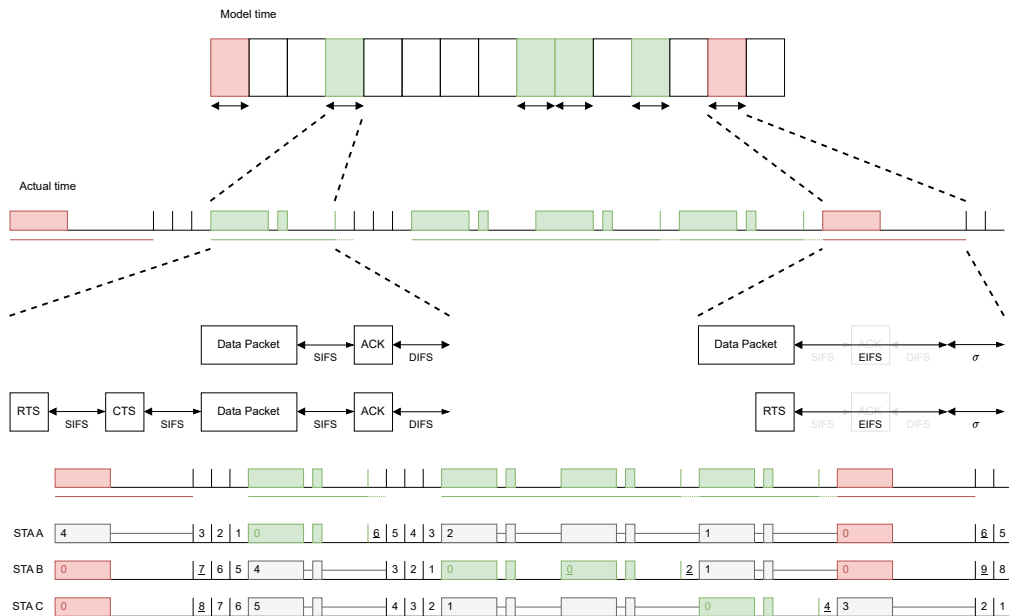




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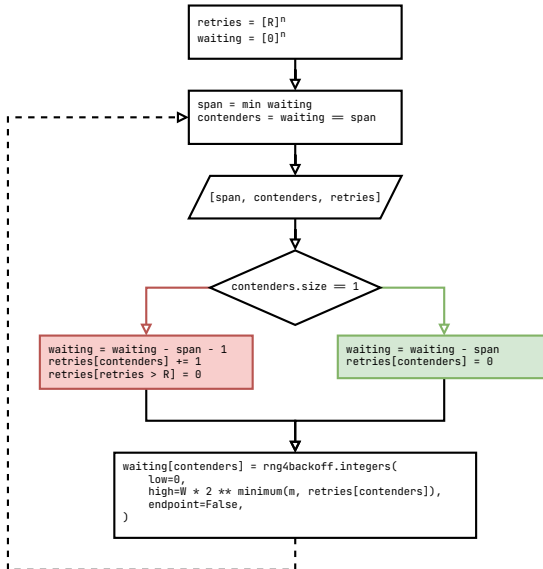
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Simulator

3 Problem definition & Setup



Characteristics:

- Discrete-Event Simulator
- Independent of any temporal specification
- Python 3.11

Assumptions:

- Constant # of STAs
- Constant payload size
- Transmission queue of each STA is always nonempty
- W/out hidden terminals
- W/out capture effect



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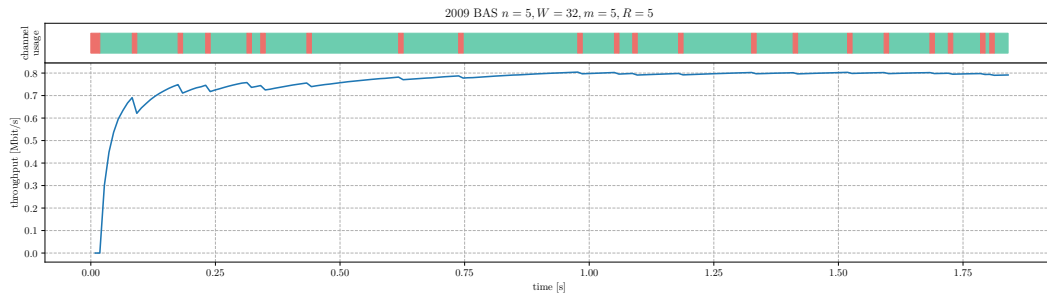
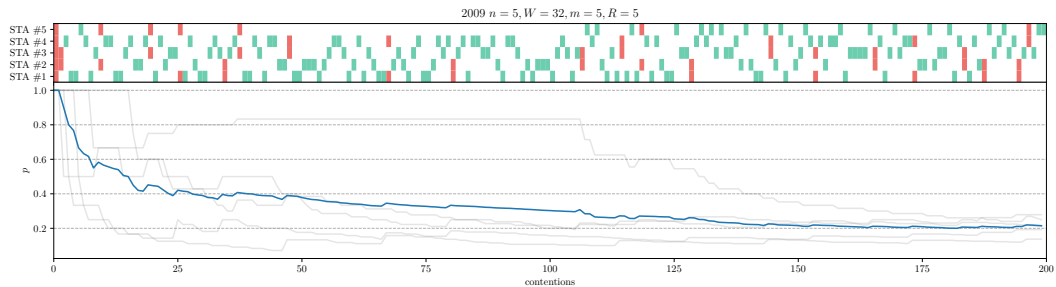
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p and S

4 Metrics





Probability of collision

4 Metrics

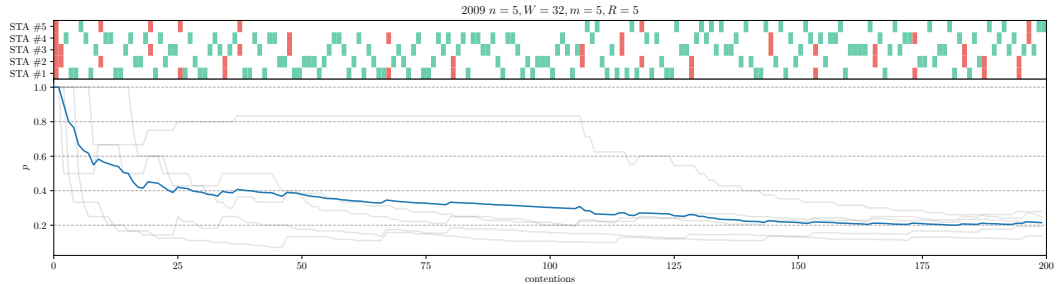


Figure: Measurement of collision probability p . Upper figure shows whether a transmission or a collision occurred in each contention round. Lower figure shows the collision probability, with grey lines describe probability of each stations while blue line describes the final mean of all probabilities.



