Simple Static Typechecking

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- Eight cats



Type Checking

The idea is to *analyze* a program *statically* to find errors that might occur when you run the program. For example:

• Ensure that exp_1 and exp_2 are both type Int in every case of

$$exp_1 + exp_2$$

ullet Ensure that exp_1 and exp_2 are both type Bool in every case of

$$exp_1$$
 Or exp_2

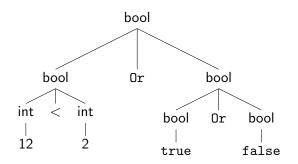
ullet Ensure that exp_1 is type Bool in every case of

If exp_1 Then exp_2 Else exp_3



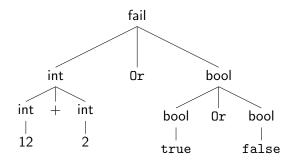
Recursive type checking

(12 < 2) Or (true Or false)



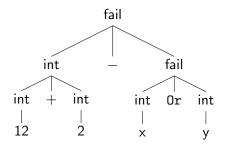
Recursive type checking

(12 + 2) Or (true Or false)



Recursive type checking

$$(12 + 2) - (x Or y)$$



The language of Homework 5 has nothing to check!

- The parser ensures that only Bool expressions occur in IF and WHILE
- Only Int expressions can occur in assignments to variables
- Only numbers, variables and arithmetic operators can occur on the RHS of assignments
- No program with a type error can even be parsed!
- My original idea was to add boolean variables, but I got a much better idea.

$$[\underbrace{\begin{array}{c} x & x + 1 \\ \text{function} \end{array}}_{\text{argument}}]$$

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• [
$$\x \{ x + 1 \} (3 * 5)] \Rightarrow$$

$$[\underbrace{\begin{array}{c} \left(x + 1\right) \\ \text{function} \end{array}}_{\text{argument}}]$$
application

• [
$$\x \{ x + 1 \} (3 * 5) \] \Rightarrow 16$$

$$[\underbrace{\begin{array}{c} \left\{\begin{array}{c} x & \left\{\begin{array}{c} x + 1 \right\} \\ \text{function} \end{array}\right]}_{\text{application}}$$

- [$\x \{ x + 1 \} (3 * 5) \] \Rightarrow 16$
- [\f { [f 3] } \x { 2 * x }] \Rightarrow

$$[\underbrace{\begin{array}{c} \left\{\begin{array}{c} x & \left\{\begin{array}{c} x + 1 \right\} \\ \text{function} \end{array}\right]}_{\text{application}}$$

- [$\x \{ x + 1 \} (3 * 5) \] \Rightarrow 16$
- [\f { [f 3] } \x { 2 * x }] \Rightarrow 6

$$[\underbrace{\begin{array}{c} x & x + 1 \\ \text{function} \end{array}}_{\text{argument}}]$$

- [$\x \{ x + 1 \} (3 * 5) \] \Rightarrow 16$
- [\f { [f 3] } \x { 2 * x }] \Rightarrow 6
- [$\x \{ \y \{ x + y \} \} 3] \Rightarrow$

$$[\underbrace{\begin{array}{c} \left\{\begin{array}{c} x & \left\{\begin{array}{c} x + 1 \right\} \\ \text{function} \end{array}\right]}_{\text{application}}$$

- [$\x \{ x + 1 \} (3 * 5) \] \Rightarrow 16$
- [\f { [f 3] } \x { 2 * x }] \Rightarrow 6
- [$\x { \y { x + y} } 3] \Rightarrow \y { 3 + y }$

$$[\underbrace{\begin{array}{c} x & x + 1 \\ \text{function} \end{array}}_{\text{argument}}]$$

- $[\ \ \ \ \ \ \ \ \ \ \ \ \ \] \Rightarrow 16$
- [\f { [f 3] } \x { 2 * x }] \Rightarrow 6
- [$\x { y { x + y} } 3] \Rightarrow \y { 3 + y }$
- \x { x + 1 } ::

$$[\underbrace{\begin{array}{c} \left(x + 1\right) \\ \text{function} \end{array}}_{\text{application}} \underbrace{(3 * 5)}_{\text{argument}}]$$

- [$\x \{ x + 1 \} (3 * 5) \] \Rightarrow 16$
- [\f { [f 3] } \x { 2 * x }] \Rightarrow 6
- [$\x \{ \y \{ x + y \} \} 3] \Rightarrow \y \{ 3 + y \}$
- \x { x + 1 } :: (int -> int)

$$[\underbrace{\begin{array}{c} x & x + 1 \\ \text{function} \end{array}}_{\text{argument}}]$$

