# A Combinatorial Presentation of the Operad of Plane Graphs

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# Diagrams for Monoidal Categories

- string diagrams graphical language for monoidal categories
- represented by graphs (Selinger, 2011)
  - · SMC: directed acyclic graphs
  - · traced: can contain cycles
  - · autonomous: wires can go the other way
- diagram equality -> graph isomorphism
- equational reasoning via rewrite rules
  - -> double pushout graph rewriting

# Diagrams for non-symmetric Monoidal Categories

### Non-symmetric case:

- printing circuits: crossings not possible
- quantum circuits: crossing not for free
- most general case: can add structure on top

### Representation:

- no crossing wires
- represented by plane graphs

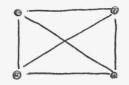
### Implementation:

This project is being implemented in Agola

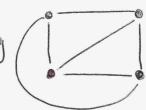
#### Definition:

- graph G = (V, E) consisting of vertices and edges
- embedding of G: drawing of G on a surface S
- G planar if there exists an embedding into the plane without crossing edges. The embedding is called plane.

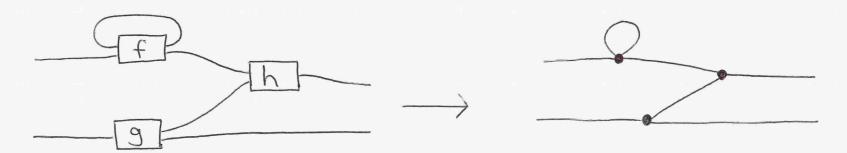
planor graph:



plane embedding of G:

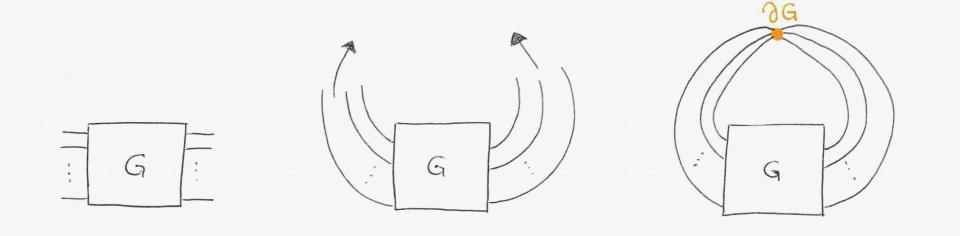


A PRO can be represented as open plane graph:



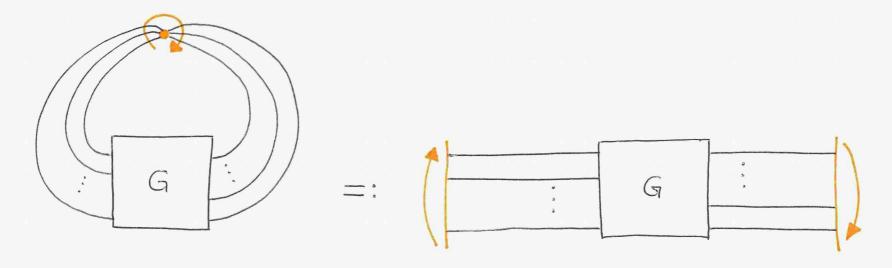
### Plane Graphs with a Boundary Vertex (1)

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Encoding the dangling input and output wires:

- introduce a new vertex  $\partial G$ , the boundary vertex
- making the boundary part of the graph



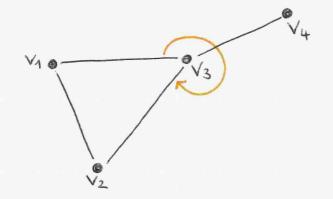
- boundary vertex part of the plane graph
- connecting parallel graphs
- nice way to represent plane graphs combinatorically

#### Definition:

- rotation of a vertex  $v \in V$ :

  cyclic ordered List of adjacent vertices
- rotation system: rotation for all vertices in the graph

Here: rotation in clockwise direction



V1: V2, V3

V2: V1, V3

V3: V1, V4, V2

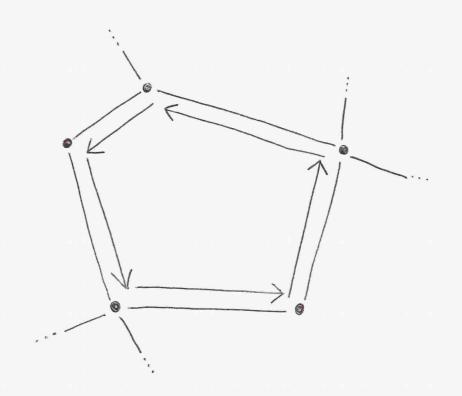
V4: V3

# Rotation Systems (2)

#### Lemma:

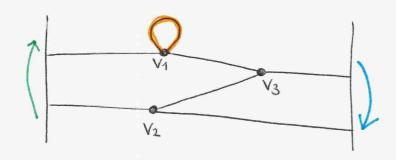
A rotation system uniquely defines a (cellular) embedding of a graph (Youngs, 1963)