Saturated throughput

4 Metrics

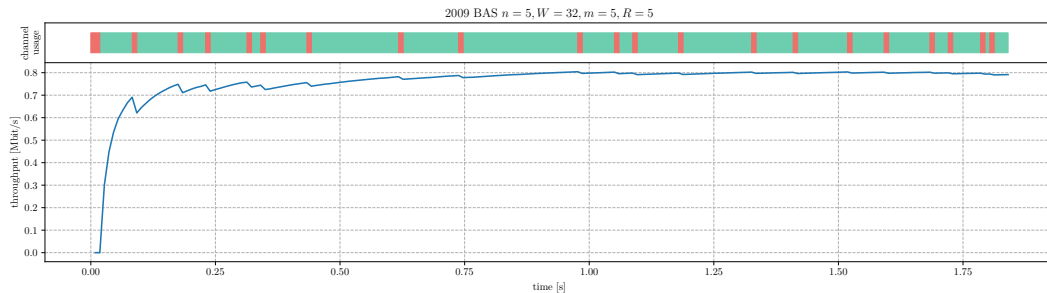
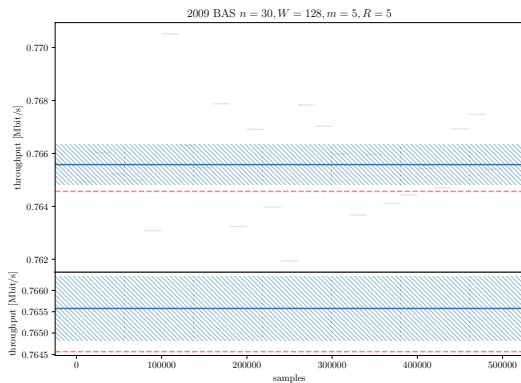


Figure: Measurement of saturated throughput S . Upper figure shows the occurrences of transmissions and collisions. Lower figure shows the saturated throughput of the channel. Each transmission improves throughput while each collision worsens throughput.



Non-overlapping batch over means

4 Metrics



$$\underbrace{Y_1, \dots, Y_s}_{B_1}, \underbrace{Y_{s+1}, \dots, Y_{2s}}_{B_2}, \dots, \underbrace{Y_{(b-1)s+1}, \dots, Y_{bs}}_{B_b}$$

$$\bar{Z}_b = \frac{1}{b} \sum_i^b Z_i \left[\bar{Z}_b \pm t_{b-1, \frac{1+\gamma}{2}} \sqrt{\frac{\hat{V}}{b}} \right]_\gamma$$

$$\hat{V} = \frac{1}{b-1} \sum_i^b (Z_i - \bar{Z}_b)^2$$



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Formulae

5 Analytical Model

σ slot time

L payload size

T_{MPDU} the time to transmit the MPDU

T_{ACK} the time to transmit an ACK

T_{RTS} the time to transmit a RTS frame

T_{CTS} the time to transmit a CTS frame

SIFS the SIFS time

DIFS the DIFS time

n number of stations

W minimum contention window

m maximum backoff stage, where maximum contention window is $2^m W$

R maximum number of retries, after which the transmission is dropped

p probability of collision

τ probability that a station will transmit

S saturation throughput



Formulae

5 Analytical Model

$$p = 1 - (1 - \tau)^{n-1}$$

$$\tau = \frac{1}{1 + \frac{1-p}{2(1-p^{R+1})} \left[\sum_{j=0}^R p^j \cdot (2^j W - 1) - (1 - p^{R+1}) \right]}$$

$$S = \frac{\mathbf{P}_s \overline{\mathbf{E}[L]}}{(1 - \mathbf{P}_b)\sigma + \mathbf{P}_s \overline{\mathbf{T}_s} + (\mathbf{P}_b - \mathbf{P}_s) \overline{\mathbf{T}_c}}$$

$$\mathbf{P}_b = 1 - (1 - \tau)^n \quad \mathbf{P}_s = n\tau(1 - \tau)^{n-1}$$

$$\overline{\mathbf{E}[L]} = L \frac{W}{W-1} \quad \overline{\mathbf{T}_s} = \mathbf{T}_s \frac{W}{W-1} + \sigma \quad \overline{\mathbf{T}_c} = \mathbf{T}_c + \sigma$$

$$\mathbf{T}_s^{\text{BAS}} = T_{\text{MPDU}} + \text{SIFS} + T_{\text{ACK}} + \text{DIFS}$$

$$\mathbf{T}_c^{\text{BAS}} = \mathbf{T}_s^{\text{BAS}}$$

$$\mathbf{T}_s^{\text{RTS/CTS}} = T_{\text{RTS}} + \text{SIFS} + T_{\text{CTS}} + \text{SIFS} + T_{\text{MPDU}} + \text{SIFS} + T_{\text{ACK}} + \text{DIFS}$$

$$\mathbf{T}_c^{\text{RTS/CTS}} = T_{\text{RTS}} + \text{SIFS} + T_{\text{ACK}} + \text{DIFS}$$



Formulae

5 Analytical Model

$$\tau = \frac{1}{1 + \frac{1-p}{2(1-p^{R+1})} \left[\sum_{j=0}^R p^j \cdot (2^j W - 1) - (1 - p^{R+1}) \right]}$$

↓

$$\tau = \frac{1}{1 + \frac{1-p}{2(1-p^{R+1})} \left[\sum_{j=0}^R p^j \cdot (\mathfrak{W}_j - 1) - (1 - p^{R+1}) \right]}$$

$$\mathfrak{W}_j = W \cdot 2^{\min(m,j)} \quad j \in [0, R]$$



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Parameters

6 Results

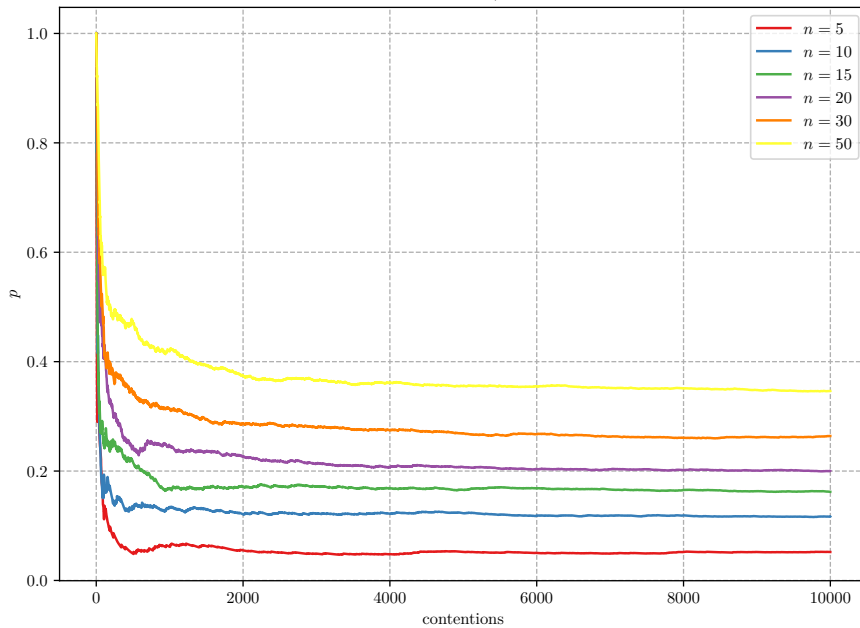
Parameter	2009 Simulator
L	1023 Byte
MAC header	34 Byte
PHY header	16 Byte
ACK	14 Byte + PHY header
RTS	20 Byte + PHY header
CTS	14 Byte + PHY header
Channel Bit Rate	1 Mbit/s
σ	50 μ s
SIFS	28 μ s
DIFS	SIFS + $2 \cdot \sigma$



