Operating Systems and Concurrency

Concurrency 3
COMP2007

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Recap

Last Lecture - Approaches to Mutual Exclusion

- Software approaches: Peterson's solution
- Hardware approaches:
 - test_and_set()
 - compare_and_swap()
- Mutexes as an abstraction providing binary locks

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Recap

Concurrency Primitives

- Recall Mutexes are a locking abstraction for providing mutual exclusion.
- Often provided by the operating system, via an API such as pthreads.
- They are binary either a thread has currently acquired the mutex or it hasn't!

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OS approaches

- Semaphores are another abstraction for mutual exclusion and process synchronisation, often provided by the operating system
 - They have a **capacity**, either a positive number or infinity.
 - We distinguish between binary (2 valued) and counting semaphores (N-valued or unbounded).
- Two functions are used to manipulate semaphores (think of the counter++ example)
 - wait () is called when a resource is acquired, the capacity is decremented
 - signal() or /post() is called when a resource is released, the capacity is incremented.
- The semaphore can only be acquired when its currently capacity is strictly positive.

• A thread calling post does not have to have previous called wait.

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OS approaches

```
typedef struct {
  int value;
  struct process * list;
} semaphore;
```

Figure: Conceptual definition of a semaphore

```
void wait(semaphore* S) {
   S->count--;
   if(S->count < 0) {
      //add process to S->list
      block(); // system call
   }
}
```

Figure: Conceptual implementation of a wait()

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OS approaches

```
void post(semaphore* S) {
   S->count++;
   if(S->count <= 0) {
      // remove process P from S->list
      wakeup(P);
   }
}
```

Figure: Conceptual implementation of post()

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Implementation

```
Thread 1
                              Thread 2
                                                            Thread 3
wait(\&s) 1 => 0
                               . . .
                                                             . . .
. . .
                               . . .
                              wait(&s)
                                                             . . .
                                                            wait(&s)
post(&s)
                               (wakeup)
                               . . .
                               . . .
                              post(&s)
                                                             (wakeup)
                               . . .
                                                             . . .
                                                            post(&s)
                               . . .
                               . . .
                                                             . . .
```

Figure: Semaphore example

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Implementation

```
Thread 1
                               Thread 2
                                                               Thread 3
wait(&s)
                                . . .
                                                               . . .
. . .
                               wait(\&s) 0 \Rightarrow -1
                                                               . . .
                                                               wait(&s)
post(&s)
                                (wakeup)
                                . . .
                                . . .
                               post(&s)
                                                               (wakeup)
                                . . .
                                                               . . .
                                                               post(&s)
                                . . .
                                . . .
                                                               . . .
```

Figure: Semaphore example

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Implementation

```
Thread 1
                               Thread 2
                                                               Thread 3
                                . . .
                                                               . . .
wait(&s)
                                . . .
                                                               . . .
. . .
                                . . .
                               wait(&s)
                                                               . . .
                                                               wait(\&s) -1 => -2
post(&s)
                                (wakeup)
                                . . .
. . .
                                                               . . .
                                . . .
                               post(&s)
                                                               (wakeup)
                                . . .
                                                               . . .
                                                               post(&s)
                                . . .
                                . . .
                                                               . . .
```

Figure: Semaphore example

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Implementation

```
Thread 1
                                Thread 2
                                                                Thread 3
wait(&s)
                                . . .
                                                                . . .
                                . . .
. . .
                                wait(&s)
                                                                . . .
                                                                wait(&s)
post(\&s) -2 \Rightarrow -1
                                (wakeup)
. . .
                                . . .
                                . . .
                                post(&s)
                                                                (wakeup)
                                . . .
                                                                . . .
                                                                post(&s)
                                . . .
                                . . .
                                                                . . .
```

Figure: Semaphore example

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Implementation

```
Thread 1
                              Thread 2
                                                            Thread 3
wait(&s)
                              . . .
                                                             . . .
. . .
                              . . .
                              wait(&s)
                                                             . . .
                                                            wait(&s)
post(&s)
                              (wakeup)
                              . . .
                                                             . . .
                              post(\&s) -1 => 0
                                                            (wakeup)
                                                            . . .
                              . . .
                                                            post(&s)
                              . . .
                              . . .
                                                             . . .
```

