Operating Systems and Concurrency

Concurrency 1 COMP2007

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Concurrency

Context

- Threads and processes execute concurrently or in parallel and can share resources (e.g., devices, memory – variables and data structures)
 - Multi-programming/multi-processing improves system utilisation
- A thread can be interrupted at any point in time (timer, I/O)
 - The process state is saved in the process control block
- The outcome of programs may become unpredictable:
 - Sharing data can lead to inconsistencies we can be interrupted part way through doing something.
 - The outcome of execution may depend on the order in which code gets to run on the CPU.

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Incrementing a counter

```
#include <stdio.h>
#include <pthread.h>
int counter = 0;
void* calc(void* param) {
  int const iterations = 500000000:
  for(int i = 0; i < iterations; i++)</pre>
    counter++;
  return 0;
int main() {
 pthread t tid1 = 0.tid2 = 0:
 pthread_create(&tid1, NULL, calc, 0);
  pthread create (&tid2, NULL, calc, 0);
  pthread_join(tid1, NULL);
  pthread join(tid2, NULL);
 printf("The value of counter is: %d\n", counter);
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```

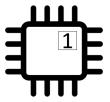
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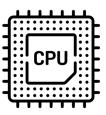
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 printf("The value of counter is: %d\n", counter);
```

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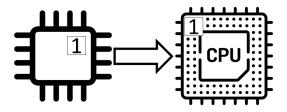
- counter++ consists of three separate actions:
 - read the value of counter from memory and store it in a register
 - 2 add one to the value in the register
 - 3 store the value of the register in counter in memory
- The above actions are NOT "atomic"





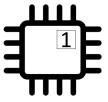
¹ Icons from https://www.flaticon.com/ @University of Nottingham

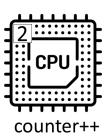
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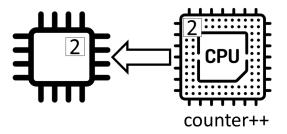
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 - read the value of counter from memory and store it in a register
 - 2 add one to the value in the register
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- The above actions are NOT "atomic"



¹ cons from https://www.flaticon.com/ ©University of Nottingham

Incrementing a counter

```
Thread 1: Thread 2: ...

Read counter -> register (= 1) ...

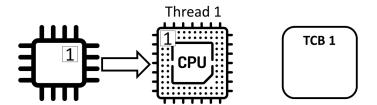
Add 1 to register value (= 2) ...

Store register in counter (= 2) ...

Read counter -> register (= 2) ...

Add 1 to register value (= 3) ...

Store register in counter (= 3)
```



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Incrementing a counter

```
Thread 1: Thread 2: ...

Read counter -> register (= 1) ...

Add 1 to register value (= 2) ...

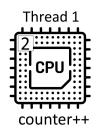
Store register in counter (= 2) ...

Read counter -> register (= 2) ...

Add 1 to register value (= 3) ...

Store register in counter (= 3)
```







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Incrementing a counter

```
Thread 1: Thread 2: ...

Read counter -> register (= 1) ...

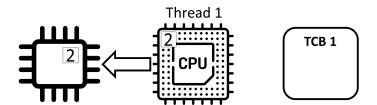
Add 1 to register value (= 2) ...

Store register in counter (= 2) ...

Read counter -> register (= 2) ...

Add 1 to register value (= 3) ...

Store register in counter (= 3)
```



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Incrementing a counter

```
Thread 1: Thread 2: ...

Read counter -> register (= 1) ...

Add 1 to register value (= 2) ...

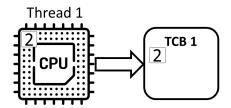
Store register in counter (= 2) ...

Read counter -> register (= 2) ...

Add 1 to register value (= 3) ...

Store register in counter (= 3)
```





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Incrementing a counter

```
Thread 1: Thread 2: ...

Read counter -> register (= 1) ...

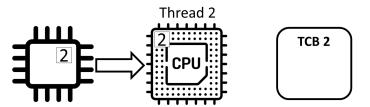
Add 1 to register value (= 2) ...

Store register in counter (= 2) ...

Read counter -> register (= 2) ...

Add 1 to register value (= 3) ...

Store register in counter (= 3)
```



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Incrementing a counter

```
Thread 1: Thread 2: ...

Read counter -> register (= 1) ...

Add 1 to register value (= 2) ...

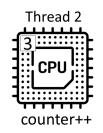
Store register in counter (= 2) ...

Read counter -> register (= 2) ...

Add 1 to register value (= 3) ...

Store register in counter (= 3)
```