Initialisation bias

6 Results

2009 $W = 128, m = 5$

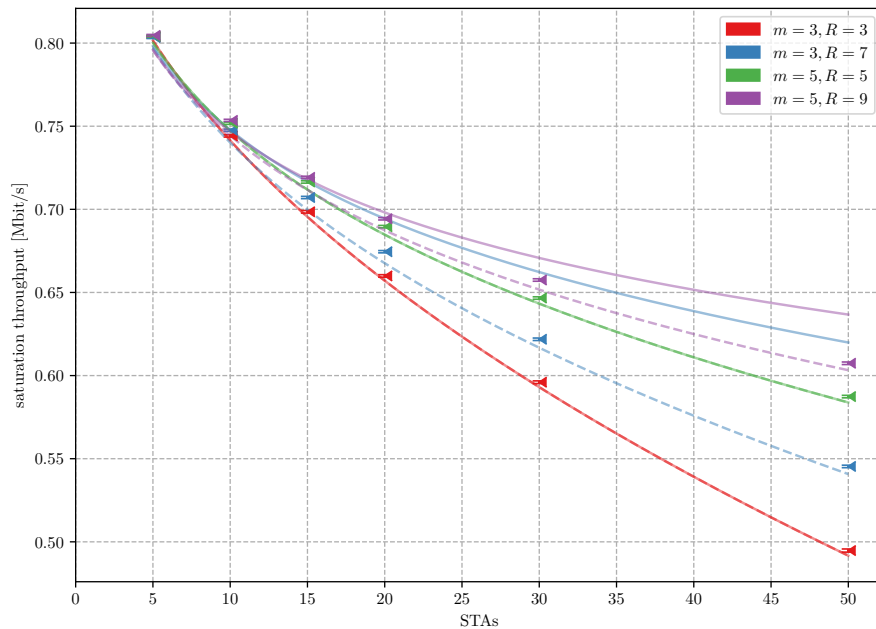




Saturation throughput

6 Results

2009 BAS $W = 32$

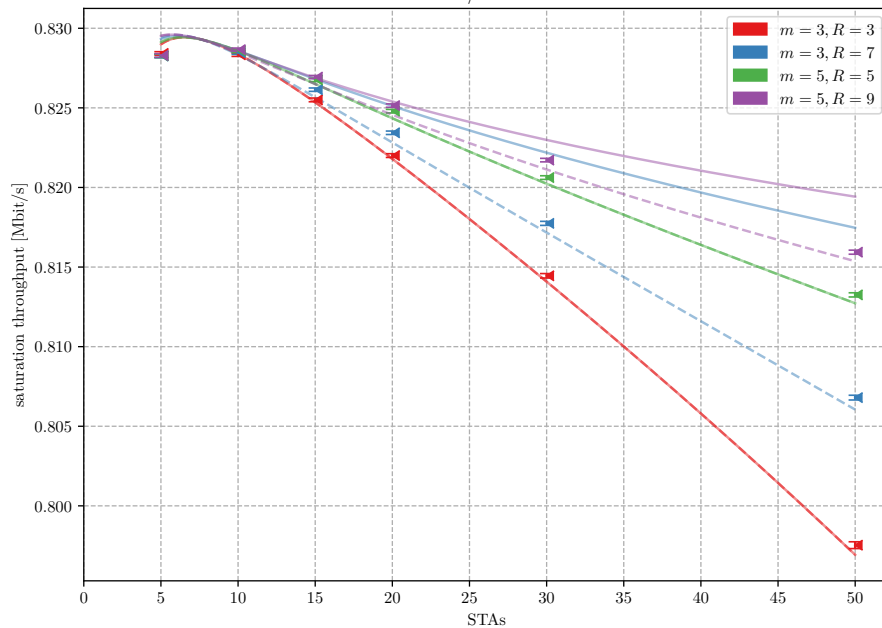




Saturation throughput

6 Results

2009 RTS/CTS $W = 32$

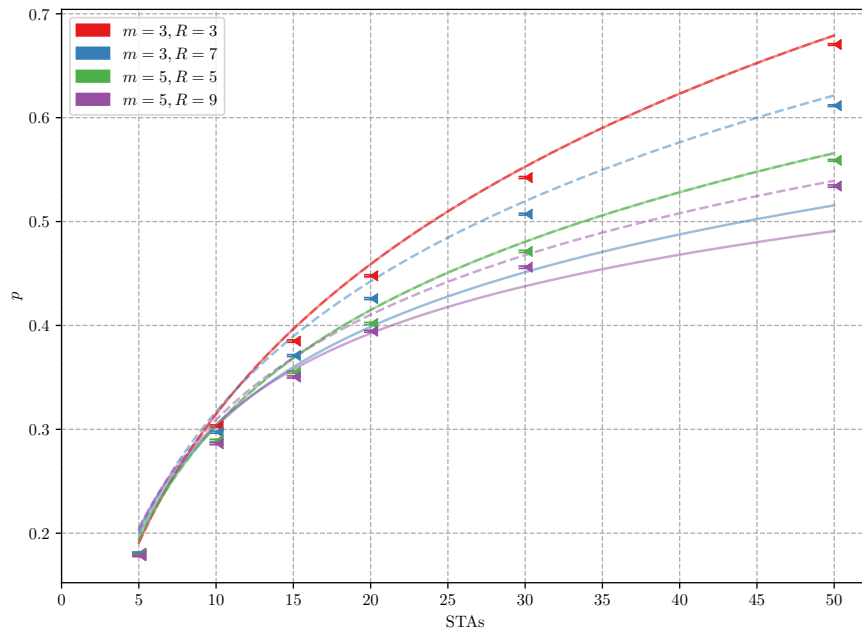




Probability of collision

6 Results

2009 $W = 32$

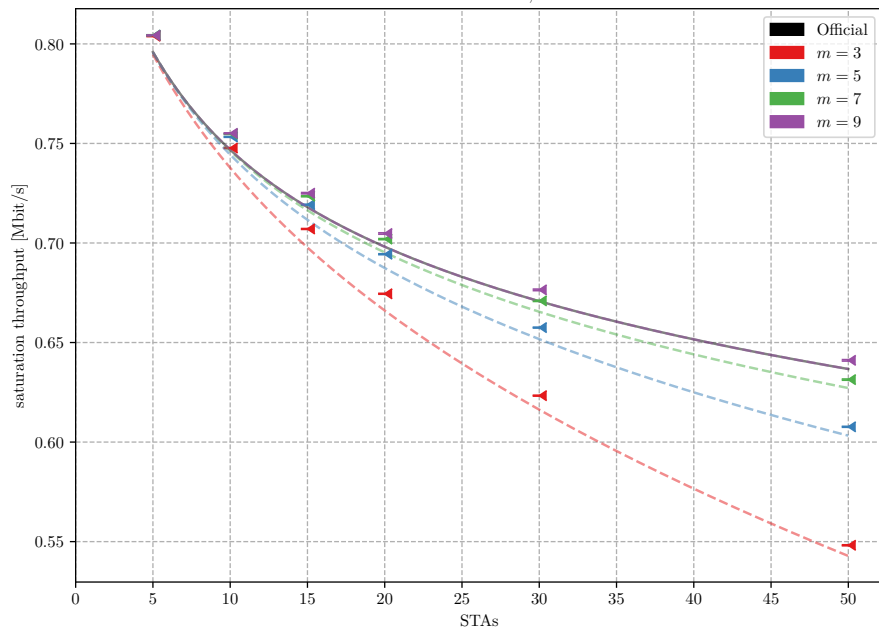




Focus

6 Results

2009 BAS $W = 32, R = 9$

