

Notes on Red-black Trees

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Red-black trees

- ▶ A variation of binary search trees.
- ▶ **Balanced:** height is $O(\lg n)$, where n is number of nodes.
- ▶ Operations will take $O(\lg n)$ in worst case.

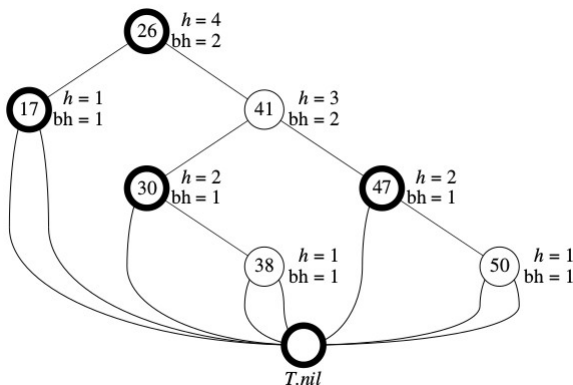
Red-black trees

- ▶ A **red-black tree** is a binary search tree.
- ▶ One bit per node stores an attribute *color*, red or black.
- ▶ All leaves are empty (*nil*) and colored black.
- ▶ We use a sentinel $T.nil$ for all the leaves of a red-black tree T .
- ▶ $T.nil.color$ is black.
- ▶ The root's parent is also $T.nil$.
- ▶ All other attributes (*key*, *left*, *right*, *p*) are inherited.
- ▶ We don't care about $T.nil.key$

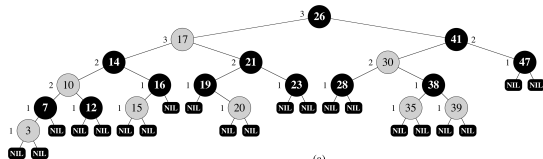
Red-black tree properties

(remember!)

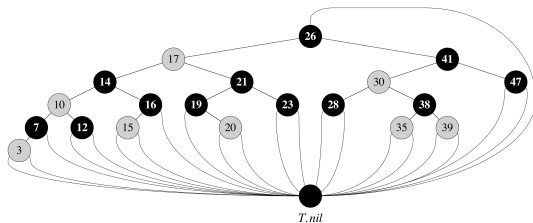
1. Every node is either red or black.
2. The root is black.
3. Every leaf ($T.nil$) is black.
4. If a node is red, then both its children are black.
5. All paths from a node to descendent leaves have same number of black nodes, called **black height**.



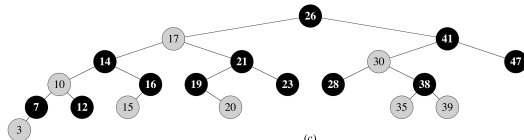
Red-black tree



(a)



(b)



(c)

Height of a red-black tree

- ▶ **Height of a node** is the number of edges in longest path to leaf.
- ▶ **Black-height** of a node x : $bh(x)$ is the number of black nodes (including $T.nil$) on a path from x to a leaf, not counting x .
 - ▶ By property 5, black-height is well defined.
 - ▶ Changing the color of a node does not change its black-height.
 - ▶ Changing the color of a node will change the black-height of its ancestors.

Claim 1:

Any node with height h has black-height $\geq h/2$.

Proof

- ▶ By property 4, $\leq h/2$ nodes on the path from node to a leaf are red.
- ▶ Hence $\geq h/2$ are black.

Claim 2:

The subtree rooted at x contains $\geq 2^{bh(x)} - 1$ internal nodes.

Proof. By induction on height of x .

Basis: Height of $x = 0 \Rightarrow x$ is a leaf and so $bh(x) = 0$, $2^0 - 1 = 0$.

