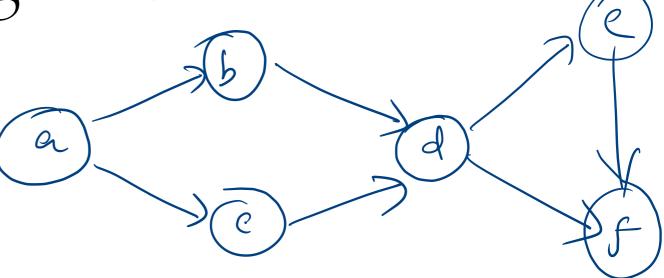
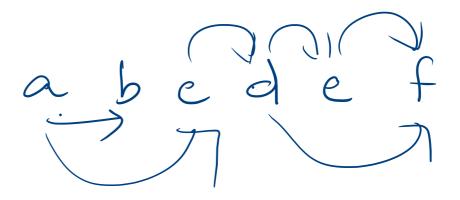
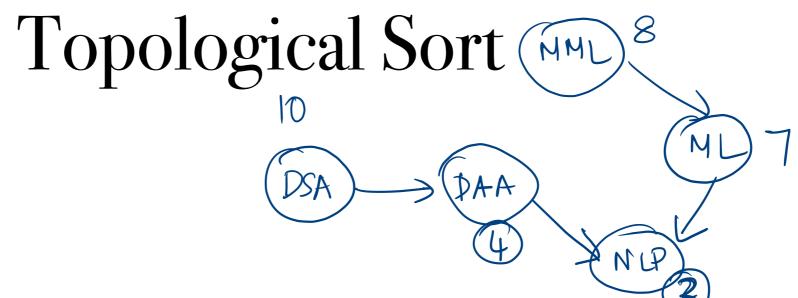
Applications of BFS and DFS

Topological Sort



A topological sort of a DAG G(V,E) is a linear ordering of vertices of G such that if G contains an edge (u,v) then u appears before v in the ordering.





A topological sort of a DAG G(V,E) is a linear ordering of vertices of G such that if G contains an edge (u,v) then u appears before v in the ordering.

- used to show precedence among events

Topological Sort u -> v





- call DFS(G) to compute f[v] for each v
- As each vertex is finished, insert it to the front of a linked list
- return the linked list

Topological Sort

A directed graph G is acyclic if an only if a DFS of G yields no back edges

E Let (u,v) be a back edger. Path from the to the in

OFS tree with (u,v) forms a cycle.

I C (V₁-V₂-···V_k-V₁). Let V₁ be the first vertex in

C to be discovered. At d[V₁], V₂···V_k one all white

I V₁ ~> V_k white path >> V_k is a descendent of V₁.

I (V_k,V₁) is a back edge.

Topological Sort

Algorithm is correct.

If
$$(u,v) \in E$$
, then $f(u) > f(v)$.

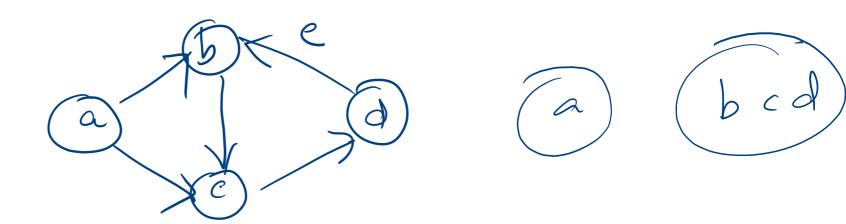
When the edge (u,v) is explosed,

Cane 1: v is white $\Rightarrow v_{-}$ is a descendent of $u \Rightarrow f(u) > f(v)$.

