

Probabilistic
Graphical
Models

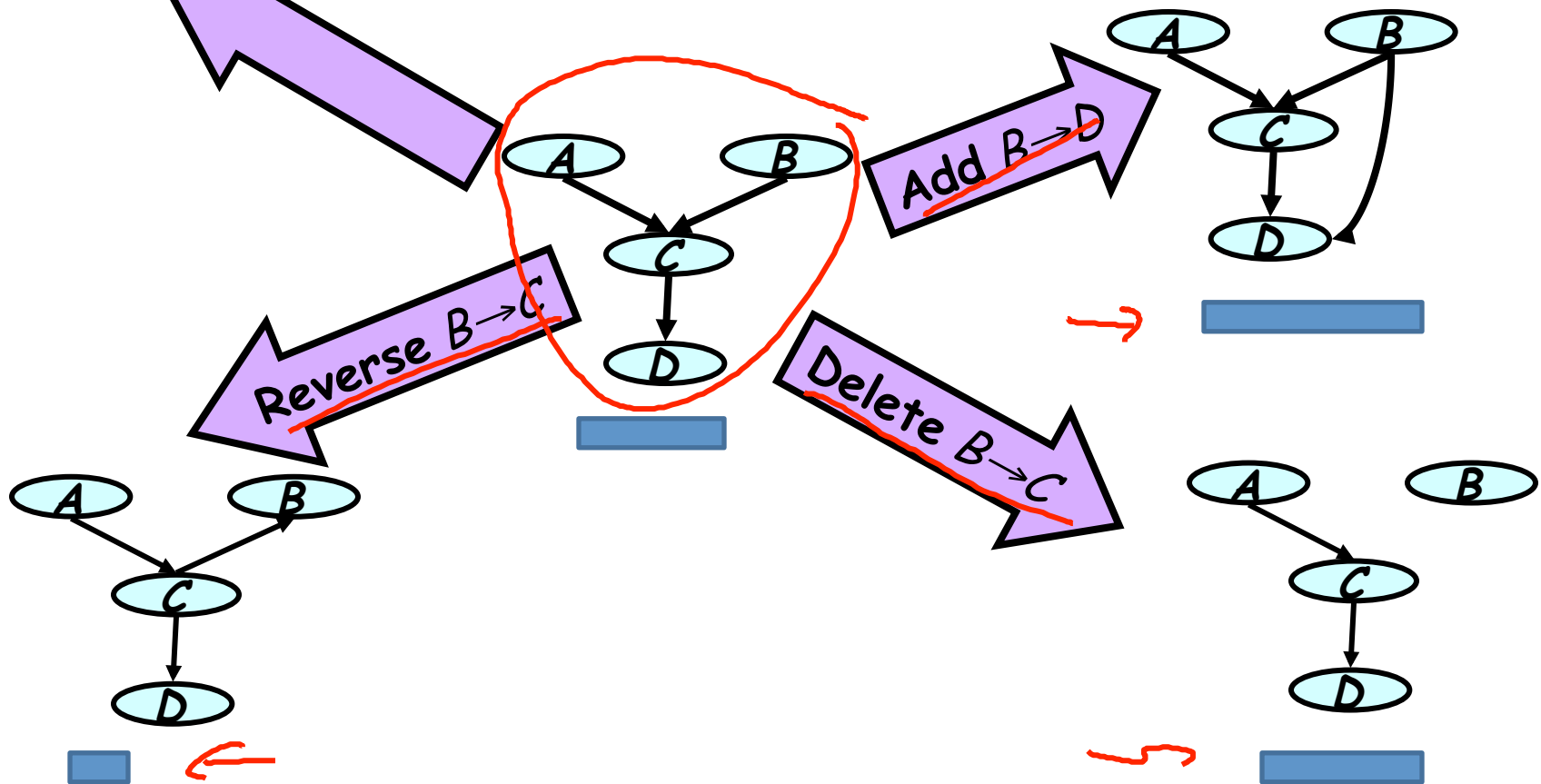


Learning

BN Structure

General Graphs:
Decomposability

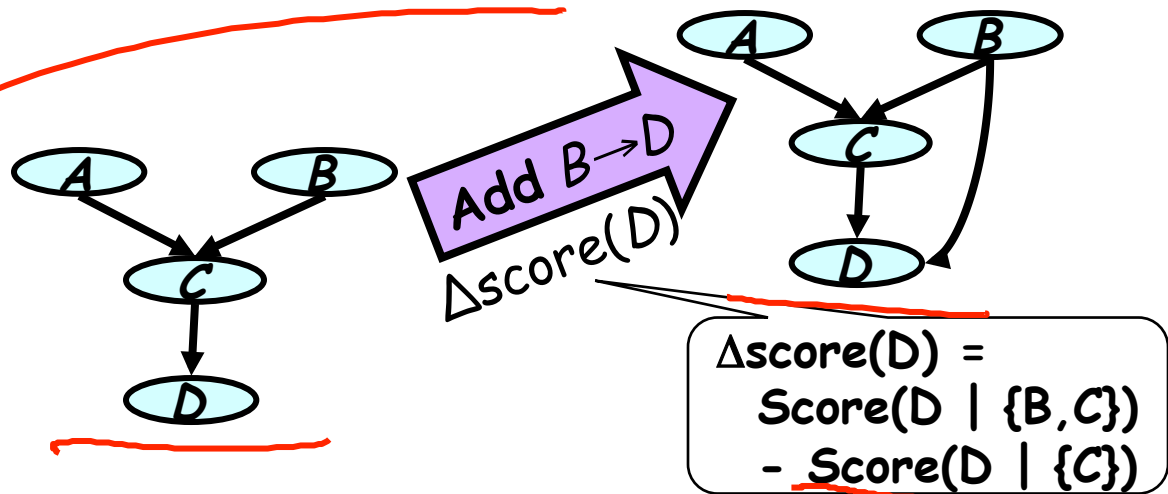
Heuristic Search



Naïve Computational Analysis

- Operators per search step:
 - $n(n-1)$ possible edges $O(n^2)$
 - present $\xrightarrow{2}$ delete $\xrightarrow{1}$ reverse
 - absent $\xrightarrow{1}$ add
- Cost per network evaluation:
 - Components in score n components
 - Compute sufficient statistics $O(m)$
 - Acyclicity check $O(m) \leftarrow \# \text{ edges}$
- Total: $O(\underline{n}^2 (\underline{M}n + \underline{m}))$ per search step

Exploiting Decomposability

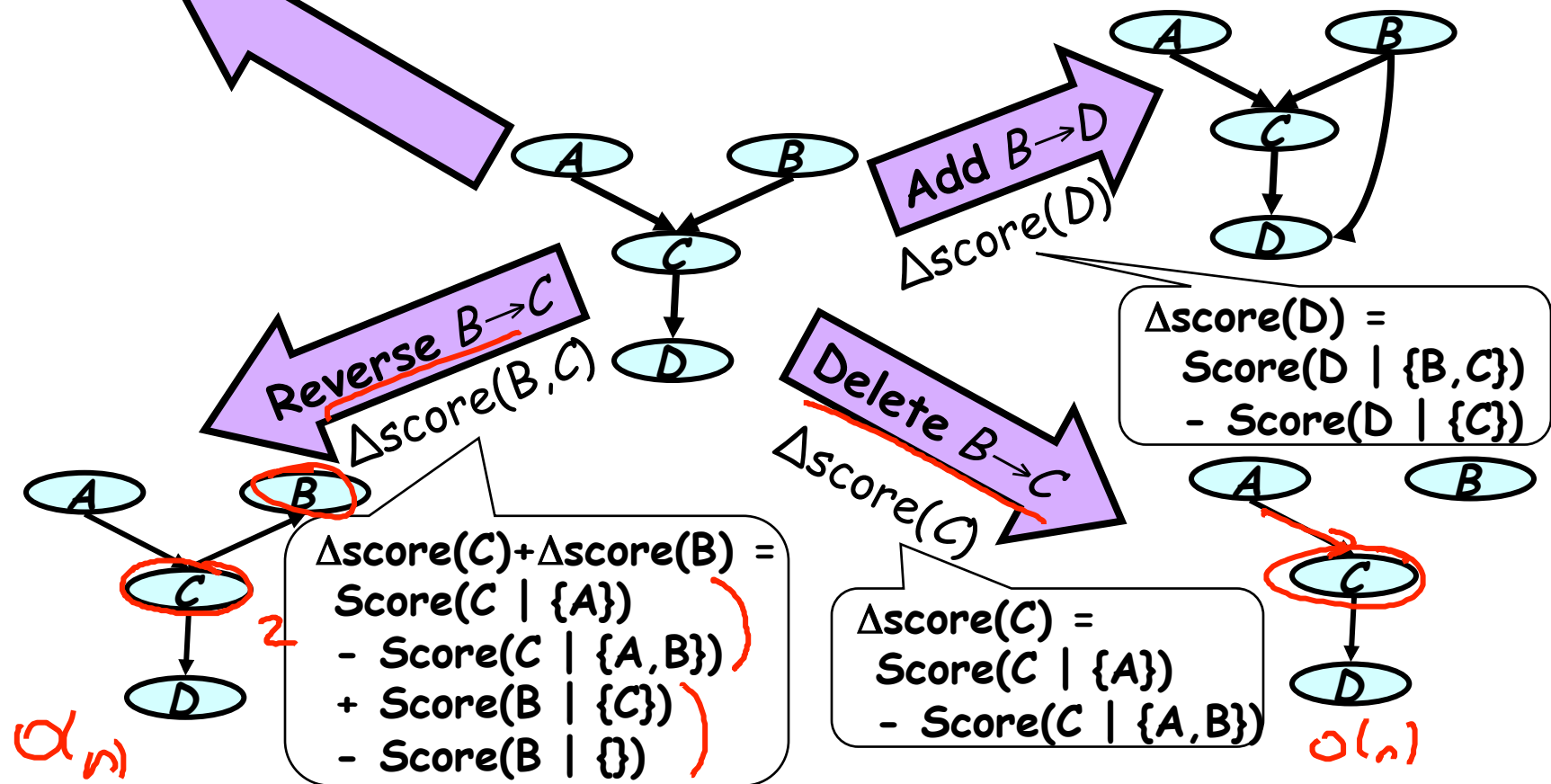


$$\text{score} = \text{Score}(A \mid \{\}) + \text{Score}(B \mid \{\}) + \text{Score}(C \mid \{A, B\}) + \text{Score}(D \mid \{C\})$$