Figure: Semaphore example

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Implementation

```
Thread 1
                               Thread 2
                                                               Thread 3
. . .
                                . . .
                                                               . . .
wait(&s)
                                . . .
                               wait(&s)
                                                               wait(&s)
                                . . .
post(&s)
                                (wakeup)
                                                               . . .
                                . . .
                                . . .
                               post(&s)
                                                               (wakeup)
                                . . .
                                                               post(\&s) 0 \Rightarrow 1
                                . . .
                                . . .
```

Figure: Semaphore example

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OS approaches

- Calling wait () will block the process when the internal counter is not positive
 - The process joins the a queue blocking on the semaphore
 - The process state is changed from running to blocked
 - Ontrol is transferred to the process scheduler
- Calling post () removes a process from the blocked queue if available:
 - The process state is changed from blocked to ready
 - ② Different queueing strategies can be employed to remove processes so avoid unjustified assumptions in your code.

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OS approaches

- The queue length is the number of processes waiting on the semaphore.
- block() and wakeup() are system calls provided by the operating system.

• post() and wait() must be atomic.

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OS approaches

```
void post(semaphore* S) {
  lock(&mutex);
  S->count++;
  if(S->count <= 0) {
    // remove process P from queue
    wakeup(P);
  }
  unlock(&mutex);
}</pre>
```

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Posix Semaphores

Counter++ revisited

- Semaphores within the same process can be declared as variables of the type sem_t
 - sem_init() initialises the value of the semaphore
 - sem_wait() decrements the value of the semaphore
 - sem_post () increments the values of the semaphore
- An explanation of any of these functions can be found in the man pages,
 e.g. by typing man sem_init on the Linux command line

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Posix Semaphores

Example

```
sem_t s;
int sum = 0;
void* calc(void* arg) {
  int const iterations = 50000000;
  for(int i = 0; i < iterations; i++) {</pre>
    sem_wait(&s);
    sum++;
    sem_post(&s);
  return 0;
int main() {
  pthread t tid1, tid2;
  sem_init(&s,0,1);
  pthread create (&tid1, NULL, calc, 0);
  pthread_create(&tid2, NULL, calc, 0);
  pthread join(tid1,NULL);
  pthread join(tid2, NULL);
  printf("The value of sum is: %d\n", sum);
```

Real-world issues

Standards support

Question

Does the previous code give the right answer on my Mac?

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Real-world issues

Standards support

Question

Does the previous code give the right answer on my Mac?

Answer

Unfortunately, running the code on my Mac gives an answer slightly below 10000000! Details:

- Compiles with compiler warnings that sem_init is deprecated.
- sem_init is always failing, returning −1!
- Using named semaphores will work see the lab for these.
- Even then, code using named semaphores must run as root on a Mac to call sem_unlink successfully.

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Real-world issues

Standards support

Lessons

- Never ignore compiler warnings.
- Always check return values slide examples don't for space reasons.
- Test code thoroughly implicit assumption Mac would work like Linux was wrong!
- Be aware of platform specific issues such as sem_unlink behaviour on Mac.
- Use the appropriate concurrency primitives the example really needed a mutex.

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Efficiency

How/when to synchronise

- Synchronising code does result in a performance penalty
 - Synchronise only when necessary.
 - Synchronise as few instructions as possible.

Carefully consider how to synchronise!

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Using Semaphores

Counter++ revisited

```
void* calc(void* increments) {
  int number_of_iterations = 50000000;
  int total = 0;
  for(int i = 0; i < number_of_iterations; i++) {
    total++; // Pretend this is non-trivial to work out
  }
  sem_wait(&s);
  sum+=total;
  sem_post(&s);
  return 0;
}</pre>
```

Figure: Fast synchronised sums

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Caveats

Potential Difficulties

- Starvation: poorly designed queueing approaches (e.g. LIFO) may result in fairness violations
- Deadlocks: two or more processes are waiting indefinitely for an event that can be caused only by one of the waiting processes
 - I.e., every process in a set is waiting for an event that can only be caused by another process in the same set
 - E.g., consider the following sequence of instructions on semaphores

```
P0 P1
wait(S); ...
wait(Q);
wait(Q); ...
wait(S);
```

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Problem Description

- Producer(s) and consumer(s) share a buffer of values this could for example be a printer queue.
 - The buffer can be of **bounded** (maximum size *N*) or **unbounded size**.
 - There can any number of producers or consumers.
- A producer attempts to add items and blocks if the buffer is full.
- A consumer attempts to remove items and blocks if the buffer is empty.

