



# Software Defined Networking

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*In this course, you will learn about software defined networking and how it is changing the way communications networks are managed, maintained, and secured.*

## **This Module: Verification**

- ⦿ **Motivation:** How do you know the network is doing the right thing?
- ⦿ Verification techniques
  - Configuration Verification: rcc (pre-SDN)
  - Control Plane Verification: Kinetic
  - Data Plane Verification
    - Header Space Analysis
    - Veriflow

## Simple Questions are Hard

- What are all the packet headers from A that can reach B?
- What will happen if I remove an entry from a firewall?
- Is Group X provably isolated from Group Y?
- Are there any loops in the network?
- Why is my network slow?

## Configuration Defines Behavior

*Provides flexibility for realizing operational goals*

- ⦿ How traffic enters and leaves the network
  - Load balance
  - Traffic engineering
  - Primary/backup paths
- ⦿ Which neighboring networks can send traffic
  - Defines business relationships and contracts
- ⦿ How routers within the network learn routes
  - Scaling and performance

**Flexibility** → **Complexity**

## **Most Important Goal: Correctness**

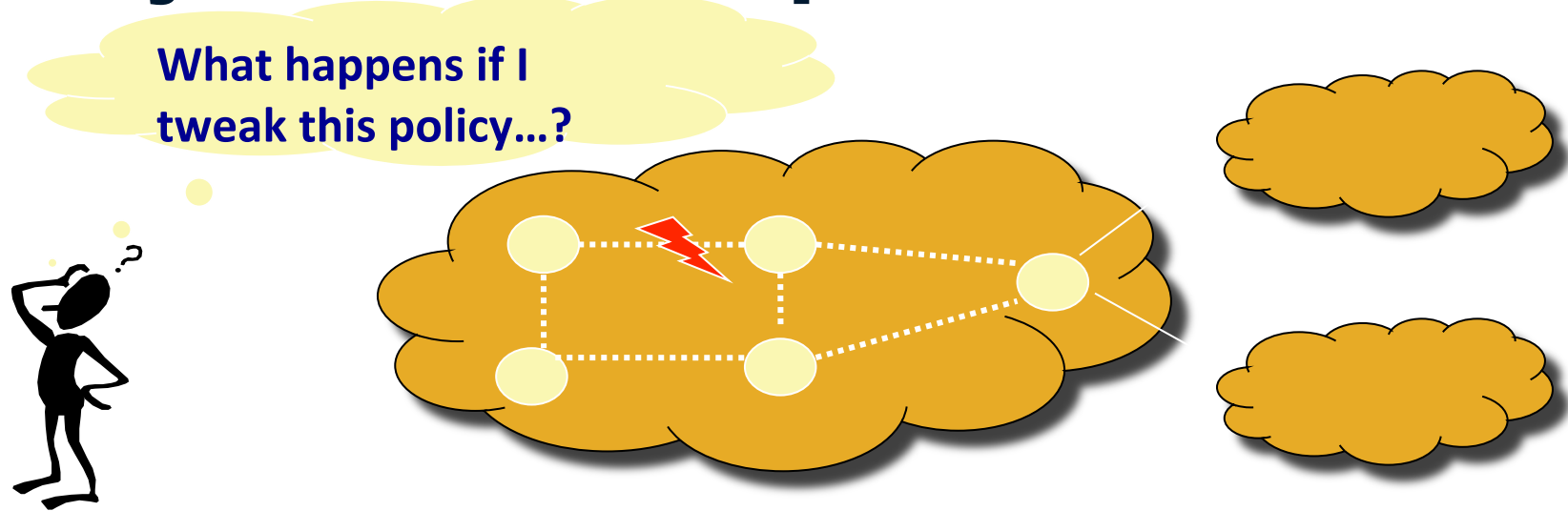
*Unfortunately...*

***Mistakes happen!***

***Why?***

- ⦿ Configuration is **difficult**. Operators make mistakes.
  - Complex policies
  - Configuration is distributed across routers
- ⦿ Each network **independently configured**
  - Unintended policy interactions