### Proof:



## Rotation Systems (3)

### Example:



V1: in, V1, V1, V3

V2: in, v3, out

V3: V1, out, V2

in: V2, V1

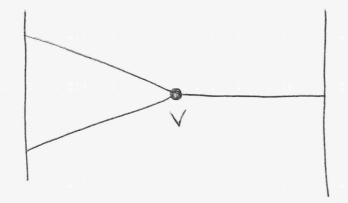
out: V3, V2

- cottegorically: inputs and outputs non-cyclic ordered lists
- combinatorically: boundary vertex as cyclic ordered list
- special case: multiple self loops (later, maybe)

boundary and inner vertices

# Building Graphs - Base Cases (1)

### Single vertex:

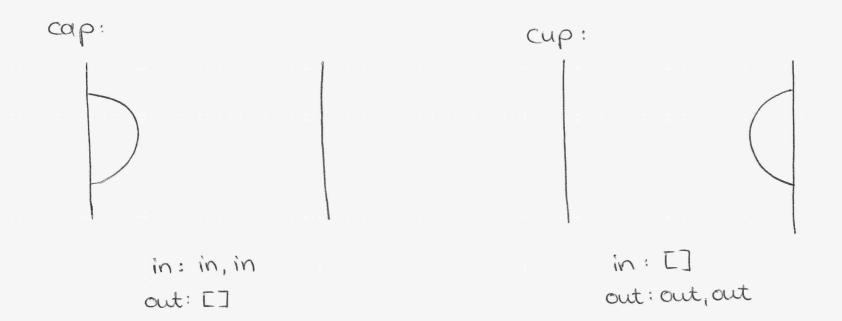


v: in, in, out

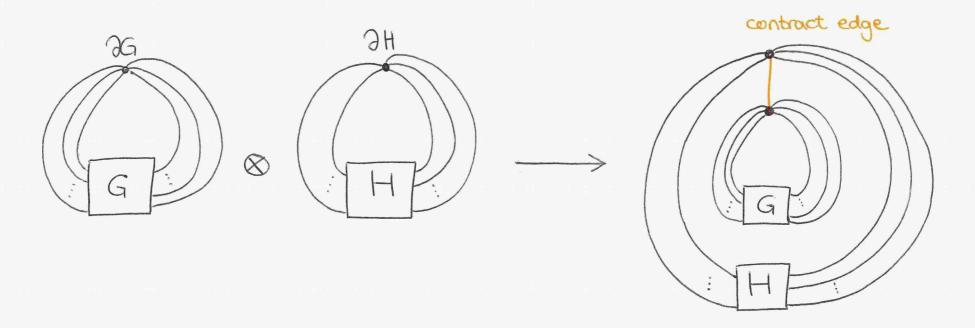
in: V, V

out: V

empty graph:			identit	identity:		
: : : : : : : : : : : : : : : : : : : :		: :::				
	in: []			in: out		
	out:[]			out: in		



cap and cup are self loops at the boundary vertex (so is the identity!)



- make names of vertices disjoint
- new rotation system: union of both rotation systems
- new boundary vertex: draw extra edge and contract it

