

AE1MCS: Mathematics for Computer Scientists

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Reading

Kenneth H. Rosen, *Discrete Mathematics and Its Applications*, 7th Edition, 2013.

- Chapter 9, Section 9.1 Relations and Their Properties
- Chapter 9, Section 9.2 ~~Database and Relations~~
- Chapter 9, Section 9.5 Equivalence Relations
- Chapter 9, Section 9.6 Partial Orderings

Relations

- Relationships between elements of sets occur in many contexts.
- Relationships between elements of sets are represented using the structure called a relation, which is just a subset of the Cartesian product of the sets.

$$\underline{A \times B = \{(a,b) \mid a \in A, b \in B\}}$$

Binary Relations

The most direct way to express a relationship between elements of two sets is to use ordered pairs made up of two related elements. Hence, sets of ordered pairs are called binary relations.

Definition

Let A and B be sets. A *binary relation* from A to B is a subset of $A \times B$.

We use $a R b$ or $R(a, b)$ to denote that $(a, b) \in R$.

$a R b$

Exercise

$$A \times A = \left\{ \begin{array}{l} (\underline{1}, \underline{1}), (\underline{1}, \underline{2}), (\underline{1}, \underline{3}), (\underline{1}, \underline{4}), \\ (\underline{2}, 1), (\underline{2}, \underline{2}), (\underline{2}, \underline{3}), (\underline{2}, \underline{4}), \\ (\underline{3}, 1), \quad \cdot \quad \cdot \quad \cdot \quad (\underline{3}, \underline{4}), \\ (\underline{4}, 1), \quad \cdot \quad \cdot \quad \cdot \quad (\underline{4}, \underline{4}) \end{array} \right\}$$

$$|A| = n$$

$$|A \times A| = n^2$$

2^{n^2} Relations

Let A be the set $\{1, 2, 3, 4\}$. Which ordered pairs are in the relation

$$R = \{(a, b) \mid a \text{ divides } b\}$$

$$R = \left\{ (1, 1), (1, 2), (1, 3), (1, 4), \overbrace{\qquad\qquad\qquad}^{\text{R}_1} \right.$$

$$\qquad\qquad\qquad (2, 2), \qquad\qquad (2, 4),$$

$$\qquad\qquad\qquad (3, 3),$$

$$\qquad\qquad\qquad \left. (4, 4) \right\}$$

$$R_1 = \{(1, 1), (1, 2)\}$$

Exercise

Let A be the set $\{1, 2, 3, 4\}$. Which ordered pairs are in the relation $R = \{(a, b) \mid a \text{ divides } b\}$?

Answer:

$$R = \{(1, 1), (1, 2), (1, 3), (1, 4), (2, 2), (2, 4), (3, 3), (4, 4)\}$$

Relations on a Set

Relations from a set A to itself are of special interest.

Definition

A *relation on a set A* is a relation from A to A .

How many relations are there on a set with n elements?

$$2^{n^2}$$

Reflexive Relations

There are several properties that are used to classify relations on a set.

Definition

A relation R on a set A is called reflexive, if $(a, a) \in R$ for every element $a \in A$.

How to use quantifiers to express it?

$$\forall a \in A, (a, a) \in R$$

$$R_1 = \{(1, 1), (1, 2), (2, 2)\} \quad A = \{1, 2\}$$

Symmetric Relations

$$R = \{(1,1), (1,2), (2,1)\}$$

Definition

A relation R on a set A is called *symmetric*, if $\underline{(b,a) \in R}$ whenever $\underline{(a,b) \in R}$, for all $a, b \in A$.

How to use quantifiers to express it?

$$\forall a \forall b \underbrace{((a,b) \in R \rightarrow (b,a) \in R)}_{\text{False.}}$$

Antisymmetric Relations

$$\forall a \forall b ((a,b) \in R \wedge (b,a) \in R \rightarrow \underline{\underline{(a=b)}}) \quad \begin{matrix} \top \\ F \end{matrix} \quad \begin{matrix} F \\ \top \end{matrix}$$

Definition

A relation R on a set A such that for all $a, b \in A$, if $\underline{(a,b) \in R}$ and $\underline{(b,a) \in R}$, then $\underline{a = b}$ is called *antisymmetric*.