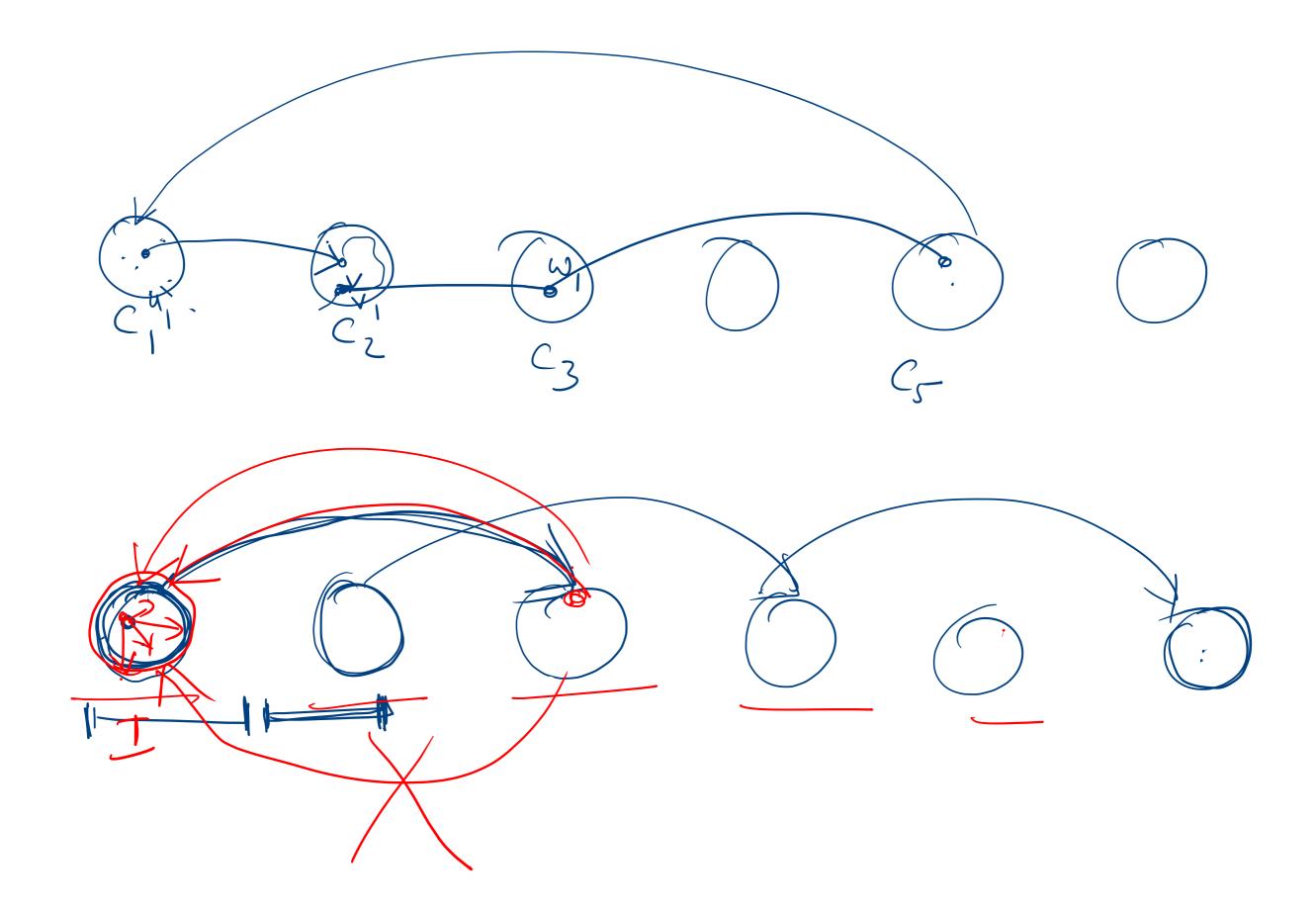
Case 2: 11 group $\Rightarrow (u,v)$ is a back edge \Rightarrow mot possible.

3: 11 black $\Rightarrow f(v) \subset f(u)$

Strongly Connected Components



• Strongly connected component of a directed graph G(V,E) is a maximal set of vertices $C\subseteq V$ such that for every pair of vertices u and v in v, v and v are reachable from each other.