### Example:



V<sub>2</sub>

V4: in, V4, V4, out

in: Va

out: V1

V2: in, out, out

in : V2

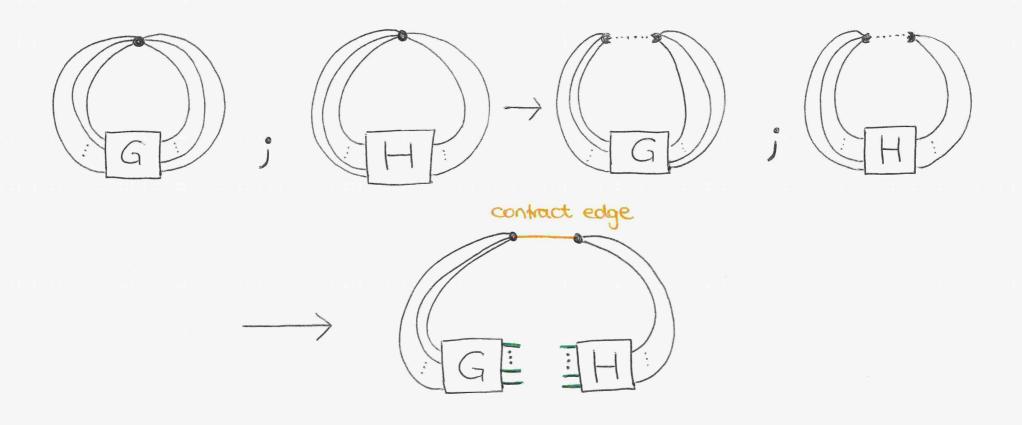
out: V2, V2

V1: in, V1, V1, out

V2: injout, out

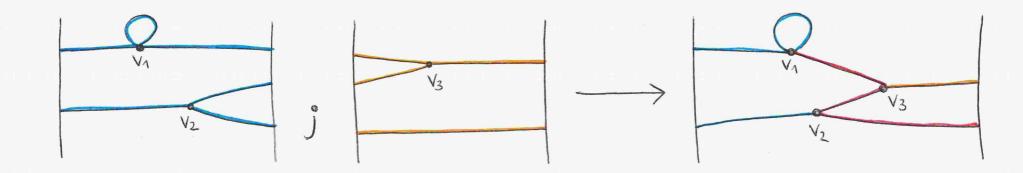
in: V2, V4

out: V1, V2, V2



- identify edges at the composition boundary
- update rotation systems on both sides
- new boundary vertex: inputs from the left outputs from the right

### Example:



V<sub>1</sub>: in, v<sub>2</sub>, out V<sub>2</sub>: in, out, out

in: V2, V1 out: V1, V2, V2 V3: in, in, out in: out, v3, V3

out: V3, in

 $V_4$ : in,  $V_4$ ,  $V_4$ 

 $V_2$ : in,  $V_3$ , out

V3: V2, V1, out

in: 1/21/4

out: V3, V2

Special cases for sequential composition:

- longer paths:



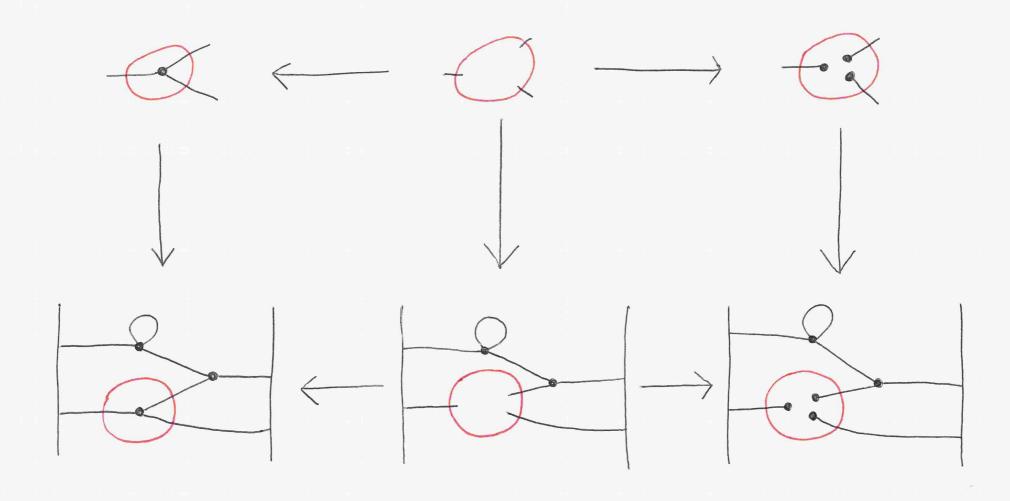
- cycles:



This representation of plane graphs with a boundary vertex defines a strict monoidal category, where

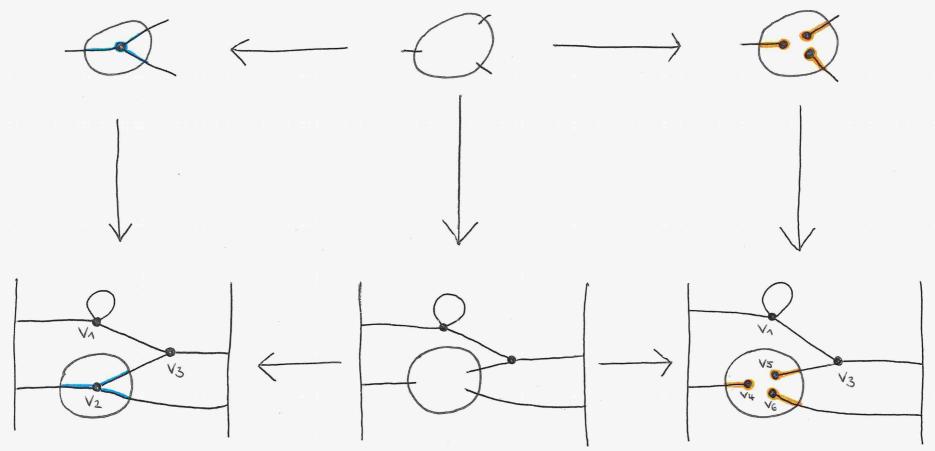
- the objects are lists of types of wires
- the morphisms are graphs
- parallel and sequential composition as defined above

Now: How does rewriting work?



