

Computer Arithmetic

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WGB 182

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Floating-Point Addition

- Consider a 4-digit decimal example
$$9.999 \times 10^1 + 1.610 \times 10^{-1}$$
- 1. Align decimal points (Shift number with smaller exponent)
$$9.999 \times 10^1 + 0.016 \times 10^1$$
- 2. Add significands
$$9.999 \times 10^1 + 0.016 \times 10^1 = 10.015 \times 10^1$$
- 3. Normalize result & check for over/underflow
 - 1.0015×10^2
- 4. Round and renormalize if necessary
 - 1.002×10^2

Floating-Point Addition

- Now consider a 4-digit binary example

- Add 0.5_{10} and -0.4375_{10}

$$1.000_2 \times 2^{-1} + -1.110_2 \times 2^{-2}$$

1. Align binary points (**Shift** number with smaller exponent)

$$1.000_2 \times 2^{-1} + -0.111_2 \times 2^{-1}$$

2. **Add** significands

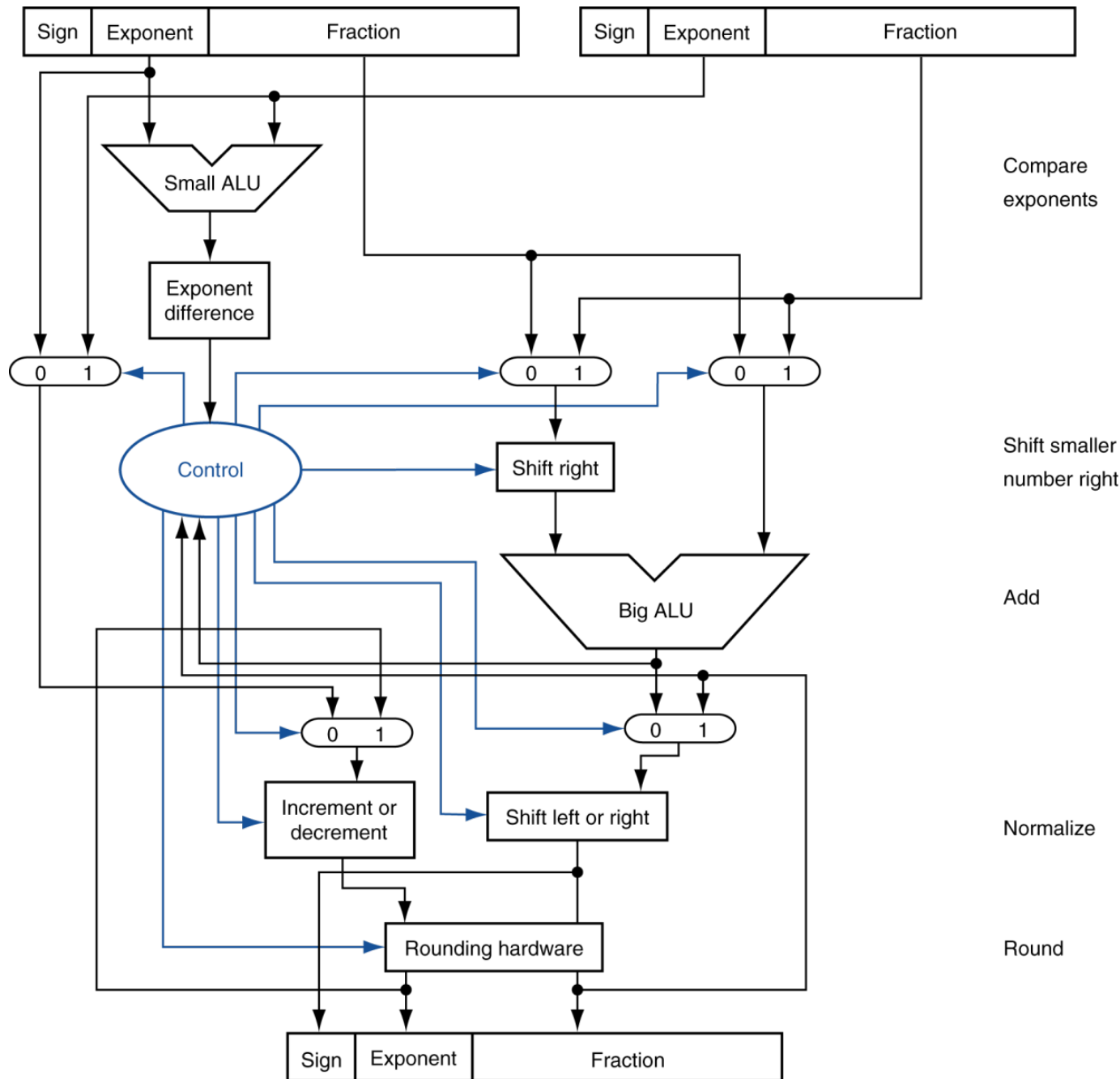
$$1.000_2 \times 2^{-1} + -0.111_2 \times 2^{-1} = 0.001_2 \times 2^{-1}$$