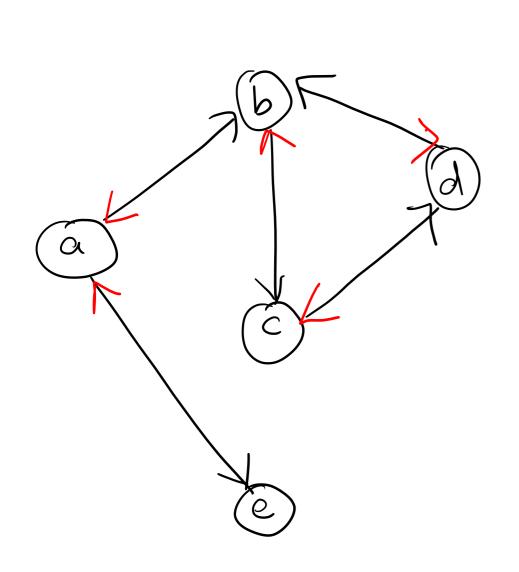


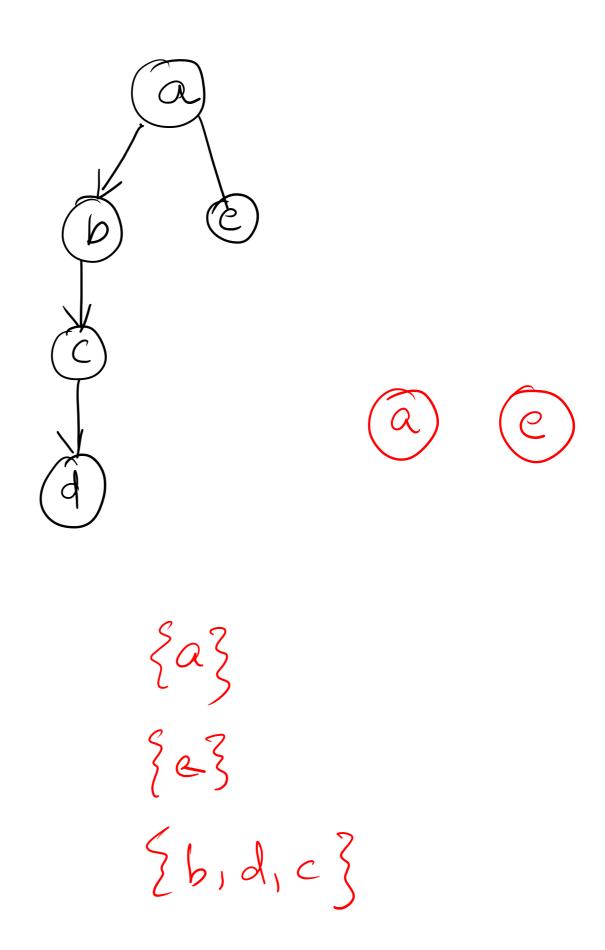
Call DFS(G) to compute f[u] for each vertex u

Compute
$$G^T = (V, E^T)$$
 $E^T = \{(u,v) \mid (v,u) \in E\}$

call DFS(G^T), consider the vertices in order of decreasing f[u]

Output the vertices in each tree in the DFS forest (formed in the previous step) as a separate strongly connected component

