Master Theorem

Let $a \ge 1$ and b > 1 be constants, let f(n) be a function, and let T(n) be defined on nonnegative integers by

$$T(n) = aT\left(\frac{n}{b}\right) + f(n)$$

Then T(n) has the following asymptotic bounds:

$$\begin{array}{c|c} f(n) & T(n) \\ \hline O(n^{\log_b a - \epsilon}) & \Theta(n^{\log_b a}) \\ \Theta(n^{\log_b a}) & \Theta(f(n) \lg n) \\ O(n^{\log_b a + \epsilon}) & \Theta(f(n)) \end{array}$$

The last one only if $af(\frac{n}{b}) \le cf(n)$ for some c < 1 and large enough n.

Master theorem examples

$$T(n) = aT\left(\frac{n}{b}\right) + f(n)$$

$$O(n^{\log_b a - \epsilon}) \Leftrightarrow \Theta(n^{\log_b a})$$

$$\Theta(n^{\log_b a}) \Leftrightarrow \Theta(f(n) | g | n)$$

$$\Omega(n^{\log_b a + \epsilon}) \Leftrightarrow \Theta(f(n))$$

$$T(n) = 2T\left(\frac{n}{2}\right) + n$$

$$\log_b a = 1$$

$$f(n) = \Theta(n^{\log_b a})$$

$$T(n) = \Theta(n \lg n)$$

Master theorem does not apply

$$T(n) = 2T(n/2) + n \lg n$$
 $n^{\log_b a} = n$
 $f(n) = n \lg n$
 $f(n) = \Omega(n^{\log_b a}) = \Omega(n)$
 $f(n) \neq \Omega(n^{\log_b a + \epsilon}) = \Omega(n^{1+\epsilon})$
 $n \lg n \neq \Omega(n^{1.000000000001})$
 $\lg n \neq \Omega(n^{0.000000000001})$

Master theorem does apply

$$T(n) = 4T(n/3) + n \lg n$$

$$n^{\log_b a} = n^{\log_c 4} = n^{1.26...}$$

$$f(n) = n \lg n$$

$$f(n) = O(n^{\log_b a}) = O(n^{1.26...})$$

$$f(n) = O(n^{\log_b a - \epsilon}) = O(n^{1.2})$$

$$n \lg n = O(n^{1.2})$$

$$\lg n = O(n^{0.2})$$