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One Consumer, One Producer, Unbounded Buffer

- The simplest version of the problem has one producer, one consumer, and a buffer of unbounded size
- A counter (index) variable keeps track of the number of items in the buffer
- It uses two binary semaphores:
 - sync synchronises access to the buffer (counter), initialised to 1
 - delay_consumer ensures that the consumer blocks when there are no items available, initialised to 0

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One Consumer, One Producer, Unbounded Buffer: First Attempt

```
void * consumer(void * p)
                                          void * producer(void * p)
 sem wait(&delay consumer); 0 => -1
 while (1)
                                            while (1)
  sem wait (&sync);
                                            sem wait (&svnc);
  items--:
                                            items++:
  printf("%d\n", items);
                                            printf("%d\n", items);
                                            if(items == 1)
  sem post (&sync);
  if(items == 0)
                                             sem post(&delay consumer);
   sem wait (&delay consumer);
                                            sem post (&sync);
```

Figure: Single producer/consumer with unbounded buffer

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One Consumer, One Producer, Unbounded Buffer: First Attempt

```
void * consumer(void * p)
{
    sem_wait(&delay_consumer);
    while(1)
    {
        sem_wait(&sync);
        items--;
        printf("%d\n", items);
        sem_post(&sync);
        if(items == 0)
        sem_wait(&delay_consumer);
    }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync); 1 => 0
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer with unbounded buffer

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One Consumer, One Producer, Unbounded Buffer: First Attempt

```
void * consumer(void * p)
{
  sem_wait(&delay_consumer);
  while(1)
  {
    sem_wait(&sync);
    items--;
    printf("%d\n", items);
    sem_post(&sync);
    if(items == 0)
    sem_wait(&delay_consumer);
  }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++; 0 => 1
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer with unbounded buffer

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One Consumer, One Producer, Unbounded Buffer: First Attempt

```
void * consumer(void * p)
{
  sem_wait(&delay_consumer);
  while(1)
  {
    sem_wait(&sync);
    items--;
    printf("%d\n", items);
    sem_post(&sync);
    if(items == 0)
    sem_wait(&delay_consumer);
  }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer with unbounded buffer

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One Consumer, One Producer, Unbounded Buffer: First Attempt

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
   {
      sem_wait(&sync);
      items--;
      printf("%d\n", items);
      sem_post(&sync);
      if(items == 0)
        sem_wait(&delay_consumer);
   }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer with unbounded buffer

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One Consumer, One Producer, Unbounded Buffer: First Attempt

```
void * consumer(void * p)
                                          void * producer(void * p)
 sem wait(&delay consumer); (wakeup)
 while (1)
                                           while (1)
  sem wait (&sync);
                                            sem wait (&sync);
  items--:
                                            items++:
  printf("%d\n", items);
                                            printf("%d\n", items);
                                            if(items == 1)
  sem post (&sync);
  if(items == 0)
                                             sem post(&delay consumer); -1 => 0
   sem wait (&delay consumer);
                                            sem post (&sync);
```

Figure: Single producer/consumer with unbounded buffer

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One Consumer, One Producer, Unbounded Buffer: First Attempt

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
   {
      sem_wait(&sync);
      items--;
      printf("%d\n", items);
      sem_post(&sync);
      if(items == 0)
        sem_wait(&delay_consumer);
   }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync); 0 => 1
  }
}
```

Figure: Single producer/consumer with unbounded buffer

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One Consumer, One Producer, Unbounded Buffer: First Attempt

