

Application for Supporting Hyperplane Theorem:
Never Best Response and Strictly Dominated
Strategies in Normal Form Game (Recitation 1 & 2)

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Additional knowledge:

Theorem 1 (Supporting Hyperplane Theorem). *Let C be a convex set in \mathbb{R}^n and x_0 be a point on the boundary of C , there exists $a \in \mathbb{R}^n$, s.t. $ax \leq ax_0, \forall x \in C$*

Here begins the application.

Definition 1 (Normal Form Game). A Normal Form Game is a triplet (N, A, u)

- $N = \{1, 2, \dots, n\}$ is a finite set of player
- $A = A_1 \times A_2 \times \dots \times A_n$ and A_i is set of actions of player i
- $u = (u_1, u_2, \dots, u_n)$ and $u_i : A \mapsto \mathbb{R}$ is pay-off function of player i

For any (measurable) space X , let $\Delta(X)$ be the set of probability measures defined on X .

Definition 2 (Strategy). A (mixed) strategy σ_i of player i is a measure defined on A_i ; i.e. $\sigma_i \in \Delta(A_i)$

Let $A_{-i} = A_1 \times A_2 \times \dots \times A_{i-1} \times A_{i+1} \times \dots \times A_n$

Definition 3 (Belief). A belief σ_{-i} of player i is a measure defined on A_{-i} ; i.e. $\sigma_{-i} \in \Delta(A_{-i})$. A belief is independent if $\sigma_{-i} \in \Delta(A_1) \times \Delta(A_2) \dots \times \Delta(A_{i-1}) \times \Delta(A_{i+1}) \times \dots \times \Delta(A_n)$.

Note: This is abuse of notation. Conventionally, $\sigma_{-i} = (\sigma_1, \dots, \sigma_{i-1}, \sigma_{i+1}, \dots, \sigma_n)$ denotes the strategy profile of all players but i . Yet here, it is just the belief of player i . They usually coincides in Nash Equilibrium.

Think of von Neumann–Morgenstern preference over $\Delta(A)$; that is, player i cares about $u_i(\sigma) = \mathbb{E}_\sigma[u_i] = \sum_{a \in A} \sigma(a)u_i(a)$ of strategy profile σ or equivalently, strategy profile σ is better than σ' to player i when $u_i(\sigma) \geq u_i(\sigma')$.

Example 1 (Prisoner Dilemma). $N = 2, A_i = \{C_i, D_i\}$ and u described by

	C_2	D_2
C_1	1,1	-1,2
D_1	2,-1	0,0

Example Strategies of player 1:

- Always cooperates: $\sigma_1(C_1) = 1$, we tend to write C_1
- Always defects: $\sigma_1(D_1) = 1$, we tend to write D_1
- Half-half: $\sigma_1(C_1) = \sigma_1(D_1) = \frac{1}{2}$, we tend to write $\frac{1}{2}C_1 + \frac{1}{2}D_1$

Example Beliefs of player 1:

- Always cooperates: $\sigma_{-1}(C_2) = 1$
- Always defects: $\sigma_{-1}(D_2) = 1$
- Half-half: $\sigma_{-1}(C_2) = \sigma_{-1}(D_2) = \frac{1}{2}$

Some strategy profiles $(C_1, C_2), (\frac{1}{2}C_1 + \frac{1}{2}D_1, D_2), \dots$

Player 1 prefers strategy profile $(\frac{1}{2}C_1 + \frac{1}{2}D_1, D_2)$ than (C_1, C_2) as

$$u_1(\frac{1}{2}C_1 + \frac{1}{2}D_1, C_2) = \frac{1}{2}u_1(C_1, C_2) + \frac{1}{2}u_1(D_1, C_2) = \frac{3}{2}$$

$$u_1(C_1, C_2) = 1$$

If someone else have a belief about what strategy profile will be played in this game, belief $\frac{1}{2}(C_1, C_2) + \frac{1}{2}(D_1, D_2)$ is not independent, while belief $\frac{1}{4}(C_1, C_2) + \frac{1}{4}(C_1, D_2) + \frac{1}{4}(D_1, C_2) + \frac{1}{4}(D_1, D_2) = (\frac{1}{2}C_1 + \frac{1}{2}D_1) \times (\frac{1}{2}C_2 + \frac{1}{2}D_2)$ is independent.

Definition 4 (Best Response). A strategy σ_i of player i is a best response to belief σ_{-i} of player i if for arbitrary strategy σ'_i of player i , $u(\sigma_i, \sigma_{-i}) \geq u(\sigma'_i, \sigma_{-i})$

Definition 5. A strategy σ_i of player i is strictly dominated if there exists a strategy σ'_i , such that $u(\sigma'_i, a_{-i}) > u(\sigma_i, a_{-i}), \forall a_{-i} \in A_{-i}$

Theorem 2. *In a finite (action) game, a strategy is never a best response to any (potentially correlated) belief if and only if it is strictly dominated.*

Proof. “If” side is trivial. If a strategy σ_i is strictly dominated by σ'_i , σ_i is never best response to any belief as σ'_i is always going to be a strictly better response.

