

## Procedure

The purpose of this lab is to be practice using assembly programming using ARMv7 in CPUlator to accomplish Tas#1, #2, and #3. The objective of this lab is to be familiarized with the layout of CPUlator, program in ARM assembly language, and understand Floating Point Addition.

### Task #1:

#### Part 1

The first part of Task #1 was to type the given program in Figure 1.1 into the CPUlator, assemble the program, make sure there is no errors, modify the code as instructed, and use the “Step Into” function to watch the values of register to answer the given questions in Figure 1.2. The code was then saved as Lab4\_Task1\_Part1.s. The answers can be found under the result section.

```
1 .text
2 .align=2
3
4 .global Start
5 Start:
6
7     mov r0, #4           @Load 4 into r0
8     mov r1, #5           @Load 5 into r1
9     add r2, r0, r1       @Add r0 to r1 and place in r2
10 S:
11     B S                 @Infinite loop ending
12 .end
```

Figure 1.1 Simple Program that Adds 2 Numbers

1. Explain what these lines mean

```
1 .text
2 .align=2
```

2. What is the value of R0, R1, R2, and PC at the start and at the end of the program?

3. Explain the S: B S line of code (lines 10 and 11)

4. Expand the program to solve  $4+5+9-3$  and save the result in the 40th word in memory. Take a screenshot of the memory for your lab report.

5. Save the current assembly code (Ctrl + S, or go to File -> Save). Change the name of the .s file to Lab4\_Task1\_Part1.s. You will need to submit this source file, along with your report. See the Deliverables section for more details.

Figure 1.2 Questions to be Answer

#### Part 2:

The second part of the task was to open a new CPUlator instance, assemble and debug the new given program in Figure 1.3. Then, follow the execution in the debugger and modify the code as instructed to answer the given question in Figure 1.4. The code was then saved as Lab4\_Task1\_Part2.s. The answers can be found under the results section.

```
1 .global _start
2 _start:
3
4     MOV r0, #3
5 LOOP:
6     PUSH {R0, LR}
7     CMP R0, #1
8     BGT ELSE
9     MOV R0, #1
10    ADD SP, SP, #8
11    MOV PC, LR
12 ELSE:
13    SUB R0, R0, #1
14    BL LOOP
15    POP {R1, LR}
16    MUL R0, R1, R0
17    MOV PC, LR
18 .end
```

Figure 1.3 Assembly Program

1. What is the value in R0 after the program ends?

2. If the value initially placed in R0 is equal to 5, what is the value in R0 when the program ends?

3. What does this program do?

4. If you replace the instructions at lines 2 and 10 in Figure 2 with PUSH {R0, R1} and POP {R1, R2} respectively, how will the program behave and why?

5. Repeat 4 with the following instructions modifications:

a. Replace the instructions at lines 2 and 10 in Figure 2 with PUSH {R3, LR} and POP {R3, LR} respectively.

b. Replace the instructions at lines 2 and 10 in Figure 2 with PUSH {LR} and POP {LR} respectively.

c. Delete the instruction at line 6

Figure 1.4 Task1 P2 Questions to Answer

## Task #2:

The second task was to design an algorithm for counting the number of 1's in a 32-bit number. The approach taken was first design a pseudocode found in Figure 1.5. Then, implement the pseudocode using ARMv7 before testing the program in CPUlator.

```
number = #
counter = 0
compared = 0x80000000
sum = 0
while (counter != 32) :
    if msb is 1
        sum++
    LSL number
    counter++
```

Figure 1.5 Pseudocode for Task#2

## Task #3:

The third task was to write an ARM assembly language function that performs floating-point addition that follows the algorithm in Figure 1.6. Note that the code will add only strictly positive numbers and the sign bits in the numbers being summed can be ignored. The sign bit of the resulting sum should be set to zero.

The approach to this task was to first study and understand the floating-point arithmetic. Then, a pseudocode was written and implemented into the CPUlator. Then, the Step In function was utilized to debug the code.