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
{
    sem_wait(&sync); 1 => 0
    items--;
    printf("%d\n", items);
    sem_post(&sync);
    if(items == 0)
        sem_wait(&delay_consumer);
   }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer with unbounded buffer

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One Consumer, One Producer, Unbounded Buffer: First Attempt

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
   {
      sem_wait(&sync);
      items--; 1 => 0
      printf("%d\n", items);
      sem_post(&sync);
      if(items == 0)
      sem_wait(&delay_consumer);
   }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer with unbounded buffer

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One Consumer, One Producer, Unbounded Buffer: First Attempt

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
   {
      sem_wait(&sync);
      items--;
      printf("%d\n", items);
      sem_post(&sync);
      if(items == 0)
        sem_wait(&delay_consumer);
   }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer with unbounded buffer

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One Consumer, One Producer, Unbounded Buffer: First Attempt

```
void * consumer(void * p)
{
    sem_wait(&delay_consumer);
    while(1)
    {
        sem_wait(&sync);
        items--;
        printf("%d\n", items);
        sem_post(&sync); 0 => 1
        if(items == 0)
        sem_wait(&delay_consumer);
    }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer with unbounded buffer

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One Consumer, One Producer, Unbounded Buffer: First Attempt

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
   {
      sem_wait(&sync);
      items--;
      printf("%d\n", items);
      sem_post(&sync);
      if(items == 0)
         sem_wait(&delay_consumer);
   }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer with unbounded buffer

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One Consumer, One Producer, Unbounded Buffer: First Attempt

```
void * consumer(void * p)
                                          void * producer(void * p)
 sem wait (&delay consumer);
 while (1)
                                           while (1)
  sem wait (&sync);
                                            sem wait (&svnc);
  items--;
                                            items++:
  printf("%d\n", items);
                                            printf("%d\n", items);
                                            if(items == 1)
  sem post (&sync);
  if(items == 0)
                                             sem post(&delay consumer);
   sem_wait(&delay_consumer); 0 => -1
                                            sem post(&sync);
```

Figure: Single producer/consumer with unbounded buffer

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One Consumer, One Producer, Unbounded Buffer: First Attempt

```
void * consumer(void * p)
{
    sem_wait(&delay_consumer);
    while(1)
    {
        sem_wait(&sync);
        items--;
        printf("%d\n", items);
        sem_post(&sync);
        if(items == 0)
        sem_wait(&delay_consumer);
    }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync); 1 => 0
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer with unbounded buffer

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One Consumer, One Producer, Unbounded Buffer: First Attempt

```
void * consumer(void * p)
{
    sem_wait(&delay_consumer);
    while(1)
    {
        sem_wait(&sync);
        items--;
        printf("%d\n", items);
        sem_post(&sync);
        if(items == 0)
        sem_wait(&delay_consumer);
    }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++; 0 => 1
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer with unbounded buffer

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One Consumer, One Producer, Unbounded Buffer: First Attempt

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
   {
      sem_wait(&sync);
      items--;
      printf("%d\n", items);
      sem_post(&sync);
      if(items == 0)
        sem_wait(&delay_consumer);
   }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer with unbounded buffer

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One Consumer, One Producer, Unbounded Buffer: First Attempt

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
   {
      sem_wait(&sync);
      items--;
      printf("%d\n", items);
      sem_post(&sync);
      if(items == 0)
         sem_wait(&delay_consumer);
   }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer with unbounded buffer

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One Consumer, One Producer, Unbounded Buffer: First Attempt

```
void * consumer(void * p)
                                          void * producer(void * p)
 sem wait (&delay consumer);
 while (1)
                                           while (1)
  sem wait (&sync);
                                            sem wait (&svnc);
  items--;
                                            items++:
  printf("%d\n", items);
                                            printf("%d\n", items);
                                            if(items == 1)
  sem post (&sync);
  if(items == 0)
                                             sem post(&delay consumer); -1 => 0
   sem_wait(&delay_consumer); (wakeup)
                                            sem post(&sync);
```

Figure: Single producer/consumer with unbounded buffer

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One Consumer, One Producer, Unbounded Buffer: First Attempt

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
   {
      sem_wait(&sync);
      items--;
      printf("%d\n", items);
      sem_post(&sync);
      if(items == 0)
         sem_wait(&delay_consumer);
   }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync); 0 => 1
  }
}
```

