

Big O definition

$O(g(n)) = \{f(n) : \text{there exists positive constants } c \text{ and } n_0 \text{ such that } f(n) \leq cg(n) \text{ for all } n \geq n_0\}$

When we say

$$f(n) = O(g(n))$$

we really mean

$$f(n) \in O(g(n))$$

For example

$$n^2 + 3n + 7 = O(n^2)$$

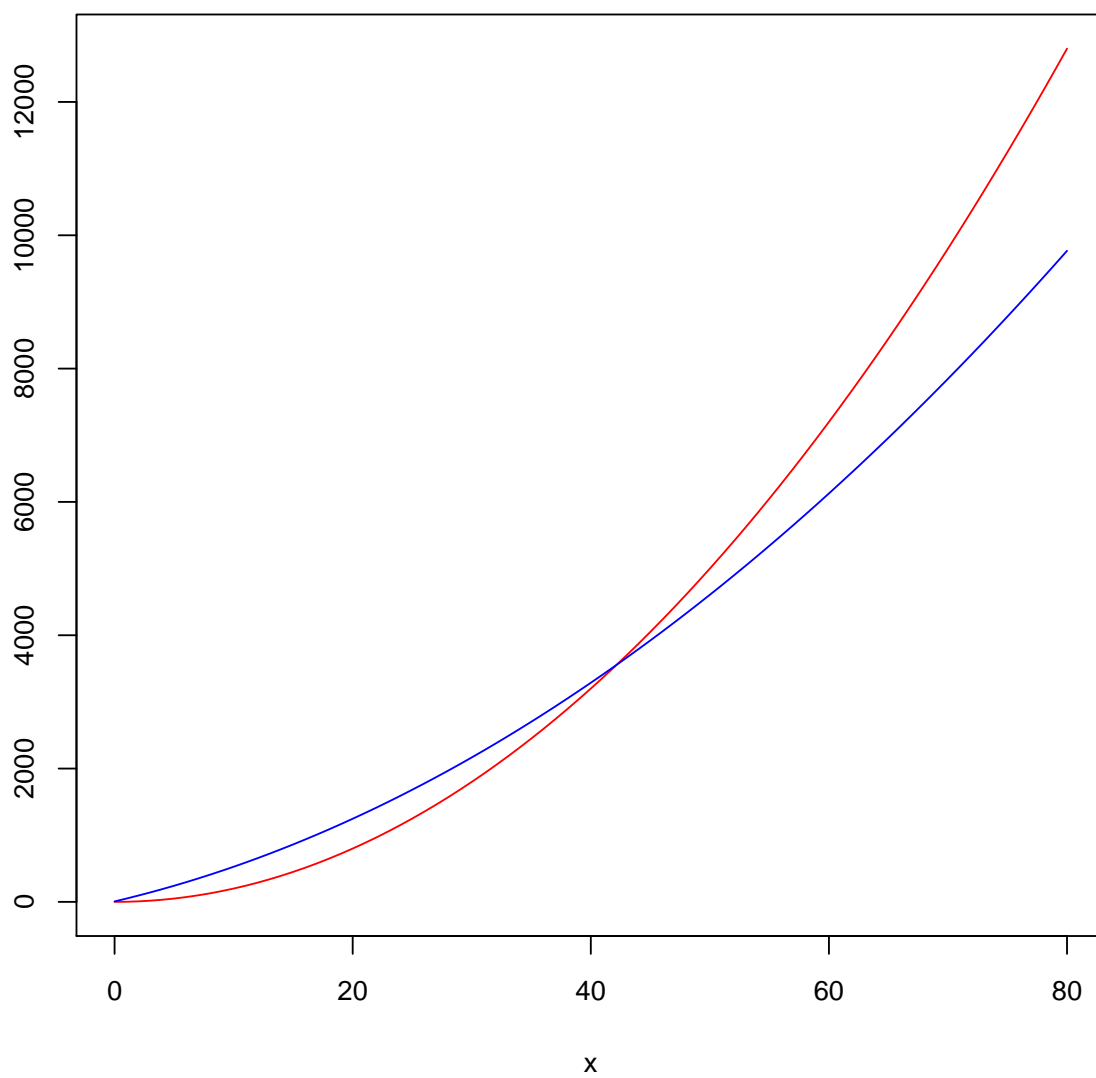
means

$$f(n) = n^2 + 3n + 7$$

is in the set

$$O(n^2)$$

$$n^2 + 42n + 7 = O(n^2)$$
$$n^2 + 42n + 7 \leq 2n^2 \text{ for all } n \geq 50$$