## Today: Stimulus-Response

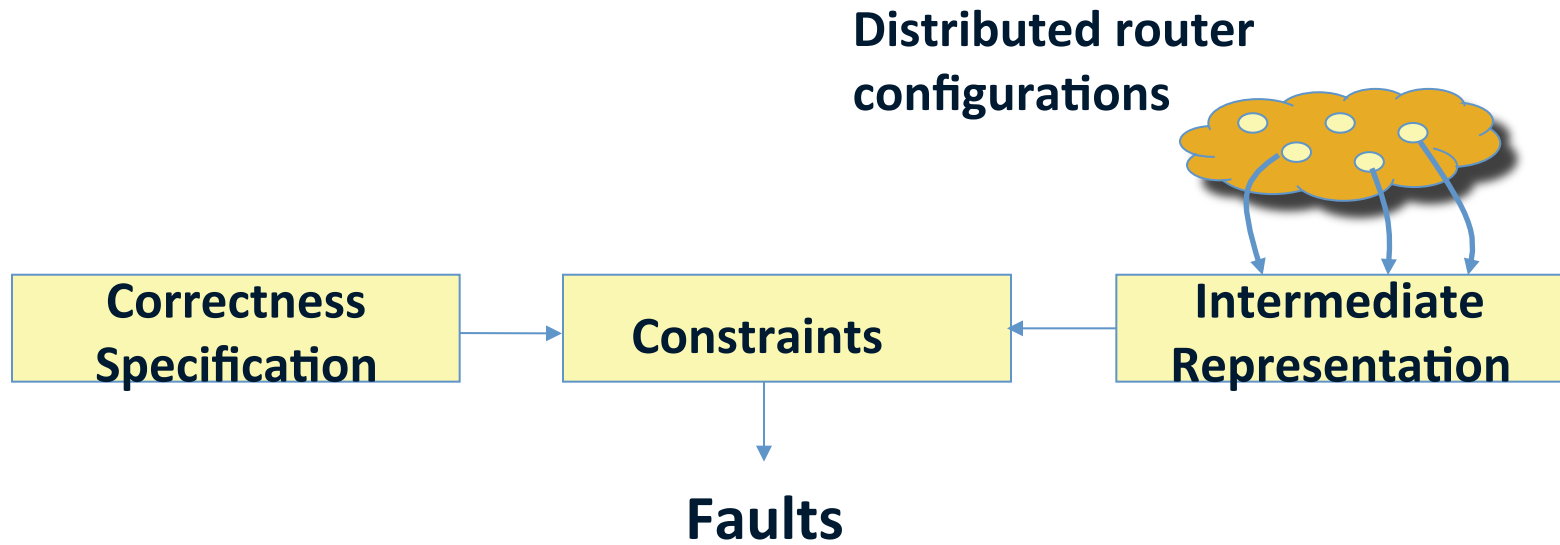


- Unacceptable programming environment
- Mistakes cause downtime
- Mistakes often not immediately apparent

## Checking Configuration

- ⦿ Correctness specification and constraints for global Internet routing
- ⦿ rcc (“router configuration checker”)
  - Static configuration analysis tool for fault detection
  - Used by network operators (including large backbone networks)
- ⦿ Analysis of real-world network configurations from 17 autonomous systems