### Graph Rewriting - Double Pushout Approach

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 $V_{1}: in, V_{1}, V_{1}, V_{3}$ 

V2: in, v3, out

V3: V1, Out, V2

in : V2, V1

out: V3, V2

 $V_{4}$ : in,  $V_{4}$ ,  $V_{4}$ ,  $V_{5}$ :  $V_{6}$ : out  $V_{3}$ :  $V_{4}$ , out,  $V_{5}$ : in :  $V_{4}$ ,  $V_{4}$ 

out: V3, V6

# Graph Rewriting

#### Lemma:

Graph rewriting (as defined above) preserves planarity.

### Proof:

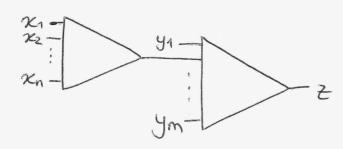
- LHS of rewrite rule is a connected graph
  - => can be contracted to a single vertex (edge contraction preserves planarity)
- substitution of a plane graph for a vertex preserves planounity

### Operads

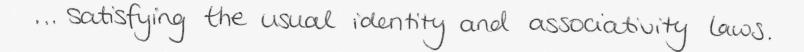
here: coloured operad (= multicategory)

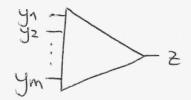
An operad consists of:

- a collection of objects
- a collection of morphisms which take multiple inputs
- Composition operation:

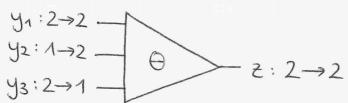


- identity x ------x





- objects: connectivity of graph variables
- morphisms: graphs

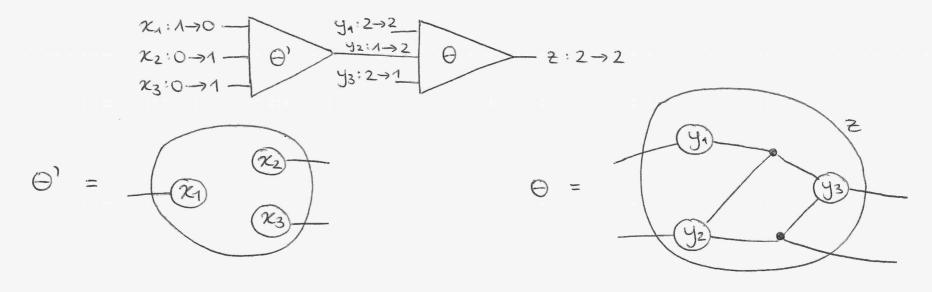


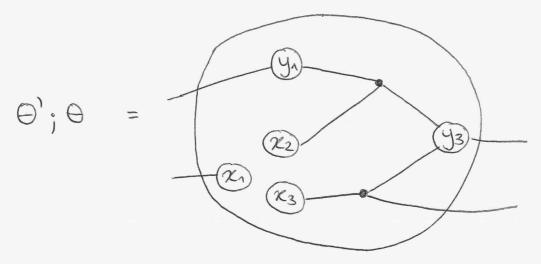
This is a symmetric operad

$$\Theta = \frac{(y_1)}{(y_2)}$$

Similar idea to the operad of wiring diagrams (Spivak, 2013)

- composition is substitution for a graph variable





Plane graphs with a boundary vertex form an operacl, where the composition operation is substitution.

- representing non-symmetric monoidal categories
- combinatorial presentation via rotation systems

#### Future work:

- more complex types of wires
- adding geometry information
- cooperads: substitution becomes patternmentahing

Thank you for your attention!

### References

- Selinger, P. (2011). A survey of Graphical languages for Monoidal Categories, pages 289-355. Springer Berlin Heidelberg.
- Spivak, D. (2013). The Operad of Wiring Diagrams: Formalizing a Graphical Language for Databases, Rocursion and Plug-and-Play Circuits.

  CORR, abs/1305,0297.
- Youngs, J.W.T. (1963). Minimal Imbeddings and the Genus of a Graph. Journal of Mouthematics and Mechanics, 12(2): 303-315.

### Extra: Self Loops

- need to distinguish



and



rotation systems [v,v,v,v,v]

[1,1,1,1,1,1]

- introduce pointers to other [YYYYYY] end of edge

[\(\lambda'\la

(validity check: well formed bracketing of pointers. LV, V, V, V, V] is not plane!)

- works for both inner vertices and the boundary