

Review 1

PROBLEM 1

- 1) Consider the following code that copies one string in the memory to another memory location

strcpy:

addi \$sp, \$sp, -4 # adjust stack for 1 item

sw \$s0, 0(\$sp) # save \$s0

add \$s0, \$zero, \$zero # i = 0

L1: add \$t1, \$s0, \$a1 # addr of y[i] in \$t1

lbu \$t2, 0(\$t1) # \$t2 = y[i]

add \$t3, \$s0, \$a0 # addr of x[i] in \$t3

sw \$t2, 0(\$t3) # x[i] = y[i]

beq \$t2, \$zero, L2 # exit loop if y[i] == 0

addi \$s0, \$s0, 1 # i = i + 1

j L1 # next iteration of loop

L2: lw \$s0, 0(\$sp) # restore saved \$s0

addi \$sp, \$sp, 4 # pop 1 item from stack

jr \$ra # and return

- Is this a leaf or non-leaf procedure?
- Which instruction format is used in every instruction?
- Assuming that the string has three characters, how many times would every instruction be executed?
- Assuming a single clock cycle processor, how many clock cycles are needed to execute this code?
- In a pipelined processor, which instructions are prone to pipelining hazards? For each of these instructions (if any)
 - Identify the type of hazard and indicate whether processor stalls can be avoided or not?
 - Which technique is used to avoid the hazard?

1) leaf procedure as it does not call external procedures or perform recursive calls.

2) I I R R I R I I I J I I R

3) 1 1 1 4 4 4 4 4 3 3 1 1 1

4) 32 cycles

- 5) sw \$s0, 0(\$sp) prone to data hazard solved by bypassing (forwarding)
 add \$t1, \$s0, \$a1 prone to data hazard solved by bypassing (forwarding)
 lbu \$t2, 0(\$t1) prone to data hazard solved by bypassing (forwarding)
 add \$t3, \$s0, \$a0 prone to data hazard solved by bypassing (forwarding)
 sw \$t2, 0(\$t3)
 beq \$t2, \$zero, L2 prone to control hazard can be avoided with proper prediction

PROBLEM 2

- 1) Consider the following C statements

```
int x1=500;
```

```
float x2=500;
```

```
float x3=500.25;
```

Write down the corresponding binary value stored for x1, x2, and x3.

1) $500 = 111110100$

2) $500 = 1.11110100 \times 2^8$

S=0, EXPONENT = 135 = 10000111, Fraction = 1111 0100 0000 0000 0000 000

3) $500.25 = 111110100.01 = 1.1111010001 \times 2^8$

S=0, EXPONENT = 135 = 10000111, Fraction = 1111 0100 0100 0000 0000 000

Double precision format: 11-bit exponent and 52-bit fraction