Additionally, a hand analysis prompts were given and answered as a guiding questions. The hand analysis and its answers can be found in Figure 1.7. To test if the program is working properly, the register memory result in R0 was tested against the expected result after the program ended, following the mentioned algorithm. The registers were viewed and the register holding the sum of two floats was examined. The final is then saved as Lab4\_Task3.s.

### The Algorithm:

In summary, your algorithm will need to do the following:

1. Mask and shift down the two exponents.
2. Mask the two fractions and append leading 1's to form the mantissas.
3. Compare the exponents by subtracting the smaller from the larger. Set the exponent of the result to be the larger of the exponents.
4. Right shift the mantissa of the smaller number by the difference between exponents to align the two mantissas.
5. Sum the mantissas.
6. Normalize the result, i.e., if the sum overflows, right shift by 1 and increment the exponent by 1.
7. Round the result (truncation is fine).
8. Strip the leading 1 off the resulting mantissa, and merge the sign, exponent, and fraction bits.

Figure 1.6 The Algorithm for Floating-Point Addition

### Hand Analysis

Before implementing floating point addition, re-familiarize yourself with the representation of floating point numbers and with carrying out addition by hand by answering the following questions. Give your answers in binary and hexadecimal. For example, 1.0 is written as an IEEE single-precision floating point number as:

1.0 = 0 01111111 000000000000000000000000 = 3F800000<sub>16</sub>

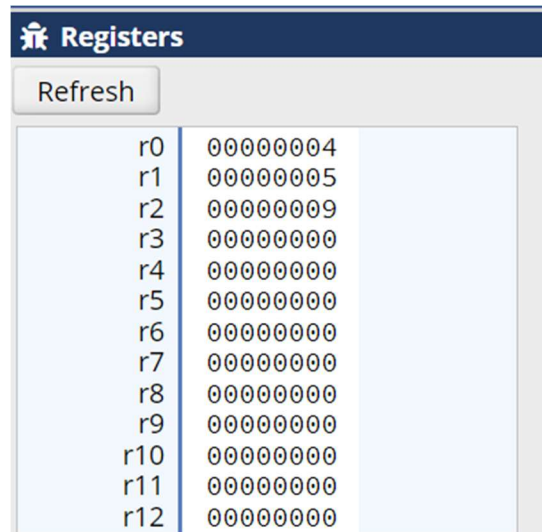
1. Write 2.0 as an IEEE single-precision floating point number.  $0 \times 40000000$
2. Write 3.5 as an IEEE single-precision floating point number.  $0 \times 40600000$
3. Write 0.50390625 as an IEEE single-precision floating point number.  $0 \times 3F010000$
4. Write 65535.6875 as an IEEE single-precision floating point number.  $0 \times 477FFB00$
5. Compute the sum of the numbers from (c) and (d) and express the result in IEEE floating point format. Truncate the sum if necessary.  $0 \times 47800018$

Figure 1.7 Hand Analysis (Answered)

## Results:

### Task #1:

#### Part 1



Register	Value
r0	00000004
r1	00000005
r2	00000009
r3	00000000
r4	00000000
r5	00000000
r6	00000000
r7	00000000
r8	00000000
r9	00000000
r10	00000000
r11	00000000
r12	00000000

Figure 2.1 Register values for Task1 Part 1

1. R0 = 0, R1=0, R2=0, PC=0  
R0=4, R1=5, R2=9, PC=0xC
2. S branches to S which then branches back to S, creating an infinite loop
- 3.