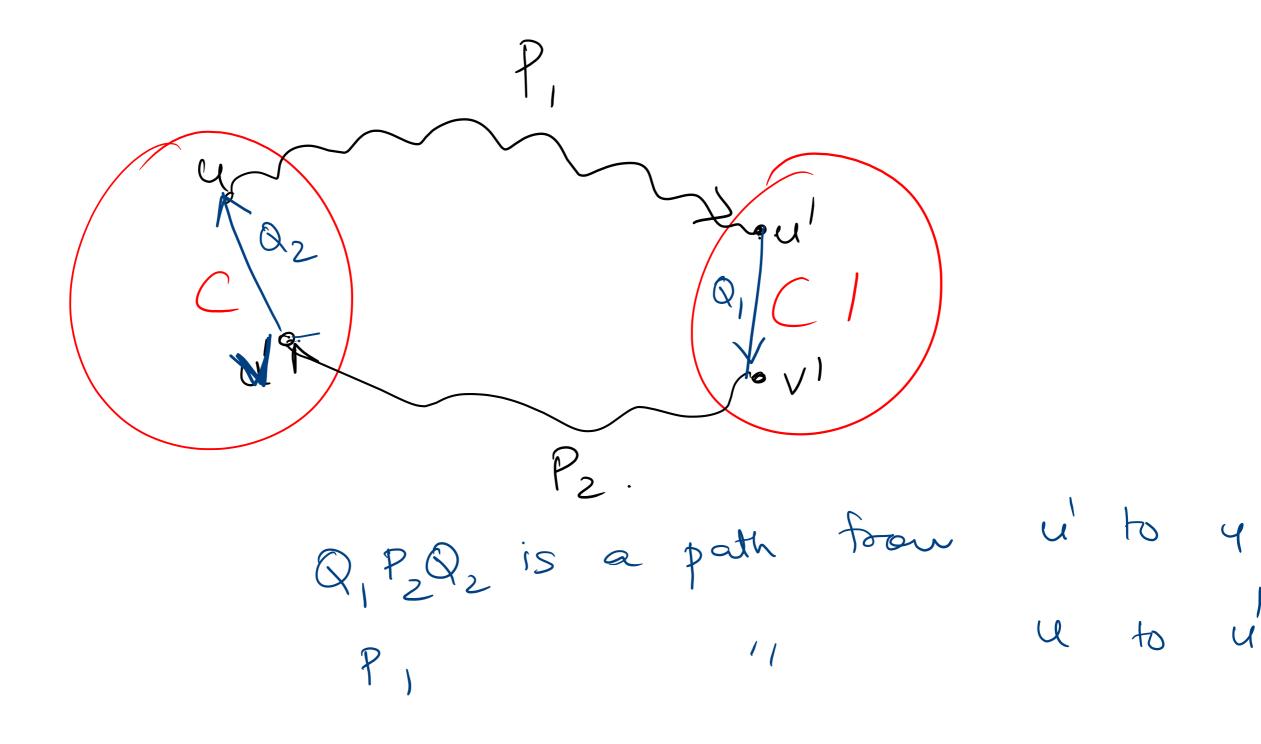


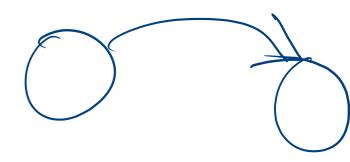


C be a:

Let C,C' be distinct strongly connected components in directed graph G(V,E).

Let $u,v\in C,u',v'\in C'$. Suppose there is a path from u to u'. Then there cannot also be a path from v





Let C,C' be distinct strongly connected components in directed graph G(V,E).

• Suppose $(u,v)\in E$ such that $u\in C$ and $v\in C'$. Then f(C)>f(C').

$$f(C) = \max_{v \in C} f(v)$$

$$d(C) = \min_{v \in C} d(v)$$

$$f(c) > f(c')$$

$$c$$

Case 1:
$$d(c) \ge d(c')$$

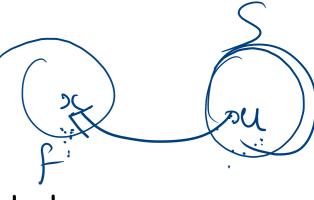
 $d(c) = d[x]$

- I white paths from x to all vertices in C + c!

> They are descendeds of x.

$$f(c) = f(x) > f(c')$$

$$\Rightarrow f(c) > f(c')$$



Let C,C' be distinct strongly connected components in directed graph G(V,E).

- Suppose $(u,v) \in E$ such that $u \in C$ and $v \in C'$. Then f(C) > f(C').
- Suppose there exists $(u, v) \in E^T, u \in C, v \in C'$. Then

Algo. is correct Proof:

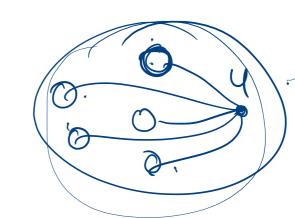
By irondu of DFS trees, k.

k=0 Base case, trivially for

k trees are retrued. Each of the correspond

S.C.C.S.

Let u be _
the root of
(k+1) th DFS tree:
het S be
the S-C. C. that I
Contain u.



V(S) = V(T)

At d'[u], all vertices in 5 au white.

By white path theorm, they are all in T.

 $\Rightarrow V(S) \subseteq V(T)$

TO prove V(T) CV(S).

Assume not.

JXET, but X&S.

x cannot belong to a S-C.C.

that is yet to be seen.

already seen.

Contradicts Industran

My 70 flusion