- The terms symmetric and antisymmetric are not opposites.
- A relation can have both of these properties or may lack both of them. Examples?

$$R = \{(1, 1)\}.$$

$$R = \{(2, 2), (3, 3), (4, 4)\}.$$

$$(a, a) \in R.$$

$$R = \{(0, 1), (1, 2), (2, 1)\}$$
$$R = \{(0, 0), (1, 5), (2, 10), (10, 2)\}$$
$$(a, b) \in R, (b, a) \in R,$$
$$(c, d) \notin R, (d, c) \in R, c \neq d$$

Transitive Relations

$\forall a \forall b \forall c, ((a,b) \in R \wedge (b,c) \in R \rightarrow (a,c) \in R).$

Definition

A relation R on a set A is called *transitive*, if whenever $\underline{(a, b) \in R}$ and $\underline{(b, c) \in R}$, then $\underline{(a, c) \in R}$, for all $a, b, c \in A$.

How to use quantifiers to express it?

$$R = \{(a, b) \mid a \mid b\}, \quad a, b \in \mathbb{Z}^+$$

Reflexive. Symmetric. Antisymmetric. transitive.

Y

N

Y

$a \mid b, b \mid a \rightarrow a = b$

Y

$a \mid b, b \mid c \rightarrow a \mid c$

Examples

Consider these relations on the set of integers:

	reflexive	Sym	anti-symm	transitive
$R_1 = \{(a, b) \mid a \leq b\},$	Y	N	Y	Y
$R_2 = \{(a, b) \mid a > b\},$	N	Y	Y	N
$R_3 = \{(a, b) \mid a = b \text{ or } a = -b\},$	Y	Y	N	Y
$R_4 = \{(a, b) \mid a = b\},$	Y	Y	Y	Y
$R_5 = \{(a, b) \mid a = b + 1\},$	N	N	Y	N
$R_6 = \{(a, b) \mid a + b \leq 3\}.$	N	Y	N	N

(2,1) (1,2); (2,2) X

Combining Relations

Because relations from A to B are subsets of $A \times B$, two relations from A to B can be combined in any way two sets can be combined.

$$R_1 \cup R_2$$

$$R_1 \cap R_2$$

$$R_1 - R_2$$

$$R_2 - R_1$$

Exercise

Let R_1 be the 'less than' relation on the set of real numbers and let R_2 be the 'greater than' relation on the set of real numbers, that is,
 $R_1 = \{(x, y) \mid x < y\}$ and $R_2 = \{(x, y) \mid x > y\}$. What are $R_1 \cup R_2$,
 $R_1 \cap R_2$, $R_1 - R_2$, $R_2 - R_1$?

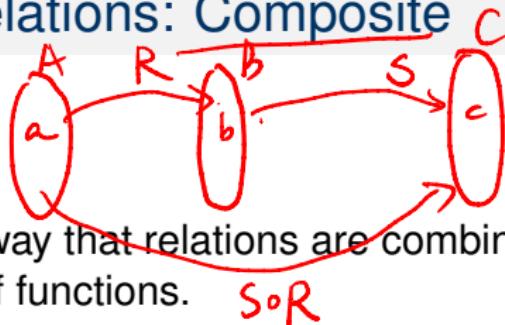
$$R_1 \cup R_2 = \{(x, y) \mid x < y \text{ or } x > y\}.$$

$$R_1 \cap R_2 = \emptyset.$$

$$R_1 - R_2 = R_1$$

$$R_2 - R_1 = R_2$$

Combining Relations: Composite



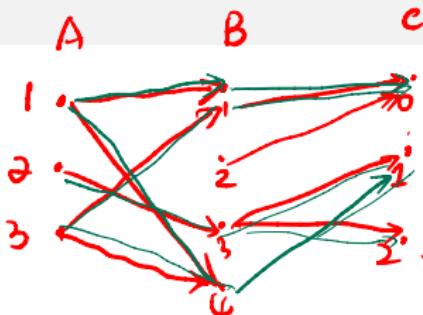
There is another way that relations are combined that is analogous to the composition of functions.

$S \circ R$

Definition

Let R be a relation from a set A to a set B and S a relation from B to a set C . The *composite* of R and S is the relation consisting of ordered pairs (a, c) , where $a \in A$, $c \in C$, and for which there exists an element $b \in B$ such that $(a, b) \in R$ and $(b, c) \in S$. We denote the composite of R and S by $S \circ R$.