3. **Normalize** result & check for over/underflow

$$1.000_2 \times 2^{-4}, \text{ with no over/underflow}$$

4. **Round** and renormalize if necessary

- $1.000_2 \times 2^{-4}$ (no change) = 0.0625_{10}



FP Adder Hardware

Step 1

Step 2

Step 3

Step 4



Floating-Point Multiplication

- Consider a 4-digit decimal example
 - $1.110 \times 10^{10} \times 9.200 \times 10^{-5}$
- 1. **Add** exponents
 - For biased exponents, subtract bias from sum
 - New exponent = $10 + -5 = 5$
- 2. **Multiply** significands
 - $1.110 \times 9.200 = 10.212 \Rightarrow 10.212 \times 10^5$
- 3. **Normalize** result & check for over/underflow
 - 1.0212×10^6
- 4. **Round** and renormalize if necessary
 - 1.021×10^6
- 5. **Determine *the sign*** of result from signs of operands
 - $+1.021 \times 10^6$

Floating-Point Multiplication

- Now consider a 4-digit binary example
 - Multiply 0.5_{10} and -0.4375_{10}
 - $1.000_2 \times 2^{-1} \times -1.110_2 \times 2^{-2}$
- 1. **Add exponents**
 - Unbiased: $-1 + -2 = -3$
 - Biased: $(-1 + 127) + (-2 + 127) = -3 + 254 - 127 = -3 + 127$
- 2. **Multiply significands**
 - $1.000_2 \times 1.110_2 = 1.110_2 \Rightarrow 1.110_2 \times 2^{-3}$
- 3. **Normalize** result & check for over/underflow
 - $1.110_2 \times 2^{-3}$ (no change) with no over/underflow
- 4. **Round** and renormalize if necessary
 - $1.110_2 \times 2^{-3}$ (no change)
- 5. **Determine sign**: $+ve \times -ve \Rightarrow -ve$
 - $-1.110_2 \times 2^{-3} = -0.21875_{10}$

FP Arithmetic Hardware

- FP multiplier is of similar complexity to FP adder
 - But uses a multiplier for significands instead of an adder
- FP arithmetic hardware usually does
 - Addition, subtraction, multiplication, division, reciprocal, square-root
 - $\text{FP} \leftrightarrow \text{integer}$ conversion
- Operations usually takes several cycles
 - Can be pipelined

Review

- Revised binary representation
- Explored basic integer operations and the [optimized] hardware implementation of multiplication and division
- Introduced floating point representation [precision and range]
- Explored the floating point adder hardware design
- Explored the procedure of multiplication and division

FP Example: °F to °C

- C code:

```
float f2c (float fahr) {  
    return ((5.0/9.0)*(fahr - 32.0));  
}
```

- fahr in \$f12, result in \$f0, literals in global memory space

- MIPS code:

```
f2c: lwc1  $f16, const5  
     lwc1  $f18, const9  
     div.s $f16, $f16, $f18  
     lwc1  $f18, const32  
     sub.s $f18, $f12, $f18  
     mul.s $f0,  $f16, $f18  
     jr   $
```

Associativity Pitfall

$$a+(b+c) = (a+b)+c$$

- Parallel programs may interleave operations in unexpected orders
 - Assumptions of associativity may fail

		$(x+y)+z$	$x+(y+z)$
x	-1.50E+38	0.00E+00	-1.50E+38
y	1.50E+38		
z	1.0	1.0	1.50E+38
		1.00E+00	0.00E+00

Floating point addition is not generally associative!

Who Cares About FP Accuracy?

- Important for scientific code
 - But for everyday consumer use?
 - “My bank balance is out by 0.0002¢!”
- The Intel Pentium FDIV bug in 1994
 - The market expects accuracy
 - Costed \$500 Million to correct
 - No Christmas bonus for Intel Engineers



Reading

- Sections 3.6, 3.7, 3.9, and 3.10