Inductive step:

- ▶ Let the height of x be h .
- ▶ Any child of x has height $h - 1$ and black-height either $bh(x)$ (if the child is red) or $bh(x) - 1$ (if the child is black).
- ▶ By inductive hypothesis, each child has $\geq 2^{bh(x)-1} - 1$ internal nodes.
- ▶ Thus, the subtree rooted at x contains $\geq 2 \cdot (2^{bh(x)-1} - 1) + 1 = 2^{bh(x)} - 1$ internal nodes.

Lemma:

A red-black tree with n internal nodes and height h has

$$h \leq 2 \lg(n + 1)$$

- ▶ Recall proven claims:
 - ▶ Any node with height h has black-height $\geq h/2$.
 - ▶ The subtree rooted at any node x contains $\geq 2^{bh(x)} - 1$ internal nodes.

Proof

Let h and b be the height and black-height of the root, respectively.

By the above two claims,

$$n \geq 2^b - 1 \geq 2^{h/2} - 1$$

Adding 1 to both sides and then taking logs gives

$$\lg(n + 1) \geq h/2$$

which implies that

$$h \leq 2 \lg(n + 1)$$

Operations on red-black trees

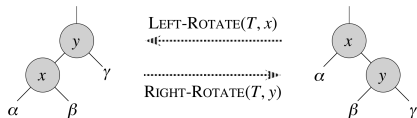
- ▶ MINIMUM, MAXIMUM, SUCCESSOR, PREDECESSOR and SEARCH all run in $O(h) = O(\lg n)$ time.
- ▶ INSERT, what color to make the new node?
 - ▶ Red?
 - ▶ Might violate property 4.
 - ▶ Black?
 - ▶ Might violate property 5.
- ▶ DELETE, what color was the old node?
 - ▶ Red?
 - ▶ Successor might be black.
 - ▶ Black?
 - ▶ Could cause two reds in a row, and violate properties 2 and 5.

Rotations

- ▶ Only pointers are changed.
- ▶ Won't upset binary-search-tree property.
- ▶ Doesn't care about red-black.

LEFT-ROTATE(T, x)

```
1   $y = x.right$ 
2   $x.right = y.left$ 
3  if  $y.left \neq T.nil$ 
4       $y.left.p = x$ 
5   $y.p = x.p$ 
6  if  $x.p == T.nil$ 
7       $T.root = y$ 
8  elseif  $x == x.p.left$ 
9       $x.p.left = y$ 
10 else  $x.p.right = y$ 
11  $y.left = x$ 
12  $x.p = y$ 
```

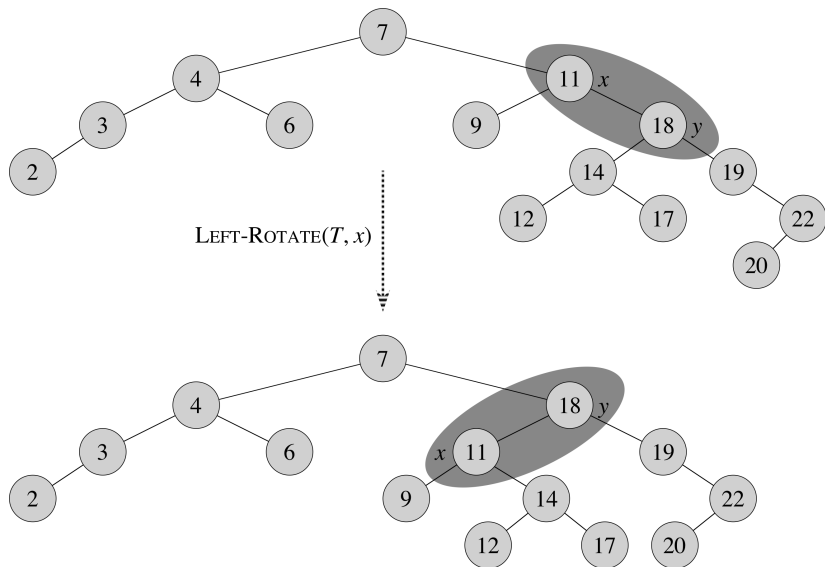


Assumes

- ▶ $x.right \neq T.nil$
- ▶ root's parent is $T.nil$

$\Theta(1)$

Rotations



TREE-INSERT(T, z)

