

1. Supposed that the following function declarations (in C++ syntax) are available in a program:

- `int pow(int, int);`
- `double pow(int, double);`
- `double pow(double, double);`

and suppose that the following code calls the `pow` function:

```
int x;  
double y;  
x = pow(2, 3);  
    C: 1  
    Java: 1  
    Ada: 1  
y = pow(2, 3);  
    C: 1  
    Java: 1  
    Ada: illegal  
x = pow(2, 3.2);  
    C: 2  
    Java: 2  
    Ada: illegal  
y = pow(2, 3.2);  
    C: 2  
    Java: 2  
    Ada: 2  
x = pow(2.1, 3);  
    C: 3  
    Java: 3  
    Ada: illegal  
y = pow(2.1, 3);  
    C: 3  
    Java: 3  
    Ada: illegal  
x = pow(2.1, 3.2);  
    C: 3  
    Java: 3
```

Ada: illegal

y = pow(2.1, 3.2);

C: 3

Java: 3

Ada: 3

Given the languages (a) C++, (b) Java, and (c) Ada, write down the number of the pow function called in each of the eight calls, or write "illegal" if a call cannot be resolved in the language, or if a data type conversion cannot be made.

2. Assume x is an int variable and y is an int * variable.
1. Which of the following C expressions are l-values? Which are not? Why?

- (1) x + 2 **not an lvalue**
- (2) &x **not an lvalue**
- (3) *&x **is an lvalue**
- (4) &x + 2 **not an lvalue**
- (5) *(&x + 2) **is an lvalue**
- (6) &*y **not an lvalue**

2. Is it possible for a C expression to be an l-value but not an r-value? Explain.

Any expression that includes a type declaration such as "int x" can only be an lvalue, never an rvalue. For instance:

int x = y is legal but

y = int x is illegal.

Therefore it is possible for a C expression to be an l-value but not an r-value.

3. Given the C declarations:

```
struct {int i; double j;} x, y;  
struct {int i; double j;} z;
```

1. The assignment x = z generates a compilation error, but the assignment x = y does not. Why? **X and Y are of the same type so the assignment is valid. C uses name equivalence for structs which does not exist between the two anonymous structs so the compiler sees x and z as different types. Without explicit rules for casting the assignment x = z illegal.**
2. Give two different ways to fix the code so that x = z works.

1) struct {int i; double j;} x, y, z;

```
x=z;
```

```
2) struct {int i; double j;} x, y;
```

```
struct {int i; double j;} z;
```

```
x.i = z.i;
```

```
x.j = z.j;
```

4. Give the output of the following program (written in C syntax) using the four parameter passing methods: Pass by Value, Pass by Reference, Pass by Value-Result, and Pass by Name.

```
int i;
```

```
int a[2];
```

```
void p(int x, int y)
```

```
{ x++;
```

```
  i++;
```

```
  y++;
```

```
}
```

```
main()
```

```
{ a[0] = 1;
```

```
  a[1] = 1;
```

```
  i = 0;
```

```
  p(a[i], a[i]);
```

```
  printf("%d %d\n", a[0], a[1]);
```

```
  return 0;
```

```
}
```

Pass by Value: 1 1

Pass by Reference: 3 1

Pass by Value-Result: 2 2

Pass by Name: 1 2

5. The following Ada program contains a function parameter.

```
with Text_IO; use Text_IO;
```

```
with Ada.Integer_Text_IO; use Ada.Integer_Text_IO;
```

```
Procedure params is
```

```
  procedure q is
```

```
    type intFunc is access function (n: integer) return integer;
```

```

m: integer := 0;

function f (n: integer) return integer is
begin
    return m + n;
end f;
procedure p (g: intFunc) is
    m: integer := 3;
begin
    put(g(2)); new_line;
end p;

begin
    p(f'access);
end q;

begin
    q;
end params;

```

1. Draw the stack of activation records after the call to g in p.

m	Activation of q
m	Activation of p
g	
n	Activation of f
m	
m	Activation of p
g	

2. What does the program print and why?

2 because f is called from inside of q so it uses q's local value of m which is 0. f returns $m + n = 0 + 2 = 2$. p prints the two and the a newline.