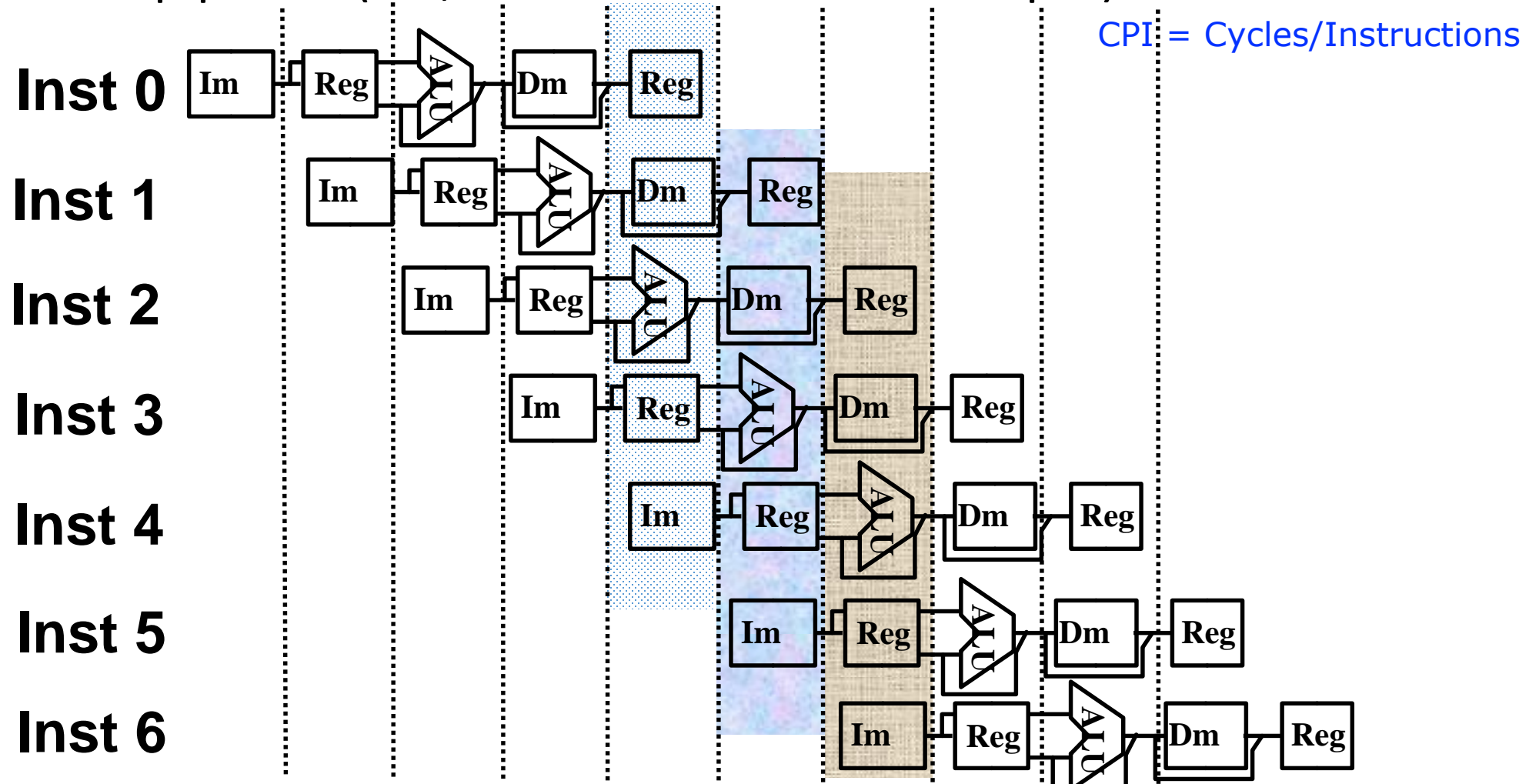


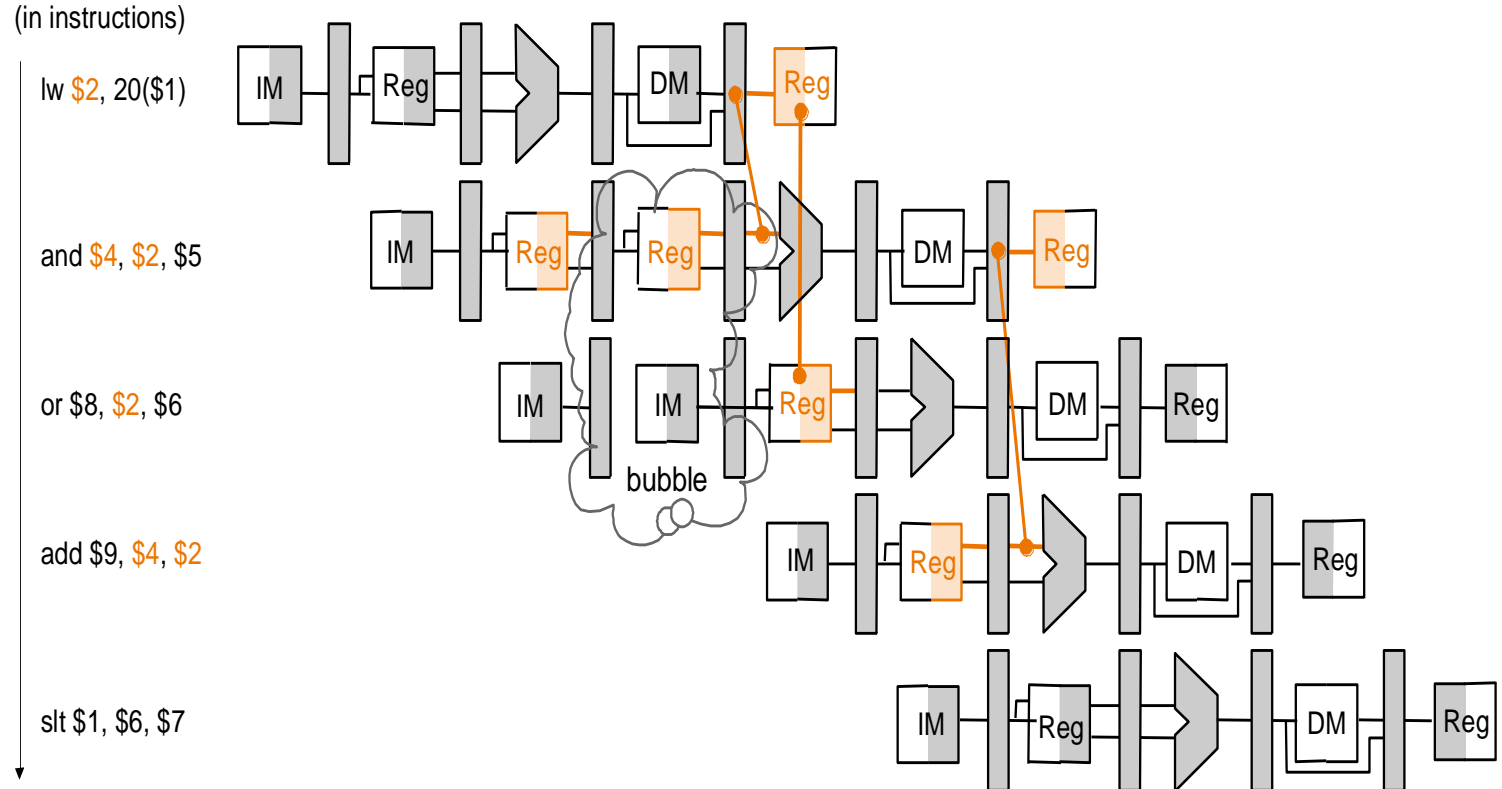
# Understanding the performance of pipeline

- Ideal CPI for pipelined processor:  $CPI = 1$ , since for each cycle, there will be one instruction taken into the pipeline and one instruction taken out of the pipeline (i.e., finished execution and quit)



# Understanding the performance of pipeline

- CPI on **data** hazard (stalling)
- When shall a **stall** happen?
  - The lw instruction and the next one has a **RAW** error. E.g.,
  - lw **\$2**, 20(\$1)  
add \$4, **\$2**, \$5
- Effects on performance?
  - Effectively, 'lw' takes 2 cycles if **RAW** happens
  - 'lw' still takes 1 cycle if no data hazard happens.



# Understanding the performance of pipeline

- CPI on **control/branch** hazard
  - **Without** early branch decision
    - **3** cycle penalty for branch taken, i.e., beq takes **4** cycles
  - **With** early branch decision
    - **1** cycle penalty for branch taken, i.e., beq takes **2** cycles
- Normally when we talk about dealing with control/branch hazard, we use flushing **with** early branch decision.
- **j** (jump) always takes **2** cycles, since branch always taken.