$y = \text{NIL}$

$x = T.\text{root}$

while $x \neq \text{NIL}$

$y = x$

if $z.\text{key} < x.\text{key}$

$x = x.\text{left}$

else $x = x.\text{right}$

$z.p = y$

if $y == \text{NIL}$

$T.\text{root} = z$ *// tree T was empty*

elseif $z.\text{key} < y.\text{key}$

$y.\text{left} = z$

else $y.\text{right} = z$

RB-INSERT(T, z)

$y = T.\text{nil}$

$x = T.\text{root}$

while $x \neq T.\text{nil}$

$y = x$

if $z.\text{key} < x.\text{key}$

$x = x.\text{left}$

else $x = x.\text{right}$

$z.p = y$

if $y == T.\text{nil}$

$T.\text{root} = z$

elseif $z.\text{key} < y.\text{key}$

$y.\text{left} = z$

else $y.\text{right} = z$

$z.\text{left} = T.\text{nil}$

$z.\text{right} = T.\text{nil}$

$z.\text{color} = \text{RED}$

RB-INSERT-FIXUP(T, z)

Insertion fixup

- ▶ Start by doing regular binary-tree insertion.
 - ▶ Color new node red.
 - ▶ May violate red-black tree properties 2 or 4:
1. Every node is either red or black.
OK.
 2. The root is black.
New node might be root.
 3. Every leaf ($T.nil$) is black.
OK.
 4. If a node is red, then both its children are black.
New node's parent might be red.
 5. For each node, all paths from the node to descendant leaves contain the same number of black nodes.
OK.

RB-INSERT-FIXUP(T, z)

while $z.p.color == \text{RED}$

if $z.p == z.p.p.left$

$y = z.p.p.right$

if $y.color == \text{RED}$

$z.p.color = \text{BLACK}$

// case 1

$y.color = \text{BLACK}$

// case 1

$z.p.p.color = \text{RED}$

// case 1

$z = z.p.p$

// case 1

else if $z == z.p.right$

$z = z.p$

// case 2

 LEFT-ROTATE(T, z)

// case 2

$z.p.color = \text{BLACK}$

// case 3

$z.p.p.color = \text{RED}$

// case 3

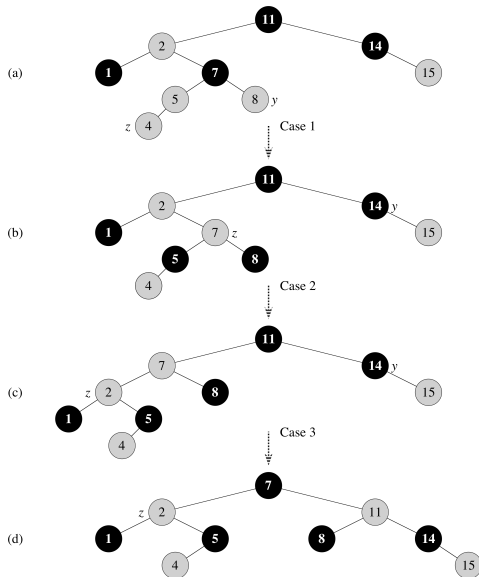
 RIGHT-ROTATE($T, z.p.p$)

// case 3

else (same as **then** clause with “right” and “left” exchanged)

$T.root.color = \text{BLACK}$

Insertion fixup



Case 1, uncle red:

- ▶ make grandpa's children black
- ▶ make grandpa red
- ▶ make grandpa new z
- ▶ loop

Case 2 (optional):

- ▶ rotate left

Case 3:

- ▶ rotate right
- ▶ make parent black

RB-INSERT-FIXUP(T, z)

```
1  while  $z.p.color == RED$ 
2      if  $z.p == z.p.p.left$ 
3           $y = z.p.p.right$ 
4          if  $y.color == RED$ 
5               $z.p.color = BLACK$ 
6               $y.color = BLACK$ 
7               $z.p.p.color = RED$ 
8               $z = z.p.p$ 
9          else if  $z == z.p.right$ 
10              $z = z.p$ 
11             LEFT-ROTATE( $T, z$ )
12              $z.p.color = BLACK$ 
13              $z.p.p.color = RED$ 
14             RIGHT-ROTATE( $T, z.p.p$ )
15      else
16          ( $right \rightleftharpoons left$ )
17   $T.root.color = BLACK$ 
```

Loop Invariant: at the start of each iteration of the **while** loop:

- ▶ z is red
- ▶ If $z.p$ is the root, then $z.p$ is black.
- ▶ There is at most one red-black violation:
 - ▶ z is a red root.
 - ▶ z and $z.p$ are both red.

Loop invariant

Loop Invariant: at the start of each iteration of the **while** loop:

1. z is red
2. If $z.p$ is the root, then $z.p$ is black.
3. There is at most one red-black violation:
 - 3.1 z is a red root.
 - 3.2 z and $z.p$ are both red.

Initialization:

- ▶ 1 because we set it that way.
- ▶ 2 because it was root to begin with.
- ▶ 3.1 possible if z is new root.
- ▶ 3.2 possible if $z.p$ was red.

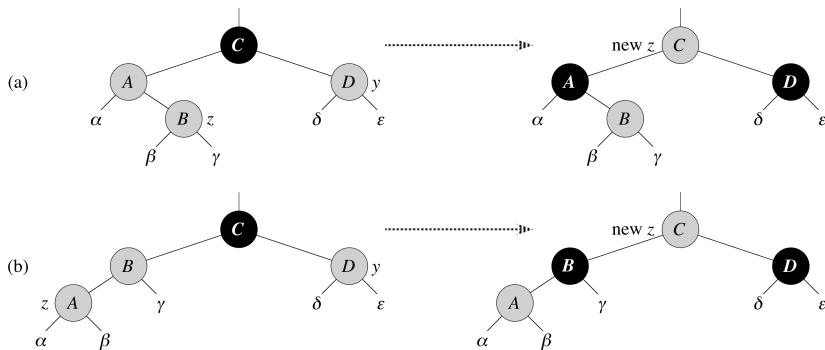
Termination:

- ▶ 3.1 last line fixes red root.
- ▶ 3.2 loop only terminates when $z.p$ is black.

Maintenance, Case 1, uncle is red

Loop Invariant: at the start of each iteration of the **while** loop:

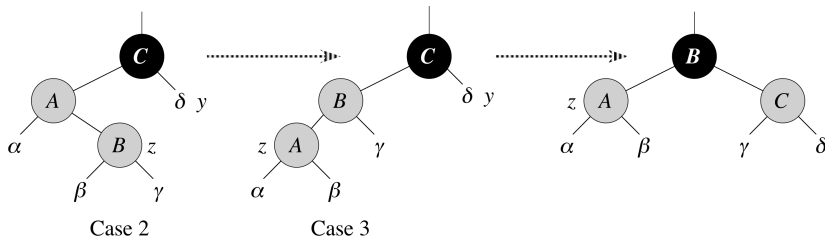
1. z is red
2. If $z.p$ is the root, then $z.p$ is black.
3. There is at most one red-black violation:
 - 3.1 z is a red root.
 - 3.2 z and $z.p$ are both red.



Maintenance, Case 2&3, uncle is black:

Loop Invariant: at the start of each iteration of the **while** loop:

1. z is red
2. If $z.p$ is the root, then $z.p$ is black.
3. There is at most one red-black violation:
 - 3.1 z is a red root.
 - 3.2 z and $z.p$ are both red.