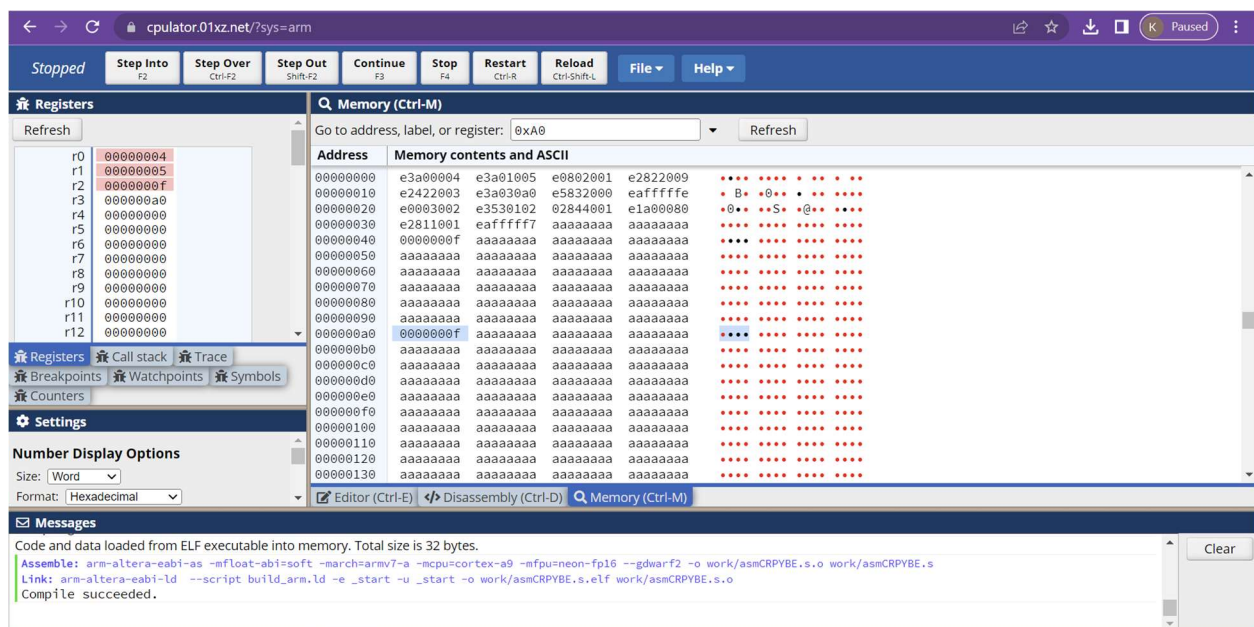


Figure 2.2 Register and Memory Results for Extended Task #1 Part 1

#### Part 2

1. 6
2. 120
3. The program takes the factorial of r0

4. R0 will become 120 because it will have more iterations multiplying for each decrement of r0

r0	00000078
r1	00000005
r2	00000028
r3	00000000
r4	00000000
r5	00000000
r6	00000000
r7	00000000
r8	00000000
r9	00000000
r10	00000000
r11	00000000

Figure 3.1 Register Results for Task #1 Part 2

## Task #2:

Registers	
Refresh	
r0	80008000
r1	00000000
r2	80000000
r3	00000000
r4	00000000
r5	00000000
r6	00000000
r7	00000000
r8	00000000
r9	00000000
r10	00000000
r11	00000000
r12	00000000

Figure 3.1 Register Results before execution

Registers		
Refresh		
r0	00000000	
r1	00000020	
r2	80000000	
r3	00000000	
r4	00000002	
r5	00000000	
r6	00000000	
r7	00000000	
r8	00000000	
r9	00000000	
r10	00000000	
r11	00000000	
r12	00000000	

Figure 3.2 Register Results after execution

### Task #3:

Task #3 did not work as expected. According to the Hand Analysis, R0 should display the result of 47800018 stored in the resistor memory as a result of the sum of the numbers found. One of the major problem encounter was in the comparing step. As a solution, more time needs to be put to study the float-point arithmiaite and more understand of assembly code must be put to study and more debugging required.

r0	0c0000ff	
r1	4f7fffb0	
r2	00000018	
r3	000000ff	
r4	e3000000	
r5	e3430f01	
r6	e30f1fb0	
r7	e344177f	
r8	e3002000	
r9	e3472f80	
r10	e0003002	
r11	e0014002	
r12	00000000	
sp	00000024	
lr	00000000	
pc	cda03bf0	
cpsr	400001d3	NZCVI SVC
spsr	00000000	NZCVI ?

Figure 4 Register Memory Result for Task #3

## **Final Product**

As a final product, we assembled the decrementing program from Task #1 and successfully familiarized ourselves with using CPUlator. For task #2, we successfully made an algorithm for counting the number of 1's in a 32-bit number using ARM assembly language. For task #3, we attempted to write an ARM assembly language function that performs floating-point addition that strictly adds positive numbers. Although unsuccessful fully throughout the lab, we learned to use CPUlator and assembly language well.