Exercise



What is the composite of the relations R and S , where R is the relation from $\{1, 2, 3\}$ to $\{1, 2, 3, 4\}$ with $\underline{\underline{R = \{(1, 1), (1, 4), (2, 3), (3, 1), (3, 4)\}}}$ and S is the relation from $\{1, 2, 3, 4\}$ to $\{0, 1, 2\}$ with $\underline{\underline{S = \{(1, 0), (2, 0), (3, 1), (3, 2), (4, 1)\}}}$?

$$S \circ R = \{(1, 0), (1, 1), (2, 1), (2, 2), (3, 0), (3, 1)\}$$

Composing a Relation with Itself

Definition

Let R be a relation on the set A . The powers R^n , $n = 1, 2, 3, \dots$, are defined recursively by $R^1 = R$ and $R^{n+1} = R^n \circ R$.

$$R^2 = R \circ R$$

$$R^3 = R \circ R \circ R$$

Exercise

Let $R = \{(1, 1), (2, 1), (3, 2), (4, 3)\}$. Find the powers R^n , $n = 2, 3, 4, \dots$.

$$R^2 = R \circ R = \{(1, 1), (2, 1), (3, 1), (4, 2)\}$$

$$R^3 = R^2 \circ R = \{(1, 1), (2, 1), (3, 1), (4, 1)\}$$

$$R^4 = R^3 \circ R = \{(1, 1), (2, 1), (3, 1), (4, 1)\}$$

⋮

$$R^n = R^3 \text{ for } n = 4, 5, 6, \dots$$

A Theorem

Theorem

The relation R on a set A is transitive if and only if $\underline{R^n \subseteq R}$ for $n = 1, 2, 3, \dots$

See Rosen's textbook, p.581

Databases and Relations

Let A_1, A_2, \dots, A_n be sets. An n-ary relation on these sets is a subset of $\underline{A_1 \times A_2 \times \dots \times A_n}$. The sets A_1, A_2, \dots, A_n are called the domains of the relation, and n is called its degree.

A **relational database model** is a model for representing databases using n-ary relations.

Primary Key: a domain of an n-ary relation such that an n-tuple is uniquely determined by its value for this domain.

Composite Key: the Cartesian product of domains of an n-ary relation such that an n-tuple is uniquely determined by its values in these domains

Primary Key and Composite Key

Major X GPA

Student_name	ID_number	Major	GPA
Ackermann	231455	Computer Science	3.88
Adams	888323	Physics	3.45
Chou	102147	Computer Science	3.49
Goodfriend	453876	Mathematics	3.45
Rao	678543	Mathematics	3.90
Stevens	786576	Psychology	2.99

- Which domains are primary keys for the n-ary relation displayed in Table 1, assuming that no n-tuples will be added in the future?
- Is the Cartesian product of the domain of major fields of study and the domain of GPAs a composite key for the n-ary relation from Table 1, assuming that no n-tuples are ever added?

Projections

The projection P_{i_1, i_2, \dots, i_m} where $i_1 < i_2 < \dots < i_m$, maps the n-tuple (a_1, a_2, \dots, a_n) to the m-tuple $(a_{i_1}, a_{i_2}, \dots, a_{i_m})$, where $m \leq n$.

What results when the projection $P_{1,3}$ is applied to the 4-tuples $(2, 3, 0, 4)$, $(\text{Jane Doe}, 234111001, \text{Geography}, 3.14)$, and (a_1, a_2, a_3, a_4) ?

$(2, 0)$. $(\text{Jane Doe}, \text{Geography})$, (a_1, a_3)

More Examples

$a \equiv b \pmod{m}$ iff $m | (a-b)$, a is congruent to b modulo m .

$$2 \equiv 7 \pmod{5} \quad 5 | (2-7)$$

$$7 \equiv 12 \pmod{5} \quad 5 | (7-12)$$

$$2 \equiv -3 \pmod{5} \quad 5 | (2 - (-3))$$

- $A = \mathbb{Z}, xRy$ if $x \equiv y \pmod{5}$ Reflexive, sym, not anti-sym, transitive
- $A = \mathbb{Z}^+$, xRy if $x|y$ Y, N, Y, Y
- $A = \mathbb{N}$, xRy if $x \leq y$ N, Y, F

Equivalence Relations

Equivalence relation: a reflexive, symmetric, and transitive relation.