Analysis

- ▶ $O(\lg n)$ time to insert into binary tree.
- ▶ Fixup also $O(\lg n)$:
 - ▶ Each pass through the loop takes $O(1)$ time.
 - ▶ Each iteration moves z up two levels or stops.
 - ▶ $O(\lg n)$ levels.
 - ▶ Also note that there are at most 2 rotations overall.
- ▶ Insertion into red-black tree is $O(\lg n)$.

RB-TRANSPLANT

TRANSPLANT(T, u, v)

```
if  $u.p == \text{NIL}$   
     $T.\text{root} = v$   
elseif  $u == u.p.\text{left}$   
     $u.p.\text{left} = v$   
else  $u.p.\text{right} = v$   
if  $v \neq \text{NIL}$   
     $v.p = u.p$ 
```

RB-TRANSPLANT(T, u, v)

```
if  $u.p == T.\text{nil}$   
     $T.\text{root} = v$   
elseif  $u == u.p.\text{left}$   
     $u.p.\text{left} = v$   
else  $u.p.\text{right} = v$   
     $v.p = u.p$ 
```

TREE-DELETE(T, z)

```

1  if  $z.left == \text{NIL}$ 
2      TRANSPLANT( $T, z, z.right$ )
3  elseif  $z.right == \text{NIL}$ 
4      TRANSPLANT( $T, z, z.left$ )
5  else
6       $y = \text{TREE-MINIMUM}(z.right)$ 
7      if  $y.p \neq z$ 
8          TRANSPLANT( $T, y, y.right$ )
9           $y.right = z.right$ 
10          $y.right.p = y$ 
11         TRANSPLANT( $T, z, y$ )
12          $y.left = z.left$ 
13          $y.left.p = y$ 

```

Differences:

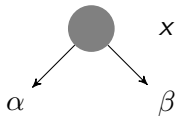
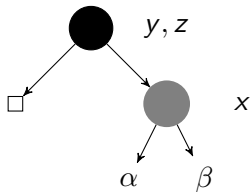
- ▶ y is either z or the node moved
- ▶ save y 's color
- ▶ x is node that moves into y 's original position
- ▶ set $x.p$ to original position of $y.p$
 - ▶ in $x.p = y$, or
 - ▶ in RB-TRANSPLANT
- ▶ RB-DELETE-FIXUP if y was black

RB-DELETE(T, z)

```

 $y = z$ 
 $y\text{-original-color} = y.color$ 
if  $z.left == T.nil$ 
     $x = z.right$ 
    RB-TRANSPLANT( $T, z, z.right$ )
elseif  $z.right == T.nil$ 
     $x = z.left$ 
    RB-TRANSPLANT( $T, z, z.left$ )
else  $y = \text{TREE-MINIMUM}(z.right)$ 
     $y\text{-original-color} = y.color$ 
     $x = y.right$ 
    if  $y.p == z$ 
         $x.p = y$ 
    else RB-TRANSPLANT( $T, y, y.right$ )
         $y.right = z.right$ 
         $y.right.p = y$ 
    RB-TRANSPLANT( $T, z, y$ )
     $y.left = z.left$ 
     $y.left.p = y$ 
     $y.color = z.color$ 
if  $y\text{-original-color} == \text{BLACK}$ 
    RB-DELETE-FIXUP( $T, x$ )

```



RB-DELETE(T, z)

$y = z$

$y\text{-original-color} = y.\text{color}$

if $z.\text{left} == T.\text{nil}$

$x = z.\text{right}$

RB-TRANSPLANT($T, z, z.\text{right}$)

elseif $z.\text{right} == T.\text{nil}$

$x = z.\text{left}$

RB-TRANSPLANT($T, z, z.\text{left}$)

else $y = \text{TREE-MINIMUM}(z.\text{right})$

$y\text{-original-color} = y.\text{color}$

$x = y.\text{right}$

if $y.p == z$

$x.p = y$

else RB-TRANSPLANT($T, y, y.\text{right}$)