Figure: Single producer/consumer with unbounded buffer

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One Consumer, One Producer, Unbounded Buffer

- It is obvious that any manipulations of items will have to be synchronised
- Race conditions still exist:
 - When the consumer has exhausted the buffer, should have blocked, but the producer increments items before the consumer checks it

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * producer(void * p)
void * consumer(void * p)
sem wait(&delay consumer); 0 => -1
while(1)
                                          while(1)
 sem wait(&sync);
                                           sem wait(&sync);
 items--;
                                           items++;
 printf("%d\n", items);
                                           printf("%d\n", items);
 sem post(&sync);
                                           if(items == 1)
  if(items == 0)
                                            sem post(&delay consumer);
  sem wait(&delay consumer);
                                           sem post(&sync);
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
  sem_wait(&delay_consumer);
  while(1)
  {
    sem_wait(&sync);
    items--;
    printf("%d\n", items);
    sem_post(&sync);
    if(items == 0)
        sem_wait(&delay_consumer);
  }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync); 1 => 0
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
  sem_wait(&delay_consumer);
  while(1)
  {
    sem_wait(&sync);
    items--;
    printf("%d\n", items);
    sem_post(&sync);
    if(items == 0)
        sem_wait(&delay_consumer);
  }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++; 0 => 1
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
  sem_wait(&delay_consumer);
  while(1)
  {
    sem_wait(&sync);
    items--;
    printf("%d\n", items);
    sem_post(&sync);
    if(items == 0)
        sem_wait(&delay_consumer);
  }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
   {
      sem_wait(&sync);
      items--;
      printf("%d\n", items);
      sem_post(&sync);
      if(items == 0)
         sem_wait(&delay_consumer);
   }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
    }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
    sem_wait(&delay_consumer); (wakeup)
    while(1)
    {
        sem_wait(&sync);
        items--;
        printf("%d\n", items);
        sem_post(&sync);
        if(items == 0)
            sem_wait(&delay_consumer);
    }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
    sem_post(&delay_consumer); -1 => 0
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
   {
      sem_wait(&sync);
      items--;
      printf("%d\n", items);
      sem_post(&sync);
      if(items == 0)
         sem_wait(&delay_consumer);
   }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync); 0 => 1
  }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
{
    sem_wait(&sync); 1 => 0
    items--;
    printf("%d\n", items);
    sem_post(&sync);
    if(items == 0)
     sem_wait(&delay_consumer);
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
  sem_wait(&delay_consumer);
  while(1)
  {
    sem_wait(&sync);
    items--; 1 => 0
    printf("%d\n", items);
    sem_post(&sync);
    if(items == 0)
    sem_wait(&delay_consumer);
  }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
   {
      sem_wait(&sync);
      items--;
      printf("%d\n", items);
      sem_post(&sync);
      if(items == 0)
         sem_wait(&delay_consumer);
   }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
   {
      sem_wait(&sync);
      items--;
      printf("%d\n", items);
      sem_post(&sync); 0 => 1
      if(items == 0)
         sem_wait(&delay_consumer);
   }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
    }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
  sem_wait(&delay_consumer);
  while(1)
  {
    sem_wait(&sync);
    items--;
    printf("%d\n", items);
    sem_post(&sync);
    if(items == 0)
        sem_wait(&delay_consumer);
  }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync); 1 => 0
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
   {
      sem_wait(&sync);
      items--;
      printf("%d\n", items);
      sem_post(&sync);
      if(items == 0)
      sem_wait(&delay_consumer);
   }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++; 0 => 1
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
   {
      sem_wait(&sync);
      items--;
      printf("%d\n", items);
      sem_post(&sync);
      if(items == 0)
         sem_wait(&delay_consumer);
   }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
  sem_wait(&delay_consumer);
  while(1)
  {
    sem_wait(&sync);
    items--;
    printf("%d\n", items);
    sem_post(&sync);
    if(items == 0)
        sem_wait(&delay_consumer);