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Incrementing a counter

```
Thread 1:
...

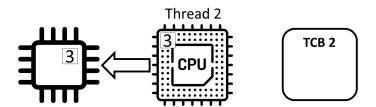
Read counter -> register (= 1)
Add 1 to register value (= 2)
...

Store register in counter (= 2)
...

Read counter -> register (= 2)
...

Add 1 to register value (= 3)
...

Store register in counter (= 3)
```



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Incrementing a counter

```
Thread 1: Thread 2: ...

Read counter -> register (= 1) ...

Add 1 to register value (= 2) ...

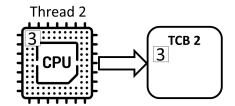
Store register in counter (= 2) ...

Read counter -> register (= 2) ...

Add 1 to register value (= 3) ...

Store register in counter (= 3)
```





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Incrementing a counter

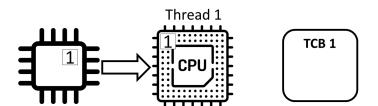
```
Thread 1:
...

Read counter -> register (= 1)
...

Add 1 register value (= 2)
Store value in counter (= 2)
...

Add 1 to register value (= 2)
...

Add 1 to register value (= 2)
Store register in counter (= 2)
```



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Incrementing a counter

```
Thread 1:
...

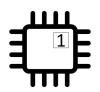
Read counter -> register (= 1)
...

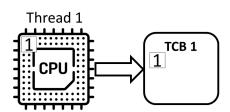
Add 1 register value (= 2)
Store value in counter (= 2)
...

Add 1 to register value (= 2)
...

Add 1 to register value (= 2)
...

Store register in counter (= 2)
```





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Incrementing a counter

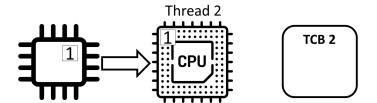
```
Thread 1:
...

Read counter -> register (= 1)
...

Add 1 register value (= 2)
Store value in counter (= 2)
...

Add 1 to register value (= 2)
...

Add 1 to register value (= 2)
Store register in counter (= 2)
```



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Incrementing a counter

```
Thread 1:
...

Read counter -> register (= 1)
...

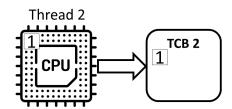
Add 1 register value (= 2)
Store value in counter (= 2)
...

Add 1 to register value (= 2)
...

Add 1 to register value (= 2)
...

Store register in counter (= 2)
```





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Incrementing a counter

```
Thread 1:
...

Read counter -> register (= 1)
...

Add 1 register value (= 2)

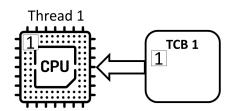
Store value in counter (= 2)
...

Add 1 to register value (= 2)
...

Add 1 to register value (= 2)
...

Store register in counter (= 2)
```





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Incrementing a counter

```
Thread 1:
...

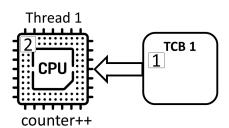
Read counter -> register (= 1)
...

Add 1 register value (= 2)
Store value in counter (= 2)
...

Add 1 to register value (= 2)
...

Add 1 to register value (= 2)
Store register in counter (= 2)
```





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Incrementing a counter

```
Thread 1:
...

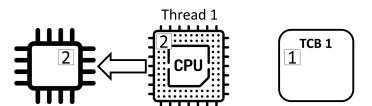
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...

Add 1 to register value (= 2)
...

Add 1 to register value (= 2)
...

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```



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Incrementing a counter

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Thread 1:
...

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...

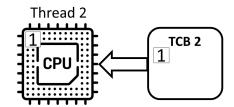
Add 1 register value (= 2)
Store value in counter (= 2)
...

Add 1 to register value (= 2)
...

Add 1 to register value (= 2)
...

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```





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Incrementing a counter

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Thread 1:
...

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...

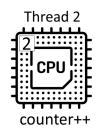
Add 1 register value (= 2)
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...

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...

Add 1 to register value (= 2)
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```







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Incrementing a counter

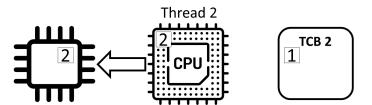
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Thread 1:
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Read counter -> register (= 1)
...

Add 1 register value (= 2)
Store value in counter (= 2)
...

Add 1 to register value (= 2)
...

Add 1 to register value (= 2)
Store register in counter (= 2)
```



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Shared procedures

- Consider the following code shared between threads/processes
- chin and chout shared global variables

```
void print()
{
  chin = getchar();
  chout = chin;
  putchar(chout);
}
```

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Shared procedures

 Consider two processes/threads and the following interleaved sequence of instructions (they do NOT interact):

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Shared procedures

 Consider two processes/threads and the following interleaved sequence of instructions (they DO interact):

```
Thread 1:
...
chin = getchar();
...
chout = chin;
putchar(chout);
...
chout = chin;
putchar(chout);
...
chout = chin;
putchar(chout);
```

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Bounded Buffers

- Consider a bounded buffer in which N items can be stored
- A counter is maintained to count the number of items currently in the buffer
 - Incremented when an item is added
 - Decremented when an item is removed
- Similar concurrency problems as with the calculation of sums happen when multiple threads read and write to the bounded buffer.