$y.\text{right} = z.\text{right}$

$y.\text{right}.p = y$

RB-TRANSPLANT(T, z, y)

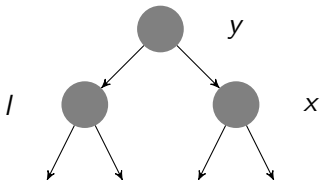
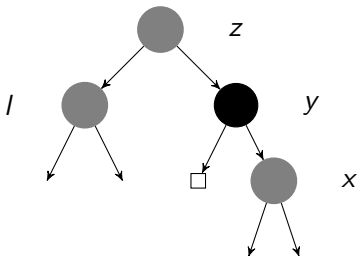
$y.\text{left} = z.\text{left}$

$y.\text{left}.p = y$

$y.\text{color} = z.\text{color}$

if $y\text{-original-color} == \text{BLACK}$

RB-DELETE-FIXUP(T, x)



RB-DELETE(T, z)

$y = z$

$y\text{-original-color} = y.\text{color}$

if $z.\text{left} == T.\text{nil}$

$x = z.\text{right}$

RB-TRANSPLANT($T, z, z.\text{right}$)

elseif $z.\text{right} == T.\text{nil}$

$x = z.\text{left}$

RB-TRANSPLANT($T, z, z.\text{left}$)

else $y = \text{TREE-MINIMUM}(z.\text{right})$

$y\text{-original-color} = y.\text{color}$

$x = y.\text{right}$

if $y.p == z$

$x.p = y$

else RB-TRANSPLANT($T, y, y.\text{right}$)

$y.\text{right} = z.\text{right}$

$y.\text{right}.p = y$

RB-TRANSPLANT(T, z, y)

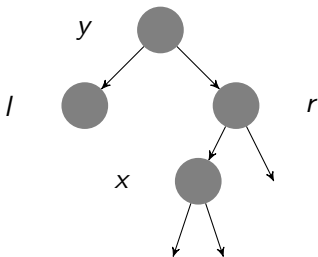
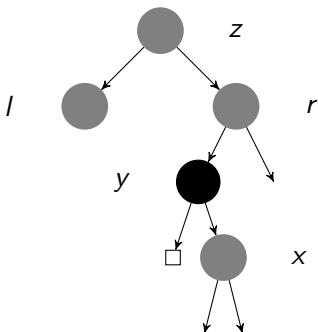
$y.\text{left} = z.\text{left}$

$y.\text{left}.p = y$

$y.\text{color} = z.\text{color}$

if $y\text{-original-color} == \text{BLACK}$

RB-DELETE-FIXUP(T, x)



RB-DELETE(T, z)

$y = z$

$y\text{-original-color} = y.\text{color}$

if $z.\text{left} == T.\text{nil}$

$x = z.\text{right}$

RB-TRANSPLANT($T, z, z.\text{right}$)

elseif $z.\text{right} == T.\text{nil}$

$x = z.\text{left}$

RB-TRANSPLANT($T, z, z.\text{left}$)

else $y = \text{TREE-MINIMUM}(z.\text{right})$

$y\text{-original-color} = y.\text{color}$

$x = y.\text{right}$

if $y.p == z$

$x.p = y$

else RB-TRANSPLANT($T, y, y.\text{right}$)

$y.\text{right} = z.\text{right}$

$y.\text{right}.p = y$

RB-TRANSPLANT(T, z, y)

$y.\text{left} = z.\text{left}$

$y.\text{left}.p = y$

$y.\text{color} = z.\text{color}$

if $y\text{-original-color} == \text{BLACK}$

RB-DELETE-FIXUP(T, x)

What violations could occur if y was red?