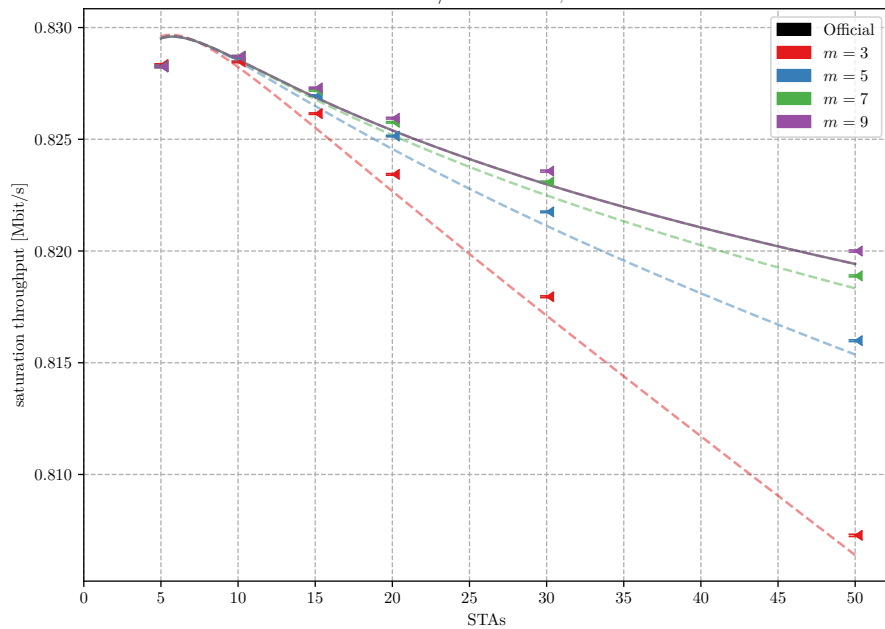




Focus

6 Results

2009 RTS/CTS $W = 32, R = 9$

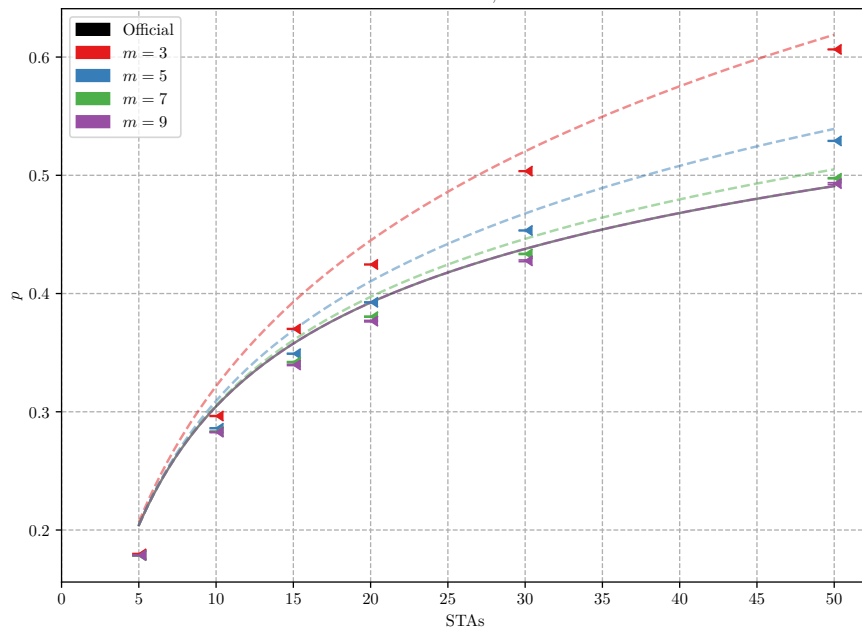




Focus

6 Results

2009 $W = 32, R = 9$





All in all

6 Results

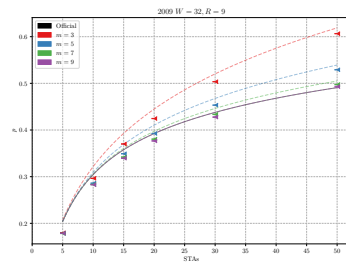
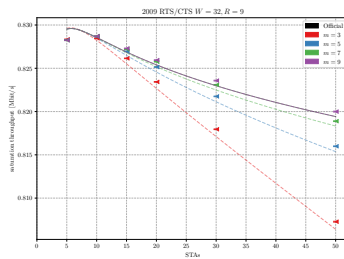
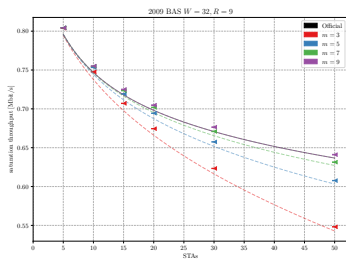
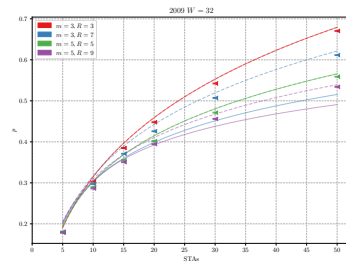
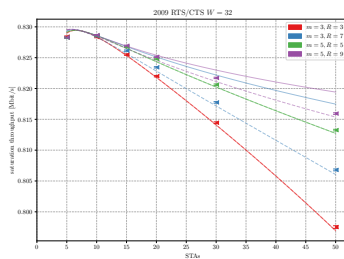
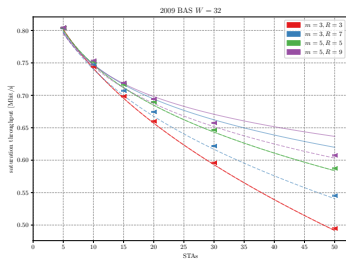