- $[\ \ \ \ \ \ \ \ \ \ \ \ \ \] \Rightarrow 16$
- [\f { [f 3] } \x { 2 * x }] \Rightarrow 6
- [$\x { y { x + y} } 3] \Rightarrow \y { 3 + y }$
- \x { x + 1 } :: (int -> int)
- \x { \y { x + y } } ::

$$\underbrace{ \left[\begin{array}{c} x & x + 1 \\ \text{function} \end{array} \right]}_{\text{argument}}$$

- [$\x \{ x + 1 \} (3 * 5) \] \Rightarrow 16$
- [\f { [f 3] } \x { 2 * x }] \Rightarrow 6
- [$\x { \y { x + y} } 3] \Rightarrow \y { 3 + y }$
- \x { x + 1 } :: (int -> int)
- \x { \y { x + y } } :: (int -> (int -> int))

$$[\underbrace{\begin{array}{c} \left\{\begin{array}{c} x & \left\{\begin{array}{c} x + 1 \right\} \\ \end{array}\right\}}_{\text{function}} \underbrace{\begin{array}{c} (3 * 5) \\ \text{argument} \end{array}}]$$

- [$\x \{ x + 1 \} (3 * 5) \] \Rightarrow 16$
- [\f { [f 3] } \x { 2 * x }] \Rightarrow 6
- [$\x { y { x + y} } 3] \Rightarrow \y { 3 + y }$
- \x { x + 1 } :: (int -> int)
- \x { \y { x + y } } :: (int -> (int -> int))
- \f { [f 3] } ::

$$[\underbrace{\begin{array}{c} \left\{\begin{array}{c} x & \left\{\begin{array}{c} x + 1 \right\} \\ \text{function} \end{array}\right]}_{\text{application}}$$

- [$\x \{ x + 1 \} (3 * 5) \] \Rightarrow 16$
- [\f { [f 3] } \x { 2 * x }] \Rightarrow 6
- [$\x { \y { x + y} } 3] \Rightarrow \y { 3 + y }$
- \x { x + 1 } :: (int -> int)
- \x { \y { x + y } } :: (int -> (int -> int))
- \f { [f 3] } :: ((int -> ??) -> ??)

My syntax:

$$\underbrace{ \left[\begin{array}{c} x & x + 1 \\ \text{function} \end{array} \right]}_{\text{argument}}$$

- [$\x \{ x + 1 \} (3 * 5) \] \Rightarrow 16$
- [\f { [f 3] } \x { 2 * x }] \Rightarrow 6
- [$\x { y { x + y} } 3] \Rightarrow \y { 3 + y }$
- \x { x + 1 } :: (int -> int)
- \x { \y { x + y } } :: (int -> (int -> int))
- \f { [f 3] } :: ((int -> ??) -> ??)

To do static recursive typechecking, we need **typed lambda expressions**.



$$[\underbrace{x:int \{x+1\}}_{\text{function}} \underbrace{(3*5)}_{\text{argument}}]$$

My syntax:

$$\underbrace{\left[\begin{array}{c} x: \text{int } \left\{\begin{array}{c} x+1 \right\} \underbrace{\left(3*5\right)}_{\text{argument}} \right]}_{\text{application}}$$

• \f:(int -> int) { [f 3] } ::

$$[\underbrace{x:int \{x+1\}}_{\text{function}} \underbrace{(3*5)}_{\text{argument}}]$$

- \f:(int -> int) { [f 3] } :: ((int -> int) -> int)
- \f:(int -> (int -> int)) { [f 3] } ::

$$[\underbrace{x:int \{x+1\}}_{\text{function}} \underbrace{(3*5)}_{\text{argument}}]$$

- \f:(int -> int) { [f 3] } :: ((int -> int) -> int)
- \f:(int -> (int -> int)) { [f 3] } :: ((int -> (int -> int)) -> (int -> int))
- \f((int -> int) -> int) { [f x] } ::

$$[\underbrace{x:int \{x+1\}}_{\text{function}} \underbrace{(3*5)}_{\text{argument}}]$$

- \f:(int -> int) { [f 3] } :: ((int -> int) -> int)
- \f:(int -> (int -> int)) { [f 3] } :: ((int -> (int -> int)) -> (int -> int))
- \f((int -> int) -> int) { [f x] } :: (((int -> int) -> int) -> int)

Typed lambda expressions and applications in our language

```
f := \x:int {x+1}; x := [f 3]

g := [\x:int {\y:int {x+y}} 4]; x := [g 5]

j := \f:(int->int){\y:int{ [f y] };

k := [j \x:int { 2*x }];

x := [k 5]

x := [[\x:int { \y:int {x + y} } 13] 12]
y := [\x:int { [\y:int {x + y} 13] } 12]
```

Typed lambda expressions and applications in our language

```
f := \x: int \{x+1\}; x := [f 3]
g := [\x:int {\y:int {x+y}} 4]; x := [g 5]
j := \f:(int->int){\y:int{ [f y] };
k := [j \x:int { 2*x }];
x := [k 5]
x := [[\x:int { \y:int {x + y} } 13] 12]
y := [\x: int { [\y: int {x + y} 13] } 12]
```

Are all of these type safe?

