

Divide & Conquer Algorithms

Divide and Conquer Algorithms

- breaks the input to several parts
- solves the problem in each part (usually recursively)
- combines the solutions to these subproblems into an overall solution

MergeSort Algorithm

- Divide the array into two equal lists
- Sort each sub-list recursively
- merge the two sorted sub-lists to become one sorted list

MergeSort Algorithm

MergeSort $A[1, \dots, n]$

if $n=1$, return A

$A_1 = \text{MergeSort } A[1, \dots, n/2]$

$A_2 = \text{MergeSort } A[n/2+1, \dots, n]$

$A = \text{Merge}(A_1, A_2)$

return A

MergeSort Algorithm

Merge(A,B)

p=1,q=1

while $p \leq n$ and $q \leq n$

 if $A[p] < B[q]$

 add $A[p]$ to C; $p++$;

 else

 add $B[q]$ to C; $q++$;

if $p > n$

 Add the remaining elements of A to C

else

 Add the remaining elements of B to C

MergeSort Algorithm

Correctness of MergeSort :

MergeSort Algorithm

Correctness of MergeSort :

- By induction on the size of the array
- Correctness of Merge()
 - After the i th iteration, C contains the i th smallest element in the i th position – By induction on i .

MergeSort Algorithm

Running Time of MergeSort :

MergeSort Algorithm

Running Time of MergeSort :

Time taken for Merge() -

- For every iteration of the while loop, an element is added to the merged list.
- The total time is of the order of size of the merged list i.e, $O(n)$
- $T(n)$ - time for MergeSort[1,, n]
- $T(n) = 2T(n/2) + O(n)$

Solving Recurrences

Unroll the recurrence :

Recursion Tree :

Substitution Method :

Master Method :

To solve recurrences of the form

$$T(n) = aT\left(\frac{n}{b}\right) + f(n)$$

where $a \geq 1, b > 1, f(n)$ is asymptotically positive.

Master Method :

1. If $f(n) = O(n^{\log_b a - \epsilon})$ for some $\epsilon > 0$, then

$$T(n) = \Theta(n^{\log_b a})$$

Master Method :

2. If $f(n) = \Theta(n^{\log_b a})$, then $T(n) = \Theta(n^{\log_b a} \log n)$

Master Method :

3. If $f(n) = \Omega(n^{\log_b a + \epsilon})$ for some $\epsilon > 0$ and

if $af(\frac{n}{b}) \leq cf(n)$ for some $c < 1$ and all large n , then

$$T(n) = \Theta(f(n))$$