## rcc Design

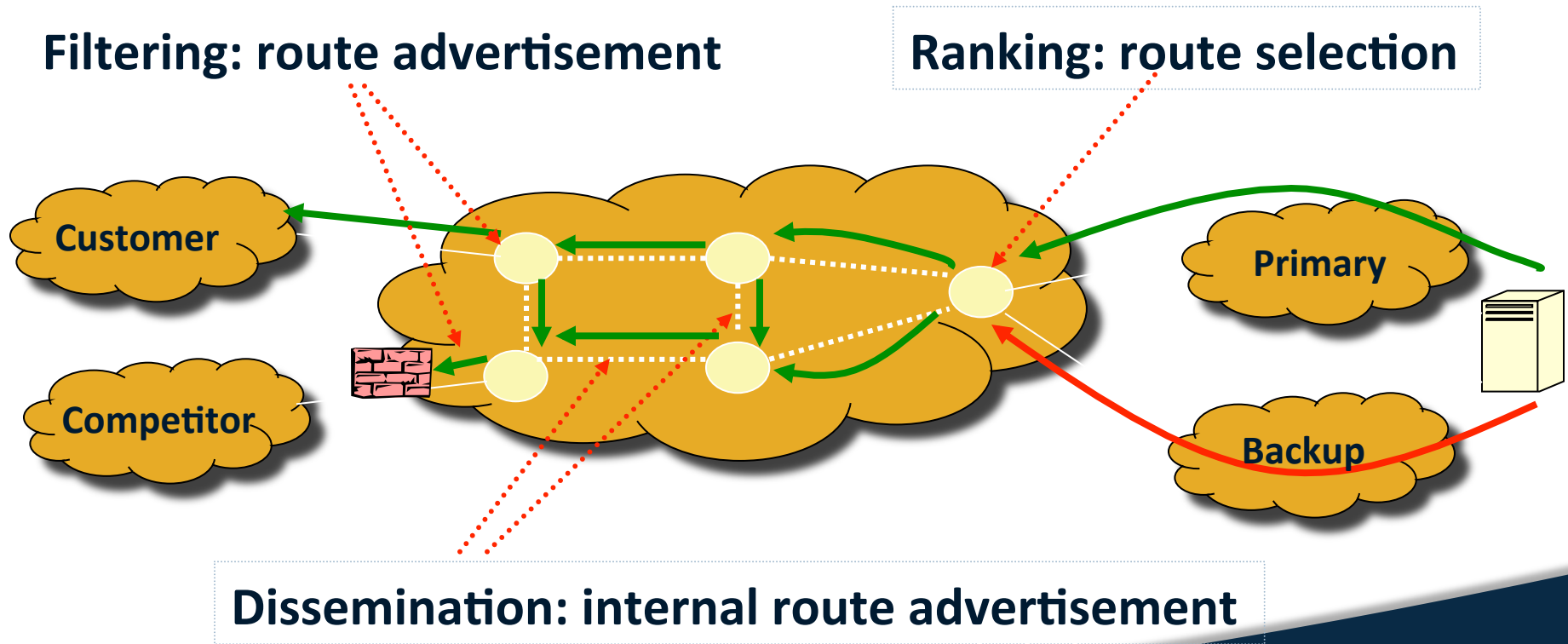




## Challenges

- ⦿ Defining a correctness specification
- ⦿ Deriving verifiable constraints from specification
- ⦿ Analyzing complex, distributed configuration

## Factoring Routing Configuration



## Path Visibility

If every router learns a route for every usable path, then path visibility is satisfied.

*A usable path:*

- Reaches the destination
- Corresponds to the path that packets take when using that route
- Conforms to the policies of the routers on that path

## Possible path visibility faults

### Dissemination

- Partition in session-level graph that disseminates routes

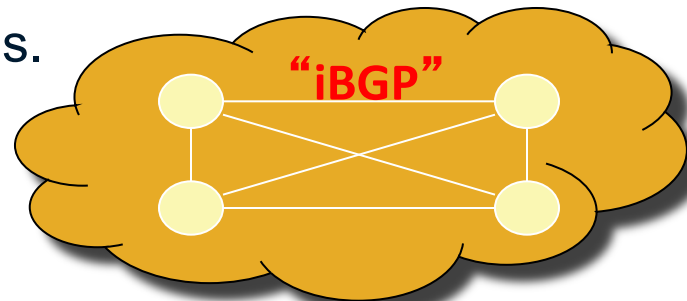
### Filtering

- Filtering routes for prefixes for usable paths

## Path Visibility: Internal BGP (iBGP)

- Default: don't re-advertise iBGP-learned routes.

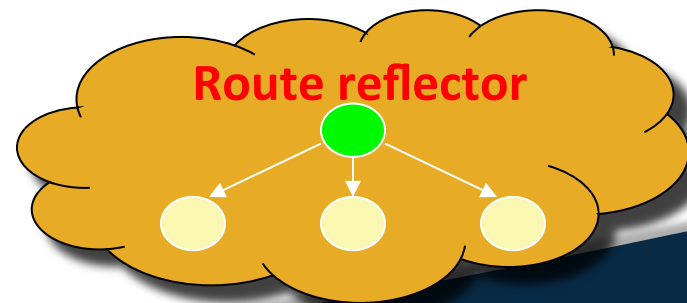
Complete propagation requires “full mesh” iBGP. *Doesn't scale.*