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Differences between analytical models

7 Appendix

$$p = 1 - (1 - \tau)^{n-1}$$

$$\tau = \frac{2}{1 + W + pW \sum_{i=0}^{m-1} (2p)^i}$$

$$\tau = \frac{1}{1 + \frac{1-p}{2(1-p^{R+1})} \left[\sum_{j=0}^R p^j \cdot (2^j W - 1) - (1 - p^{R+1}) \right]}$$

$$S = \frac{P_s P_{tr} L}{(1 - P_{tr})\sigma + P_{tr} P_s T_s + P_{tr}(1 - P_s) T_c}$$

$$S = \frac{P_s \overline{E[L]}}{(1 - P_b)\sigma + P_s \overline{T_s} + (P_b - P_s) \overline{T_c}}$$

$$P_{tr} = 1 - (1 - \tau)^n \quad P_s = \frac{n\tau(1 - \tau)^{n-1}}{P_{tr}}$$

$$P_b = 1 - (1 - \tau)^n \quad P_s = n\tau(1 - \tau)^{n-1}$$

$$\overline{E[L]} = L \frac{W}{W - 1} \quad \overline{T_s} = T_s \frac{W}{W - 1} + \sigma \quad \overline{T_c} = T_c + \sigma$$



Differences between analytical models

7 Appendix

$$T_s^{\text{BAS}} = T_{\text{MPDU}} + \delta + \text{SIFS} + T_{\text{ACK}} + \delta + \text{DIFS}$$

$$T_c^{\text{BAS}} = T_{\text{MPDU}} + \delta + \text{DIFS}$$

$$T_s^{\text{RTS/CTS}} = T_{\text{RTS}} + \delta + \text{SIFS} + T_{\text{CTS}} + \delta + \text{SIFS} \\ + T_{\text{MPDU}} + \delta + \text{SIFS} + T_{\text{ACK}} + \delta + \text{DIFS}$$

$$T_c^{\text{RTS/CTS}} = T_{\text{RTS}} + \delta + \text{DIFS}$$

$$T_s^{\text{BAS}} = T_{\text{MPDU}} + \text{SIFS} + T_{\text{ACK}} + \text{DIFS}$$

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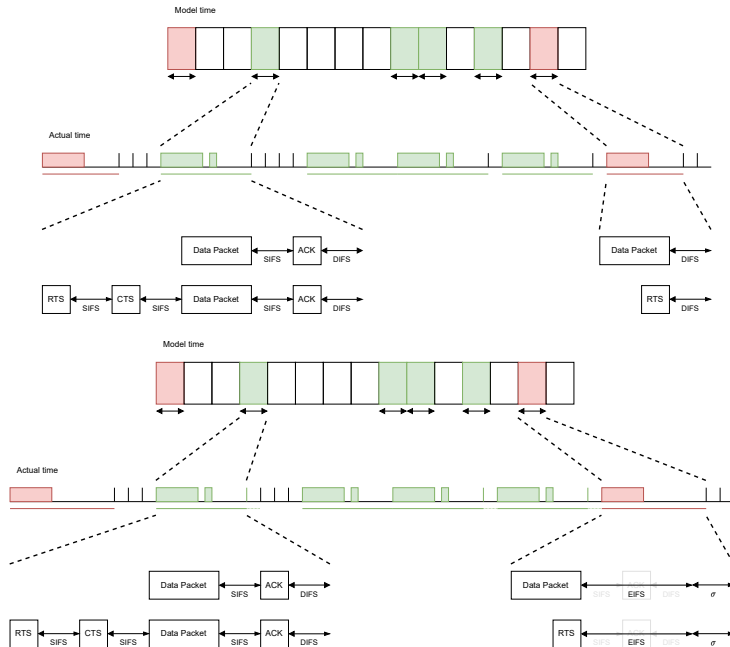
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$$T_c^{\text{RTS/CTS}} = T_{\text{RTS}} + \text{SIFS} + T_{\text{ACK}} + \text{DIFS}$$



Differences between simulators

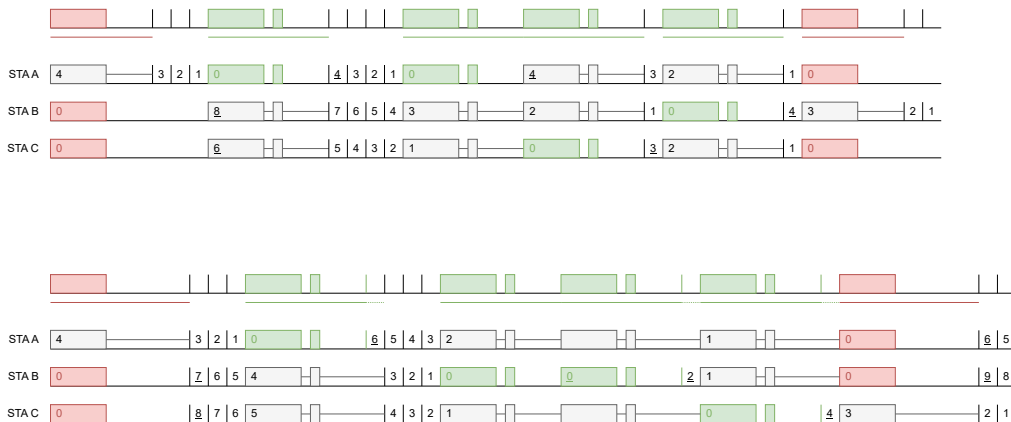
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Differences between simulators

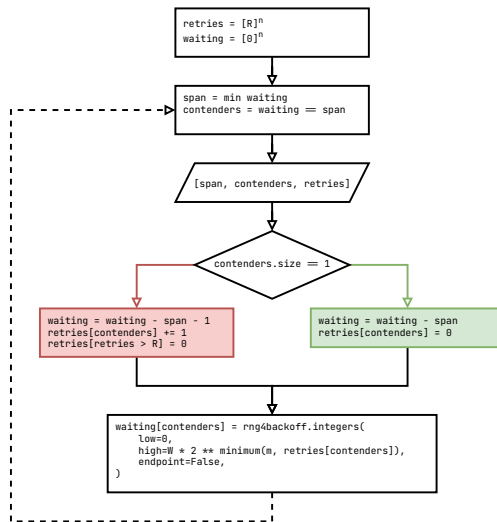
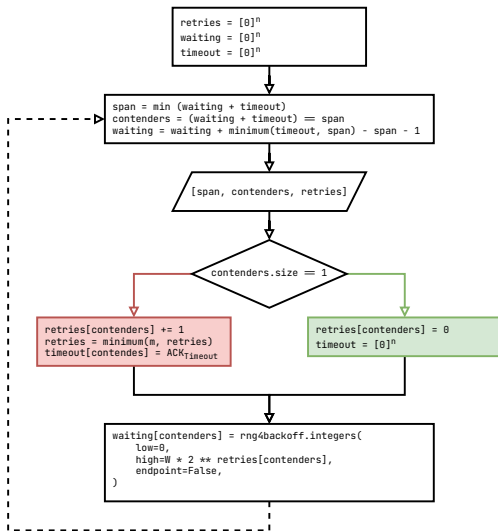
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Differences between simulators