## Appendix

```
Editor (Ctrl-E)
Compile and Load (F5)  Language: [ARMv7 v]  untitled.s [changed since save] [changed since compile]

1 .text
2 .align=2
3
4 .global _start
5 _start:
6
7     mov r0, #4           @Load4 into r0
8     mov r1, #5           @Load5 into r1
9     add r2, r0, r1       @Add r0 to r1 and place into r2
10
11    add r2, r2, #9        @Adds 9 to r2
12    sub r2, r2, #3        @Subtracts 3 from r2
13
14    mov r3, #0xA0         @Stores word 40 into r3
15    str r2, [r3]          @Stores the value of r2 into word 40 (Address 160)
16 S:
17    B S                   @Infinite loop ending
18 .end
19
```

### Appendix 1 Expanded Task1 Part 1

```
1 .global _start
2     PUSH {LR}
3
4     MOV r0, #5
5 LOOP:
6     PUSH {R0, LR}
7     CMP R0, #1
8     BGT ELSE
9     MOV R0, #1
10    POP {LR}
11    MOV PC, LR
12 ELSE:
13    SUB R0, R0, #1
14    BL LOOP
15    POP {R1, LR}
16    MUL R0, R1, R0
17    MOV PC, LR
18    .end
```

### Appendix 2 Expanded Task1 Part 2



```

.global _start
_start:
    MOVW R0, #0x8000 // Right half of number
    MOVT R0, #0x8000 // Left half of number
    MOV R1, #0        // Counter
    MOVW R2, #0x0000 // Right half of compared
    MOVT R2, #0x8000 // Left half of compared
    MOV R4, #0        // Sum

WHILE:
    CMP R1, #32
    BEQ END            // Branch to end if counter equal to 32 (done counting)
    AND R3, R0, R2      // Compare msb to number (if 1 then and will be the equivalent of msb being 1)
    CMP R3, #0x80000000
    ADDEQ R4, R4, #1    // Increment sum if the result is 1
    LSL R0, R0, #1      // Shift the given number bits to the left
    ADD R1, R1, #1      // Increment counter
    B WHILE            // Re-enter the while loop

END:|

```

## Appendix 3 Expanded Task2



```

// Kiana Peterson and Nattapon Oonlamom
// 05/19/2023
// EE 469
// Lab 4: Assembly Programming

/* Overview:
 *          This code perform floating-point addition.
 *          The implementation assumes that the inputs are
 *          strictly positive single-precision floating-
point.
 * Inputs:
 *          R0 = The first floating-point number
 *          R1 = The second floating-point number
 * Output:
 *          R0 = Result of the addition
 */
.global _start
_start:

    // Initialize the floating-point number
    MOVW R0, #0x0000 // Initialize R0 with the first
floating-point number
    MOVT R0, #0x3F01
    MOVW R1, #0xFFB0 // Initialize R1 with the second
floating-point number
    MOVT R1, #0x477F

    // Step 1: Mask and shift down the two exponents
    MOVW R2, #0x0000 // Initialize exponent mask
    MOVT R2, #0x7F80
    AND R3, R0, R2 // Extract the exponent
    AND R4, R1, R2
    LSR R3, R2, #23 // Shift the exponent to the right
    LSR R4, R3, #23

    // Step 2: Mask the two fractions and append leading 1's
to form the mantissas
    MOVW R2, #0x007F // Initialize fraction mask
    MOVT R2, #0x0000
    AND R5, R0, R2 // Extract the fraction
    AND R6, R1, R2
    ORR R0, R0, #0x80000000 // Append leading 1 to memory
    ORR R1, R1, #0x80000000

    // Step 3: Compare the exponents by subtraction and
// Set the result exponent to be the larger
    CMP R3, R4
    MOVGE R7, R3 // If R3 > R4
    MOVLt R7, R4 // If R3 < R4

    // Step 4: Right shift the smaller mantissa to align the two
mantissas

```