“Only if” side is equivalent to the statement “any strategy that is not strictly dominated is a best response to some (potentially correlated) belief”. Suppose a strategy σ_i^0 is not strictly dominated, we are going to construct such a belief to which σ_i^0 is best response.

Consider the following set

$$C = \{x \in \mathbb{R}^{|A_{-i}|} : \exists \sigma_i \in \Delta(A_i), x \leq u_i(\sigma_i) = (u(\sigma_i, a_{-i}))_{a_{-i} \in A_{-i}}\}$$

C is convex. If $x_1, x_2 \in C$, then there exists $\sigma_i^{(1)}, \sigma_i^{(2)} \in \Delta(A_i)$, such that $x_1 \leq u(\sigma_i^{(1)})$, $x_2 \leq u(\sigma_i^{(2)})$. Then,

$$\alpha x_1 + (1 - \alpha)x_2 \leq \alpha u(\sigma_i^{(1)}) + (1 - \alpha)u(\sigma_i^{(2)}) = u(\alpha \sigma_i^{(1)} + (1 - \alpha)\sigma_i^{(2)}), \forall \alpha \in [0, 1]$$

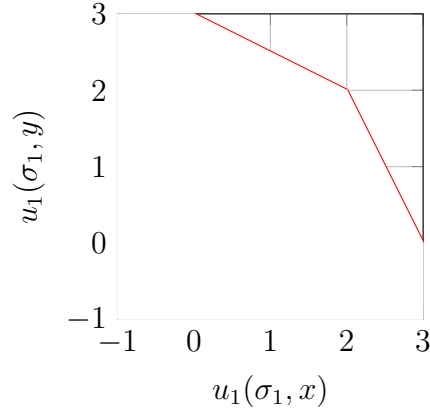
and $\alpha \sigma_i^{(1)} + (1 - \alpha)\sigma_i^{(2)} \in \Delta(A_i)$. That is to say, $\alpha x_1 + (1 - \alpha)x_2 \in C$.

$u(\sigma_i^0)$ is on the boundary of C . Otherwise, $u(\sigma_i^0)$ is in the interior of C . But if it is in the interior of C , then there exists $\varepsilon, x \in \mathbb{R}^{|A_{-i}|}$, where $x \gg 0$ and $u_i(\sigma_i^0) = x + \varepsilon \leq u_i(\sigma_i^x) + \varepsilon$ for some $\sigma_i^x \in \Delta(A_i)$ determined by x . That is to say, σ_i^0 is strictly dominated by σ_i^x . This conflicts with σ_i^0 is not strictly dominated.

Combining the previous two observations, by Supporting Hyperplane Theorem, we can find $a \in \mathbb{R}^{|A_{-i}|}$, such that $au_i(\sigma_i^0) \geq ax, \forall x \in C$. As in every direction, $-\infty$ is in C , we have $a > 0$. The belief can then be constructed by $\frac{a}{\|a\|}$. σ_i^0 is a best response to $\frac{a}{\|a\|}$. \square

Example 2 (Numerical Example for the Proof). Consider a case where $n = 2$, $A_1 = \{a, b, c, d\}$, $A_2 = \{x, y\}$

$u_1(a_1, a_2)$	x	y
a	3	0
b	0	3
c	2	2
d	1	1



White area (extending outside the graph) depicts C . Payoff vectors of strategies that are not strictly dominated lie on the red curve. Belief is a direction on the graph.

Example 3 (Strategies Never Best Response to Independent Belief may not be Strictly Dominated). $N = 2, A_1 = \{U, D\}, A_2 = \{L, R\}, A_3 = \{a, b, c, d\}$ and u_3 is described by

a	L	R	b	L	R	c	L	R	d	L	R
U	9	0	U	0	9	U	0	0	U	6	0
D	0	0	D	9	0	D	0	9	D	0	6

d is not best response to any **independent** belief. Suppose Player 1 uses $pU + (1 - p)D$ and player 2 uses $qL + (1 - q)D$. Need to show d cannot be better than a, b, c at the same time. Expected payoff of using d is $6(pq + (1 - p)(1 - q))$; expected

payoff of using a is $9pq$; expected payoff of using b is $9p(1 - q) + 9q(1 - p)$; expected payoff of using c is $9(1 - p)(1 - q)$. d better than a requires

$$6(pq + (1 - p)(1 - q)) \geq 9pq \Rightarrow \frac{pq}{(1 - p)(1 - q)} \leq 2$$

d better than c requires

$$6(pq + (1 - p)(1 - q)) \geq 9(1 - p)(1 - q) \Rightarrow \frac{pq}{(1 - p)(1 - q)} \geq \frac{1}{2}$$

d better than b requires

$$6(pq + (1 - p)(1 - q)) \geq 9p(1 - q) + 9q(1 - p) \Rightarrow 3\frac{p}{1 - p} + 3\frac{q}{1 - q} - 2\frac{pq}{(1 - p)(1 - q)} \leq 2$$

The last one implies

$$6\sqrt{\frac{pq}{(1 - p)(1 - q)}} - 2\frac{pq}{(1 - p)(1 - q)} \leq 2$$

It can be easily verified the three inequalities cannot hold at the same time.

And d is not strictly dominated. To strictly dominate d , a strategy needs to put more than $\frac{2}{3}$ probability on both a and c , but that is more than 1.