  }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
   {
      sem_wait(&sync);
      items--;
      printf("%d\n", items);
      sem_post(&sync);
      if(items == 0)
         sem_wait(&delay_consumer);
   }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
    sem_post(&delay_consumer); 0 => 1
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
   {
      sem_wait(&sync);
      items--;
      printf("%d\n", items);
      sem_post(&sync);
      if(items == 0)
         sem_wait(&delay_consumer);
   }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync); 0 => 1
  }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
  sem_wait(&delay_consumer);
  while(1)
  {
    sem_wait(&sync);
    items--;
    printf("%d\n", items);
    sem_post(&sync);
    if(items == 0)
    sem_wait(&delay_consumer);
  }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
{
    sem_wait(&sync); 1 => 0
    items--;
    printf("%d\n", items);
    sem_post(&sync);
    if(items == 0)
     sem_wait(&delay_consumer);
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
    }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
   {
      sem_wait(&sync);
      items--; 1 => 0
      printf("%d\n", items);
      sem_post(&sync);
      if(items == 0)
         sem_wait(&delay_consumer);
   }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
    }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
  sem_wait(&delay_consumer);
  while(1)
  {
    sem_wait(&sync);
    items--;
    printf("%d\n", items);
    sem_post(&sync);
    if(items == 0)
        sem_wait(&delay_consumer);
  }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
    }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
  sem_wait(&delay_consumer);
  while(1)
  {
    sem_wait(&sync);
    items--;
    printf("%d\n", items);
    sem_post(&sync); 0 => 1
    if(items == 0)
        sem_wait(&delay_consumer);
  }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
  sem_wait(&delay_consumer);
  while(1)
  {
    sem_wait(&sync);
    items--;
    printf("%d\n", items);
    sem_post(&sync);
    if(items == 0)
    sem_wait(&delay_consumer);
  }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
                                         void * producer(void * p)
sem wait(&delay consumer);
while(1)
                                          while(1)
 sem wait(&sync);
                                           sem wait(&sync);
 items--;
                                           items++;
 printf("%d\n", items);
                                           printf("%d\n", items);
                                           if(items == 1)
 sem post(&sync);
  if(items == 0)
                                            sem post(&delay consumer);
   sem wait(&delay consumer); 1 => 0
                                           sem post(&sync);
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
{
    sem_wait(&sync); 1 => 0
    items--;
    printf("%d\n", items);
    sem_post(&sync);
    if(items == 0)
    sem_wait(&delay_consumer);
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
    }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
  sem_wait(&delay_consumer);
  while(1)
  {
    sem_wait(&sync);
    items--; 0 => -1
    printf("%d\n", items);
    sem_post(&sync);
    if(items == 0)
    sem_wait(&delay_consumer);
  }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
   sem_wait(&delay_consumer);
   while(1)
   {
      sem_wait(&sync);
      items--;
      printf("%d\n", items);
      sem_post(&sync);
      if(items == 0)
         sem_wait(&delay_consumer);
   }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
  sem_wait(&delay_consumer);
  while(1)
  {
    sem_wait(&sync);
    items--;
    printf("%d\n", items);
    sem_post(&sync); 0 => 1
    if(items == 0)
    sem_wait(&delay_consumer);
  }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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One Consumer, One Producer, Unbounded Buffer: Non-Existing Items

```
void * consumer(void * p)
{
  sem_wait(&delay_consumer);
  while(1)
{
    sem_wait(&sync);
    items--;
    printf("%d\n", items);
    sem_post(&sync);
    if(items == 0)
        sem_wait(&delay_consumer);
    }
}
```

```
void * producer(void * p)
{
  while(1)
  {
    sem_wait(&sync);
    items++;
    printf("%d\n", items);
    if(items == 1)
        sem_post(&delay_consumer);
    sem_post(&sync);
  }
}
```

Figure: Single producer/consumer and an unbounded buffer: Race condition (non-existing element => items = -1)

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Test your understanding

- Is a binary semaphore the same thing as a mutex?
- When should you prefer a mutex rather than a binary semaphore?
- Is there a straightforward way to check concurrent code is correct?

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