Typed lambda expressions and applications in our language

```
f := \x:int {x+1}; x := [f 3]

g := [\x:int {\y:int {x+y}} 4]; x := [g 5]

j := \f:(int->int){\y:int{ [f y] };

k := [j \x:int { 2*x }];

x := [k 5]

x := [[\x:int { \y:int {x + y} } 13] 12]
y := [\x:int { [\y:int {x + y} 13] } 12]
```

Are all of these type safe?

You will learn how to add lambdas to your interpreters very soon.

Today we'll just do a typechecker for them.



Typed recursive functions

What more type information do you need to typecheck recursive functions?

REC f =
$$\x$$
 { [f (x+1)] } IN [f 3] END

Typed recursive functions

What more type information do you need to typecheck recursive functions?

REC f =
$$\x$$
 { [f (x+1)] } IN [f 3] END

REC f = int
$$x:int { [f (x+1)] } IN [f 3] END$$

Two new types of Arithmetic Expressions for Parsing

```
data AExp =
   Var String
| Num Int
| Plus AExp AExp
| Times AExp AExp
| Neg AExp
| Div AExp AExp
| Lambda String Type AExp --NEW
| App AExp AExp
deriving (Show, Eq)
```

- We won't talk about the parsing today.
- AExp are not always numbers anymore!



Example typechecking

Expressions

```
[x:int {99} 22] + 3
                                                     --accept
[x:(int->int) \{99\} 22]
                                                     --reject
[\x:(int->int) \{[x 2]\} 9]
                                                     --reject
\lceil x:(int->int) \rceil \lceil x 2\rceil \rceil \ x:int \lceil 9\rceil \rceil
                                                     --accept
[\x:(int->int) \{x + 2\} \x:int\{9\}]
                                                     --reject
Programs
x := \x: int{x+1}:
IF x < 2 THEN y := 3 ELSE y := 4;
                                                     --reject
x := 3;
z := [\y:int \{x + y\} 5];
w := x + y
                                                     --reject
```

A type for types

```
x := 5; y := \x:int \{x+x\}; z := \x:(int->int) \{ [x 4] \}
```

```
typeCheckStmt :: (Stmt AExp BExp) -> TypeStore
                              -> (Bool, TypeStore)
typeCheckStmt program store =
 case program of
  Seq s1 s2 ->
    let (s1Good, store') = typeCheckStmt s1 store
    in if not s1Good
       then (False, store')
       else let (s2Good, store'') = typeCheckStmt s2 store'
            in (s2Good, store'')
```

```
x := 5; y := \w:int \{x+w\}; z := [y x]
```



```
typeCheckStmt :: (Stmt AExp BExp) -> TypeStore
                               -> (Bool, TypeStore)
typeCheckStmt program store =
 case program of
  If b s1 s2 \rightarrow
    let (bType, store') = findTypeBExp b store in
    if bType /= BoolType
    then (False, store')
    else let (s1Good, store'') = typeCheckStmt s1 store' in
       if not s1Good
       then (False, store'')
       else let (s2Good, store'') = typeCheckStmt s1 store'
       in (s2Good, store',')
       -- Do you see a problem here?
```

```
typeCheckStmt :: (Stmt AExp BExp) -> TypeStore
                               -> (Bool, TypeStore)
typeCheckStmt program store =
 case program of
  If b s1 s2 \rightarrow
    let (bType, store') = findTypeBExp b store in
    if bType /= BoolType
    then (False, store')
    else let (s1Good, store'') = typeCheckStmt s1 store' in
       if not s1Good
       then (False, store'')
       else let (s2Good, store'') = typeCheckStmt s1 store'
       in (s2Good, store'')
       -- Do you see a problem here?
IF a < b THEN z := 3 ELSE z := \x: int\{x+1\} END; w := z + 1
```

```
typeCheckStmt :: (Stmt AExp BExp) -> TypeStore
                              -> (Bool, TypeStore)
typeCheckStmt program store =
 case program of
    While b s ->
      let (bType, store') = findTypeBExp b store in
      if bType /= BoolType
      then (False, store')
      else let (sGood, store'') = typeCheckStmt s store' in
          (sGood, store'')
```

```
typeCheckStmt :: (Stmt AExp BExp) -> TypeStore
                              -> (Bool, TypeStore)
typeCheckStmt program store =
 case program of
    Assign x val ->
      let (valType, store') = findTypeAExp val store
      in if valType /= FailureType
         then (True, Map.insert x valType store')
         else (False, store)
```

