Rima: Building Math Models for Reuse

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Composing Models from Reusable Parts

Reuse: Someone has written a model or model part, and someone would like to use the model or part in a different context

In this talk: You have a model for a *single* knapsack and you would like to extend it to a *multiple* knapsack model (generalised assignment)

We would like to fill a single sack with items of the highest value:

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maximize(sum(i in ITEMS) take(i) * value(i))
sum(i in ITEMS) take(i) * size(i) <= CAPACITY
forall(i in ITEMS) take(i) is_binary</pre>
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Now we'd like to extend the model to cover several sacks:

```
maximize(
   sum(s in SACKS, i in ITEMS) take(s, i) * value(i))
forall(s in SACKS)
   sum(i in ITEMS) take(s, i) * size(i) <= CAPACITY(s)
forall(s in SACKS, i in ITEMS) take(s, i) is_binary
forall(i in ITEMS) sum(s in SACKS) take(s, i) <= 1</pre>
```

Why do we care?

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- We might have a more complex model than a knapsack
- Altering the text of the model makes it hard to re-use
 - to re-use it, we have to understand it enough to modify it
 - it is hard to share improvements between the two models

Why program by hand in five days what you can spend five years of your life automating?

- Terence Parr

What is Rima?

Rima:

- is Yet-Another Math Programming Modelling Language
- focuses on making it easy to construct and re-use models

- is MIT licensed and available at http://rima.googlecode.com/
- is implemented in Lua: http://www.lua.org/
 - a small, fast "scripting" language
- currently binds to CLP, CBC and Ipsolve
- has been submitted to COIN for review

Contents

Algorithms + Data Structures = Programs

- Niklaus Wirth

 $Symbolic\ Expressions\ +\ Structured\ Data\ =\ Reusable\ Model\ Components$

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All the rima.R's are cumbersome, so there is a shortcut:

```
rima.define("a, x, b, y")
e = a * x + b * y
print(e) --> a*x + b*y
```

You can manipulate expressions like references:

Evaluating Expressions

rima. E evaluates expressions by matching references to a table of values:

```
rima.define("a, x, b, y")
e = a * x + b * y
print(rima.E(e, {a=2,x=3,b=4,y=5}))--> 26
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The values you provide as data to rima. E can be other expressions:

```
rima.define("xpos, xneg")
print(rima.E(e, {x=xpos - xneg}) --> a*(xpos - xneg) + b*y
```

A Simple LP (1)

rima.mp.new creates a model, and rima.mp.C builds a constraint:

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rima.define("a, b, x, y")

M = rima.mp.new({
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 x = rima.positive(),
  y = rima.positive()
})
```

A Simple LP (2)

As with expressions, you can print the model:

```
print(M)
--> Maximise:
--> a*x + b*y
--> Subject to:
--> C1: x + 2*y <= 3
--> C2: 2*x + y <= 3
--> 0 <= x <= inf, x real
--> 0 <= y <= inf, x real</pre>
```

A Simple LP (3)

rima.mp.solve takes the model and a table of data and solves:

M encapsulates a complete, symbolic representation of the model

You can index references as if they were arrays:

```
rima.define("X")
e = X[1] + X[2] + X[3]
print(e) --> X[1] + X[2] + X[3]
```

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```
rima.define("X")
e = X[1] + X[2] + X[3]
print(e) --> X[1] + X[2] + X[3]
print(rima.E(e, {X={1,2,3}})) --> 6
```

rima.sum sums an expression over a set:

```
rima.define("x, X")
e = rima.sum{x=X}(x^2)
print(rima.E(e, {X={1,2,3}})) --> 14
```

You can index references as if they were arrays:

rima.sum sums an expression over a set:

You can assign to a whole array at once:

Structures

As well as using arrays, you can also index references as if they were structures:

```
rima.define("item")
mass = item.volume * item.density
print(mass)
   --> item.volume * item.density
print(rima.E(mass, {item={volume=10, density=1.032}}))
   --> 10.32
```

As usual, rima can write the model out for us:

```
print(rima.repr(knapsack, {format="latex"}))
```

In LATEX:

$$\mathbf{maximise} \sum_{i \in \mathsf{items}} i_{\mathsf{take}} i_{\mathsf{value}}$$

subject to

$$\begin{aligned} \text{capacity_limit} : & \sum_{i \in \text{items}} i_{\text{size}} i_{\text{take}} \leq \text{capacity} \\ & \text{items}_{i, \text{take}} \in \{0, 1\} \forall i \in \text{items} \end{aligned}$$

```
ITEMS = {
 camera = { value = 15, size = 2 },
 necklace = { value = 100, size = 20 },
 vase = { value = 15, size = 20 },
 picture = { value = 15, size = 30 },
 tv = { value = 15, size = 40 },
 video = { value = 15, size = 30 },
 chest = { value = 15, size = 60 },
 brick = { value = 1, size = 10 }}
primal = rima.mp.solve("cbc", knapsack,
 {items=ITEMS, capacity=102})
print(primal.objective)
                                 --> 160
print(primal.items.camera.take)
                                --> 1
print(primal.items.vase.take)
                                 --> 1
print(primal.items.brick.take)
                                 --> 0
```

Constraints are Data Too

Suppose, for example, you can only pick one of the camera or the vase.

Constraints, like expressions, are just data, so modelling this is easy:

```
primal = rima.mp.solve("cbc", knapsack,
    {items=ITEMS, capacity=102,
    camera_xor_vase =
        rima.mp.C(items.camera.take + items.vase.take, "<=", 1)})

print(primal.objective) --> 146
print(primal.items.camera.take) --> 1
print(primal.items.vase.take) --> 0
```

What if we want to reuse this constrained model?

Extensible Models

rima.mp.new can take two arguments, the model you want to extend and any extensions to the model:

```
side_constrained_knapsack = rima.mp.new(knapsack, {
   camera_xor_vase =
        rima.mp.C(items.camera.take + items.vase.take, "<=", 1)})

primal = rima.mp.solve("cbc", side_constrained_knapsack,
   {items=ITEMS, capacity=102})

print(primal.objective) --> 146
```

Multiple Sacks

Now we are ready to try a multiple sack model:

```
rima.define("s, sacks")

multiple_sack = rima.mp.new({
    sense = "maximise",
    objective = rima.sum{s=sacks}(s.objective),
    only_take_once[{i=items}] =
        rima.mp.C(rima.sum{s=sacks}(s.items[i].take), "<=", 1)
})</pre>
```

Note that:

- we are treating sacks like a "substructure"
- we have not said anything about what sacks is

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Multiple Knapsacks

We can specify what the knapsack submodel is when we solve:

```
primal = rima.mp.solve("cbc", multiple_sack, {
  items = ITEMS,
  [sacks[s].items] = items,
  sacks = {{capacity=51}, {capacity=51}},
  [sacks[s]] = knapsack})

print(primal.objective) --> 146
```

Sack 1: camera, vase, brick Sack 2: necklace, video

Multiple Constrained Knapsacks

What if we can't carry the camera and vase in the same sack?

```
primal = rima.mp.solve("cbc", multiple_sack, {
  items = ITEMS,
  [sacks[s].items] = items,
  sacks = {{capacity=51}, {capacity=51}},
  [sacks[s]] = side_constrained_knapsack})

print(primal.objective) --> 146
```

Sack 1: camera, picture, brick Sack 2: vase, necklace

Conclusion

We wrote a knapsack model and reused it without any modification:

- in a side-constrained knapsack
- as a part of a multiple knapsack problem
- as a part of a multiple side-constrained knapsack problem

Structured symbolic models enable reuse without alteration:

- we only need to understand the model interface to reuse it
- it is easy to share improvements

Can we "compose" complex models, and is it worthwhile?