1. Every node red or black. OK.
2. Root is black. y could not have been the root, OK.
3. All leaves are black. OK.
4. No red nodes became adjacent (check). OK.
5. No black heights have changed. OK.

What violations could occur if y was black?

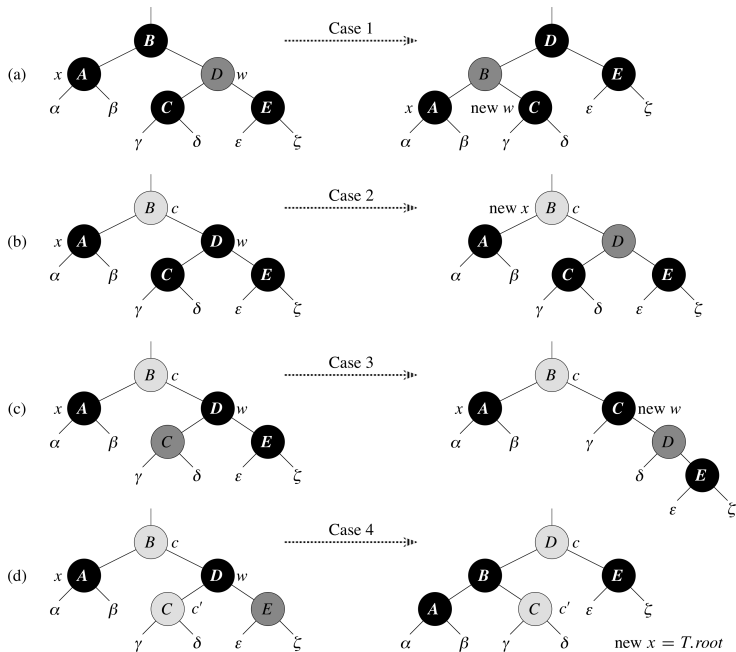
1. Every node is either red or black. Still OK.
2. The root is black.
Violation if y is the root and x is red.
3. Every leaf ($T.nil$) is black. Still OK.
4. If a node is red, then both its children are black.
Violation if $x.p$ and x are both red.
5. All paths from a node to descendent leaves have same number of black nodes.

Any paths containing y now have 1 fewer black.

- ▶ Correct by giving x an “extra black.”
- ▶ Add 1 to count of black nodes on paths containing x .
- ▶ Now 5 is OK, but 1 is not, and possibly 3.
- ▶ x is either **black&black** or **red&black**
- ▶ $x.color$ is still either RED or BLACK
- ▶ x just marks the node with a black to move somewhere.

RB-DELETE-FIXUP

- ▶ Idea: move the extra black up the tree until
 - ▶ x points to a **red&black** node \implies turn it black,
 - ▶ This also handles problem of x and $x.p$ both red.
 - ▶ x points to the root \implies just remove the extra black, or
 - ▶ we can do rotations and recolorings to finish.
- ▶ Within the **while** loop:
 - ▶ x always points to a nonroot doubly black node
 - ▶ w is x 's sibling
 - ▶ w cannot be $T.nil$ since that would violate 5 at $x.p$.
- ▶ Four cases when x is a left child:
 - ▶ Case 1: w is red
 - ▶ Case 2: w 's children both black
 - ▶ Case 3: w 's left child red, right child black
 - ▶ Case 4: w 's right child red



