Divide & Conquer - II

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- Output: Product of x and y

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- Bruteforce method $O(n^2)$

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$$T(n) = 3T(\frac{n}{2}) + n$$

$$= n^{\log_2 3} = n^{1.59}$$

Recursive_Multiply(x,y)

$$x = x_1 2^{\frac{n}{2}} + x_0$$

$$y = y_1 2^{\frac{n}{2}} + y_0$$

$$p = \text{Recursive_Multiply}(x_1 + x_0, y_1 + y_0)$$

$$q = \text{Recursive_Multiply}(x_1, y_1)$$

$$r = \text{Recursive_Multiply}(x_0, y_0)$$

return $q2^{n} + (p - q - r)2^{\frac{n}{2}} + r$

- Input: P, set of n points in the plane
- Output: p,q in P that minimises dist(p,q)

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Points on the real line?

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Brute force?

- Input: P, set of n points in the plane
- Output: p,q in P that minimises dist(p,q)

Divide and Conquer

Solve for n/2 points on the left and n/2 points on the right

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• How is the solution related to the solutions of the sub-problems?

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Divide and Conquer

Solve for n/2 points on the left and n/2 points on the right

• How is the solution related to the solutions of the sub-problems?

 (q_0,r_0) - closest pair of points among n/2 points in the left (q_1,r_1) -closest pair of points among n/2 points in the right d= min (dist (q_0,r_0) , dist (q_1,r_1))

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 (q_1, r_1) -closest pair of points among n/2 points in the right

d= min (dist (q_0, r_0) , dist (q_1, r_1))

The solution is either (q_0,r_0) OR (q_1,r_1) OR some (q',r') such that q' is in left and r' is in right and $dist(q',r') \leq d$

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Let L be the line that divides P into two

Let S be the set of points in P that lie within distance d from L

$$q', r' \in S$$

Let S_{y} be S sorted by y co-ordinates

Claim : If there exists s and s' in S such that dist(s, s') < d, then s and s' are within 15 positions of each other in S_y

Closest_Pair(P)

 $P_{\scriptscriptstyle \chi}$ - P sorted by x-co-ordinate

 P_{y} - P sorted by y-co-ordinate

return Closest_Pair_Rec (P_x, P_y)

• Closest_Pair_Rec(P_x , P_y)

If |P|<3 then solve by brute force

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Construct Q_x, Q_y, R_x, R_y
(q_0, q_1) = Closest_Pair_Rec(Q_x, Q_y)
(r_0, r_1) = Closest_Pair_Rec(R_x, R_v)
d= min (dist(q_0, q_1), dist(r_0, r_1))
Let S be the set of points in P that lie within distance d from L
Construct S_{v}
For each point in S_{
m y}, compute distance from next 15 points in S_{
m y}
Find the minimum of all distances - (s, s')
If min < d, return (s, s')</pre>
else
    if \operatorname{dist}(q_0,q_1)=\operatorname{d}, then return (q_0,q_1)
   else return (r_0, r_1)
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