Typechecking Boolean Expressions

Typechecking Boolean Expressions

```
findTypeBExp :: (BExp AExp) -> TypeStore
                    -> (Type, TypeStore)
findTypeBExp b store =
  case b of
    . . .
    Or x y ->
    let (t, store') = (findTypeBExp x store) in
    if t /= BoolType
    then (FailureType, store')
    else let (t', store'') = (findTypeBExp y store') in
                if t' /= BoolType
                then (FailureType, store'')
                else (BoolType, store'')
```

Typechecking Boolean Expressions

```
findTypeBExp :: (BExp AExp) -> TypeStore
                         -> (Type, TypeStore)
findTypeBExp b store =
  case b of
    . . .
    Lt x y ->
     let (t, store') = (findTypeAExp x store) in
     if t /= IntType
     then (FailureType, store')
     else let (t', store'') = (findTypeAExp y store') in
                if t' /= IntType
                then (FailureType, store'')
                else (BoolType, store'')
```

```
findTypeAExp :: AExp -> TypeStore -> (Type, TypeStore)
findTypeAExp a store =
  case a of
   ...
  Num x -> (IntType, store)
  ...
```

```
findTypeAExp :: AExp -> TypeStore -> (Type, TypeStore)
findTypeAExp a store =
 case a of
  Num x -> (IntType, store)
  . . .
Next: What if it's a variable?
data AExp =
  Var String
```

```
findTypeAExp :: AExp -> TypeStore -> (Type, TypeStore)
findTypeAExp a store =
  case a of
   ...
  Var x -> (Map.findWithDefault FailureType x store, store)
  ...
```

```
findTypeAExp :: AExp -> TypeStore -> (Type, TypeStore)
findTypeAExp a store =
 case a of
  Var x -> (Map.findWithDefault FailureType x store, store)
  . . .
Next: What if it's a Plus?
data AExp =
  ... Plus AExp AExp
```

```
findTypeAExp :: AExp -> TypeStore -> (Type, TypeStore)
findTypeAExp a store =
 case a of
    Plus x y ->
    let (t, store') = (findTypeAExp x store) in
      if t /= IntType
      then (FailureType, store')
      else let (t', store'') = (findTypeAExp y store') in
                     if t' /= IntType
                     then (FailureType, store'')
                     else (IntType, store'')
```

```
findTypeAExp :: AExp -> TypeStore -> (Type, TypeStore)
findTypeAExp a store =
 case a of
    Plus x y ->
    let (t, store') = (findTypeAExp x store) in
      if t /= IntType
      then (FailureType, store')
      else let (t', store'') = (findTypeAExp y store') in
                     if t' /= IntType
                      then (FailureType, store'')
                      else (IntType, store'')
  . . .
Next: What if it's a Lambda?
data AExp = ... Lambda String Type AExp
```

```
findTypeAExp :: AExp -> TypeStore -> (Type, TypeStore)
findTypeAExp a store =
  case a of
  ...
  Lambda s t a ->
  let (t', store') =
     findTypeAExp a (Map.insert s t store)
     in (LambdaType t t', store)
     -- Why not store' ?
  ...
```

```
findTypeAExp :: AExp -> TypeStore -> (Type, TypeStore)
findTypeAExp a store =
 case a of
    Lambda s t a \rightarrow
    let (t', store') =
       findTypeAExp a (Map.insert s t store)
        in (LambdaType t t', store)
                       -- Why not store'?
  . . .
Next: What if it's a application of a function?
data AExp =
  App AExp AExp
```

```
findTypeAExp :: AExp -> TypeStore -> (Type, TypeStore)
findTypeAExp a store =
 case a of
    App f x \rightarrow
    let (t,store') = (findTypeAExp f store) in
      case t. of
        LambdaType t1 t2 ->
           let (t',store'') = (findTypeAExp x store) in
           if t1 == t'
           then (t2, store)
           else (FailureType, store)
        _ (FailureType, store)
                  -- what about store, or store,,?
```

Your Turn! Write typecheckers for some new expressions.

Let's Do IfExp!

```
IF x < 4 THEN 8 + x ELSE x + 9 END
```

Let's Do Let!

```
LET x = 2 + 2 IN x + x END
```

Let's Do Recursion!!

```
REC f = int x:int { [f x+1] } IN [f 3] END
```

I've done the parsing for you, just add clauses to the findTypeAExp function.