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Bounded Buffers – Producer/Consumer

```
// producer
while (true) {
 while (counter == BUFFER SIZE);
 buffer[in] = new item;
 in = (in + 1) % BUFFER_SIZE;
 counter++;
// consumer
while (true) {
 while (counter == 0);
 consumed = buffer[out];
 out = (out + 1) % BUFFER SIZE;
 counter --;
```

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Race Conditions

Definition

- Code has a race condition if its behaviour is dependent on the timing of when computation is performed.
- A race condition typically occurs when multiple threads access shared data and the result is dependent on the order in which the instructions are interleaved.
- We will be interested in mechanisms to provide synchronised access to data and avoid race conditions.

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Concurrency within the OS

Data Structures

- Kernels are preemptive these days
 - Multiple threads are running in the kernel.
 - Kernel code can be interrupted at any point.
- The kernel maintains data structures such as process tables and open file lists:
 - These data structures may be accessed concurrently.
 - These can be subject to concurrency issues.
- The OS must make sure that interactions within the OS do not result in race conditions.

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Critical Sections and Mutual Exclusion

- As race conditions can lead to wild, unpredictable and plain wrong behaviour, we want software abstractions help to prevent them.
- A critical section is a section of code that can only be run by one thread at a time. This property is referred to as mutual exclusion.
- The question then becomes how to enforce mutual exclusion? Potentially:
 - The O/S and compiler could provide direct support for critical sections as a programming primitive.
 - The O/S and compiler provide locks which can be held by at most one thread at a time.

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Mutual Exclusion

Definition

```
do {
    ...
    // ENTRY to critical section
    critical section, e.g.counter++;

// EXIT critical section

// remaining code
    ...
} while (...);
```

Critical Sections, Mutual Exclusion

Definition

- A solution to the critical section problem should satisfy the following requirements:
 - Mutual exclusion: only one process can be in its critical section at any one point in time
 - Progress: any process must be able to enter its critical section at some point in time
 - Processes/threads in the "remaining code" do not influence access to critical sections
 - Fairness/bounded waiting: fairly distributed waiting times/processes cannot be made to wait indefinitely
- These requirements have to be satisfied, independent of the order in which computations are executed

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Enforcing Mutual Exclusion

Approaches

- A standard approach to enforcing mutual exclusion is via locks known as mutexes. These can be implemented in various ways:
 - Software based: Peterson's solution
 - Hardware based: test_and_set(), swap_and_compare()
 - O/S based the operating system blocks processes waiting for the lock.

 Unfortunately, mutexes and other concurrency primitives such as semaphore introduce a new problem - deadlocks.

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Example

- Assume that X and Y are mutexes.
- Thread A and B need to acquire both mutexes, and request them in opposite orders.
- The following sequence of events could occur in a multi-programmed system:

```
THREAD A:

request mutex X

...

...

request mutex X

...

request mutex Y

...

request mutex Y

...

request mutex Y

...

request mutex X
```

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Definition

Tanenbaum

"A set of threads is deadlocked if **each thread** in the set is waiting for **an event** that only the **other thread** in the set can cause"

- Each deadlocked thread is waiting for a resource held by an other deadlocked thread (which cannot run and hence cannot release the resources)
- This can happen between any number of threads and for any number of resources



Minimum Conditions

- Four conditions must hold for deadlocks to occur (Coffman et al. (1971)):
 - Mutual exclusion: a resource can be assigned to at most one process at a time
 - Hold and wait condition: a resource can be held whilst requesting new resources
 - No preemption: resources cannot be forcefully taken away from a process
 - Circular wait: there is a circular chain of two or more processes, waiting for a resource held by the other processes

No deadlocks can occur if one of the conditions is not satisfied

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 - Hold and wait condition: a resource can be held whilst requesting new resources
 - No preemption: resources cannot be forcefully taken away from a process
 - Circular wait: there is a circular chain of two or more processes, waiting for a resource held by the other processes
- No deadlocks can occur if one of the conditions is not satisfied
- If your coursework solution deadlocks, check for the order in which resources are requested

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Test your understanding

- ullet The code x := y doesn't modify anything. Is it certain to occur atomically?
- Can race conditions or deadlocks occur in practice on a machine with a single hardware thread?
- Can two threads running the same function deadlock against each other?

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