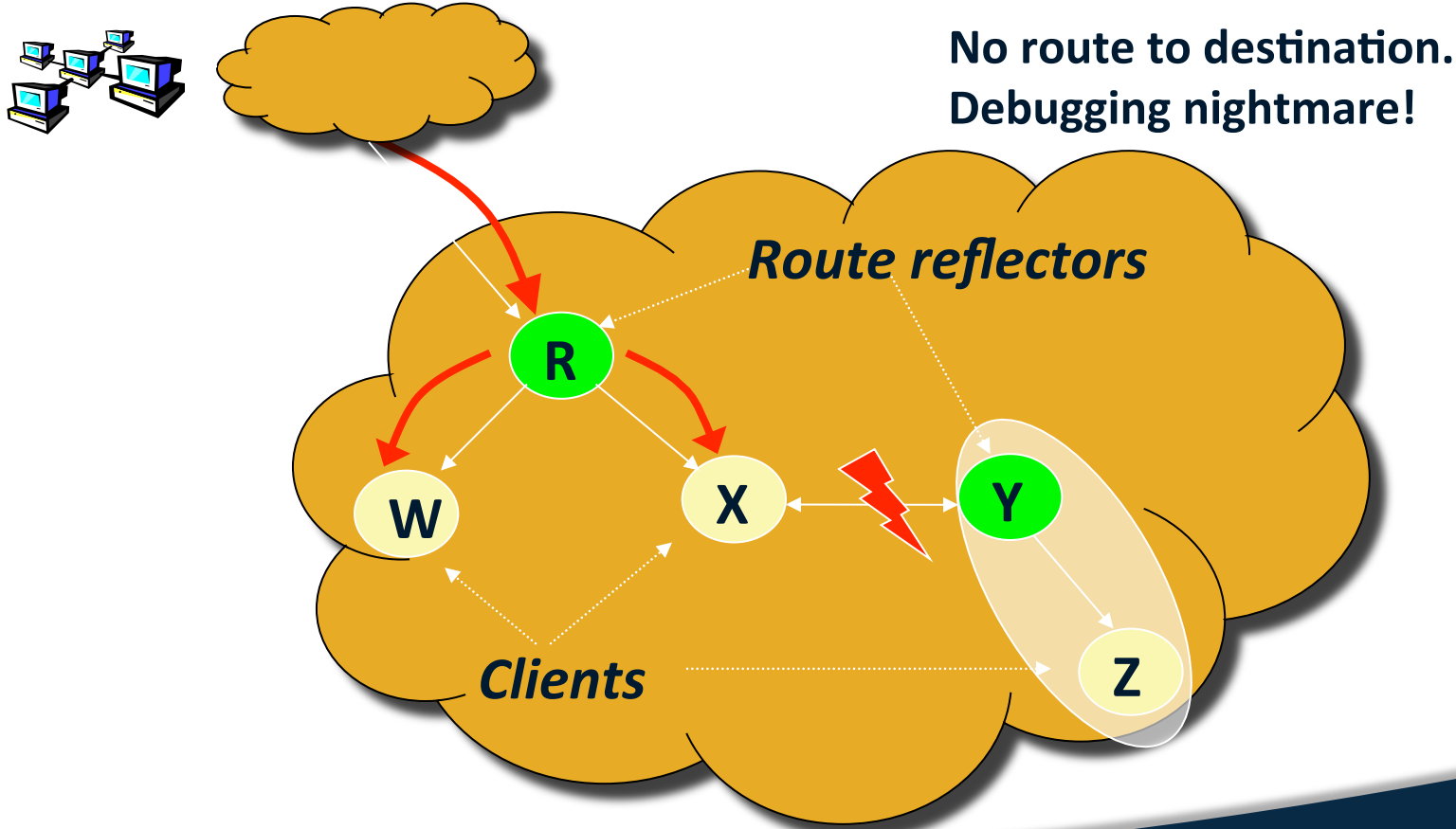
- “Route reflection” improves scaling.

