The Concurrent Computing Landscape

Andrews Chapter 1

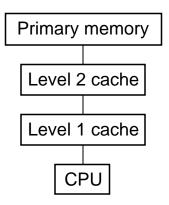


Figure 1.1 Processors, cache, and memory in a modern machine.

- Temporal locality
- Cache hit
- Cache miss
- Write through cache
- Write back cache
- Cache lines
- Spatial locality

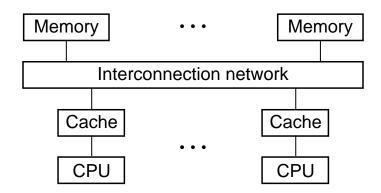


Figure 1.2 Structure of Shared-Memory Multiprocessors.

- Cache consistency problem Snooping
- Memory consistency problem Compilers and libraries
- False sharing Padding

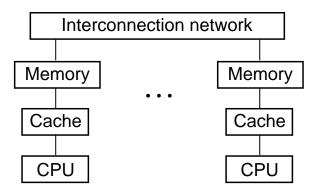


Figure 1.3 Structure of distributed-memory machines.

- Clusters
- Networks
- Message passing

Paradigms

- Iterative parallelism
 - Each process is an iterative program. Scientific computation.
- Recursive parallelism
 - Recursive procedure calls are independent.
- Producers and consumers
 - Each process is a filter. Pipelines.
- Clients and servers
 - World Wide Web. Graphics cards.
- Interacting peers
 - Decentralized decision making.

```
double a[n,n], b[n,n], c[n,n];

for [i = 0 to n-1] {
   for [j = 0 to n-1] {
     # compute inner product of a[i,*] and b[*,j]
     c[i,j] = 0.0;
   for [k = 0 to n-1]
     c[i,j] = c[i,j] + a[i,k]*b[k,j];
   }
}
```

Sequential Matrix Multiplication

Embarrassingly Parallel

- Read set: set of variables read by a process
- Write set: set of variables written to by a process
- Two operations are **independent** if the write set of each is disjoint from both the read and write sets of the other.

```
co [i = 0 to n-1] { # compute rows in parallel
  for [j = 0 to n-1] {
    c[i,j] = 0.0;
    for [k = 0 to n-1]
       c[i,j] = c[i,j] + a[i,k]*b[k,j];
  }
}
```

Parallel Matrix Multiplication by Rows

```
co [j = 0 to n-1] { # compute columns in parallel
  for [i = 0 to n-1] {
    c[i,j] = 0.0;
    for [k = 0 to n-1]
       c[i,j] = c[i,j] + a[i,k]*b[k,j];
  }
}
```

Parallel Matrix Multiplication by Columns

Parallel Matrix Multiplication by Rows and Columns

```
co [i = 0 to n-1] {  # rows in parallel then
  co [j = 0 to n-1] {  # columns in parallel
    c[i,j] = 0.0;
    for [k = 0 to n-1]
       c[i,j] = c[i,j] + a[i,k]*b[k,j];
  }
}
```

Parallel Matrix Multiplication Using Nested co Statements

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• Can we parallelize the loop over k?

```
process row[i = 0 to n-1] { # rows in parallel
  for [j = 0 to n-1] {
    c[i,j] = 0.0;
    for [k = 0 to n-1]
        c[i,j] = c[ij] + a[i,k]*b[k,j];
  }
}
```

Parallel Matrix Multiplication Using a Process Declaration

- Processes are named.
- Processes cannot be nested.
- Program does not wait for processes to end.

```
process worker[w = 1 to P] {  # strips in parallel
  int first = (w-1) * n/P;  # first row of strip
  int last = first + n/P - 1;  # last row of strip
  for [i = first to last] {
    for [j = 0 to n-1] {
       c[i,j] = 0.0;
       for [k = 0 to n-1]
        c[i,j] = c[i,j] + a[i,k]*b[k,j];
    }
}
```

Parallel Matrix Multiplication by Strips (Blocks)

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• Useful because we often have only 4 or 8 cores.

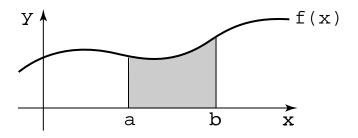


Figure 1.4 The quadrature problem.

- Find the area under a function.
- Iterative approach:

Break region into many equal intervals.