7 Appendix



simulation()		simulation()
seeds: t.Iterator[NDArray[np.uint32]],	27 27	seeds: t.Iterator[NDArray[np.uint32]],
n: int,	28 28	n: int,
W: int,	29 29	W: int,
m: int,	30 30	m: int,
ack_timeout: int,	>> 31 31 <<	R: int,
) → t.Iterator[Log]:	32 32) → t.Iterator[Log]:
rng4backoff: np.random.Generator = np.random.default_rng(next(seeds))	33 33	rng4backoff: np.random.Generator = np.random.default_rng(next(seeds))
	34 34	
timeouts: NDArray[bool] = np.full(n, fill_value=False)	>> 35 35 <<	retries: NDArray[int] = np.full(n, R, dtype=int)
retries: NDArray[int] = np.zeros(n, dtype=int)	36 36	waiting: NDArray[int] = np.zeros(n, dtype=int)
waiting: NDArray[int] = np.zeros(n, dtype=int)	37 37	
	38 38	
while True:	39 39	while True:
	40 40 <<	span: int = np.amin(waiting)
span: int = np.amin(waiting + timeouts * ack_timeout)	>> 41 41	mask: NDArray[bool] = waiting == span
mask: NDArray[bool] = (waiting + timeouts * ack_timeout) == span	42 42	contenders: NDArray[int] = np.flatnonzero(mask)
contenders: NDArray[int] = np.flatnonzero(mask)	43 43	
	44 44	
yield Log(45 45	yield Log(
span=span,	46 46	span=span,
contenders=mask,	47 47	contenders=mask,
attempt=retries.copy(),	48 48	attempt=retries.copy(),
)	49 49)
	50 50 <<	waiting = waiting + span
waiting = waiting + np.minimum(timeouts * ack_timeout, span) - span - 1	>> 51 51	
	52 52	match contenders.size:
match contenders.size:	53 53	case 1:
case 1:	54 54 <<	retries[contenders] = 0
timeouts = np.full(n, fill_value=False)	>> 55 55	case _:
retries[contenders] = 0	56 56 <<	waiting -= 1
case _:	57 57	retries[contenders] += 1
timeouts = mask	>> 58 58 <<	retries[retries > R] = 0
retries[contenders] += 1	59 59	
retries.clip(max=m, out=retries)	>> 60 60	waiting[contenders] = rng4backoff.integers(
	61 61	low=0,
waiting[contenders] = rng4backoff.integers(62 62 <<	high=W * 2 ** np.minimum(n, retries[contenders]),
low=0,	63 63	endpoint=False,
high=W * 2 ** retries[contenders],	>> 64 64)
endpoint=False,	65 65	
)	66 66	



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Parameters

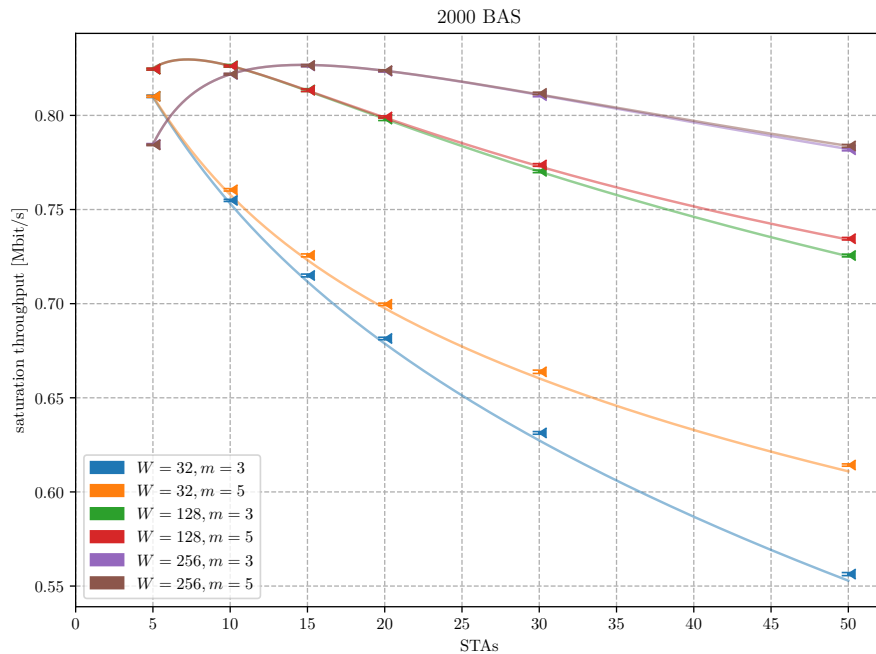
8 Results

Parameter	2000 Simulator	2009 Simulator
L	1023 Byte	1023 Byte
MAC header	34 Byte	34 Byte
PHY header	16 Byte	16 Byte
ACK	14 Byte + PHY header	14 Byte + PHY header
RTS	20 Byte + PHY header	20 Byte + PHY header
CTS	14 Byte + PHY header	14 Byte + PHY header
Channel Bit Rate	1 Mbit/s	1 Mbit/s
σ	$50 \mu s$	$50 \mu s$
SIFS	$28 \mu s$	$28 \mu s$
DIFS	$SIFS + 2 \cdot \sigma$	$SIFS + 2 \cdot \sigma$
δ	$1 \mu s$	—
$ACK_{Timeout}$	$300 \mu s$	—



Saturation throughput

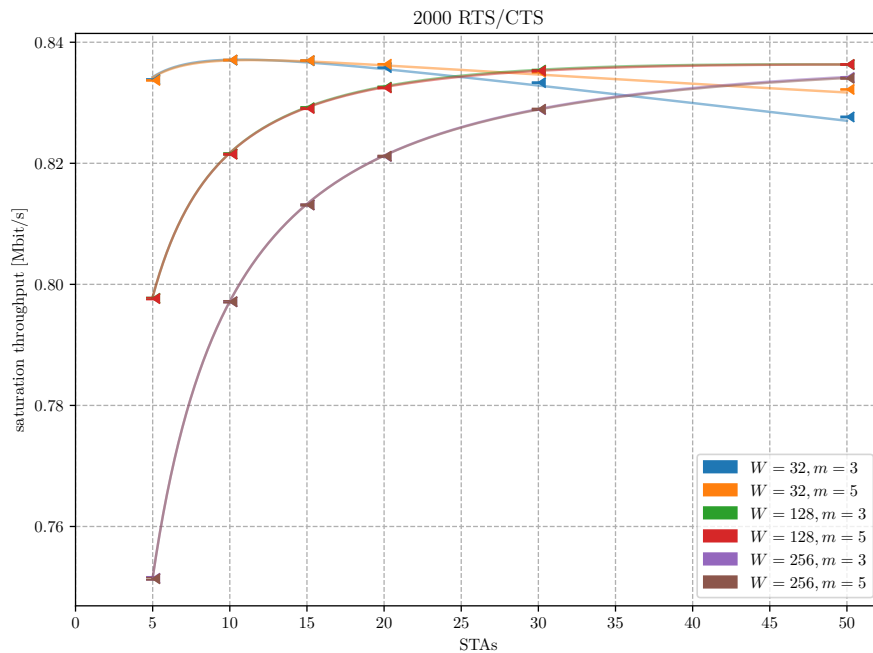
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Saturation throughput