Proof of $n^2 + 42n + 7 = O(n^2)$

$$\begin{aligned} n^2 + 42n + 7 &\leq n^2 + 42n^2 + 7n^2 && \text{for } n \geq 1 \\ &= 50n^2 \end{aligned}$$

So $n^2 + 42n + 7 = O(n^2)$, with $c = 50$ and $n_0 = 1$

Proof of $4n^2 + 5n + 3 = O(n^2)$

$$\begin{aligned} 4n^2 + 5n + 3 &\leq 4n^2 + 5n^2 + 3n^2 & n \geq 1 \\ &= 12n^2 \end{aligned}$$

so $4n^2 + 5n + 3 = O(n^2)$ with $c = 12$ and $n_0 = 1$

Proof of $5n \lg n + 8n - 200 = O(n \lg n)$

Note: if $n \geq 2$ then $\lg n \geq 1$.

$$\begin{aligned} 5n \lg n + 8n - 200 &\leq 5n \lg n + 8n \\ &\leq 5n \lg n + 8n \lg n \quad \text{for } n \geq 2 \\ &\leq 13n \lg n \end{aligned}$$

So

$$5n \lg n + 8n - 200 = O(n \lg n)$$

with $c = 13$ and $n_0 = 2$

Proof of $(n + 5) \lg(3n^2 + 7) = O(n \lg n)$

$$\begin{aligned}(n + 5) \lg(3n^2 + 7) &\leq (n + 5n) \lg(3n^2 + 7n^2) \quad n \geq 1 \\&= 6n \lg(10n^2) \\&\leq 6n \lg(n^3) \quad n \geq 10 \\&= 6n(3 \lg(n)) \\&= 18n \lg(n)\end{aligned}$$

So $(n + 5) \lg(3n^2 + 7) = O(n \lg n)$ for $c = 18$ and $n_0 = 10$

Proof of $(n^2 + 5 \lg n)/(2n + 1) = O(n)$

$$\begin{aligned} \frac{n^2 + 5 \lg n}{2n + 1} &\leq \frac{n^2 + 5n^2}{2n + 1} & n \geq 1 \\ &\leq \frac{n^2 + 5n^2}{2n} \\ &= 3n \end{aligned}$$

So $(n^2 + 5 \lg n)/(2n + 1) = O(n)$ for $c = 3$ and $n_0 = 1$

Useful facts

- For any $a < b$:

$$O(n^a) \subset O(n^b)$$

- For any $a, b > 0, c > 1$:

$$O(a) \subset O(\lg n) \subset O(n^b) \subset O(c^n)$$

You can multiply to find

$$O(an) = O(n) \subset O(n \lg n) \subset O(n^{b+1}) \subset O(nc^n)$$

Asymptotic notation in equations

Right hand side:

$$2n^2 + 3n + 1 = 2n^2 + \Theta(n)$$

means

$$2n^2 + 3n + 1 = 2n^2 + f(n)$$

for *some* $f(n) \in \Theta(n)$.

Left hand side:

$$2n^2 + \Theta(n) = \Theta(n^2)$$

means for *all* $f(n) \in \Theta(n)$,

$$2n^2 + f(n) = \Theta(n^2)$$

We can chain them:

$$\begin{aligned} 2n^2 + 3n + 1 &= 2n^2 + \Theta(n) \\ &= \Theta(n^2) \end{aligned}$$

Relational properties

Reflexive:

$$f(n) = \Theta(f(n))$$

Symmetric:

$$f(n) = \Theta(g(n)) \Leftrightarrow g(n) = \Theta(f(n))$$

Transitive:

$$(f(n) = \Theta(g(n)) \wedge g(n) = \Theta(h(n))) \Rightarrow f(n) = \Theta(h(n))$$

Transpose symmetry:

$$f(n) = O(g(n)) \Leftrightarrow g(n) = \Omega(f(n))$$

$$f(n) = o(g(n)) \Leftrightarrow g(n) = \omega(f(n))$$

Note: We may have two functions f and g such that

$$f(n) \neq O(g(n))$$

$$f(n) \neq \Theta(g(n))$$

$$f(n) \neq \Omega(g(n))$$