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## Time complexity of each sort and some words on the constant

**Bubble sort** - best case  $O(n)$ , worst case  $O(n^2)$ ; average case  $O(n^2)$

When  $n > 50$ , the constant is about 0.75

When  $n < 50$ , the constant is more sporadic, mostly because some best case behaviour is more likely to occur. If we assume it is behaving with a complexity of  $O(n^2)$ , the constant fluctuates from 0.5 to 1

**Shell sort** - best case  $\Omega(n)$   $O(n \log n)$ , average case  $\Omega(n \log n)$   $O(n^2)$

It is difficult to determine the constant, Shell sort appears to get worse than quicksort as  $n$  increases, and better than bubble sort as  $n$  increases.

Its average asymptotic behavior is probably faster than  $n^2$  and slower than  $n \log n$

**Quick sort** - best case  $O(n \log n)$ , worst case  $O(n^2)$ , average case  $O(n \log n)$

When  $n > 20$ , the constant is about 2

When  $n < 20$ , the sort is more likely to exhibit worst case behavior. Under the assumption that it is exhibiting average case behavior, the constant can range from 2 to 2.5

**Binary insertion sort** - best case  $O(n)$ , worst case  $O(n^2)$ , average case  $O(n^2)$

Empirical note: As far as I have seen the number of moves for binary insertion sort *exactly* equals the number of moves for bubble sort. Here's why I think this may be the case:

**Claim:** For both algorithms( binary insertion sort and bubble sort)

For each element A

The number of times it is swapped is equal to # of elements greater than it that are below it + # of elements lesser than it that are above it

Here's why:

**I:** Both of these algorithms only swap adjacent elements.

**II:** Both algorithms will only swap adjacent elements if the lower element is the greater of the two.

**III:** From **I** and **II** it follows that for each element A:

The elements greater than A yet below A must swap with A once to surpass it

The elements greater than A yet above A will not swap with A

The elements lesser than A yet below A will not swap with A

The elements lesser than A yet above A must swap with A once to subceed it

The elements equal to A will not swap with A

**IV:** For every element A, the other elements are exactly in one of these five sets

From **I, II, III** and **IV**, the **Claim** follows