**Client:** re-advertise as usual.

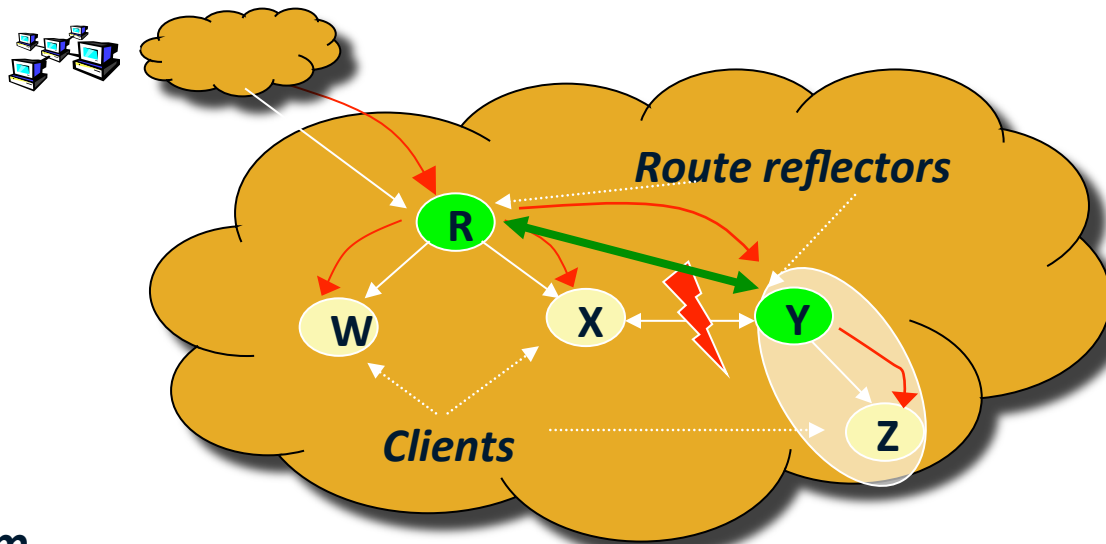
**Route reflector:** reflect non-client routes to all clients, client routes to non-clients and other clients.



## Path Visibility: iBGP Signaling



## Path Visibility: iBGP Signaling



### Theorem.

Suppose the iBGP reflector-client relationship graph contains no cycles. Then, path visibility is satisfied if, and only if, *the set of routers that are not route reflector clients forms a full mesh.*

*Condition is easy to check with static analysis.*

## Route Validity

**If every route that a router learns corresponds to a usable path, then route validity is satisfied.**

*A usable path:*

- Reaches the destination
- Corresponds to the path that packets take when using that route
- Conforms to the policies of the routers on that

## Possible route validity faults

### Filtering

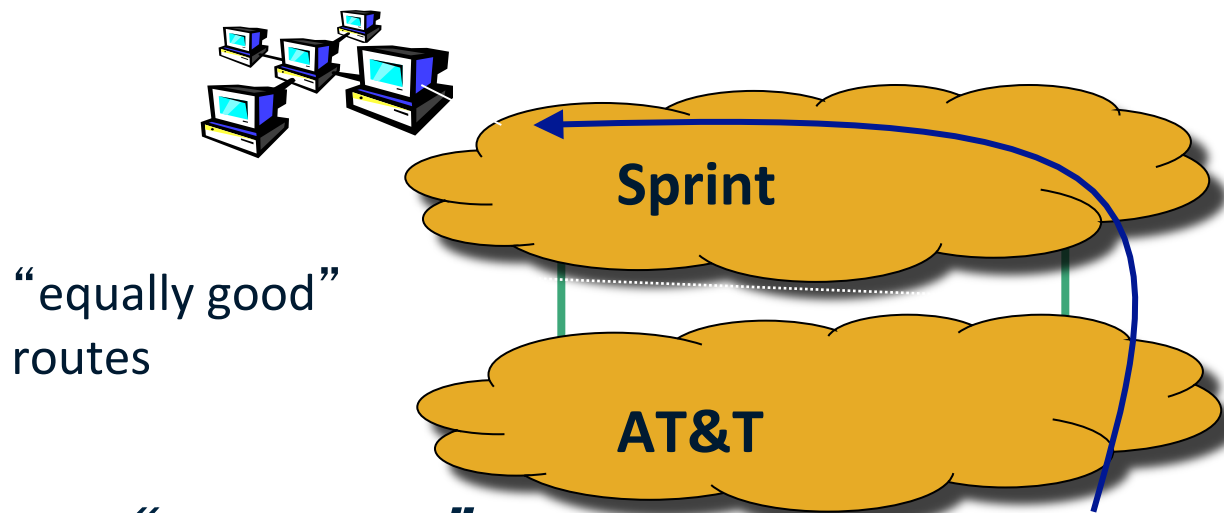
- Unintentionally providing transit service
- Advertising routes that violate higher-level policy
- Originating routes for private (or unowned) address space

### Dissemination

- Loops and “deflections

## Route Validity: Consistent Export

- Rules of *settlement-free peering*:
  - Advertise routes at all peering points
  - Advertised routes must have equal “AS path length”



*Enables “hot potato” routing.*



## **Summary: Verifying Config is Hard**

- ⦿ Distributed configuration is a bad idea.
- ⦿ SDN can allow us to treat verification as a distributed program, applying concepts from software engineering, testing, formal verification...

**SDN to the rescue...**