Two elements a and b that are related by an equivalence relation are called equivalent. The notation $a \sim b$ is often used to denote that a and b are equivalent elements with respect to a particular equivalence relation.

$$\{(a, b) \mid a \equiv b \pmod{5}\}. \quad \underline{2 \sim 7}$$

$$\{(a, b) \mid a = b, \text{ or } a = -b\}$$

Equivalence Classes

$[a]_R$ (equivalence class of a with respect to R): the set of all elements of A that are equivalent to a

$$R = \{ (a, b) \mid a \equiv b \pmod{4} \}$$

$\begin{matrix} -4, 0, 4 \end{matrix}$ $\overbrace{\hspace{1cm}}$ $[a]_R = \{ s \mid (a, s) \in R \}$.

Example

What are the equivalence classes of 0 and 1 for congruence modulo 4?

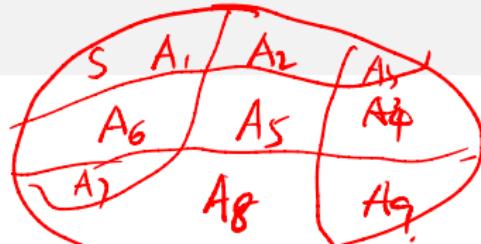
$$[0]_R = \{ \dots, -8, \underline{-4}, \underline{0}, \underline{4}, \underline{8}, 12, 16, \dots \}$$

$$[1]_R = \{ \dots, \underline{-7}, \underline{-3}, \underline{1}, \underline{5}, \underline{9}, 13, 17, \dots \}$$

$[a]_m$ the congruence class of an integer a modulo m .
equivalence.

$$[a]_m = \{ \dots, a-2m, a-m, a, a+m, a+2m, \dots \}$$

Partition



A partition of a set S is a collection of disjoint nonempty subsets of S that have S as their union.

- $A_i \neq \emptyset$ for $i \in I$
- $\underline{A_i \cap A_j = \emptyset}$ when $i \neq j$
- $\underline{\cup_{i \in I} A_i = S}$

What are the sets in the partition of the integers arising from congruence modulo 4?

$$\left\{ \begin{array}{l} [0]_4 \\ [1]_4 \\ [2]_4 = \{ \dots -6, -2, 2, 6, 10, 14, \dots \} \\ [3]_4 = \{ \dots -5, -1, 3, 7, 11, 15, \dots \} \end{array} \right.$$

Partial Orderings

$R = \{(x, y) \mid x \geq y\}$, $x, y \in \mathbb{Z}$. Partial ordering

\geq is a partial ordering on the set of \mathbb{Z} .

$$\begin{array}{c} (\mathbb{Z}, \geq) \\ \hline (\mathbb{Z}, \leq) \\ \hline (\mathbb{Z}, >) \end{array}$$

Partial ordering: a relation that is reflexive, antisymmetric, and transitive

A set S together with a partial ordering R is called a partially ordered set, or poset, and is denoted by (S, R) . Members of S are called elements of the poset.

$$\underline{(\mathbb{Z}^+, \mid)}$$

$$(\underline{\wp(S)}, \subseteq)$$

$$\underline{\{(A, B) \mid A \subseteq B\}}, A, B \in \wp(S)$$

Partial Orderings

In different posets different symbols such as \leq , \sqsubseteq , and $|$, are used for a partial ordering.

- \leq
- $|$
- \sqsubseteq

Comparable V.S. Incomparable

$$\underline{A = \{1, 2\}}, \quad \underline{B = \{1, 3\}}$$

$$(A, B) \notin R'$$

Comparable: the elements a and b in the poset (A, \preceq) are comparable if $a \preceq b$ or $b \preceq a$

Incomparable: When a and b are elements of S such that neither $a \preceq b$ nor $b \preceq a$, a and b are called incomparable.

- $(P(\mathbb{Z}), \subseteq)$

- $(\mathbb{Z}^+, |)$

$$a=3, \quad b=9 \quad a|b = \frac{b}{a} = \frac{9}{3} = 3.$$
$$(a, b) \in R = \{(x, y) \mid x|y\}, \quad a \leq b, \quad b \not\leq a$$

$$c=5, \quad d=7.$$

$$(c, d) \notin R$$

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