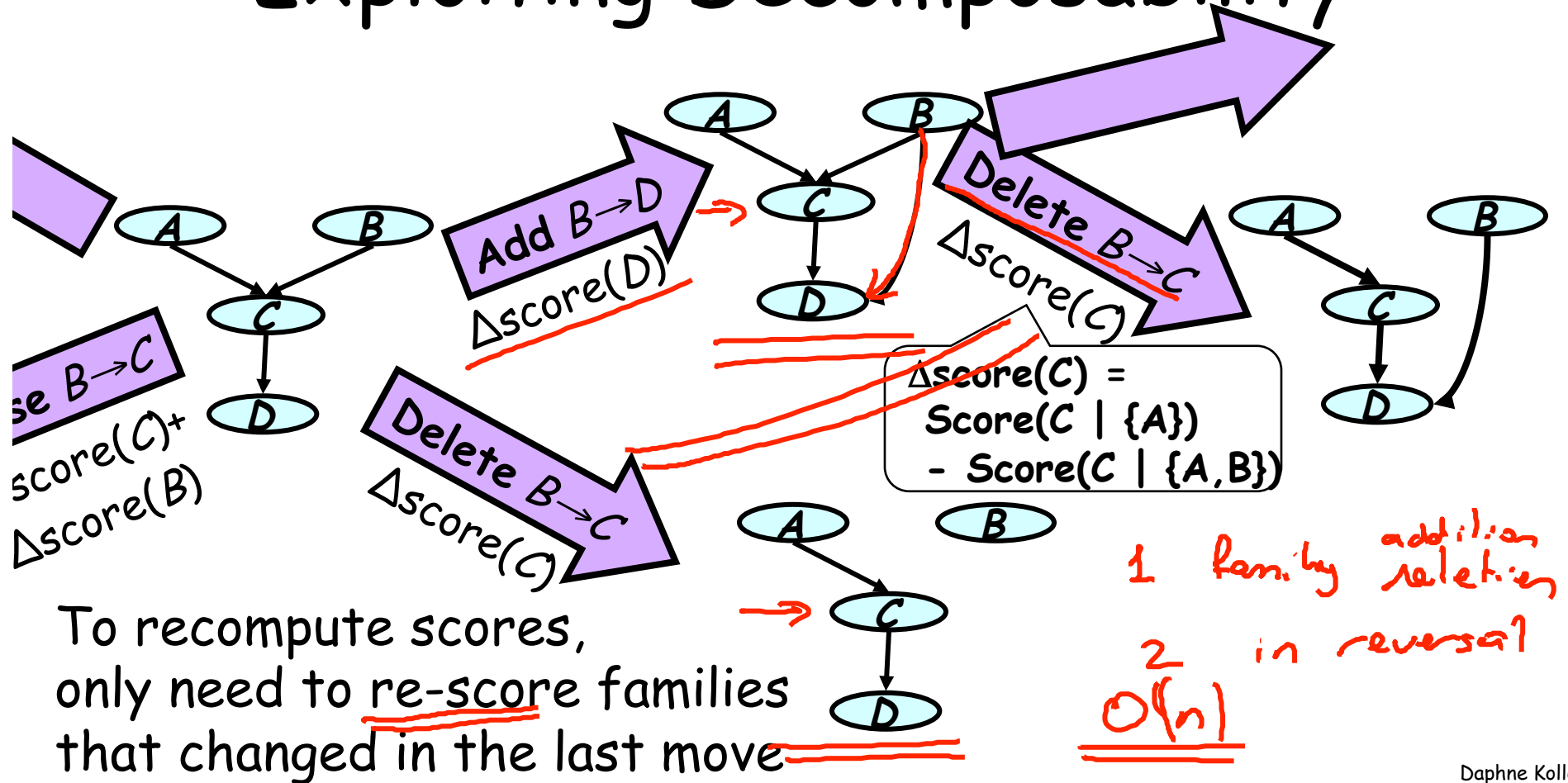
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$O(n)$ savings

Exploiting Decomposability



Exploiting Decomposability



Computational Cost

- Cost per move
 - Compute $O(n)$ delta-scores damaged by move
 - Each one takes $O(M)$ time
- Keep priority queue of operators sorted by delta-score - $O(n \log n)$

More Computational Efficiency

- Reuse and adapt previously computed sufficient statistics
- Restrict in advance the set of operators considered in the search

$$\underbrace{A, B, C}_{\text{Sufficient Statistics}} \quad m[A, B, C]$$
$$m[A, B] = \sum_c m[A, B, c]$$

Summary

- Even heuristic structure search can get expensive for large n
- Can exploit decomposability to get orders of magnitude reduction in cost
- Other tricks are also used for scaling