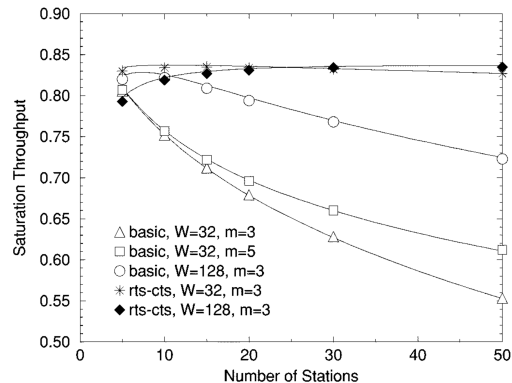
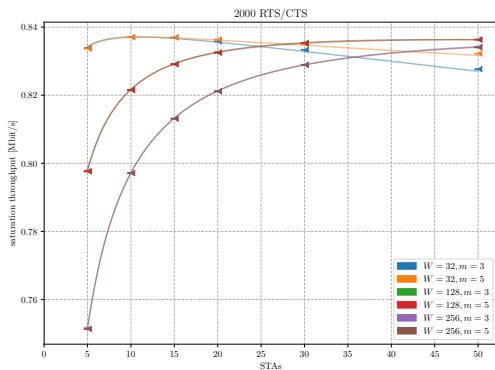
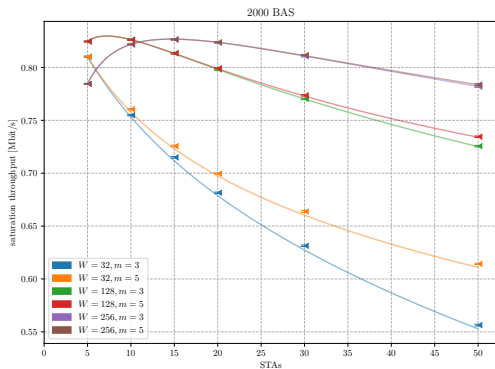
8 Results





Saturation throughput

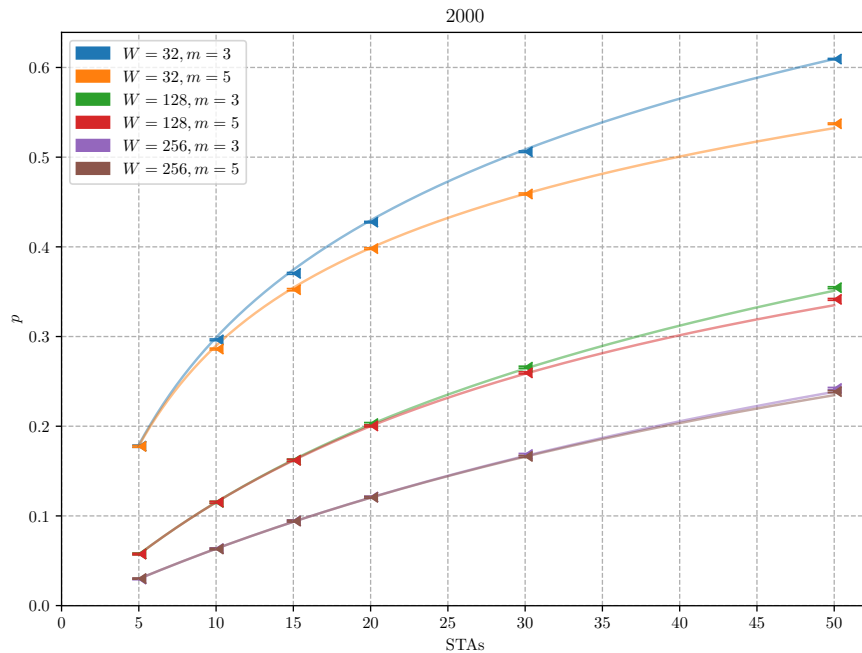
8 Results





Probability of collision

8 Results

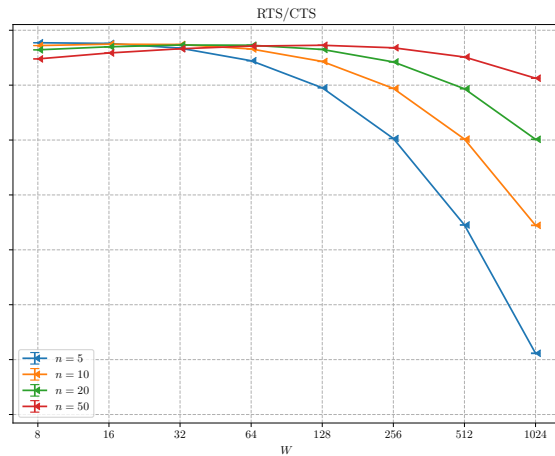
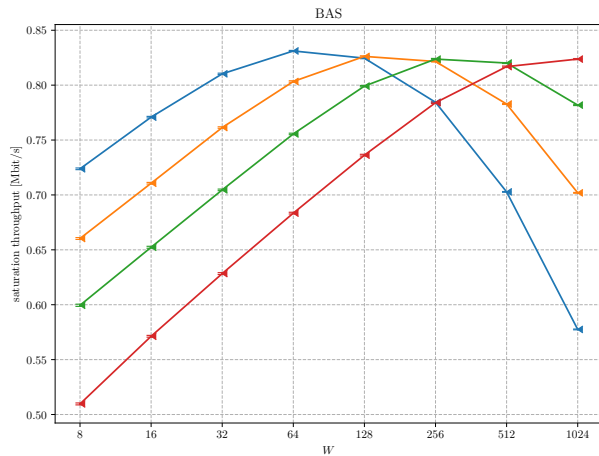