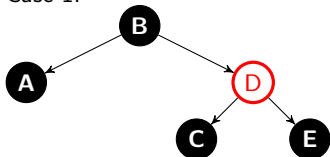
RB-DELETE-FIXUP(T, x)

```

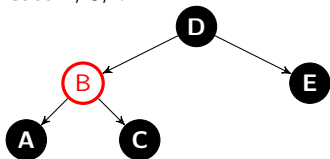
while  $x \neq T.root$  and  $x.color == BLACK$ 
    if  $x == x.p.left$ 
         $w = x.p.right$ 
        if  $w.color == RED$ 
             $w.color = BLACK$  // case 1
             $x.p.color = RED$  // case 1
            LEFT-ROTATE( $T, x.p$ ) // case 1
             $w = x.p.right$  // case 1
        if  $w.left.color == BLACK$  and  $w.right.color == BLACK$ 
             $w.color = RED$  // case 2
             $x = x.p$  // case 2
        else if  $w.right.color == BLACK$ 
             $w.left.color = BLACK$  // case 3
             $w.color = RED$  // case 3
            RIGHT-ROTATE( $T, w$ ) // case 3
             $w = x.p.right$  // case 3
             $w.color = x.p.color$  // case 4
             $x.p.color = BLACK$  // case 4
             $w.right.color = BLACK$  // case 4
            LEFT-ROTATE( $T, x.p$ ) // case 4
             $x = T.root$  // case 4
    else (same as then clause with “right” and “left” exchanged)
 $x.color = BLACK$ 

```

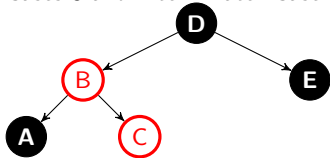
Case 1:



Case 2, 3, or 4:



Cases 3 and 4 terminate. Case 2:



Also terminates, because x is now red.

RB-DELETE-FIXUP(T, x)

while $x \neq T.root$ and $x.color == BLACK$

if $x == x.p.left$

$w = x.p.right$

if $w.color == RED$

$w.color = BLACK$

$x.p.color = RED$

 LEFT-ROTATE($T, x.p$)

$w = x.p.right$

if $w.left.color == BLACK$ and $w.right.color == BLACK$

$w.color = RED$

$x = x.p$

else if $w.right.color == BLACK$

$w.left.color = BLACK$

$w.color = RED$

 RIGHT-ROTATE(T, w)

$w = x.p.right$

$w.color = x.p.color$

$x.p.color = BLACK$

$w.right.color = BLACK$

 LEFT-ROTATE($T, x.p$)

$x = T.root$

else (same as **then** clause with “right” and “left” exchanged)

$x.color = BLACK$

Possible paths

- ▶ Case 1
Case 2
Terminate
- ▶ Case 1
Case 3
Case 4
Terminate
- ▶ Case 1
Case 4
Terminate
- ▶ Case 2
Loop
- ▶ Case 3
Case 4
Terminate
- ▶ Case 4
Terminate

RB-DELETE-FIXUP(T, x)

while $x \neq T.root$ and $x.color == BLACK$

if $x == x.p.left$

$w = x.p.right$

if $w.color == RED$

$w.color = BLACK$

$x.p.color = RED$

 LEFT-ROTATE($T, x.p$)

$w = x.p.right$

if $w.left.color == BLACK$ and $w.right.color == BLACK$

$w.color = RED$

$x = x.p$

else if $w.right.color == BLACK$

$w.left.color = BLACK$

$w.color = RED$

 RIGHT-ROTATE(T, w)

$w = x.p.right$

$w.color = x.p.color$

$x.p.color = BLACK$

$w.right.color = BLACK$

 LEFT-ROTATE($T, x.p$)

$x = T.root$

else (same as **then** clause with “right” and “left” exchanged)

$x.color = BLACK$

Analysis

- ▶ $O(\lg n)$ time for RB-DELETE up to the call of RB-DELETE-FIXUP
- ▶ Within RB-DELETE-FIXUP:
 - ▶ Case 2 is the only case in which more iterations occur:
 - ▶ x moves up 1 level
 - ▶ Hence $O(\lg n)$ iterations
 - ▶ Each of cases 1, 3, and 4 has 1 rotation $\implies \leq 3$ rotations in all
 - ▶ Hence, $O(\lg n)$ time.
- ▶ Note:
- ▶ Red-black trees use at most a constant number of rotations.
- ▶ AVL trees in worst case use $\Omega(\lg n)$ rotations.