• Recursive approach:

Break region in half, recursion on each half.

```
double fleft = f(a), fright, area = 0.0;
double width = (b-a) / INTERVALS;
for [x = (a + width) to b by width] {
  fright = f(x);
  area = area + (fleft + fright) * width / 2;
  fleft = fright;
}
```

Iterative Quadrature Program

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• Cannot be parrallelized because of area.

```
double quad(double left,right,fleft,fright,lrarea) {
   double mid = (left + right) / 2;
   double fmid = f(mid);
   double larea = (fleft+fmid) * (mid-left) / 2;
   double rarea = (fmid+fright) * (right-mid) / 2;
   if (abs((larea+rarea) - lrarea) > EPSILON) {
        # recurse to integrate both halves
        larea = quad(left, mid, fleft, fmid, larea);
        rarea = quad(mid, right, fmid, fright, rarea);
   }
   return (larea + rarea);
}
```

Recursive Procedure for Quadrature Problem

- Only recurse when area changing. Copyright © 2000 by Addison Wesley Longman, Inc.
- Change the recursion to:

```
co larea = quad(left, mid, fleft, fmid, larea);
// rarea = quad(mid, right, fmid, fright, rarea);
oc
```

• co statements do not end until all branches end.

Independent procedure calls

- If a procedure does not reference global variables and has only value parameters, then every call of the procedure will be independent.
- Functional programming has these features.
- For example, quicksort.

```
double quad(double left,right,fleft,fright,lrarea) {
  double mid = (left + right) / 2;
  double fmid = f(mid);
  double larea = (fleft+fmid) * (mid-left) / 2;
  double rarea = (fmid+fright) * (right-mid) / 2;
  if (abs((larea+rarea) - lrarea) > EPSILON) {
    # recurse to integrate both halves in parallel
    co larea = quad(left, mid, fleft, fmid, larea);
    // rarea = quad(mid, right, fmid, fright, rarea);
    oc
  }
  return (larea + rarea);
}
```

Recursive Parallel Adaptive Quadrature

Source files
$$\rightarrow$$
 (sed) \rightarrow (tbl) \rightarrow $(groff)$ \rightarrow $(proff)$ \rightarrow $($

Figure 1.5 A pipeline of processes.

• Producers and consumers.

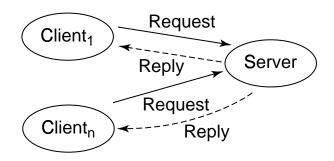


Figure 1.6 Clients and servers.

- File servers
- Web pages
- Graphics cards

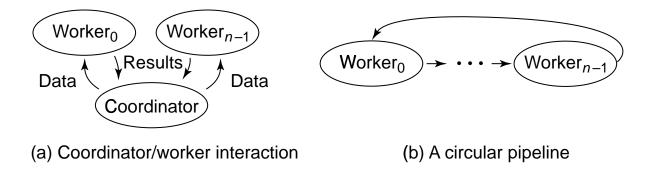


Figure 1.7 Matrix multiplications using message passing.

• Interacting peers

```
process worker[i = 0 to n-1] {
  double a[n]; # row i of matrix a
  double b[n,n]; # all of matrix b
  double c[n]; # row i of matrix c
  receive initial values for vector a and matrix b:
  for [j = 0 \text{ to } n-1] {
    c[j] = 0.0;
    for [k = 0 \text{ to } n-1]
      c[j] = c[j] + a[k] * b[k,j];
  send result vector c to the coordinator process;
process coordinator {
  double a[n,n]; # source matrix a
  double b[n,n]; # source matrix b
  double c[n,n]; # result matrix c
  initialize a and b:
  for [i = 0 to n-1] {
    send row i of a to worker[i];
    send all of b to worker[i];
  for [i = 0 \text{ to } n-1]
    receive row i of c from worker[i];
  print the results, which are now in matrix c;
```

```
process worker[i = 0 to n-1] {
  double a[n];  # row i of matrix a
  double b[n]; # one column of matrix b
  double c[n];  # row i of matrix c
double sum = 0.0; # storage for inner products
  int nextCol = i; # next column of results
  receive row i of matrix a and column i of matrix b;
  # compute c[i,i] = a[i,*] \times b[*,i]
  for [k = 0 \text{ to } n-1]
    sum = sum + a[k] * b[k];
  c[nextCol] = sum;
  # circulate columns and compute rest of c[i,*]
  for [j = 1 to n-1] {
    send my column of b to the next worker;
    receive a new column of b from the previous worker;
    sum = 0.0;
    for [k = 0 \text{ to } n-1]
      sum = sum + a[k] * b[k];
    if (nextCol == 0)
      nextCol = n-1;
    else
      nextCol = nextCol-1;
    c[nextColl = sum;
  send result vector c to coordinator process;
```