

Problem 1.  $(S, T), c: E \rightarrow \mathbb{R}, f: E \rightarrow \mathbb{R}, C: \text{maximum weight of base.}$

$$\left\{ \begin{array}{l} \sum_{e \in S} c(e)x_e = C \\ \sum_{e \in U} x_e \leq f(U) \quad \forall U \subseteq S \\ \sum_{e \in S} x_e = f(S) \\ x_e \geq 0 \end{array} \right\} \quad \text{let } = P'$$

<Solution> Observe that the polyhedron  $\{x \mid \sum_{e \in U} x_e \leq f(U), \forall U \subseteq S, \sum_{e \in S} x_e = f(S), x_e \geq 0\} = P$   
 $\dots P = P'$  // network graph.

Since  $\{\sum_{e \in S} c(e)x_e = C\}$  defines a supporting hyperplane of  $P$ ,  
 We have that the polyhedron defined,  $P'$  also is an integer polyhedron  
 since a face of a face is a face.

( $\otimes$  since  $C$  exists,  $P'$  is nonempty).

Problem 2.  $(G = (V_L \cup V_R, E))$ , Prove Hall's Theorem.

Given  $G$ ,  $G$  has a matching of size  $|V_L|$  if and only if for every  $S \subseteq V_L$ ,  
 We have  $|N(S)| \geq |S|$ , where  $N(S) = \{b \in V_R \mid \exists a \in S \text{ with } (a,b) \in E\}$

<Proof>  $\Rightarrow$  Trivial.

$\Leftarrow$  Let  $U := V_L \setminus C$  for an arbitrary vertex cover  $C \Rightarrow N(U) \subseteq C$

$\otimes$  Supp not.  $\Rightarrow \exists (x,y) \in E$  s.t.  $x \in U$  and  $y \notin C$ . therefore  $x \notin C$ .

$\Rightarrow G(x,y)$  is not covered by  $C \dots$  contradicts our choice of  $C$

• We thus have  $|C| = |C \cap V_L| + |C \cap V_R|$

$$= |V_L| - |U| + |C \cap V_R|$$

$$= |V_L| - |U| + |N(U)| \quad \text{since } N(U) \subseteq C \cap V_R$$

$$\geq |V_L| - |U| + |U| \quad \text{since we supposed } |N(U)| \geq |U|$$

$$\geq |V_L| - |U| + |U|$$

$$= |V_L|$$

• Maximum Cardinality of Matching in  $G := |M^*| = |C^*|$  by König's Theorem

$\Rightarrow |M^*| \geq |V_L| \dots$  Therefore  $G$  has a matching of size  $|V_L|$

Problem 3. (b)  $S := V_L$ ,  $\mathcal{I} := \{U \subseteq V_L \mid \exists \text{ matching where every vertex in } U \text{ is matched}\}$   
 Prove that  $(S, \mathcal{I})$  is a matroid.

proof). If  $I \in \mathcal{I}$ , then for any  $J \subseteq I$ , the matching that certifies  $I$  also satisfies that  $J \in \mathcal{I}$ . ... ① satisfied

• Suppose that  $I, J \in \mathcal{I}$  and  $|I| < |J|$

Let  $M_I, M_J \in \mathcal{E}$  be the matching in which every vertex in  $I$  (and  $J$ ) is matched. (respectively)

(WLOG) assume that  $|M_I| = |I|$  and  $|M_J| = |J|$

• Now Consider a directed bipartite graph on  $(V_L \cup V_R)$ ,

where we have an arc  $\langle \vec{x}, \vec{y} \rangle$  iff  $(x, y) \in M_J$   
 and  $\langle \vec{y}, \vec{x} \rangle$  iff  $(x, y) \in M_I$

$\Rightarrow$  Since every vertex in this graph  $\rightarrow (\text{indeg} \leq 1, \text{outdeg} \leq 1)$

The graph is a union of disjoint directed paths and directed cycles.

• Observe that the number of vertices in  $V_L$  whose indeg is 0 is  $|J \setminus I|$  and the number of vertices in  $V_L$  whose outdeg is 0 is  $|I \setminus J|$

• Since  $|J \setminus I| > |I \setminus J|$ , among the  $|J \setminus I|$  directed paths that starts from a vertex in  $|J \setminus I|$ , at least one of them ends at a vertex in  $V_R$

• Let  $P$  this path, and  $M_I \Delta P$  shows that  $\exists x \in J \setminus I, I + x \in \mathcal{I}$ .

Problem 5 (a) given  $(S, \mathcal{I})$ , The intersection of two flats is a flat.

proof). Let  $F_1 = \text{span}(F_1)$ ,  $F_2 = \text{span}(F_2)$  and  $x \in \text{span}(F_1 \cap F_2)$

We then have  $r(F_1 \cap F_2 + x) = r(F_1 \cap F_2)$

• and also.  $r(F_1 + x) = r(F_1)$   
 $r(F_2 + x) = r(F_2)$  ( $\because$  submodularity)

• This shows that  $x \in \text{span}(F_1)$  and  $x \in \text{span}(F_2) \therefore x \in \underline{F_1 \cap F_2}$

(b) Let  $F = \text{span}(F)$ ,  $t \in S \setminus F$  and  $F'$  be a flat with the maximum cardinality that contains  $(F + t) \Rightarrow$  then  $\nexists$  a flat  $F''$  s.t.  $F \subsetneq F'' \subsetneq F'$

proof) First, we claim that  $F' = \text{span}(F + t)$

$\rightarrow$  proof)  $(F + t) \subseteq F'$  and suppose that  $y \in \text{span}(F + t)$

$\Rightarrow$  then  $r(F + t + y) = r(F + t) \dots$  implying  $r(F + y) = r(F)$

This shows that  $\text{span}(F + t) \subseteq F'$  and hence  $F' = \text{span}(F + t)$

- Note that  $\text{span}(F+t)$  is a flat and the unique set that contains  $\text{span}(F+t)$  has cardinality  $\leq |\text{span}(F+t)|$ .
- We have  $r(F') = r(\text{span}(F+t))$   
 $= r(F) + 1$  since  $t \notin F = \text{span}(F)$ .

This implies that  $r(F'') \in \{r(F), r(F) + 1\}$

Case 1)  $r(F'') = r(F)$

Suppose toward contradiction that  $\exists x \in F'' \setminus F$

and then we have  $r(F'') \geq r(F+x) > r(F)$ . ✗

Case 2)  $r(F'') = r(F')$

Suppose toward contradiction that  $\exists$  some  $x \in F' \setminus F''$

and then we have  $r(F') \geq r(F''+x) > r(F'')$ . ✗