Experiment

8 Results

2000 $m = 6$

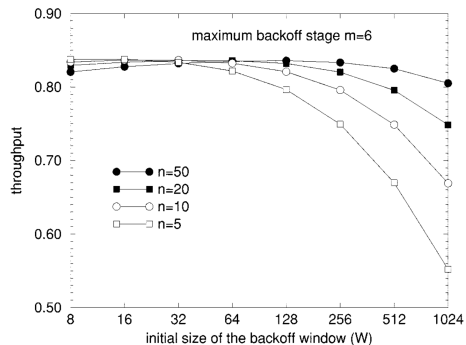
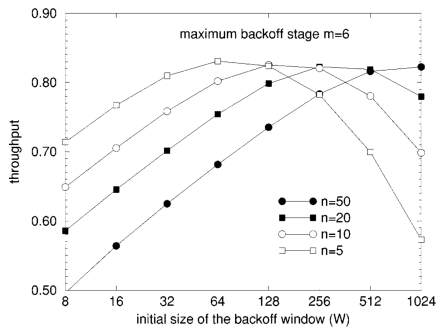
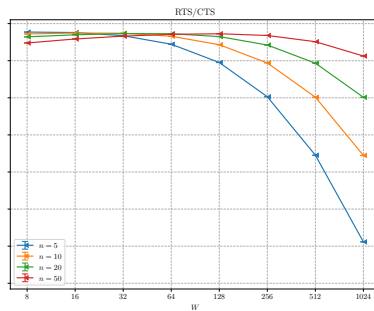
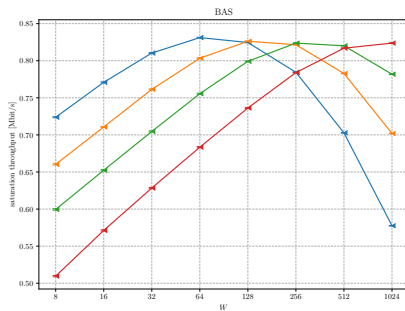




Experiment

8 Results

2000 $m=6$

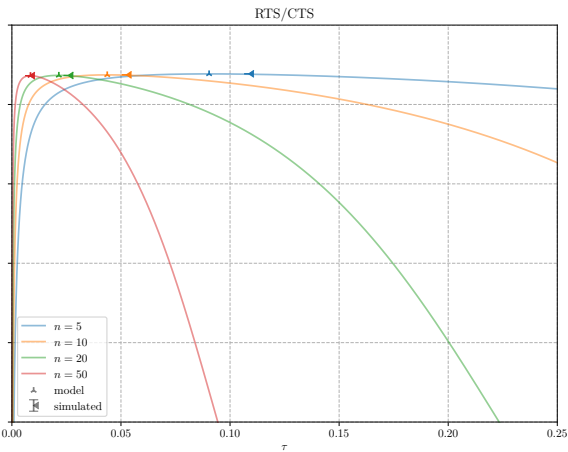
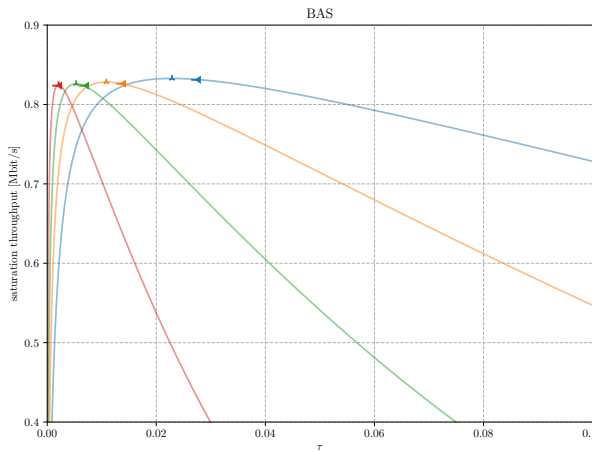




Experiment

8 Results

2000 $m = 6$

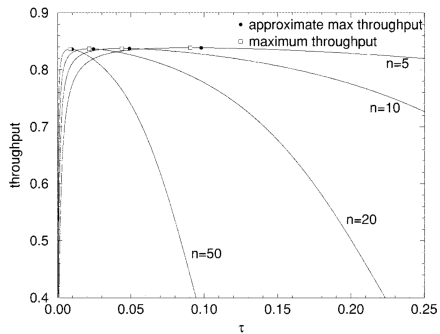
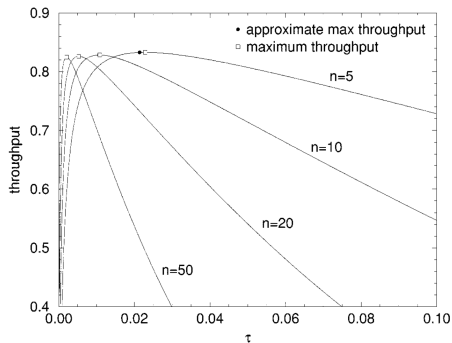
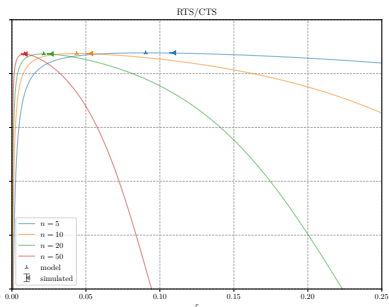
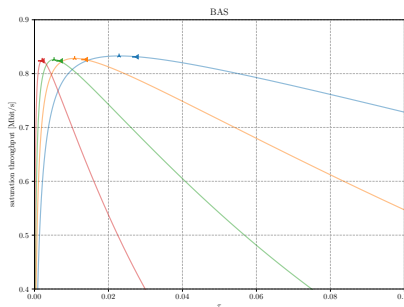




Experiment

8 Results

2000 $m = 6$

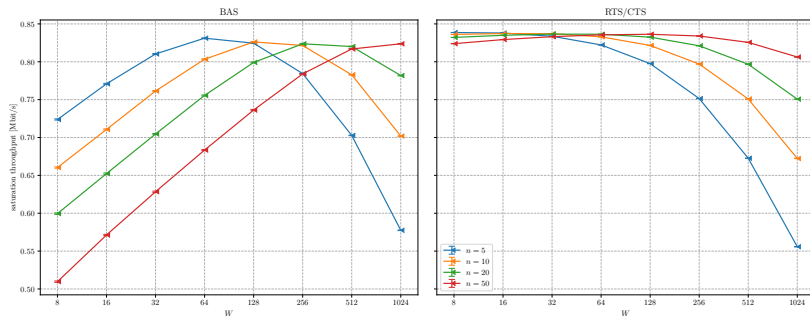




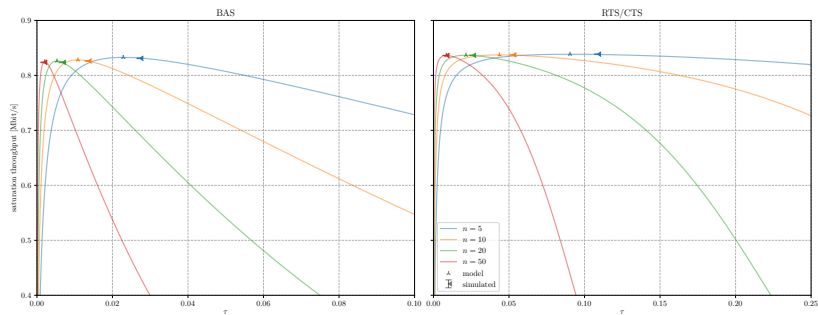
Experiment

8 Results

2000 $m = 6$



2000 $m = 6$





Q&A